# WFF Guide

There are three types of cubes in WFF N Proof: Sentence variables, which are the main part of well formed formulas and are similar to (but not exactly like) nouns in English; Connectives, which hold the sentence variables together; and Blanks which are combined with connectives to form words.

Sentence Variables: p, q, r, and s Connectives: C, A, K, E, and N,

Blanks: R, i, and o

The sentence variables (p, q, r, and s) are all well formed formulas (WFFs) by themselves. Two WFFs joined by a connective (excluding N) **in front** of them is also a WFF.

N is a leech letter. Any WFF with N in front of it is still a WFF. There is no way to take an N out of a WFF or put an N into a WFF, in Basic WFF N' Proof.

Examples of WFFs: p, q, Kpq, Asr, Ars, KArsErp, ECCrpqp, KKKKKKKAAAAAAApppppsssssssrrr, Np, Ns, NNKKpqs, NNNNNNNNKKKrspq

Not WFFs: pq, qps, Eqps, ssEErr, KKKKssqqKs, KKAECKKKsq

The goal of WFF N Proof is to transform a supposition(sometimes known as a "Said-So"), your starting point or given, into the goal, which has to be a WFF. You do this by using rules, which help you change around WFFs. These rules are a connective followed by an i or an o. The rules used in basic WFF are, Ko, Ki, Ai, Co, Eo, Ei, Rp. o stands for out and i stands for in. The R cube can stand for any rule (CAKE letter and blank) and R doesn't have to be the same thing everyplace you use it. R can also stand for the R in Rp.

**Ki**: Ki says if you have two WFFs, then you can connect those two WFFs with a K. This rule can be stated s,  $r \rightarrow Ksr$  or Krs. Example proofs and solutions: $\rightarrow$ 

$$\begin{array}{c|cccc} & s, p \rightarrow Kps \\ \hline 1 & s & s & s, p/Ki \\ \hline 2 & p & s & \\ \hline 3 & Kps & Ki 1,2 & \end{array}$$

- 1.  $p, q \rightarrow Kpq$
- 2.  $s, p, q \rightarrow KKqsKKspKpq$
- 3. s, r, p,  $q \rightarrow KKKspqr$
- 4.  $r, p \rightarrow KKrpKKrpr$
- 5. s, r,  $q \rightarrow KKsrKKqsr$

**Rp:** Repeat. Rp lets you repeat a WFF. The rule can be stated as  $p \rightarrow p$  Example proofs and solutions:

- 1.  $s \rightarrow Kss$
- 2.  $p \rightarrow KKKpppKpp$
- 3. s,  $q \rightarrow KsKKssKqqq$
- 4.  $r, p \rightarrow KKrpKpp$
- 5. s, p,  $r \rightarrow KpKKpsKKrrKss$

**Ko**: "When in doubt, Kout." Ko says that if you have a K WFF (any WFF in which K is the left most connective) you can get either of the two smaller WFFs contained in the K WFF. This rule can be stated as  $Ksr \rightarrow r$  or s. Example proofs and solution:

	$Krs \rightarrow Ksr$		
1_	Krs	S	
2 3	S	Ko 1	Krs/Ko, Ki
3	r	Ko 1	
4	Ksr	Ki 1,2	

_	KKspq → KKKsqKqp	KKspq	
1	KKspq	S	
$2^{-}$	q	Ko 1	
3	Ksp	Ko 1	KKspq/Ki, Ko
4	S	Ko 3	
5	p	Ko 3	
5	Ksq	Ki 2,3	
6	Kqp	Ki 2,4	
7	KKsqKqp	Ki 5,6	
8	KKKsqKqpKKspq	Ki 1,7	

	$Ksp, KAprq \rightarrow$	KpApr	
1	Ksp	S	
2_	<u>K</u> Aprq	S	
3	p	Ko 1	Ksp, KAprq/Ko, Ki
4	Apr	Ko 2	
5	KpApr	Ki 3,4	

- 1.  $KsApr \rightarrow KAprs$
- 2. KCrpq, KEqps  $\rightarrow$  Ksq
- 3.  $KECrpsp \rightarrow KKECrpspp$
- 4. KKpsr  $\rightarrow$  KKrps
- 5. KpKrs, Kqs  $\rightarrow$  KKpqKsr

Ai: "Buy one, get one free." Ai says if you have a WFF, you can get an A WFF with the WFF you had and a new WFF. Ai can be stated  $r \rightarrow Arp$  or Apr.

Example proofs and solutions:

1 2 3	KpNs → Arp KpNs p Arp	s Ko 1 Ai 2	KpNs/Ko, Ai
1 2 3 4	p → AKppr p p Kpp AKppr	s Rp 1 Ki 1,2 Ai 3	p/Rp, Ki, Ai
4	KKsrNCsr → A	<u>EprAsr</u>	
1	_KKsrNCsr	S	
2	Ksr	Ko 1	
3	S	Ko 2	KKsrNcsr/Ko, Ai
4	Asr	Ai 3	
5	AEprAsr	Ai 4	

- 1.  $s \rightarrow Ars$
- 2.  $p \rightarrow ArKKppp$
- 3. Krs  $\rightarrow$  Arp
- 4. KKprs  $\rightarrow$  AKrsq
- 5.  $Kqp \rightarrow KKppAKpps$

**Co**: "Cancel Out" Co says that if you have a C WFF and the first little WFF of the C WFF, then you can get the second little WFF of the C WFF. This can be stated as Cqs,  $q \rightarrow s$  Example proofs and solutions:

KCrpr → p 1 KCrpr 2 Crp 3 r 4 p	s Ko 1 Ko 1 Co 3,4	KCrpr/Ko, Co
KCCrpsCrp →  KCCrpsCrp  CCrps  Crp  Crp  Srp  KCCrpsCrp	s Ko 1 Ko 1 Co 3,4	KCCrpsCrp/Ko, Co
KCKqpsp, Krq KCKqpsp Krq CKqps CKqps p Kqps Krq Kqps	→ s s Ko 1 Ko 1 Ko 2 Ki 4,5 Co 3,6	KCKqpsp, Krq/Ko, Ki, Co
CArsp, Kqs → 1 CArsp s 2 Kqs 3 s 4 Ars 5 p 6 q 7 Kqp	s Ko 2 Ai 3 Co 1,4 Ko 2 Ki 5,6	CArsp, Kqs/Ko, Ai, Co, Ki

- 1. r, CAAssrq  $\rightarrow$  AqNp
- 2. Krp, CAsKprEps  $\rightarrow$  KEpsp
- 3. CArps, Cqr, Krq  $\rightarrow$  Ksr
- 4. Crp, Csr, Ksq  $\rightarrow$  KApEqrq
- 5. CpCrq, Csp, Ksr  $\rightarrow$  KKsrq

**Ei**: E in says Cqs, Csq → Eqs or Esq. Some other examples of Ei are CrErq, CErqr → CrErq or CErqr and Cpq, Cqp → Cqp or Cpq. Both C WFFs are needed to get the E WFF. Example Proofs and Solutions:

Cqr, KCrqs 1 Cqr 2 KCrqs 3 Crq 4 Erq	→ Erq s s Ko 2 Ei 1,3	Cqr, KCrqs/Ko, Ei
KCsqq, Cqs KCsqq KCsqq Cqs Cqs Cqs Csq Esq Ksq Ksq Ksq Ksq	→ KEsqKsq s s Ko 1 Ei 2,3 Ko 1 Co 2,5 Ki 5,6 Ki 4,7	KCsqq, Cqs/Ko, Ki, Co, Ei
KCEsqrCsq  KCEsqrCsq  CpCqs Cqs Cqs Csq Esq CEsqr ACEsqr ACEsqr ANrr	Ko Ei Ko Co	KCEsqrCsq, p, CpCqs/ co 2 and 3 co 1 d and 5 co 1 do 6 and 7 i 8

- 1.  $KCrrCrr \rightarrow Err$
- 2. Cps, KpCsp →Esp
- 3. Cps,  $KpCsp \rightarrow KsEps$
- 4. Csr, Crs →Esr
- 5. KEsrCqp, Cpq → KEqpEsr

**Eo**: Eo says if you have an E WFF, then you can get a C WFF with the two little WFFs of the E WFF reversed. This rule can be stated as  $\text{Erp} \rightarrow \text{Cpr}$  or Crp.

Example proofs and solutions:

Eps, $s \rightarrow p$ 1 Eps 2 s 3 Csp 4 p	s s Eo 1 Co 2, 3	Eps, s/Eo, Co
Erp → KCrpCpr  1 Erp  2 Crp  3 Cpr  4 KCrpCpr	s Eo 1 Eo 1 Ki 2,3	Erp/Eo, Ki
KEspp → ANKr  KEspp  Esp  Cps  p  ANKr  A	s Ko 1 Eo 2 Ko 1 Co 3,4 Ai 5	KEspp/Ko, Eo, Co, Ai
KEsqApr → AK  KEsqApr  Esq  Csq  Cqs  Eqs  Apr  KEqsApr  KEqsApr	EqsAprK s Ko 1 Eo 2 Eo 2 Ei 3, 4 Ko 1 Ki 5,6 Ai 7	rp KEsqApr/Ko, Eo, Ei, Ko, Ki, Ai

- 1.  $Erp \rightarrow KCrpCpr$
- 2. Esq  $\rightarrow$  Eqs
- 3. EqKrs,  $q \rightarrow s$
- 4.  $Erp \rightarrow KCrpAEprAss$
- 5. KCprEpr  $\rightarrow$  Erp

#### **How To Present A Solution**

A solution consists of your proof, and the cubes you used (including the rules). In addition to a solution, you must also write a proof:

Ш	$Given \rightarrow Goal$		
1	Given	S	Given/RuleA, RuleB
2	~~~	RuleA 1	
3	~~	RuleB 1,2	

When writing a solution you should

- 1) Number your lines (this will make it easier for the checker to check)
- 2) Include the line numbers with which you used the rule (Ko uses one line, Ki uses two lines, Ai uses one line, Co uses to lines, Ei uses two lines, and Eo uses one line)
- 3) Put a line after your given in your proof
- 4) Put arrows under the cubes in essentials in your solution. Remember, each proof must have at least one rule, or it's wrong, and if you use a rule twice, you only need to list it once in your solution.

Setting the Goal: You have two minutes to set the goal. The maximum number of cubes in the goal is seven. A goal must be a WFF, or it is a Never.

## The Playing Mats:

The Equations playing mat, with Required, Permitted and Forbidden is used to play WFF N Proof. The permitted and forbidden section operate exactly as they do in Equations. The required section, however, is slightly different. When a cube is place in the required section in WFF, it must be an essential part of any acceptable solution. If you can take a cube that is in required out of a solution and the solution still works, that cube is non-essential, and therefore the solution is incorrect. The cube(s) you take out must be a WFF or a rule, not a part of a WFF or part of a rule. You are allowed to take out as many rules and WFFs out of a solution as you want, but you can not add any rules. To prove non-essentiality, you must write a proof showing that the parts you took out are non-essential. If you can prove a cube that is in essentials is non-essential, then the solution is wrong. Here is an example of this:

Someone calls Now on you and presents a solution like this

_	$Krp \rightarrow KKrpKrp$		
1	_Krp	S	
2	r	Ko 1	Krp/Ko, Ki, Rp
3	p	Ko 1	_
4	Krp	Ki 2,3	
5	Krp	Rp 1	
6	KKrpKrp	Ki 1,5	

You say that that the K in essentials is non-essential. Here is your proof:

$$\begin{array}{c|cccc} & Krp \rightarrow KKrpKrp \\ \hline 1 & Krp & s & Krp/Rp, Ki \\ \hline 2 & Krp & Rp 1 \\ 3 & KKrpKrp & Ki 1,2 \\ \end{array}$$

You have proven your opponent is wrong, and they lose the shake.