OS Homework 3 Report Team21

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contribution: 一起

1-1. New -> Ready

Kernel::ExecAll()

呼叫Kernel::Exec()去個別生成thread來執行kernel中所有的execfile, 執行完後呼叫currentThread的Finish(), 也就是結束掉main thread

Kernel::Exec(char*)

新增一個thread給要執行的user program, 再make一個新的pageTable給這個thread, 接著呼叫 Thread::Fork來做allocate stack, 把thread放進ready queue等操作, 最後回傳這個thread的thread number

Thread::Fork(VoidFunctionPtr, void*)

先做StackAllocate(), allocate空間給要執行的thread, 接下來disable interrupt, 把thread放進ready queue, 回復interrupt原本的狀態

Thread::StackAllocate(VoidFunctionPtr, void*)

```
306 void
307 void
308 {
309 stack = (int *) AllocBoundedArray(StackSize * sizeof(int));
```

create一塊pagesize*2+size的空間,保護前後兩個page,來偵測有沒有overflow,回傳第二個 page,也就是第一個可以用的page

把stackTop移動到stack+StackSize-4的位置,然後--stackTop後指向threadRoot,當我們switch到這個process的時候,就會從Threadroot開始run

```
#else
machineState[PCState] = (void*)ThreadRoot;
machineState[StartupPCState] = (void*)ThreadBegin;
machineState[InitialPCState] = (void*)func;
machineState[InitialArgState] = (void*)arg;
machineState[WhenDonePCState] = (void*)ThreadFinish;
#endif
#endif
```

將ThreadRoot、ThreadBegin、func procedure、arg等pointer,存進machineState這個register,之後要呼叫這些procedure或使用argument就從這些register找

Scheduler::ReadyToRun(Thread*)

thread的status設為READY, 把thread丟進ready queue, 表示他準備好執行

1-2. Running -> Ready

Machine::Run()

```
void

Machine::Run()

Instruction *instr = new Instruction; // storage for decoded instruction

if (debug→IsEnabled('m')) {

cout < "Starting program in thread: " < kernel→currentThread→getName();

cout < ", at time: " < kernel→stats→totalTicks < "\n";

kernel→interrupt→setStatus(UserMode);

for (;;) {

DEBUG(dbgTraCode, "In Machine::Run(), into OneInstruction " < "= Tick " < kernel→stats→totalTicks < " = ");

OneInstruction(instr);

DEBUG(dbgTraCode, "In Machine::Run(), return from OneInstruction " < "= Tick " < kernel→stats→totalTicks < " = ");

kernel→interrupt→OneTick();

DEBUG(dbgTraCode, "In Machine::Run(), into OneTick " < "= Tick " < kernel→stats→totalTicks < " = ");

kernel→interrupt→OneTick();

DEBUG(dbgTraCode, "In Machine::Run(), return from OneTick " < "= Tick " < kernel→stats→totalTicks < " = ");

if (singleStep 66 (runUntilTime ≤ kernel→stats→totalTicks))

Debugger();

}

75

}
```

當user program開始執行的時候,準備好一個Instruction的data structure, 並把status設成 user mode, 在for迴圈裡 OneInstruction一個一個讀並執行指令, 每個指令執行完後呼叫 OneTick()增加simulated time, 同時檢查有沒有interrupt需要處理

Interrupt::OneTick()

看現在是user mode還是kernel mode增加simulated time, 先disable interrupt, 用 CheckIfDue()去檢查並處理interrupt, 再enable interrupt回來, 最後檢查yieldOnReturn是否為TRUE, 若為TRUE就呼叫Yield(), 做完再把yieldOnReturn設回FALSE, yieldOnReturn會根據以下的步驟被設為TRUE。

```
alarm = new Alarm(randomSlice); // start up time slicing
```

在kernel::Initialize會create一個alarm object

alarm會create一個time, 並把doRandom和自己當作參數

Time把傳進來的alarm object設為callBack object callPeriodically, doRandom決定是否random arrange interrupt, 把disable設為FALSE, 最後呼叫SetInterrupt()

接下來根據random決定delay interval,然後schedule這個timer device的interrupt,傳入timer自己和 delay的時間

create一個toOccur interrupt,會傳入我們前面傳的timer object(toCall)和interrupt發生的時間(when),然後insert到PendingInterrupt中

當OneTick或Idle呼叫CheckIfDue, 而且時間到了interrupt該執行的時間時, 會在359行呼叫interrupt的CallBack(), 也就是timer的callback

呼叫callPeriodically(Alarm)的CallBack, 同時再set一個新的interrupt

如果Machine的status不是IdleMode, 呼叫interrupt的YieldOnReturn

把yieldOnReturn設為TRUE

Thread::Yield()

當currentThread要把CPU的使用權交給其他thread, 並且currentThread還沒finish, 之後還要執行, 所以用FindNextToRun找有沒有其他thread要run, 有的話呼叫 ReadyToRun()把currentThread放回 ready queue, 然後呼叫scheduler run去做context switch

Scheduler::FindNextToRun()

回傳下一個要做的thread, 若無則回傳NULL

Scheduler::ReadyToRun(Thread*)

thread的status設為READY, 把thread丟進ready queue, 表示他準備好執行

Scheduler::Run(Thread*, bool)

如果old Thread是個user program, save他目前的state, 然後check 他有沒有stack overflow, 將 currentThread設為nextThread並把nextThread的status設為RUNNING

133 SWITCH(oldThread, nextThread);

呼叫SWITCH停止oldThread開始nextThread. 完成switch

```
// we're back, running oldThread

// interrupts are off when we return from switch!

ASSERT(kernel→interrupt→getLevel() = IntOff);

DEBUG(dbgThread, "Now in thread: " << oldThread→getName());

CheckToBeDestroyed(); // check if thread we were running

// before this one has finished

// and needs to be cleaned up

if (oldThread→space ≠ NULL) { // if there is an address space

oldThread→RestoreUserState(); // to restore, do it.

oldThread→space→RestoreState();

}

oldThread→space→RestoreState();

}
```

當又回到了oldThread之後,CheckToBeDestroyed判斷是否刪除前一個執行的Thread,然後確認有沒有state要restore,有的話在把之前oldThread存下來的state還原回去

1-3. Running -> Waiting

SynchConsoleOutput::PutChar(char)

- 1. lock -> Acquire(): 避免和其他thread同時output, 透過lock::Acquire和semaphore::P去檢查和搶資源
- 2. consoleOutput -> PutChar(): 把char寫進display
- 3. waitFor->P(): 等待後面呼叫了CallBack()中的waitFor->V()表示output結束
- 4. lock->Release(): thread釋出lock

Semaphore::P()

先disable interrupt, 因為85到89行確認value跟減value的動作不能被中斷, 若value=0, 代表有其他thread在用, 所以這個thread被block住了, 所以把它放進waiting queue, 然後讓他 sleep, 若value>0, 表示resource空下來了, 所以用value--去搶這個resource, 然後再回復 interrupt的狀態。

List<T>::Append(T)

```
template <class T>
void
template <class T>
void

List<T>::Append(T item)

{
    ListElement<T> *element = new ListElement<T>(item);

ASSERT(!IsInList(item));

if (IsEmpty()) {    // list is empty
    first = element;

last = element;

last = element;

last = element;

last = element;

ast = element;

and multist = element;

ASSERT(IsInList(item));

ASSERT(IsInList(item));

ASSERT(IsInList(item));

and multist = element;

and mult
```

List是一個link list的結構, append會把新的item接在last的尾端, numInList++Thread::Sleep(bool)

把thread status設為BLOCKED,如果ready queue是空的,呼叫Interrupt::Idle去檢查有沒有 interrupt要處理,如果ready queue有thread在等待,把CPU從A thread交給ready queue的第一個 thread, A thread可能是因為已經做完或在等synchronization variable,若是前者, A thread之後會被刪掉;若是後者之後會把A thread重新放回ready queue,未來會叫醒A thread

Scheduler::FindNextToRun()

```
Thread *
Thread
```

回傳下一個要做的thread, 若無則回傳NULL Scheduler::Run(Thread*, bool)

如果old Thread是個user program, save他目前的state, 然後check 他有沒有stack overflow, 將 currentThread設為nextThread並把nextThread的status設為RUNNING

133 SWITCH(oldThread, nextThread);

呼叫SWITCH停止oldThread開始nextThread. 完成switch

```
// we're back, running oldThread

// interrupts are off when we return from switch!

ASSERT(kernel→interrupt→getLevel() = IntOff);

DEBUG(dbgThread, "Now in thread: " << oldThread→getName());

CheckToBeDestroyed(); // check if thread we were running

// before this one has finished

// and needs to be cleaned up

if (oldThread→space ≠ NULL) { // if there is an address space oldThread→RestoreUserState(); // to restore, do it.

oldThread→space→RestoreState(); }

oldThread→space→RestoreState(); }

}
```

當又回到了oldThread之後,CheckToBeDestroyed判斷是否刪除前一個執行的Thread,然後確認有沒有state要restore,有的話在把之前oldThread存下來的state還原回去

1-4. Waiting -> Ready

Semaphore::V()

先disable interrupt,因為85到89行確認value跟加value的動作不能被中斷,檢查waiting queue有沒有waiter,有的話把waiter放進ready queue,增加value來釋出資源,然後回復 interrupt狀態

Scheduler::ReadyToRun(Thread*)

thread的status設為READY, 把thread丟進ready queue, 表示他準備好執行

1-5. Running -> Terminated

ExceptionHandler(ExceptionType) case SC Exit

```
case SC_Exit:

DEBUG(dbgAddr, "Program exit\n");

val=kernel→machine→ReadRegister(4);

cout << "return value:" << val << endl;

kernel→currentThread→Finish();

break;
```

CPU收到SC_exit的exception type, 故結束currentThread

Thread::Finish()

當一個thread做完就會呼叫,先Disable interrupt,實作方式是呼叫Sleep裡面傳TRUE,讓currentThread Sleep,因為現在還沒完成switch,所以還不能刪除這個thread,Sleep的參數TRUE代表這個Thread已經結束了,之後才會在Scheduler::CheckToBeDestroyed將這個Thread刪除

Thread::Sleep(bool)

```
void
Thread::Sleep (bool finishing)

{

Thread *nextThread;

ASSERT(this = kernel→currentThread);

ASSERT(kernel→interrupt→getLevel() = IntOff);

DEBUG(dbgThread, "Sleeping thread: " < name);

DEBUG(dbgTraCode, "In Thread::Sleep, Sleeping thread: " < name < ", " < kernel→stats→totalTicks);

status = BLOCKED;

//cout < "debug Thread::Sleep " < name < "wait for Idle\n";

while ((nextThread = kernel→scheduler→FindNextToRun()) = NULL) {
    kernel→interrupt→Idle(); // no one to run, wait for an interrupt
}

// returns when it's time for us to run
kernel→scheduler→Run(nextThread, finishing);
}
</pre>
```

把thread status設為BLOCKED,如果ready queue是空的,呼叫Interrupt::Idle去檢查有沒有 interrupt要處理,如果ready queue有thread在等待,把CPU從A thread交給ready queue的第一個 thread, A thread可能是因為已經做完或在等synchronization variable,若是前者, A thread之後會被刪掉;若是後者之後會把A thread重新放回ready queue,未來會叫醒A thread

Scheduler::FindNextToRun()

回傳下一個要做的thread. 若無則回傳NULL

Scheduler::Run(Thread*, bool)

如果old Thread是個user program, save他目前的state, 然後check 他有沒有stack overflow, 將 currentThread設為nextThread並把nextThread的status設為RUNNING

133 SWITCH(oldThread, nextThread);

呼叫SWITCH停止oldThread開始nextThread. 完成switch

```
// we're back, running oldThread

// interrupts are off when we return from switch!

ASSERT(kernel→interrupt→getLevel() = IntOff);

DEBUG(dbgThread, "Now in thread: " << oldThread→getName());

CheckToBeDestroyed(); // check if thread we were running

// before this one has finished

// and needs to be cleaned up

if (oldThread→space ≠ NULL) { // if there is an address space oldThread→RestoreUserState(); // to restore, do it.

oldThread→space→RestoreState(); }

oldThread→space→RestoreState(); }

150 }
```

當又回到了oldThread之後,CheckToBeDestroyed判斷是否刪除前一個執行的Thread,然後確認有沒有state要restore,有的話在把之前oldThread存下來的state還原回去

1-6. Ready -> Running

Scheduler::FindNextToRun()

回傳下一個要做的thread. 若無則回傳NULL

Scheduler::Run(Thread*, bool)

如果old Thread是個user program, save他目前的state, 然後check 他有沒有stack overflow, 將 currentThread設為nextThread並把nextThread的status設為RUNNING

呼叫SWITCH停止oldThread開始nextThread. 完成switch

```
// we're back, running oldThread

// interrupts are off when we return from switch!

ASSERT(kernel→interrupt→getLevel() = IntOff);

DEBUG(dbgThread, "Now in thread: " << oldThread→getName());

CheckToBeDestroyed(); // check if thread we were running

// before this one has finished
// and needs to be cleaned up

// if there is an address space

oldThread→space ≠ NULL) { // if there is an address space

oldThread→RestoreUserState(); // to restore, do it.

oldThread→space→RestoreState();

}

150 }
```

當又回到了oldThread之後,CheckToBeDestroyed判斷是否刪除前一個執行的Thread,然後確認有沒有state要restore,有的話在把之前oldThread存下來的state還原回去

SWITCH(Thread*, Thread*)

```
329 ** on entry, stack looks like this:

330 ** 8(esp) \rightarrow thread *t2

331 ** 4(esp) \rightarrow thread *t1

332 ** (esp) \rightarrow return address
```

```
SWITCH:

√ SWITCH:

         movl
                 %eax, eax save
                  4(%esp),%eax
         movl
         movl
                  %ebx,_EBX(%eax)
                  %ecx,_ECX(%eax)
         movl
         movl
                 %edx,_EDX(%eax)
                 %esi,_ESI(%eax)
         movl
         movl
                 %edi,_EDI(%eax)
                 %ebp,_EBP(%eax)
         movl
                 %esp,_ESP(%eax)
         movl
         movl
                  eax save,%ebx
                 %ebx,_EAX(%eax)
         movl
                  0(%esp),%ebx
         movl
                 %ebx, PC(%eax)
         movl
         movl
                 8(%esp),%eax
         movl
                  _EAX(%eax),%ebx
         movl
                 %ebx,_eax_save
         movl
                  _EBX(%eax),%ebx
         movl
                 _ECX(%eax),%ecx
         movl
                  EDX(%eax),%edx
         movl
                 ESI(%eax),%esi
         movl
                  _EDI(%eax),%edi
                 _EBP(%eax),%ebp
         movl
         movl
                  ESP(%eax),%esp
                  _PC(%eax),%eax
         movl
         movl
                 %eax,4(%esp)
         movl
                  _eax_save,%eax
 #endif // x86
```

先把%eax暫存到_eax_save, 然後把t1 pointer存到%eax, 346-352行把register存到t1的stack中, 353-354行把%eax的restore後也存進t1的stack, 355-356行把return address存進pc storage中。 358行把t2 pointer存到%eax。

360-361行把t2的%eax放到%ebx後暫存到_eax_save, 362-369行從t2的stack restore register的值。 370把t2 return address存到%esp+4

371 restore t2 %eax的值。

373 ret 會回到%esp, 下一個執行的指令是esp+4

depends on the previous process state, [New,Running,Waiting]→Ready

```
/* void ThreadRoot( void )

**

** expects the following registers to be initialized:

** eax points to startup function (interrupt enable)

** edx contains inital argument to thread function

** esi points to thread function

** edi point to Thread::Finish()

*/
```

New:會從Thread root開始做

```
for (;;) {

DEBUG(dbgTraCode, "In Machine::Run(), into OneInstruction " ≪ "= Tick " ≪ kernel→stats→totalTicks ≪ " =");

OneInstruction(instr);

DEBUG(dbgTraCode, "In Machine::Run(), return from OneInstruction " ≪ "= Tick " ≪ kernel→stats→totalTicks ≪ " =");

DEBUG(dbgTraCode, "In Machine::Run(), into OneTick " ≪ "= Tick " ≪ kernel→stats→totalTicks ≪ " =");

kernel→interrupt→OneTick();

DEBUG(dbgTraCode, "In Machine::Run(), return from OneTick " ≪ "= Tick " ≪ kernel→stats→totalTicks ≪ " =");

if (singleStep 56 (runUntilTime ≤ kernel→stats→totalTicks))

Debugger();

}
```

Running:代表t2原本是在跑Machine::Run的OneTick, 因為yield被放進了ready queue, context switch回來後會從之前的program counter的instruction繼續run

Waiting:t2原本是在Semaphore::P等待資源而被放進了waiting queue,接下來又因為資源可以用了而被放進ready queue,所以context switch回來後,會從Semaphore::P的while繼續run

Implementation

debug.h

```
33 // MP3
34 const char z = 'z';
35
```

新增一個debug flag 'z'

stats.h

```
// MP3
// const int ConsoleTime = 100; // time to read or write one character
const int ConsoleTime = 1;
// MP3
// const int ConsoleTime = 1;
```

kernel.cc

```
#endif // FILESYS_STUB

// postOfficeIn = new PostOfficeInput(10);

// postOfficeOut = new PostOfficeOutput(reliability);

// delete postOfficeIn;

// delete postOfficeOut;

// MP3

else if (strcmp(argv[i], "-ep") = 0){

execfile[+execfileNum] = argv[++i];

priorities[execfileNum] = atoi(argv[++i]);

}
```

把要執行的檔案放進execfile中, 並給它priority

create好thread之後, 設定他的initital priority

Kernel.h

```
// MP3
int Exec(char* name, int priority);

// MP3
You, 1 s
int priorities[10];
```

Scheduler.h

```
// MP3
int Aging();
// Thread* CheckPreemptive();

bool CheckPreemptive();

// SelfTest for scheduler is implemented in class Thread

private:

List<Thread *> *readyList; // queue of threads that are ready to run,

// but not running
// MP3

SortedList<Thread *> *L1;
SortedList<Thread *> *L2;
List<Thread *> *L3;

List<Thread *> *L3;
```

Scheduler.cc

```
Scheduler::Scheduler()
69
70
         readyList = new List<Thread *>;
71
72
         // MP3
         L1 = new SortedList<Thread *>(L1Compare);
73
         L2 = new SortedList<Thread *>(L2Compare);
74
75
         L3 = new List<Thread *>;
76
         toBeDestroyed = NULL;
77
78
```

在scheduler新增L1, L2, L3三個queue, L1和L2是sortedList, 因為要根據remain time和priority來排序,在create sortedlist時,要把compare的function傳進去(L1Compare, L2Compare)。

L1Compare比較的是approxRemainTime, approxRemainTime較小的優先

L2Compare比較的是priority, priority較大的優先

```
Scheduler::ReadyToRun (Thread *thread)
          ASSERT(kernel->interrupt->getLevel() == IntOff);
          DEBUG(dbgThread, "Putting thread on ready list: " << thread->getName());
          //cout << "Putting thread on ready list: " << thread->getName() << endl;</pre>
          thread->setStatus(READY);
110
          // readyList->Append(thread);
          // MP3
          int threadPriority = thread->getPriority();
112
113
          thread->setStartWaitingTime(kernel->stats->totalTicks);
115
          if (threadPriority > 149) {
116
              DEBUG(z, "Thread priority is out of range");
118
119
120
          else if (threadPriority >= 100) {
121
122
              L1->Insert(thread);
              thread->InsertedIntoQueue(1);
124
          else if (threadPriority >= 50) {
125
              // L2 (50 - 99)
126
127
              L2->Insert(thread);
              thread->InsertedIntoQueue(2);
128
129
          else if (threadPriority >= 0) {
130
132
              L3->Append(thread);
              thread->InsertedIntoQueue(3);
135
              // out of range
136
137
              DEBUG(z, "Thread priority is out of range");
138
```

當要將thread放進ready queue時,我們先將start waiting time設為現在,然後根據thread的priority用 Insert(thread)把thread放進queue, InsertedIntoQueue()則設定thread的queuelevel並print出debug訊 息

```
Thread *
150
151
      Scheduler::FindNextToRun ()
152
153
          ASSERT(kernel->interrupt->getLevel() == IntOff);
154
          // if (readyList->IsEmpty()) {
155
          // return NULL;
156
157
          // return readyList->RemoveFront();
158
159
          // MP3
          if (!L1->IsEmpty()) { ···
          else if (!L2->IsEmpty()) { ···
          else if (!L3->IsEmpty()) { ···
          else {
              return NULL;
200
```

FindNextToRun依序從L1, L2, L3尋找ready thread

```
if (!L1->IsEmpty()) {
    Thread *currThread = kernel->currentThread;
    Thread *nextThread = L1->RemoveFront();

    nextThread->setCpuStartTime(kernel->stats->totalTicks);
    nextThread->setTotalWaitingTime(0);

    nextThread->RemovedFromQueue();
    currThread->ContextSwitch(nextThread->getID());

    return nextThread;
}
```

假設L1 queue中有next thread, 則把它從L1 queue中remove, 然後將cpuStartTime設為現在, totalWaitingTime重置為0, RemovedFromQueue()和ContextSwitch()會print出debug訊息, 最後回傳 next thread

```
if (!L1->IsEmpty()) {
   iterator = new ListIterator<Thread *>(L1);
   for (; !iterator->IsDone(); iterator->Next()) {
      thread = iterator->Item();
      // Increase total waiting time
      bool doAging = thread->IncreaseTotalWaitingTime();
      thread->setStartWaitingTime(kernel->stats->totalTicks);
      // Check whether over 1500
      if (doAging) {
            thread->ChangePriority();
      }
}
```

```
if (!L2->IsEmpty()) {
   iterator = new ListIterator<Thread *>(L2);
   for (; !iterator->IsDone(); iterator->Next()) {
        thread = iterator->Item();

        // Increase total waiting time
        bool doAging = thread->IncreaseTotalWaitingTime();
        thread->setStartWaitingTime(kernel->stats->totalTicks);

        // Check whether over 1500
        if (doAging) {
              thread->ChangePriority();
        }

        if (thread->getPriority() > 99) {
              L2->Remove(thread);
             L1->Insert(thread);
              thread->RemovedFromQueue();
              thread->InsertedIntoQueue(1);
        }
    }
}
```

```
if (!L3->IsEmpty()) {
    iterator = new ListIterator<Thread *>(L3);
    for (; !iterator->IsDone(); iterator->Next()) {
        thread = iterator->Item();

        // Increase total waiting time
        bool doAging = thread->IncreaseTotalWaitingTime();
        thread->setStartWaitingTime(kernel->stats->totalTicks);

        // Check whether over 1500
        if (doAging) {
             thread->ChangePriority();
        }

        if (thread->getPriority() > 49) {
             L3->Remove(thread);
            L2->Insert(thread);
            thread->RemovedFromQueue();
            thread->InsertedIntoQueue(2);
        }

}
```

Aging會對所有queue中的thread去做檢查看是否要等待超過1500 ticks,有的話就增加10 priority, L2和L3中的thread需要再確認queue level是否可以升級了

```
// MP3
     bool
     Scheduler::CheckPreemptive()
          Thread *currThread = kernel→currentThread;
          int currThreadLevel = currThread→getQueueLevel();
          if (currThreadLevel = 1) {
              if (!L1 \rightarrow IsEmpty()) {
                  if (L1Compare(currThread, L1→Front()) = 1) {
                       return TRUE;
          else if (currThreadLevel = 2) {
              if (!L1→IsEmpty()) {
                  return TRUE;
          else {
              if (!L1 \rightarrow IsEmpty()){
                  return TRUE;
              if (!L2 \rightarrow IsEmpty()) {
                   return TRUE;
              if (!L3→IsEmpty()) {
378
                   return TRUE;
          return FALSE;
```

- 1. 如果現在的thread是L1, 只有L1的thread可以preemptive, 所以拿他的remain time去和現在L1的第一個thread比. 決定是否可以preemptive
- 2. 如果現在的thread是L2, 只有L1的thread可以preemptive
- 3. 如果現在的thread是L3, 所有thread都可以preempt

alarm.cc

在time slice到時, 會呼叫CallBack(), 在裡面計算accuTicks、cpuBurstTime、approxRemainTime, 接著呼叫aging。最後確認thread是否可被其他thread搶, 若可以則會呼叫YieldOnReturn()會把yieldOnReturn設為true, 然後在OneTick()就會呼叫Yield()

thread.h

```
156
           // MP3
           int priority;
157
           int queueLevel;
158
159
           int startWaitingTime;
           int totalWaitingTime;
161
162
           double cpuStartTime;
163
           double cpuBurstTime;
164
           double approxBurstTime;
165
           double approxRemainTime;
167
           double accuTicks;
168
```

在class Thread中新增幾項data, 如下:

int priority:該thread的priority

int queueLevel:該thread在哪個ready queue

int startWaitingTime: 進入Waiting state的時間

int totalWaitingTime: 在waiting等了多少時間

double cpuStartTime: 進入Ready state的時間

double cpuBurstTime: 進入RuningState的時間

double approxBurstTime: 估計會進入RuningState的時間

double approxRemainTime: 估計所需要的剩餘時間

double accuTicks: 最後一次的連續cpuBurstTime

```
int getPriority() { return priority; }
int getQueueLevel() { return queueLevel; }
int getStartWaitingTime() { return startWaitingTime; }
int getTotalWaitingTime() { return totalWaitingTime; }

double getCpuStartTime() { return cpuStartTime; }
double getCpuBurstTime() { return approxBurstTime; }
double getApproxBurstTime() { return approxBurstTime; }
double getApproxRemainTime() { return approxRemainTime; }

void setPriority(int initPriority) { priority = initPriority; }
void setStartWaitingTime(int newTime) { startWaitingTime = newTime; }

void setCpuStartTime(double newTime) { cpuStartTime = newTime; }
void setCpuBurstTime(double newTime) { cpuBurstTime = newTime; }
void setCpuBurstTime(double newTime) { cpuBurstTime = newTime; }
void setAccuTicks(double newTick) { accuTicks = newTick; }
```

定義一些get和set function

Thread.cc

```
Thread::Thread(char* threadName, int threadID)

ID = threadID;
name = threadName;
// MP3
priority = 0;
queueLevel = 0;

totalWaitingTime = 0;
startWaitingTime = 0;
cpuStartTime = 0;
approxBurstTime = 0;
approxRemainTime = 0;
accuTicks = 0;
accuTicks = 0;
```

Thread::Thread把剛才新加的data初始化為0

```
void
Thread::Yield ()

If thread *nextThread;
IntStatus oldLevel = kernel->interrupt->SetLevel(IntOff);

DEBUG(dbgThread, "Yielding thread: " << name);

Calculate burst and remain time

// // nextThread = kernel->scheduler->FindNextToRun();

// this->IncreaseCpuBurstTime();

// this->increaseCpuBurstTime();

// this->setCpuStartTime(kernel->stats->totalTicks); // Because we have to keep accumulate burst time

// // do aging and then check preemptive

// kernel->scheduler->Aging();

// kernel->scheduler->FindNextToRun();

// nextThread = kernel->scheduler->CheckPreemptive();
nextThread = kernel->scheduler->FindNextToRun();

// kernel->scheduler->ReadyToRun(this);
k
```

Thread::Yield()

current thread要將讓出cpu給其他thread, 呼叫FindNextToRun()去找下一個thread

Thread::Sleep()

先計算accuTick、cpuBurstTime、approxBurstTime,然後因為thread從runnning移到了waiting,故要重設cpuBurstTime為0

Thread::InsertedIntoQueue()當有thread priority改變, 以致從一個queue跑到另一個 queue時呼叫(如 L2->L1, 就會insert到L1)

Thread::RemovedFromQueue()當有thread priority改變, 以致從一個queue跑到另一個 queue時呼叫(如L2->L1, 就會從L2移出)

```
310  void
311  Thread::ChangePriority()
312  {
    int oldPriority = this->priority;
    int newPriority = min(oldPriority + 10, 149);
314    priority = newPriority;
315
316    priority = newPriority;
317
318    DEBUG(z, "[c] Tick [" << kernel->stats->totalTicks << "]: Thread [" << ID << "] changes its priority from [" << oldPriority << "] to [" << newPages]
319  }</pre>
```

將priority增加10,上限是149

```
341  // MP3
342  void
343  Thread::IncreaseCpuBurstTime()
344  {
345  | cpuBurstTime = cpuBurstTime + (kernel->stats->totalTicks - cpuStartTime);
346  }
```

累積cpuBurstTime

累積totalWaitingTime, 如果等了超過1500 Ticks, 則減去1500並回傳true, 代表要增加priority

```
Thread::UpdateApproxBurstTime()

double newApproxBurstTime = 0.5 * cpuBurstTime + 0.5 * approxBurstTime;

DEBUG(z, "[D] Tick [" << kernel->stats->totalTicks << "]: Thread [" << ID << "

// update approxBurstTime and reset approxRemainTime
approxBurstTime = newApproxBurstTime;
approxRemainTime = newApproxBurstTime;

approxRemainTime = newApproxBurstTime;
```

當thread從running to waiting,用公式計算新的approxBurstTime和approxRemainTime

更新approxRemainTime

```
370  // MP3
371  void
372  Thread::IncreaseAccuTicks()
373  {
374  | accuTicks = accuTicks + (kernel->stats->totalTicks - cpuStartTime);
375  }
```

累積AccuTicks

心得

part 1 因為跟前兩次有蠻多地方重疊,所以code trace比較快,也比較快理解thread是如何在各個state切換。

part 2 則讓我們學會multy queueLevel scheduling是如何在OS上實作的,像是non-preemtive 跟preemptive在實作上的差異,以及SJF如何實現的等等。