

## Intuition for new memory hierarchy

### V2P address translation becomes a bottleneck

- Memory capacity significantly grows in the past decade
- Entries in Translation Lookaside Buffers are limited



### Adding a new layer in current memory hierarchy

- New memory architecture — Midgard<sup>[1]</sup> proposes a new layer to provide translation/protection at VMA granularity
- New cache hierarchy uses Midgard addresses

### Explore the reasonable LLC size for new hierarchy

- Cache miss ratio decreases as LLC size increases
- Large LLC brings additional cost

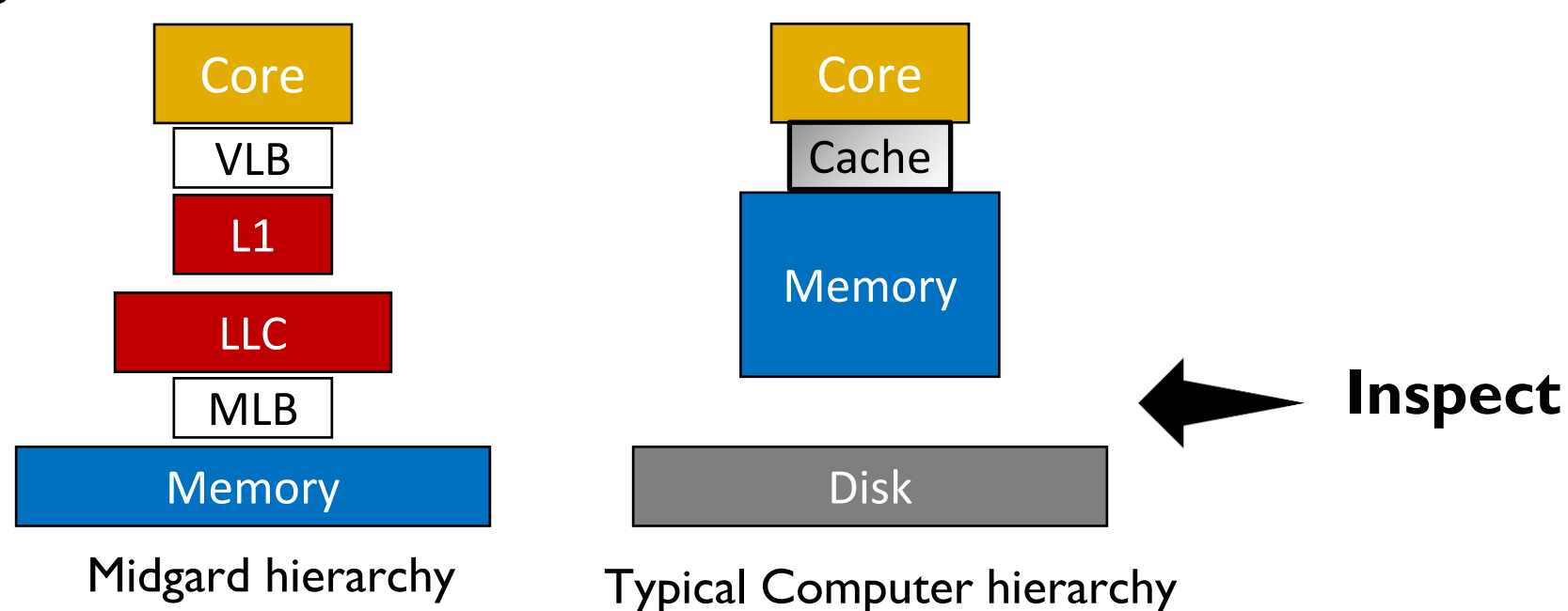


## Address translation drives memory hierarchy design

## Utilize typical computer architecture

### Memory hierarchy in testing

- Use container to control the resources a program can access
- Assume hard disk as the memory abstraction
- Memory can be regarded as the cache abstraction for the memory in Midgard



Cache miss rates in Midgard can be evaluated by page faults rate in traditional hierarchy

### Page faults rate as metrics

## Characterize the cache behavior

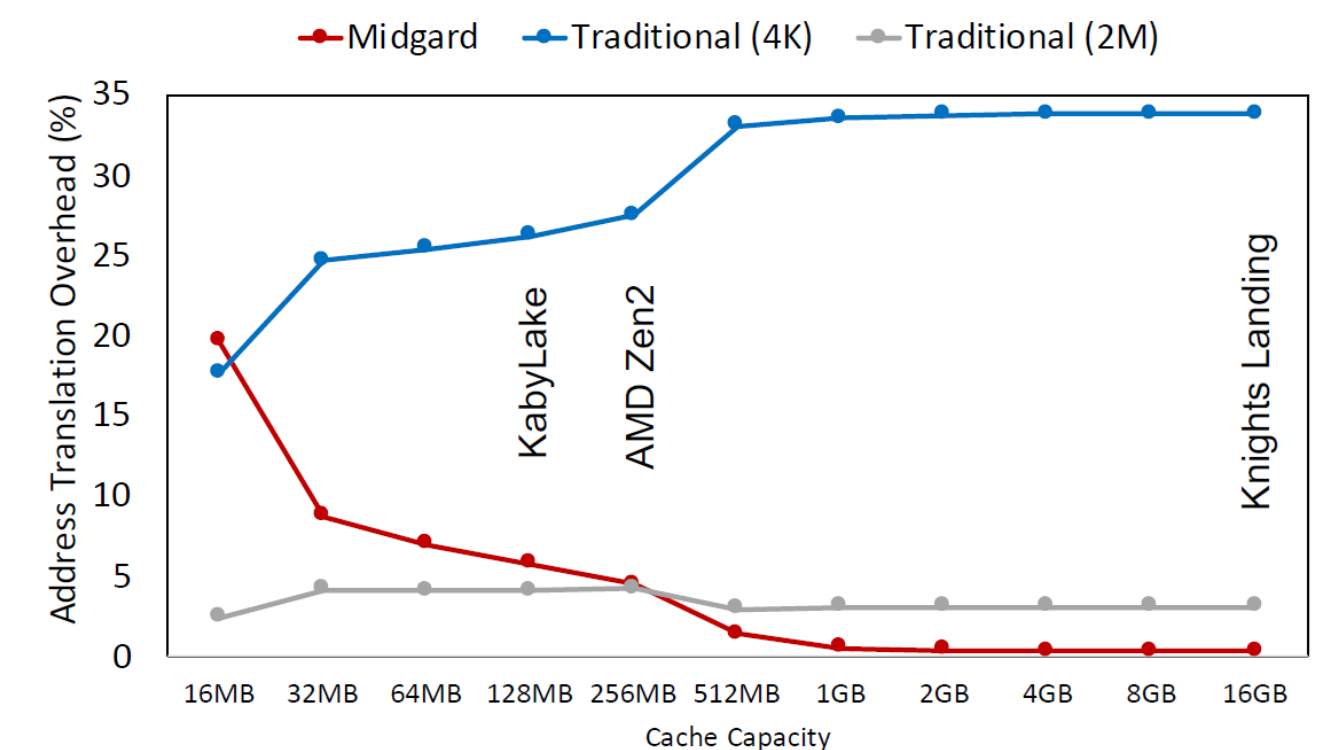
### Evaluation based on LLC size for Midgard needs testbench

- Tested under real-world and widely-used workloads

**CloudSuite:** Benchmark suite contains applications which are popular among today's datacenters.



Under the dataset of 200GB, Midgard can provide nearly 0% Address Translation Overhead when Cache Capacity is 4GB (2% of the dataset size).



## Characterizing under real-world workloads

## Metrics monitor tools

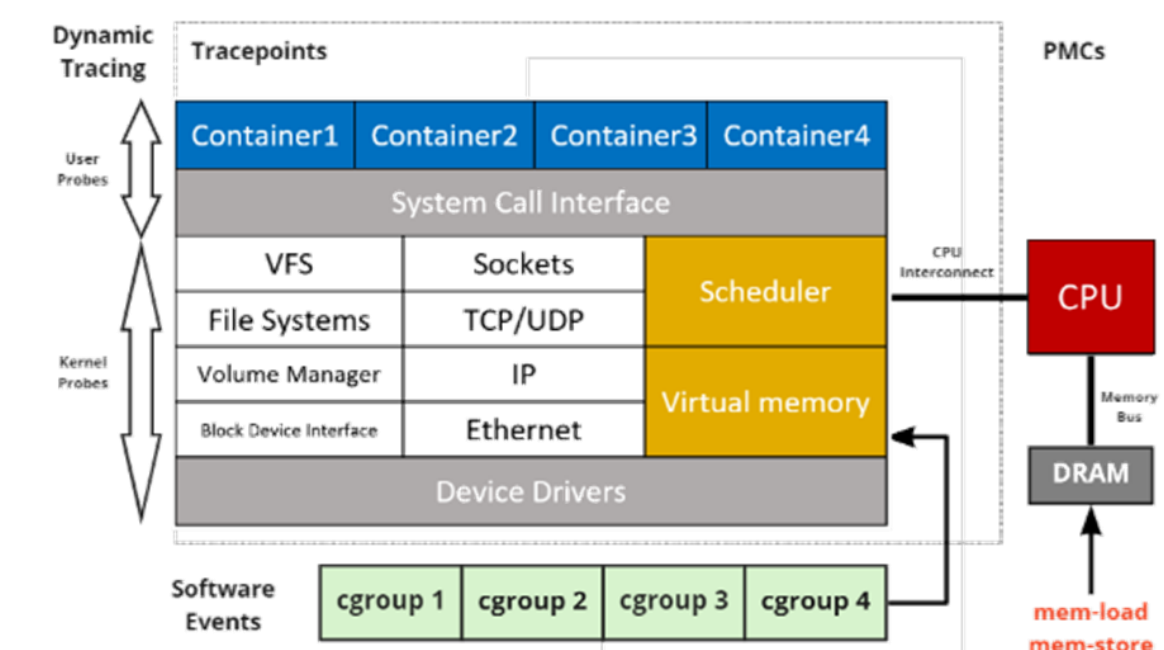
### Docker runtime metrics

- Docker provides tools to limit the memory usage
- Control groups in Linux containers expose memory measurements including major page faults



### Perf tools monitor system events

- Perf records the total memory accesses during the server computation



### Perf & cgroup works together

## Evaluation

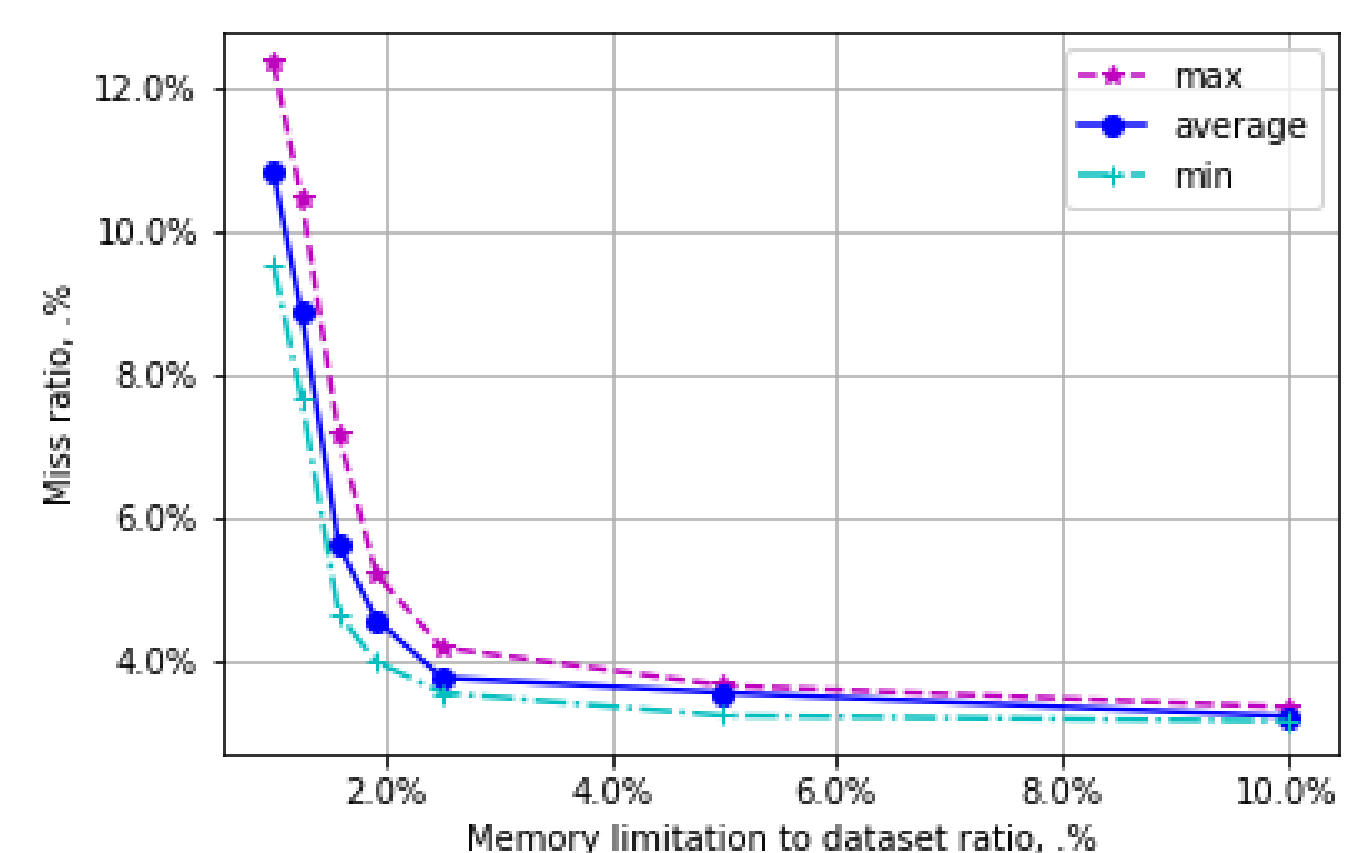
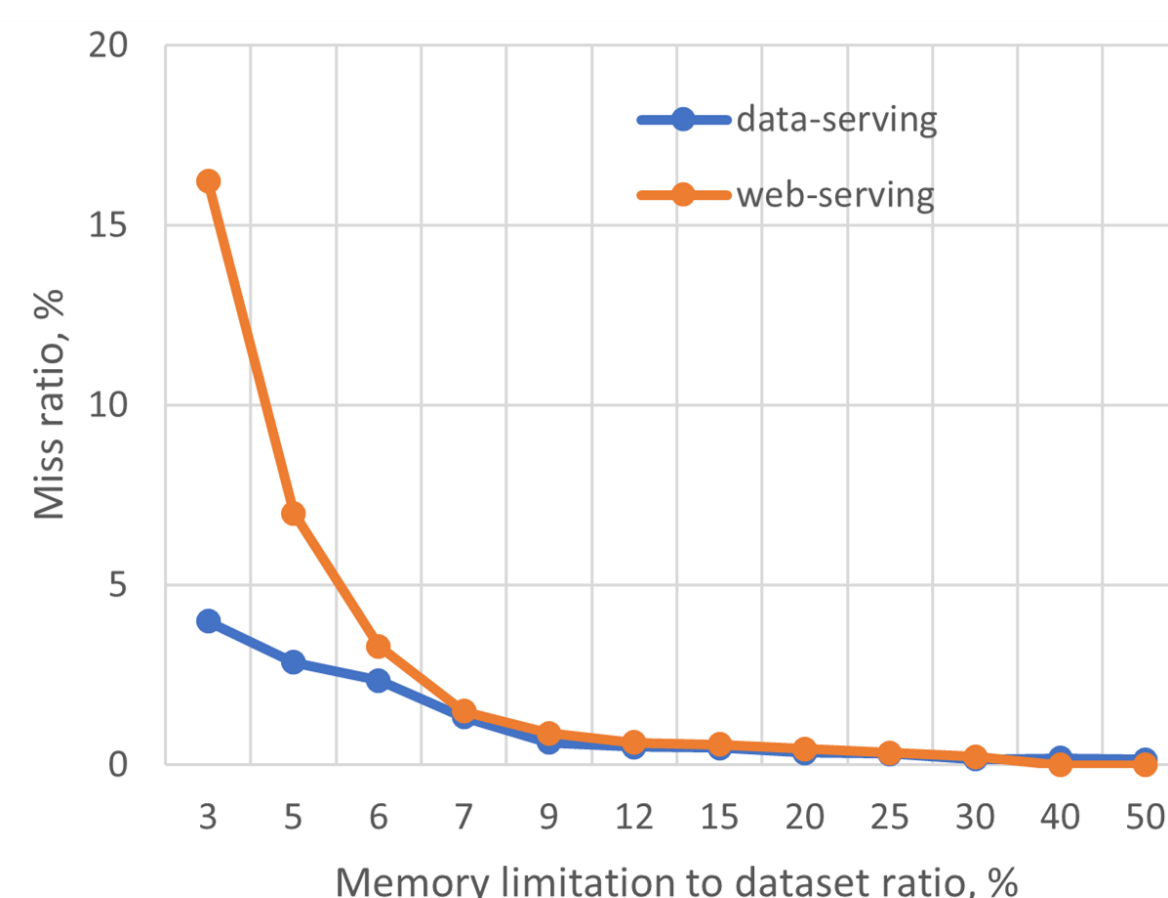
### Methodology

- Use of a scientific infrastructure designed for cloud-related experiments: CloudLab.
- Use of multi-core Intel processors with 128GB of RAM on average and ~1 TB of storage for the datasets.



### Characterisation of memory usage by workloads

- Temporal locality is observed at the simulated cache level
- An inflection point separates a memory-intensive region from one where few new data are requested  
⇒ Around 5% of the considered dataset size.



### Conclusions

- Empirical results over multiple workloads compatible with Midgard claims.
- Costlier LLC misses in Midgard system can therefore be compensated by faster virtual to midgard addresses translations thanks to a new level of indirection (Virtual Memory Areas of unrestricted size).

Under real-world workloads, LLC of ~5% of dataset size provides <3% miss ratio