

CS 173: Discrete Structures, Summer 2014

Homework 2

This homework contains 5 problems, each worth 10 points. It is due in class on Wednesday, July 2nd. **Please follow the guidelines on the class web page about homework format and style.**

In all questions, you must explain how you get your answers. Stating the answer with no supporting work will not receive full credit.

1. Prime modulus

- (a) Prove that if x is an integer, p is a prime, and $x^2 \equiv 1 \pmod{p}$, then either $x \equiv 1 \pmod{p}$ or $x \equiv -1 \pmod{p}$ or both. (Hint: Use Euclid's lemma.)

Solution. Suppose $x \in \mathbb{Z}$ and p is prime. Then,

$$x^2 \equiv 1 \pmod{p} \implies p \mid x^2 - 1 \tag{1}$$

$$\implies p \mid (x+1)(x-1). \tag{2}$$

By Euclid's lemma, either $p \mid (x+1)$ or $p \mid (x-1)$.

Case 1: $p \mid (x+1)$. In this case,

$$x+1 \equiv 0 \pmod{p} \implies x \equiv -1 \pmod{p}.$$

Case 2: $p \mid (x-1)$. In this case,

$$x-1 \equiv 0 \pmod{p} \implies x \equiv 1 \pmod{p}.$$

In both cases, either $x \equiv 1 \pmod{p}$ or $x \equiv -1 \pmod{p}$, as desired. \square

- (b) Show that this statement is false if p is allowed to be composite. (In other words, disprove the following statement: If x and p are integers, p is positive, and $x^2 \equiv 1 \pmod{p}$, then either $x \equiv 1 \pmod{p}$ or $x \equiv -1 \pmod{p}$ or both.)

Solution. The statement we want to disprove is

$$\forall x \in \mathbb{Z}, \forall p \in \mathbb{Z}^+, [x^2 \equiv 1 \pmod{p}] \rightarrow [x \equiv 1 \pmod{p} \vee x \equiv -1 \pmod{p}]$$

The negation of this is:

$$\exists x \in \mathbb{Z}, \exists p \in \mathbb{Z}^+, x^2 \equiv 1 \pmod{p} \wedge x \not\equiv 1 \pmod{p} \wedge x \not\equiv -1 \pmod{p}.$$

To prove this negation, let $x = 3$ and $p = 8$. Then,

$$3^2 \equiv 1 \pmod{8}, \tag{3}$$

$$3 \not\equiv 1 \pmod{8}, \tag{4}$$

$$3 \not\equiv -1 \pmod{8}, \tag{5}$$

as desired. \square

2. Extended Euclidean algorithm

- (a) Let $a = 826$ and $b = 470$. Find a pair of integers (s_1, t_1) such that $s_1a + t_1b = \gcd(a, b)$.

Solution. Running Euclid's algorithm, we get

$$826 = 1(470) + 356 \quad (6)$$

$$470 = 1(356) + 114 \quad (7)$$

$$356 = 3(114) + 14 \quad (8)$$

$$114 = 8(14) + 2 \quad (9)$$

$$14 = 7(2). \quad (10)$$

Thus, $\gcd(a, b) = 2$ and

$$2 = 114 - 8(14) \quad (11)$$

$$= (470 - 356) - 8(356 - 3(114)) \quad (12)$$

$$= 470 - 9(356) + 24(114) \quad (13)$$

$$= 470 - 9(826 - 470) + 24(470 - 356) \quad (14)$$

$$= 34(470) - 9(826) - 24(356) \quad (15)$$

$$= 34(470) - 9(826) - 24(826 - 470) \quad (16)$$

$$= 58(470) - 33(826), \quad (17)$$

so $(s_1, t_1) = \boxed{(-33, 58)}$ works. \square

- (b) Find a different pair of integers (s_2, t_2) such that $s_2a + t_2b = \gcd(a, b)$. (Hint: Find integers s and t such that $sa + tb = 0$. What is $(s_1 + s)a + (t_1 + t)b$?)

Solution. We have

$$(470 - 33)826 + (58 - 826)470 \quad (18)$$

$$= [(-33)826 + 58(470)] + [(470)(826) + (-826)(470)] \quad (19)$$

$$= 2 + 0 = 2, \quad (20)$$

so $(s_2, t_2) = \boxed{(437, -768)}$ works. $\boxed{(-503, 884)}$ also works. \square

3. Modular arithmetic

- (a) Evaluate $7^{500} \bmod 17$.

Solution. By repeated squaring, we have

$$7^2 \equiv 49 \equiv -2 \pmod{17} \quad (21)$$

$$7^4 \equiv 4 \pmod{17} \quad (22)$$

$$7^8 \equiv 16 \equiv -1 \pmod{17} \quad (23)$$

$$7^{16} \equiv 1 \pmod{17} \quad (24)$$

$$7^{32} \equiv 1 \pmod{17} \quad (25)$$

$$\vdots \quad (26)$$

Thus,

$$7^{500} \equiv 7^{256+128+64+32+16+4} \equiv 7^{256} 7^{128} 7^{64} 7^{32} 7^{16} 7^4 \equiv 4 \pmod{17}$$

and $7^{500} \bmod 17 = \boxed{4}$. Alternatively, Fermat's theorem tells us that $7^{16} \equiv 1 \pmod{17}$, so

$$7^{500} \equiv 7^{500 \bmod 16} \equiv 7^4 \equiv 4 \pmod{17}.$$

□

- (b) Use Fermat's theorem to find a modulo 17 multiplicative inverse of 7.

Solution. By Fermat's theorem, we want $7^{15} \bmod 17$. Since

$$7^{15} \equiv 7^{8+4+2+1} \equiv 7^8 7^4 7^2 7 \equiv (-1)(4)(-2)(7) \equiv 56 \equiv 5 \pmod{17},$$

the answer is $\boxed{5}$. Indeed $(7)(5) \equiv 35 \equiv 1 \pmod{17}$. □

4. Sets

- (a) Prove that if A , B , and C are sets, then $A - (B \cap C) \subseteq (A - B) \cup (A - C)$. (Note: To show that $S \subseteq T$, you must pick an arbitrary element of S and show that it is in T . A Venn diagram does not count as a proof.)

Solution. Suppose that $x \in A - (B \cap C)$. Then, $x \in A$ and $x \notin B \cap C$. By De Morgan's Law, this means that $x \notin B$ or $x \notin C$.

Case 1: $x \notin B$. Then, $x \in A - B$, so $x \in (A - B) \cup (A - C)$.

Case 2: $x \notin C$. Then, $x \in A - C$, so $x \in (A - B) \cup (A - C)$.

In both cases, $x \in (A - B) \cup (A - C)$, so $A - (B \cap C) \subseteq (A - B) \cup (A - C)$. □

- (b) Let A and B be sets. Prove that $A \subseteq B$ if and only if $A - B = \emptyset$.

Solution. (\Rightarrow): We will prove the contrapositive. The contrapositive is: If $A - B \neq \emptyset$, then $A \not\subseteq B$. If $A - B \neq \emptyset$, then we can let $x \in A - B$. This means that $x \in A$ but $x \notin B$. This means that not every element in A is in B also, so $A \not\subseteq B$.

(\Leftarrow): Again, we will prove the contrapositive: If $A \not\subseteq B$, then $A - B \neq \emptyset$. If $A \not\subseteq B$, then not every element in A is also in B , so there must be some $x \in A$ such that $x \notin B$. Then, $x \in A - B$, so $A - B \neq \emptyset$. □

5. Partial Order

Let $A = \mathbb{Z}^2$, and define the relation R as follows: $(x_1, y_1) R (x_2, y_2)$ if

- (a) $x_1 < x_2$, or
- (b) $x_1 = x_2$ and $y_1 \leq y_2$.

Show that R is a partial order.

Solution. We must show that R is reflexive, antisymmetric, and transitive.

- Reflexive: Suppose that $(x, y) \in A$. Since $x = x$ and $y \leq y$, we have $(x, y) R (x, y)$, so R is reflexive.
- Antisymmetric: Suppose that $(x_1, y_1) R (x_2, y_2)$ and $(x_2, y_2) R (x_1, y_1)$. Then, $x_1 \leq x_2$ and $x_2 \leq x_1$. The only way this is possible is if $x_1 = x_2$. This means that $y_1 \leq y_2$ and $y_2 \leq y_1$. The only way this is possible is if $y_1 = y_2$. Thus, $(x_1, y_1) = (x_2, y_2)$, and R is antisymmetric.
- Transitive: Suppose that $(x_1, y_1) R (x_2, y_2)$ and $(x_2, y_2) R (x_3, y_3)$. Then, $x_1 \leq x_2$ and $x_2 \leq x_3$, so $x_1 \leq x_3$. There are two cases:
Case 1: $x_1 < x_3$. Then, $(x_1, y_1) R (x_3, y_3)$.
Case 2: $x_1 = x_3$. Since $x_1 \leq x_2 \leq x_3$, this means that $x_1 = x_2 = x_3$. Hence, $y_1 \leq y_2$ and $y_2 \leq y_3$, so $y_1 \leq y_3$. This means that $(x_1, y_1) R (x_3, y_3)$.
In both cases, we have $(x_1, y_1) R (x_3, y_3)$, so R is transitive.

□