

Today: Simply connectedness, complex logarithm, complex roots.

Pf. of $\text{mero. fun on } \hat{\mathbb{C}}$ are rat^l poly. $\frac{P(z)}{Q(z)}$ (Sketch):

- $\hat{\mathbb{C}}$ is cpt set.
- the sing. pts of a mero. fun are isolated, set of isolated pts in a cpt set is always finite.
⇒ set of poles is a finite set

- Suppose f has a pole at ∞

i.e. $f\left(\frac{1}{w}\right)$ has a pole at $w=0$

$$\Rightarrow \frac{a_{-n}}{w^n} + \frac{a_{-n+1}}{w^{n-1}} + \dots + \frac{a_{-1}}{w} + \text{holo. fr. (near } w=0\text{)}$$

• $f\left(\frac{1}{w}\right) - \left(\frac{a_{-n}}{w^n} + \dots + \frac{a_{-1}}{w} \right)$ holo. near $w=0$.

• $f(z) - (a_{-n}z^n + \dots + a_{-1}z)$ holo. near $z=\infty$.

Similarly, for other poles $\alpha_1, \dots, \alpha_n \in \mathbb{C}$

locally near each α_i ,

$$f(z) = \frac{* \text{ const.}}{(z-\alpha_i)^k} + \dots + \frac{*}{z-\alpha_i} + \text{holo. ...}$$

$$\Rightarrow f(z) - \left(\frac{*}{(z-\alpha_i)^k} + \dots + \frac{*}{z-\alpha_i} \right) \text{ holo. near } \alpha_i$$

$$\Rightarrow f(z) - \left(f_{\text{prin}, \alpha_1}(z) + \dots + f_{\text{prin}, \alpha_n}(z) \right) \text{ holo. in } \hat{\mathbb{C}} \equiv \text{const.}$$

rat^l fun

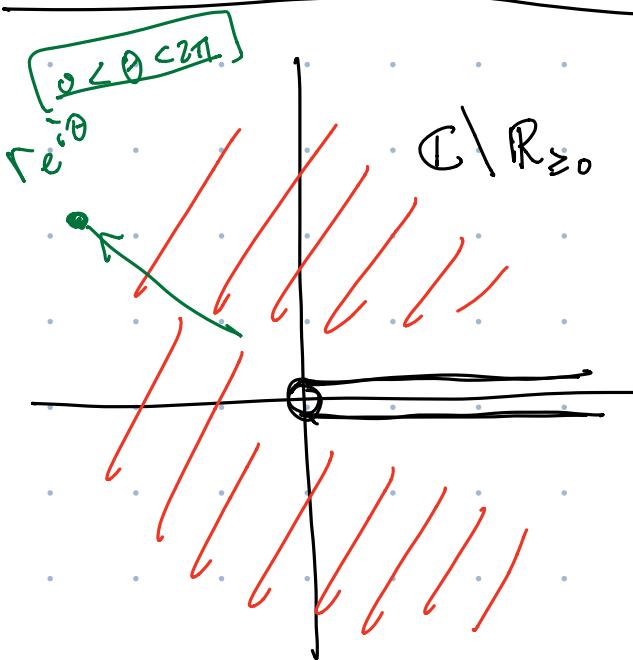
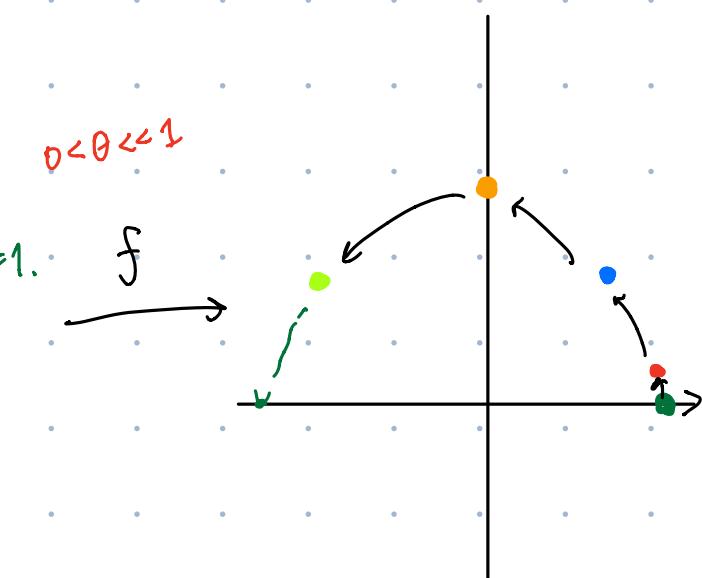
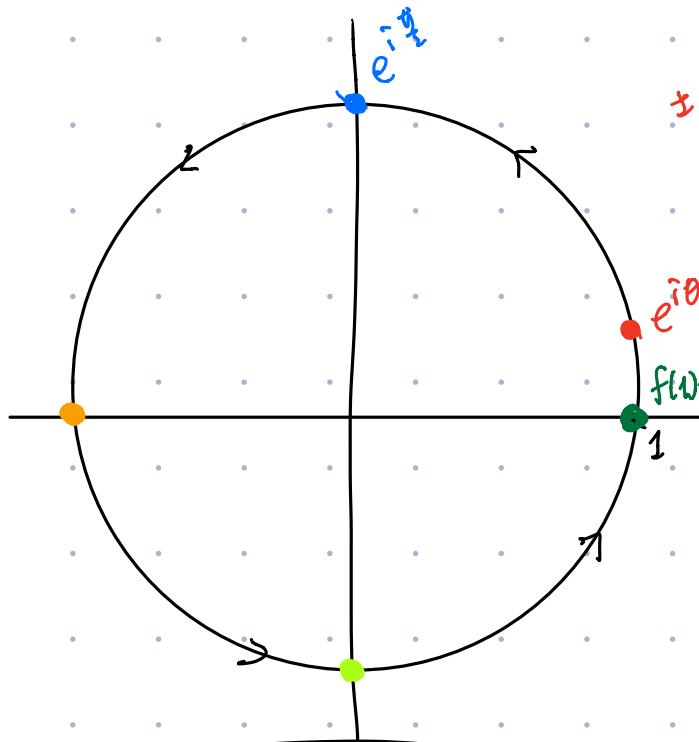


"Complex square root"

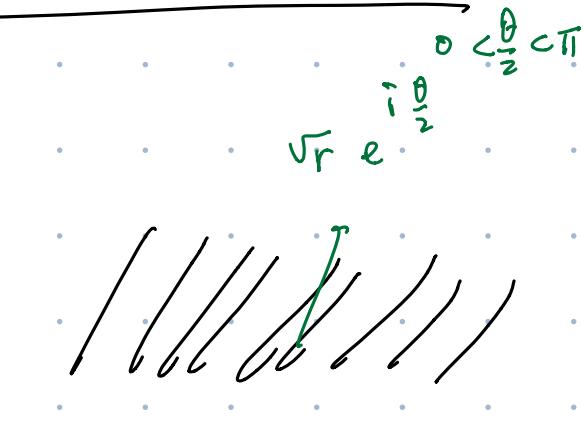
Q: Does there exist an entire fun: $f: \mathbb{C} \rightarrow \mathbb{C}$

a1. $f(z)^2 = z \quad \forall z \in \mathbb{C}$

if $z = r e^{i\theta}$, $f(z) = \sqrt{r} e^{\frac{i\theta}{2}}, -\sqrt{r} e^{\frac{i\theta}{2}}$
 the roots of $\frac{f^2}{r^2} = z$



$$\frac{\sqrt{z}}{2\pi i}$$



More generally, the square root fun can be defined on
any simply connected domain not containing 0.

Recall: Ω is simply connected

\Leftrightarrow any hol. fun on Ω has a primitive. F ($F' = f$)

$$(F' = f)$$

Def: $\Omega \subseteq \mathbb{C}$ open.

γ_0, γ_1 curves in Ω w/ the same end points.

$$\gamma_0, \gamma_1: [a, b] \rightarrow \Omega, \quad \gamma_0(a) = \gamma_1(a) = \alpha \in \Omega$$

$$\gamma_0(b) = \gamma_1(b) = \beta \in \Omega$$

Say γ_0, γ_1 are homotopic in Ω if

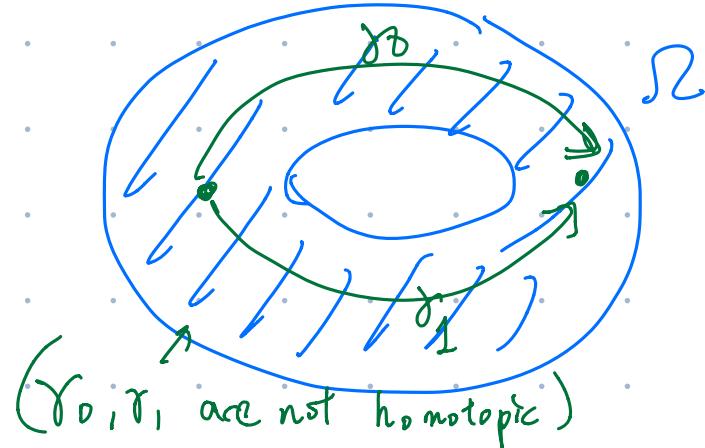
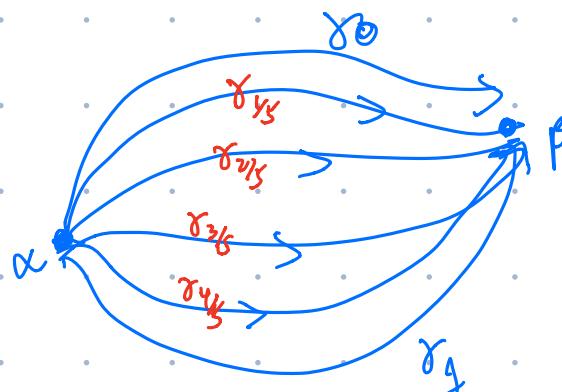
$$\exists \gamma: [a, b] \times [0, 1] \rightarrow \Omega \quad \text{cont. fun.}$$

$$(t, s) \mapsto \gamma(t, s).$$

at.

- $\gamma(0, s) = \alpha, \quad \gamma(1, s) = \beta \quad \forall s \in [0, 1],$
- $\gamma(t, 0) = \gamma_0(t), \quad \gamma(t, 1) = \gamma_1(t) \quad \forall t \in [a, b].$

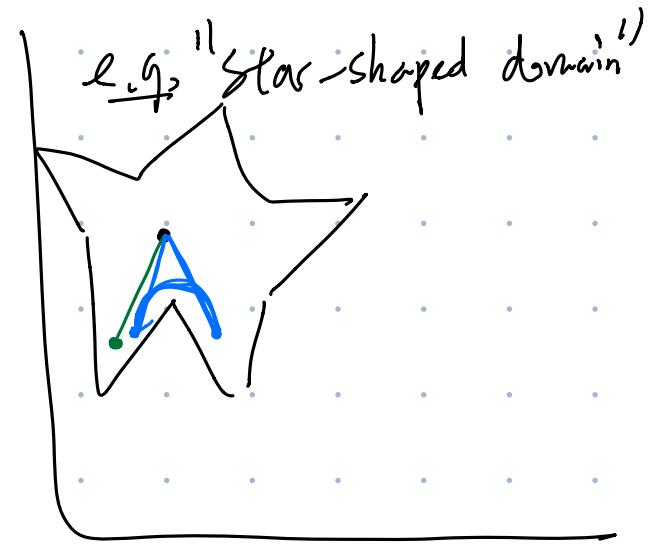
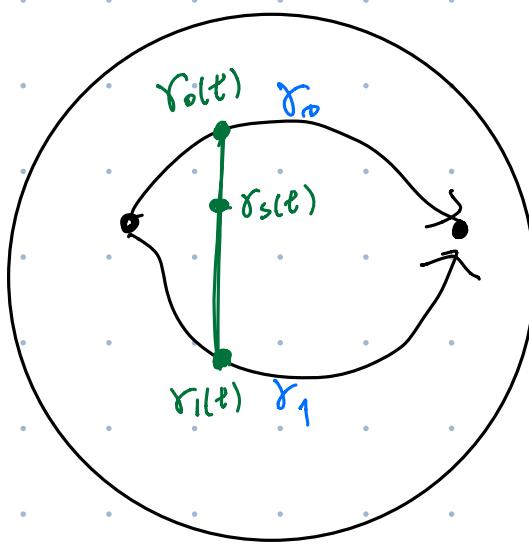
(think of γ as a family of parametrized curves γ_s .)



Def. A domain is simply connected if any two curves w/
the same endpoints are homotopic.

e.g. \mathbb{D}

/
or any
convex set
 $\subset \mathbb{C}$



$$\gamma(t, s) := (1-s)\gamma_0(t) + s\gamma_1(t) \quad \forall s \in [0, 1] \\ t \in [\alpha, \beta]$$

Thm (Contour integral is deformation invariant)

$f: \Omega \rightarrow \mathbb{C}$ holomorphic, $\gamma_0, \gamma_1 \subseteq \Omega$ homotopic.

Then $\int_{\gamma_0} f(z) dz = \int_{\gamma_1} f(z) dz.$

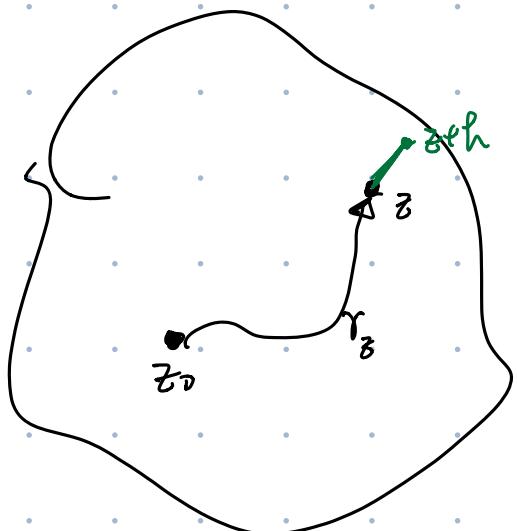
Thm $\Omega = \overline{\text{simply connected}}$, $f: \Omega \rightarrow \mathbb{C}$ holomorphic.

Then f has a primitive, F , in Ω ,

(i.e. $\exists F: \Omega \rightarrow \mathbb{C}$ holomorphic s.t. $F' = f$)

Pf: Choose any $z_0 \in \Omega$,

$\forall z \in \Omega$, choose any path γ_z in Ω connecting z_0 & z



Define $F(z) := \int_{\gamma_z} f(z) dz,$

(it's a well-defined fun on \mathbb{C} b/c
 $S2$ is simply connected & the deformation
 inv. of contour integrals)

Goal:
 $\lim_{h \rightarrow 0} \frac{F(z+h) - F(z)}{h}$ exists & $= f(z)$

$$\left[\frac{1}{h} \int_{z}^{z+h} f(z) dz \right]$$

we proved before
 (only uses f is conti.)

□

Complex logarithm:

Q: Does there exist a hol. fun. $f: \mathbb{C}^* \rightarrow \mathbb{C}$.

$$e^f = z \quad \forall z \in \mathbb{C}^*$$

No

Thm. Ω simply connected, $0 \notin \Omega$.

- $z_0 \in \Omega$, $z_0 = r e^{i\theta}$, $r > 0$, $\theta \in \mathbb{R}$

Then: $\exists!$ $F: \Omega \rightarrow \mathbb{C}$ holo.

at.

- $e^{F(z)} = z \quad \forall z \in \Omega$.
- $F(z_0) = \log r + i\theta$

$$\boxed{e^F = z_0 = r e^{i\theta} = e^{\log r + i\theta}}$$

$$\boxed{\log r + i(\theta + 2\pi k)}$$

pf: $\frac{1}{z}$ holo. in Ω $\Rightarrow \exists$ primitive F on Ω

↑
simply connected

- $F'(z) = \frac{1}{z} \quad \forall z \in \Omega$
- $F(z_0) = \log r + i\theta$

$$(z e^{-F(z)})' = e^{-F(z)} - z F'(z) e^{-F(z)} \\ = e^{-F(z)} \cdot (1 - z F'(z)) = 0$$

$$\Rightarrow z \cdot e^{-F(z)} = \text{const.} = 1$$

$$z_0 \cdot e^{-F(z_0)} = r e^{i\theta} e^{-(\log r + i\theta)} = 1.$$

Uniqueness. $\boxed{F, \tilde{F}}$

$$\boxed{F(z) - \tilde{F}(z) = G(z)} : \Omega \rightarrow \mathbb{C}$$

- $G(z_0) = 0$
- $e^{\underline{G(z)}} = e^{F(z) - \tilde{F}(z)} \\ = e^{F(z)} / e^{\tilde{F}(z)} = 1$
- $\Rightarrow G(z) = 2\pi i z \quad \forall z$

$$\boxed{e^{F(z)} = z}$$

□

Thm: Ω simply connected, $f: \Omega \rightarrow \mathbb{C}$ nonvanishing holo.

Then $\exists g: \Omega \rightarrow \mathbb{C}$ holo s.t. $f(z) = e^{g(z)} \quad \forall z \in \Omega$

Idea: $(f \cdot e^{-g})' = f' e^{-g} - fg' e^{-g}$

$$f' = fg' \quad g' = \boxed{\frac{f'}{f}}$$

pf: $\frac{f'}{f}$ hol. in Ω $\Rightarrow \exists$ primitive g on Ω
 \uparrow
 Simply conn. $g' = \frac{f'}{f}$

- $g' = \frac{f'}{f}$
- Choose $g(z_0)$ s.t. $e^{g(z_0)} = f(z_0)$

$$\Rightarrow (f e^{-g})' = 0$$

$$\Rightarrow f(z) e^{-g(z)} = \text{const.} \equiv 1$$

$$\textcircled{B} \Rightarrow e^{g(z)} = \underline{f(z)} \quad \square$$

(Complex n-th root)

Thm. Ω simply conn. $0 \notin \Omega$.

$$z_0 \in \Omega. \quad z_0 = r e^{i\theta}, \quad r > 0, \quad \theta \in \mathbb{R}.$$

Then $\exists!$ $F: \Omega \rightarrow \mathbb{C}$ hol. s.t.

- $F(z)^n = z$

- $F(z_0) = r^{\frac{1}{n}} e^{i\frac{\theta}{n}}$

RF: Define

$$F(z) = e^{\frac{\log_2(z)}{n}}$$

$$F(z)^n = e^{\log_2(z)} = z \quad \forall z \in \Omega.$$

$$F(z_0) = e^{\frac{\log_2(z_0)}{n}} = e^{\frac{\log r + i\theta}{n}} = r^{\frac{1}{n}} e^{i\frac{\theta}{n}}$$
