Threshold Bandit, With and Without Censored Feedback

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Previous work

- A sequential decision making
- Select an arm i and receive a stochastic reward R_i^t
 - Independently and identically distribution
- Action is a function of past observed (action, reward) pairs

Previous work

- Optimal solution
 - Pull the arm with highest expected reward
 - Identical in each round
- UCB(Upper Confidence Bound) algorithm
 - Maximization over rewards estimated from previous data
 - Bias each estimation according to its uncertainty
 - Mean + confidence interval

Threshold bandit

- A piece of side information c_t
- Choose one arm $i \in [K]$, produce a value X_i^t
 - Survival function $F_i(x) \sim \Pr(X_i^t \ge x)$
- Binary reward $R_i^t = \mathbb{I}[X_i^t \ge c^t]$
- Expected payoff $\mathbb{E}[R_i^t] = F_i(c^t)$
- The optimal arm is related the value of c^t
- c^t is drawn from $M = \{1, 2, ..., m\}$
 - Observe at the beginning of each round

Applications

- Packet Delivery with Deadlines
 - Ship a stream of packages from source to destination
 - Each package is supplied with a delivery deadline
 - Select a transportation route with the highest probability of on-time arrival
 - Uncertain: faster with higher volatility

Applications

- Supplier Selection
 - Produce a product satisfying specific quality demands
 - Select one of several suppliers to contract out the work
 - Uncertain: capabilities and variabilities of the products from each supplier

Applications

- Dark Pool Brokerage
 - Buy/sell various sized bundles of shares
 - Establish financial exchanges that match buyers and sellers
 - Execute the transaction if there is suitable liquidity
 - Choose the dark pool with highest probability of completion
 - Uncertain: the successful rate of transaction

Benchmark

- Compare with the best policy instead of the best arm
 - Optimal policy may incorporate with the threshold
- Regret
 - $Reg(T) = \mathbb{E}\left[\sum_{t=1}^{T} \left(\max_{i \in [n]} R_i^t R_{I^t}^t\right)\right]$
- Previous work
 - Static optimal solution
 - E.g., adversary bandit

Feedback

- Uncensored feedback
 - Observe the sample X_i^t regardless of the threshold
 - E.g., learn the travel time of a package regardless of the deadline having been met
- Censored feedback
 - Observe a null value when exceeding the threshold
 - Observe X_i^t otherwise
 - E.g., learn the product quality when it is rejected by the customer

A New Perspective on UCB

- Potential function
 - The current number of plays of each arm
- UCB algorithm
 - Redefine the score
 - Play arm $I^t = \operatorname{argmax}_i \hat{\mu}_i^t + \zeta(N_i^t, \delta)$
- The regret is bounded by potential functions
- Bound the number of pulls N_i^t following the standard techniques

Uncensored feedback

Estimate the survival function via empirical distribution

$$\bullet \ \widehat{F}_i^t(j) = \frac{\sum_{\tau=1}^{t-1} \mathbb{I}[X_{I^\tau}^\tau \ge j, I^\tau = i]}{N_i^t}$$

- The number of times that arm i is selected and exceeds the threshold j
- Converge to F_t^i at the exponential rate
- DKWUCB algorithm
 - Play arm $I^t = \operatorname{argmax}_i \hat{F}_i^t(c^t) + \zeta(N_i^t, \delta)$

Uncensored feedback

- Redefine potential function
 - The number of pulls of each arm when it is not optimal
 - Track the accumulation of the estimation gap
- Previous analysis in UCB
 - The number of pulls of a bad arm i is $O(\frac{\log(T)}{\Delta_i^2})$
 - The regret suffered by any such pull is $O(\Delta_i)$
 - The contribution of arm i to total regret is $O(\frac{\log(T)}{\Delta_i})$

Uncensored feedback

- Current analysis
 - Bound the difference between the empirical mean estimator and the true mean of the selected arm
 - Uniform convergence rate
 - For any error ϵ , there exists N such that $|f_n(x) f(x)| < \epsilon, \forall n > N$
 - Accuracy after enough rounds
 - Improve the bound with a constant factor

Censored feedback

 For a given threshold value, all the feedbacks from larger threshold values are useful

•
$$\hat{F}_i^t(j) = \frac{\sum_{\tau=1}^{t-1} \mathbb{I}[\min\{X_{I^{\tau}}^{\tau}, c^t\} \ge j, I^{\tau} = i]}{N_i^t(j)}$$

- ullet Realization and threshold exceed the threshold j
- Error bound estimator
 - Define the martingale sequence
 - The summation
 - The gap between expected reward and real reward for each arm
 - By Azuma's inequality

Martingale Sequence

- $\mathbb{E}[Y_{n+1}|Y_1,\ldots,Y_n]=Y_n$
- The conditional expected value of the next observation = the value of last observation
 - Given all the past observations
- Knowledge of past events never helps predict the mean of the future winnings
- Only the current event matters

Martingale Sequence

- Betting game
- Y_n : A gambler's cumulative fortune after n tosses of a fair coin
 - Win 1 if the coin comes up heads
 - Lose 1 if it's tails
- The gambler's conditional expected fortune after the next trial
 - Given the history
 - Equal to his present fortune

Azuma's inequality

- When the values of martingales that have bounded differences
- High probability event
- $\Pr[X_n > t] \le \exp(\frac{-t^2}{2n})$
- Provide a bound similar in form to the Chernoff bound
 - Without assuming independence

The order of threshold values

- Adversarial setting
 - Threshold values arrive in a non-decreasing order
 - 1,1 ..., 1,2, ... 2,3, ..., *m*
 - Only the samples observed within the same stage can help to inform decision making
 - Play m independent copies of bandits
 - Reduce to stochastic MAB problems
 - Regret scales with m

The order of threshold values

- Optimistic setting
 - Threshold values arrive in a non-increasing order
 - m, ..., m, m 1, ..., m 1, ..., 1
 - Make full use of samples
 - Every sample observed in the previous rounds is useful in later rounds
 - Reduce to the problem with uncensored feedback

Cyclic Permutation Setting

• Threshold values are a cyclic permutation order of 1,2,...,m

- Divide the time horizon into epochs of length Km
- Initialize: pull each arm once for each threshold value
- Arm elimination process
 - Leverage information across threshold values
 - Utilize observations from higher thresholds to estimate mean payoffs for lower thresholds
 - Information does not flow in the other direction

Arm elimination process

- Once all but one arm has been eliminated, proceed to pull the single remaining arm for the given threshold value
- For any threshold values below j
 - The best arm has been determined
 - Play that arm constantly
- For any threshold values greater or equal to j
 - Explores all arms uniformly
- All sub-optimal arms can be detected after $O(\log T)$ epochs with high probability

Algorithm

```
Algorithm 1 KMUCB
 1: Input: A set of arms 1, 2, \dots, K.
 2: Initialization: L_j \leftarrow [K] \ \forall j \in M, k \leftarrow 1, j \leftarrow 1
3: for epoch k = 1, 2, ..., T/Km do
          \operatorname{count}[j'] \leftarrow 0 \ \forall j' \in M
 4:
                                                        Number of pulls for threshold j
          for t from (t_{k-1}+1) to t_k do
 5:
              Observe c^t = j' and set count[j'] \leftarrow \text{count}[j'] + 1
 6:
              if j' < j then
 7:
                                                                                           Stick to the single best arm
                    I^t \leftarrow \text{index of the single arm remaining in } L_{i'}
 8:
 9:
              else
                                                                            Uniform explore
                    I^t \leftarrow \text{count}[j'].
10:
              end if
11:
                                                             guaranteed by cyclic permutation
          end for
12:
         if j \le m and \max_{i' \in [K]} \hat{F}_{i'}^{t_k}(j) - \hat{F}_i^{t_k}(j) \ge \sqrt{\frac{16\log(Tk)}{(m-j+1)k}} \ \forall i \in L_j \setminus \{\arg\max_{i' \in [K]} \hat{F}_{i'}^{t_k}(j)\} then
13:
14:
                                                                                  Estimation is accurate enough
                                                                          Remove all suboptimal arms and
          end if
15:
16: end for
                                                                                     increment threshold
```

Proof Sketch

- The probability that optimal arm for threshold j is eliminated
 - $O\left(\frac{1}{T}\right)$
- The number of times a suboptimal arm is pulled for threshold j
 - Under the condition that real optimal arm is not eliminated
 - Within $O(\log T)$ epochs
- Bounded by criteria in elimination

Summary

- Using potential function in analysis may be of independent interest
- The elimination process is different from the standard UCB work
- Similar to current problem
 - Only get feedback when reward does not exceed the threshold
 - A reverse direction
 - Threshold (demand) is uncertain
 - No side information

Thank you!