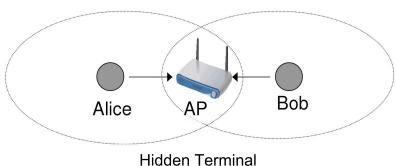
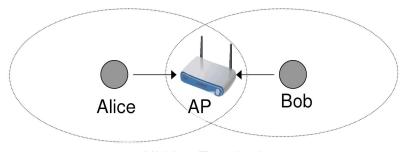
# A Brief Introduction to Interference Cancellation in Wireless Networks

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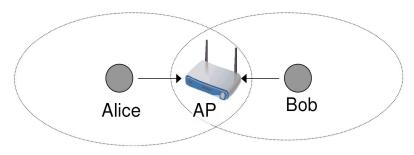
- 1 Introduction
- 2 Successive Interference Cancellation
- 3 ZigZag Decoding
- 4 Conclusion and Discussions





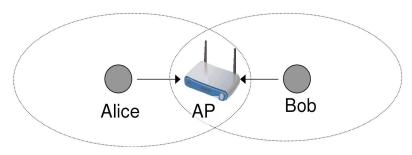
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- Engineering Perspective: With CSMA/CD, if Alice and Bob conduct transmission to the AP concurrently, they will suffer from packet loss and long delay.



Hidden Terminal

- Theory Perspective: Graph interference model is unrealistic. Is it possible to schedule Alice and Bob simultaneously in the SINR model?
- **Engineering Perspective**: Is it possible to reduce the packet loss rate?

### What is Interference Cancellation?

- No formal definition.
- A class of techniques.

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Interference Cancellation techniques are any technique or combination of techniques that allow an existing receiver to operate with co-channel interference and recover the data.

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$$SINR_i = \frac{P_i/d_{ii}^{\alpha}}{N_0 + \sum_{e_k \in \Lambda} P_k/d_{ki}^{\alpha}} \ge \beta$$

Here,  $\Lambda$  denotes the set of links that transmit simultaneously with  $e_i$ .  $P_i$  and  $P_k$  denote the transmission power at the transmitters of link  $e_i$  and  $e_k$ , respectively.  $d_{ii}$  ( $d_{ki}$ ) is the distance between transmitters of link  $e_i$  ( $e_k$ ) and the receiver of link  $e_i$ .  $N_0$  is ambient background noise,  $\alpha$  is path loss ratio,  $\beta$  is SINR threshold.

Successively recover the data and cancel its interference to other co-channel signals.

Suppose Alice's distance to AP is  $d_1$  and its transmission power is  $P_1$ , Bob's distance to AP is  $d_2$  and its transmission power is  $P_2$ .

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Without Successive Interference Cancellation

- Time slot 1: Alice to AP
- Time slot 2: Bob to AP

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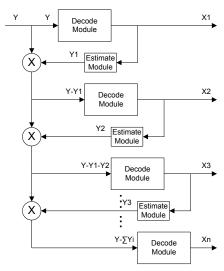
Without Successive Interference Cancellation

- Time slot 1: Alice to AP
- Time slot 2: Bob to AP

With Successive Interference Cancellation

■ Time slot 1: Alice to AP and Bob to AP

#### Execution details of SIC



### Condition for a successful SIC

$$SINR_i = \frac{P_i/d_{ii}^{\alpha}}{N_0 + \sum_{e_j \in \Lambda_i} P_j/d_{ji}^{\alpha} + \sum_{e_k \in \Lambda - \Lambda_i} P_k/d_{ki}^{\alpha}} \ge \beta$$

Here,  $\Lambda_i$  denotes the set of links that share the same receiver with  $e_i$  but canceled after  $e_i$ .

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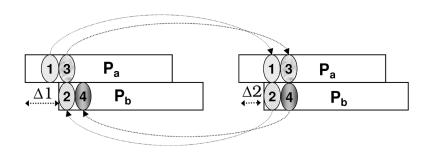
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**Feature**: Chunk-level decoding instead of packet-level decoding **Performance**: In a testbed of 14 USRP nodes, ZigZag reduces the average packet loss rate at hidden terminals from 72.6% to about 0.7%.

## Execution details of ZigZag



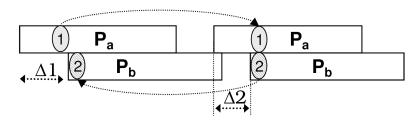


Figure: Overlapped collisions.

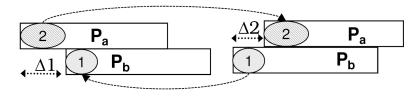


Figure: Flipped Order.

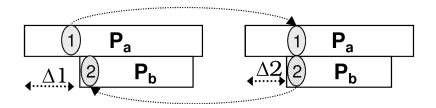


Figure: Different Packet Sizes.

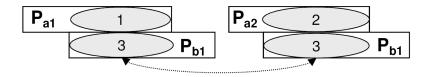


Figure: Alice's Packets Enjoy the Capture Effect.

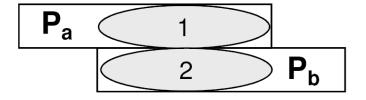


Figure: Single Decodable Collision; Inefficient Choice of Bit Rates.

└─ ZigZag Decoding

## **Examples**

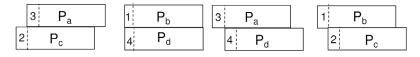


Figure: Nodes A and B are hidden from C and D.

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  - Power control and Exponential power increase: for the  $i^{th}$  canceled signal, we should have  $P_i \geq (1+\beta)^{n-i}\beta N_0 d_{ii}^{\alpha}$  (proved by induction).

Conclusion and Discussions

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- No requirement for scheduling, power control, synchronization or coding.
- No changes to 802.11 protocol.
- Drawbacks:
  - Cannot reduce the scheduling latency.

Conclusion and Discussions

# Thank You!