Modeling and Analysis of Peer-topeer Video on Demand Streaming

Outline

- Characteristics of P2P VoD streaming
- Design Problems of P2P VoD streaming
- Modeling and analysis of P2P VoD streaming

Asynchronous playback

- Due to random arrival time, peers watch different parts of a video at the same time
 - In P2P live streaming, the peers have a shared temporal content focus
 - P2P-VoD has a greater temporal diversity of content requests
 - Require each peer to contribute an amount of storage (such as 1GB) to compensate the problem

VCR functions

- Users can offer control over the video streams with VCR operations: pause & resume; random seek
- VCR operations introduce more dynamics to overlay structures.
- Need mechanisms to support VCR

Design Problems

- Segment Sizes
 - Tradeoff between overhead and utility of upload capacity
- Replication Strategy
 - Increase content availability
 - Replication strategies: LRU, LFU; MVC, SVC; weight-based evaluation process
- Content Discovery and Peer Overlay Management
- Piece selection
 - Sequential
 - rarest first
- Transmission strategy

Models on P2P VoD

- Analysis of Piece selection
 - C. Williamson et al., "Analysis of BitTorrent-like Protocols for On-Demand Stored Media Streaming", SIGMETRICS'08
- Analysis of the effect of VCR functions
 - Zhen Wei Zhao et al., "Modeling the Effect of User Interactions on Mesh-based P2P VOD Streaming Systems", To be published in INFOCOM'11

□ Assumptions:

- Each peer is allowed U concurrent upload connections, D concurrent download connections, D>U.
- Each connection achieves mean throughput C.
- Not consider peer selection strategy. Assume peers can find needed contents.
- Peers are cooperative

- □ Baseline model: Rarest-First
 - Fluid modeling approach
 - System efficiency η is taken as 1.
 - The change of downloader and seed population is:

$$\frac{dx}{dt} = \lambda - (x + y)UC,$$

$$\frac{dy}{dt} = (x + y)UC - \mu y.$$

The average number in steady-state:

$$\overline{x} = \lambda \left[\frac{1}{UC} - \frac{1}{\mu} \right] \qquad \overline{y} = \frac{\lambda}{\mu} \qquad T = \frac{\overline{x}}{\lambda} = \frac{1}{UC} - \frac{1}{\mu}.$$

- Strict In-Order Piece Selection (Random)
 - Peer relationships are asymetric: a given peer can only download from older peers, and can only provide content to younger peers
 - An uploader that receives more than U requests chooses at random U recipients for service. The remaining unsatisfied requests are purged from the system.

- Considering this, calculate the average downloading rate for a downloader.
 - \Box The probability for a peer been in system for time t to obtain a download connection:

$$p(t) = \frac{\tilde{U}(t)}{\tilde{D}(t)} = \alpha \frac{(x + y - \lambda t)U}{xD}.$$

The average downloading rate:

$$\begin{split} \gamma &=& \frac{1}{T} \int_0^T Dp(t) C dt \\ &=& \frac{1}{T} \int_0^T \alpha \left(\left(\frac{x+y}{x} \right) U C - \frac{\lambda t U C}{x} \right) dt \\ &=& \alpha \left(\frac{x+y}{x} \right) U C - \alpha \frac{\lambda T U C}{2x} \\ &=& \alpha U C \left[\left(1 + \frac{y}{x} \right) - \frac{\lambda T}{2x} \right]. \end{split}$$

The differential equation for number of downloaders and seeds is:

$$\frac{dx}{dt} = \lambda - \gamma x = \lambda - \alpha \left(\frac{1}{2}x + y\right) UC,$$

$$\frac{dy}{dt} = \gamma x - \mu y = \alpha \left(\frac{1}{2} x + y \right) UC - \mu y.$$

The number can be obtained:

$$\overline{x} = 2\lambda \left[\frac{1}{\alpha UC} - \frac{1}{\mu} \right] \qquad \overline{y} = \frac{\lambda}{\mu} \qquad T = \frac{\overline{x}}{\lambda} = 2 \left[\frac{1}{\alpha UC} - \frac{1}{\mu} \right]$$

Insights:

- The average download time almost doubles compared to Rarest-First. This is a large price to pay for the benefit of In-Order retrieval.
- The number of downloaders in steady-state almost doubles compared to Rarest-First. The number of seeds remains the same as in Rarest-First.
- Unlike the Rarest-First policy, the total swarm population depends on the seed residence time.

- Strict In-Order Piece Selection (FCFS)
 - A variant of the foregoing In-Order piece selection model: The requests for downloading are queued until they are serviced. The uploading peers don't purge the unfilled requests after each uploading round.
 - Finite queue + FCFS service model at uploading.
 - "Daisy-chain effect" is possible: peers of age t download their needed pieces from peers of age t+E, and those peers in turn download their needed pieces from peers of age t+2E, and so on.

- The differential equation for number of downloaders and seeds in this strategy:
 - Regardless of the actual peer relationships formed, all peers are busy uploading, therefor, the system efficiency is close to 1.
 - The steady-state swarm population and average download time is the same with the rarest-first model:

$$\overline{x} = \lambda \left[\frac{1}{UC} - \frac{1}{\mu} \right] \qquad \overline{y} = \frac{\lambda}{\mu} \qquad T = \frac{1}{UC} - \frac{1}{\mu}.$$

Insights:

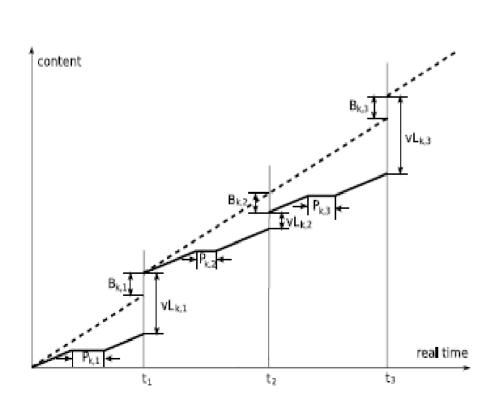
- In-Order (FCFS) piece selection policy may achieve near optimal system efficiency and achieve download time identical to that of the rarest-first policy
- In-Order (FCFS) provide ideal sequential progress at the peers.

□ Assumptions:

- Ignore the peer overlay structure and piece selection strategy. Assume download rate is larger than video rate, D>v.
- Peers adopt sequential prefetching.

- Characterizing Seek and Pause
 - Seek distance (L): the amount of video time that a peer skips during a seek
 - Inter-seek distance (I): the amount of video time between two consecutive seeks.
 - Pauses: Type A: pauses that occur between two seeks; Type B: pauses that occur after the last seek.
 - All pauses that occur between two successive seeks into one aggregated pause

Calculate the downloaded content under user interactions:



$$b_s = \lambda(z - \alpha b_{up}t_s)$$

$$z = vl_v - \bar{N}\bar{G}$$

$$t_s = l_v - \bar{N}\bar{L} + \bar{N}\bar{P} + \bar{Q} + \xi c$$

$$G_{k,i} = -min(B_{k,i}, 0)$$

$$\begin{split} B_{k,1} &= (D_k - v)t_1 - vL_{k,1} + vP_{k,1} \\ B_{k,2} &= \begin{cases} (t_2 - t_1)(D_k - v) - vL_{k,2} + vP_{k,2} & \text{if } B_{k,1} < 0 \\ (t_2 - t_1)(D_k - v) - vL_{k,2} + vP_{k,2} + B_{k,1} & \text{if } B_{k,1} \ge 0 \end{cases} \\ &= (D_k - v)t_2 - \min(B_{k,1}, 0) - v\sum_{j=1}^2 (L_{k,j} - P_{k,j}) \\ B_{k,3} &= (D_k - v)t_3 - \sum_{j=1}^2 \min(B_{k,j}, 0) - v\sum_{j=1}^3 (L_{k,j} - P_{k,j}) \end{split}$$

By generalization, we can get

$$B_{k,i} = (D_k - v)t_i - \sum_{j=1}^{i-1} \min(B_{k,j}, 0) - v \sum_{j=1}^{i} (L_{k,j} - P_{k,j})$$

With

$$t_i = \sum_{j=1}^{i} (I_{k,j} + P_{k,j})$$

$$G_{k,i} = -min(B_{k,i}, 0)$$

□ The generalized equation can be written as:

$$B_{k,i} = (D_k - v) \sum_{j=1}^i I_{k,j} + \sum_{j=1}^{i-1} G_{k,j} - v \sum_{j=1}^i L_{k,j} + D_k \sum_{j=1}^i P_{k,j}$$

Extend the Equation to a large number of peers:

$$\begin{split} \sum_{k=1}^{n_p} B_{k,N_k+1} &= \sum_{k=1}^{n_p} \sum_{j=1}^{N_k+1} (D_k - v) I_{k,j} + \sum_{k=1}^{n_p} \sum_{j=1}^{N_k} G_{k,j} \\ &- v \sum_{k=1}^{n_p} \sum_{j=1}^{N_k+1} L_{k,j} + \sum_{k=1}^{n_p} \sum_{j=1}^{N_k+1} D_k P_{k,j} \end{split}$$

Simplify it by taking mean values:

$$\sum_{k=1}^{n_p} B_{k,N_k+1} = (n_s + n_p)(\bar{D} - v)\bar{I} + n_s\bar{G}$$
$$- (n_s + n_p)v\bar{L} + (n_s + n_p)\bar{D}\bar{P}$$

■ We can get

$$\begin{split} \bar{G} &= (1 + \frac{n_p}{n_s})(v\bar{L} - (\bar{D} - v)\bar{I} - \bar{D}\bar{P}) + \frac{n_p}{n_s}\bar{B}_{\bar{N}+1} \\ &= v\bar{L} - (\bar{D} - v)\bar{I} - \bar{D}\bar{P} + \delta \end{split}$$

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$$\delta = \frac{1}{\bar{N}}(\bar{B}_{\bar{N}+1} - \bar{B}_1)$$

THANK YOU