

# Efficiency and Nash Equilibria in a Scrip System for P2P Networks

Eric J. Friedman

Joseph Y. Halpern

Ian Kash

# The Free Rider Problem

Almost 70% of Gnutella users share no files and nearly 50% of responses are from the top 1% of sharing hosts

- Adar and Huberman '00

# Approaches Taken

- Reputation Systems (Xiong and Liu '02, Eigentrust: Kamvar et al. '03, Gupta et al. '03, Guha et al. 04, ...)
  - Assume “good” vs “bad” agents
- Barter-like Systems (BitTorrent, Anagnostakis and Greenwald '04, ...)
  - No analogue of money – can only do swaps
  - Very good for popular files, not so good for rarer ones
- Scrip Systems (Karma: Vishnumurthy et al. '03)
  - How to find the right amount of money
  - How to deal with new players

# Outline

- **Model**
- Non-Strategic Play
- Strategic Play
- Designing a System

# Model

- $n$  agents
- In round  $t$ , agent  $\mathbf{p}_t$  is chosen at random to make a request
- With probability  $\beta$ , an agent can satisfy the request, then decides whether to volunteer
- One volunteer  $\mathbf{v}_t$  is chosen at random to satisfy the request
- For round  $t$ ,  $\mathbf{p}_t$  gets a payoff of  $\mathbf{1}$  (if someone volunteered),  $\mathbf{v}_t$  get a payoff of  $-\alpha$  (a small cost)
- Total utility for a player is the discounted sum of round payoffs ( $\mathbf{U}_t$ ):

$$\sum_{t=0}^{\infty} \delta^t u_t$$

# Adding Scrip

- M dollars in the system
- Requests cost 1 dollar
- How to decide when I should satisfy a request?

# Why Do I Want to Satisfy, or not?



# Threshold Strategies

**$S_k$** : Volunteer if the requester can pay me and I have less than  $k$  dollars

**$k$**  is your “comfort level,” how much you want to have saved up for future requests

**$S_0$**  corresponds to never volunteering



# Main Results

- This game has a Nash equilibrium in which all agents play threshold strategies
- The optimal amount of money  $M$  (to maximize agent utilities) is  $mn$  for some average amount of money  $m$ .

# Outline

- Model
- **Non-Strategic Play**
- Strategic Play
- Designing a System

What happens if everyone uses the same strategy  $S_k$ ?

- Models the system as a Markov chain
- States are vectors of the amount of money each agent has
- Transition probabilities are the probabilities of the relevant agents being chosen as requester and volunteer respectively

# Theorem: Concentration Phenomenon

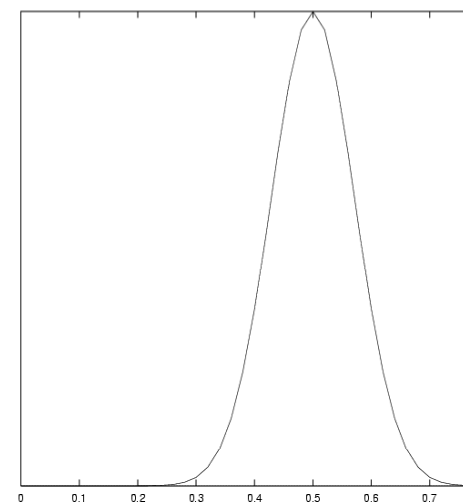
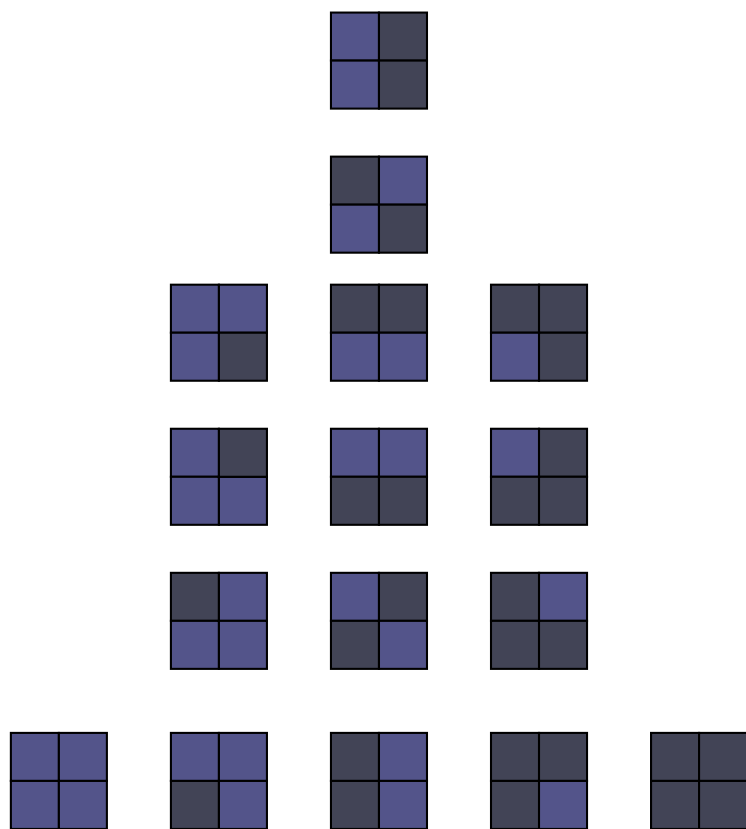
*Remarkable property of this Markov Chain:* For any amount of money, the number of people with that amount of money is essentially constant.

More formally, there exists a distribution:

$$(d(0), \dots, d(k)) \in \Delta^{k+1}$$

(dependent on the parameters of the game) such that the fraction of people in the system with  $k$  dollars is within  $\varepsilon$  of  $d(k)$  with high probability.

# Maximum Entropy



# Maximum Entropy

We can express our constraints as:

$$1 = \sum_{j=0}^k d(j) \qquad M = \sum_{j=0}^k n \cdot d(j) j$$

There is a **unique** solution that maximizes entropy among all solutions that satisfy the constraints. This solution characterizes the number of people with each amount of money.

The system is **stable**: with high probability it will be in a state close to this distribution.

The system **stays away from bad states**: entropy is maximized when money is divided “as evenly as possible” subject to the constraints.

# Outline

- Model
- Non-Strategic Play
- **Strategic Play**
- Designing a System

# Equilibria

Trivially,  $S_0$  is always an equilibrium.

We would like to say:

“There is **always** a Nash equilibrium where all agents play  $S_k$  for some  $k > 0$ .”

Two Problems:

- If  $\delta$  is small, agents will be unwilling to do any work now for future payoff
- If  $n$  is small, the theorems do not apply (because each transaction causes a big shift, the system could get quite far away from the maximum entropy distribution)



**THEOREM 4.1.** *Fix a strategy  $S_\gamma$  and an agent  $i$ . There exists  $\delta^* < 1$  and  $n^*$  such that if  $\delta > \delta^*$ ,  $n > n^*$ , and every agent other than  $i$  is playing  $S_\gamma$  in game  $G(n, \delta)$ , then there is an integer  $k'$  such that the best response for agent  $i$  is  $S_{k'}$ . Either  $k'$  is unique (that is, there is a unique best response that is also a threshold strategy), or there exists an integer  $k''$  such that  $S_{\gamma'}$  is a best response for agent  $i$  for all  $\gamma'$  in the interval  $[k'', k'' + 1]$  (and these are the only best responses among threshold strategies).*

There exist  $\delta^* < 1$  and  $n^*$  such that for all systems with discount factor  $\delta > \delta^*$  and  $n > n^*$  agents, if every agent but  $i$  is playing  $S_k$  then there exists a best response for  $i$  of the form  $S_{k'}$ ,

## Proof Sketch

If  $n$  is large enough, the actions of a single agent have essentially no impact on the system. Thus our calculations before make the problem a **Markov Decision Problem** (MDP).

## Theorem: Existence of Equilibria

For all  $M$ , there exist  $\delta^* < 1$  and  $n^*$  such that for all systems with discount factor  $\delta > \delta^*$  and  $n > n^*$  agents, there is a Nash equilibrium where all agents play  $S_k$  for some  $k > 0$ .

Experiments show that the system converges quickly to equilibrium when agents play equilibrium strategies.

# Outline

- Model
- Non-Strategic Play
- Strategic Play
- **Designing a System**

## Choosing The Price of a Job

- The maximum entropy distribution only depends on the average amount of money  $\mathbf{m} = \mathbf{M} / \mathbf{n}$  and the strategy  $\mathbf{k}$ .
- The equilibrium solution depends only on the parameters of the game and the maximum entropy distribution.
- Therefore, the optimal value of  $\mathbf{M}$  (to maximize expected utility of the agents) is  $\mathbf{mn}$  for some average amount of money  $m$ .

# Handling New Players

- Let new players join with 0 dollars
- Adjust the amount of money to maintain the optimal ratio  $M/n$

# Recap

- The long run behavior of a simple scrip system is stable.
- There is a nontrivial Nash equilibrium where all agents play a threshold strategy.
- There is an optimal average amount of money, which leads to maximum total utility. It depends only on the discount factor  $\delta$ .
- The system is scalable; we can deal with new agents by adjusting the price of a job.

# Open Problems

- Can we characterize the set of equilibria and the conditions under which there is a unique nontrivial equilibrium?
- Can we find an analytic solution for the best response function and the optimal pricing?
- What are the effects of sybils and collusion?