

# A Tale of Two Metrics: Simultaneous Bounds on Competitiveness and Regret

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- Realistic scenario is rare
- Idea is inspiring

# Two communities

- Online convex optimization (OCO)  $\sum_{t=1}^T c^t(x^t)$
- Metric task system (MTS)

$$\sum_{t=1}^T (c^t(s^t) + d(s^t, s^{t-1}))$$

A constant number of states, a cost  $c^t$  for each given state, a fixed switching cost  $d$  for given  $s^t, s^{t-1}$ . At each  $t$ ,  $c^t$  is revealed. Then we choose a state  $s^t$  for this time  $t$ .

# Example:dynamic datacenter right sizing

$$\text{minimize } \sum_{t=1}^T \sum_{i=1}^{x^t} f(\lambda_i^t) + \beta \sum_{t=1}^T (x^t - x^{t-1})^+$$

$$\text{subject to: } 0 \leq \lambda_i^t \leq 1 \text{ and } \sum_{i=1}^{x_t} \lambda_i^t = \lambda_t$$

- $f(\cdot)$  is convex,  $x^t, \lambda_i^t$  are fractional variables.
- $\beta$  is fixed. At each  $t$ ,  $\lambda^t$  arrives first. Algorithm makes decisions then.

# Example:dynamic datacenter right sizing

$$\text{minimize } \sum_{t=1}^T \underline{x^t f(\lambda^t / x^t)} + \beta \sum_{t=1}^T (x^t - x^{t-1})^+$$

subject to:  $x^t \geq \lambda^t$

- $f(\cdot)$  is convex,  $x^t, \lambda_i^t$  are fractional variables.
- $\beta$  is fixed. At each  $t$ ,  $\lambda^t$  arrives first. Algorithm makes decisions then.

# Smoothed online convex optimization (SOCO)

- OCO + switching cost

At each  $t$ , algorithm chooses  $x^t$ , then experiences loss function  $c^t(x^t) + ||x^t - x^{t-1}||$ .  $c^t(x^t)$  is convex.

How to connect the two problems

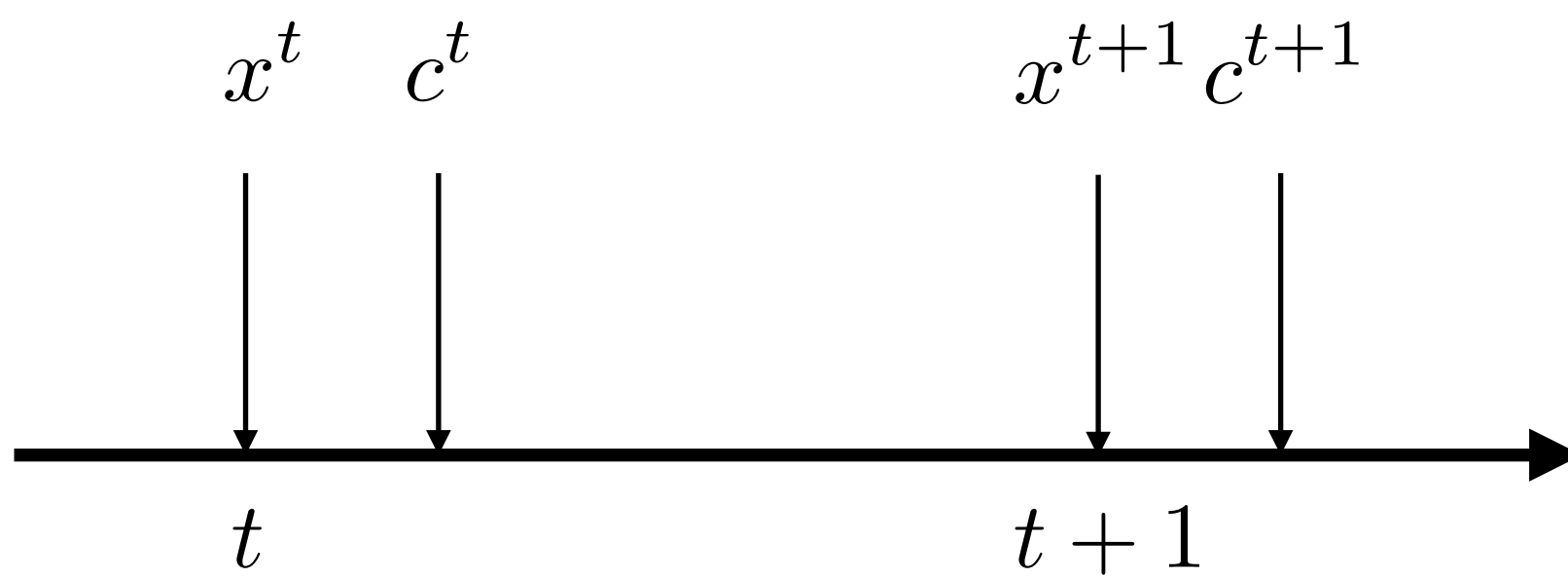
# Problem formulation

$$C_i^\alpha(A, T) = \sum_{t=1}^T c^t(x^{t+i}) + \alpha ||x^{t+i} - x^{t+i-1}||$$

- $\alpha$ -penalized cost with lookahead  $i$
- Euclidean norms of subgradients of  $c^t$  are bounded

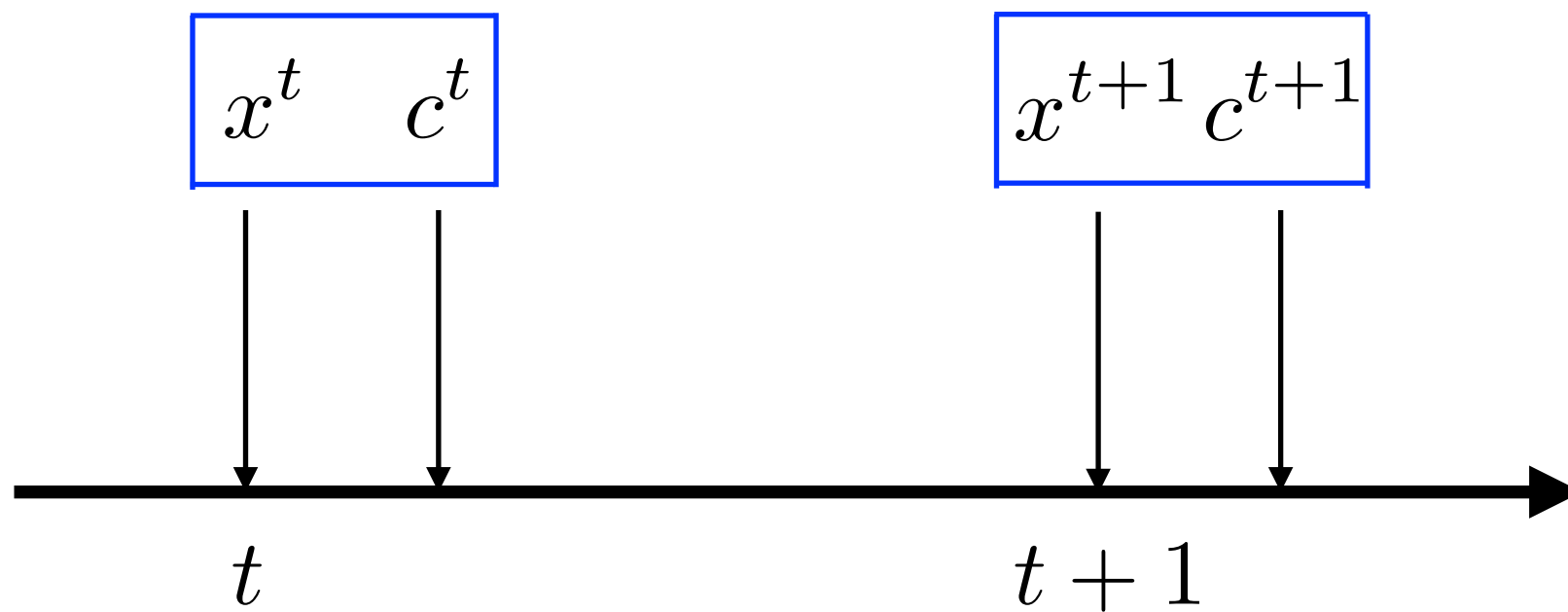
$$x^i = 0$$





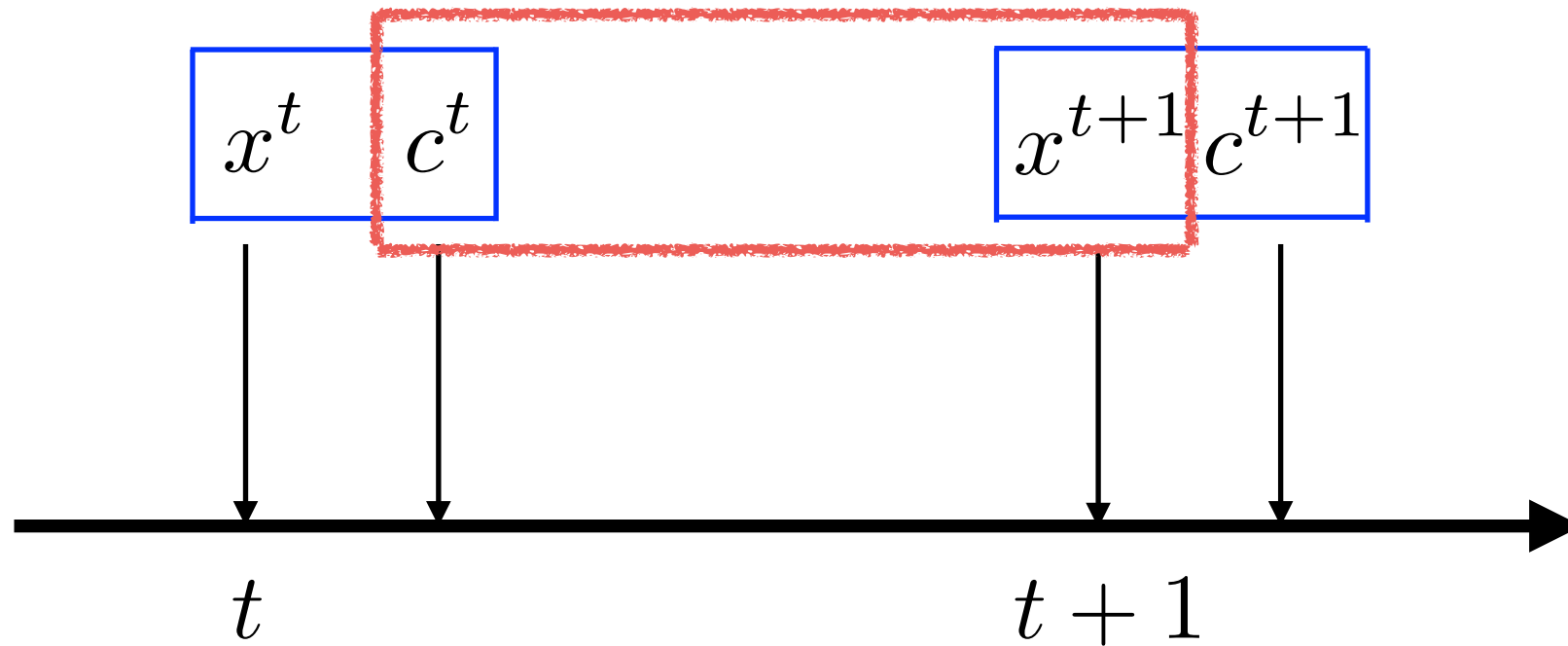
$$x^0 = 0$$

$$i=0: \quad C_0^\alpha(A, T) = \sum_{t=1}^T c^t(x^t) + \alpha ||x^t - x^{t-1}||$$



$$x^0 = 0$$

$$i=0: \quad C_0^\alpha(A, T) = \sum_{t=1}^T c^t(x^t) + \alpha ||x^t - x^{t-1}||$$



$$i=1: \quad C_1^\alpha(A, T) = \sum_{t=1}^T c^t(x^{t+1}) + \alpha ||x^{t+1} - x^t||$$

$$x^0 = 0 \quad x^1 = 0$$

# Traditional performance metrics

$$C_i^\alpha(A, T) = \sum_{t=1}^T c^t(x^{t+i}) + \alpha ||x^{t+i} - x^{t+i-1}||$$

- OCO:  $R_0^0$
- MTS:  $CR_1^1$

# Traditional performance metrics

- In OCO literature:  $OPT_s = \min_{x \in \mathcal{F}} \sum_{t=1}^T c^t(x)$

$$R_0^0(A, T) = \max_{\vec{c}} (C_0^0(A) - OPT_s)$$

- In MTS literature:

$$OPT_d = \min_{x \in \mathcal{F}^T} \sum_{t=1}^T c^t(x^{t+1}) + ||x^{t+1} - x^t||$$

$$CR_1^1(A, T) = \max_{\vec{c}} (C_1^1(A) / OPT_s)$$

# Traditional algorithms

- Online context optimization (OCO)  $\sum_{t=1}^T c^t(x^t)$

Regret sub-linear with  $T$

- Metric task system (MTS)  $\sum_{t=1}^T (c^t(s^t) + d(s^t, s^{t-1}))$

Competitive ratio is independent with  $T$

# In the unified problem

$$C_i^\alpha(A, T) = \sum_{t=1}^T c^t(x^{t+i}) + \alpha ||x^{t+i} - x^{t+i-1}||$$

- $i$  connects online algorithm and online learning
- $\alpha$  connects dynamic offline optimal and static offline optimal

# Contributions

- Connecting OCO and (convex)MTS into one general problem
- Finding the incompatibility of regret and competitive ratio
- Designing an algorithm trading off the performance between the two metrics



# Traditional algorithm

Online Gradient Decent (OGD)

- works for OCO
- works for SOCO (contribution of this paper)
- has unbounded competitive ratio for (convex)MTS

- Is there fundamental incompatibility between these two problems?

# With window i

- In OCO literature:  $OPT_s = \min_{x \in \mathcal{F}} \sum_{t=1}^T c^t(x)$

$$R_i^0(A) = \max_{\vec{c}} (C_i^0(A) - OPT_s)$$

- In MTS literature:

$$OPT_d = \min_{x \in \mathcal{F}^T} \sum_{t=1}^T c^t(x^{t+i}) + \alpha ||x^{t+i} - x^{t+i-1}||$$

$$CR_{i+1}^\alpha(A) = \max_{\vec{c}} (C_{i+1}^\alpha(A) / OPT_s)$$

# Incompatibility of regret and competitive ratio

$$R_i^0(A) = \max_{\vec{c}}(C_i^0(A) - OPT_s)$$

$$CR_{i+1}^\alpha(A) = \max_{\vec{c}}(C_{i+1}^\alpha(A)/OPT_s)$$

- For any algorithm  $A$  ,  $\alpha \geq 1$  ,  $\gamma \geq 0$  , either  $R_i^0(A) = O(T)$  or  $CR_{i+1}^\alpha(A) \geq \gamma$

Both with switching cost,  
Both without lookahead window

$$R_0^1(A) = \max_{\vec{c}}(C_0^1(A) - OPT_s)$$

$$CR_0^\alpha(A) = \max_{\vec{c}}(C_0^\alpha(A)/OPT_s)$$

- For any algorithm  $A$ ,  $\alpha \geq 1$ ,  $\gamma \geq 0$ , either  $R_0^1(A, T) = O(T)$  or  $CR_0^\alpha(A, T) \geq \lambda$

- Regret =  $O(T^{1-\epsilon})$ , Competitive ratio =  $O(T^\epsilon)$

$$\epsilon \rightarrow 0$$

# A unified algorithm

## **Algorithm 2 (Randomly Biased Greedy, RBG( $N$ ))**

*Given a norm  $N$ , define  $w^0(x) = N(x)$  for all  $x$  and  $w^t(x) = \min_y \{w^{t-1}(y) + c^t(y) + N(x - y)\}$ . Generate a random number  $r$  uniformly in  $(-1, 1)$ . For each time step  $t$ , go to the state  $x^t$  which minimizes  $Y^t(x^t) = w^{t-1}(x^t) + rN(x^t)$ .*

**Theorem 7** *For a SOCO problem in a one-dimensional normed space  $\|\cdot\|$ , running RBG( $N$ ) with a one-dimensional norm having  $N(1) = \theta\|1\|$  as input (where  $\theta \geq 1$ ) attains an  $\alpha$ -unfair competitive ratio  $CR_1^\alpha$  of  $(1 + \theta)/\min\{\theta, \alpha\}$  and a regret  $R'_0$  of  $O(\max\{T/\theta, \theta\})$ .*

# A unified algorithm

## Algorithm 2 (Randomly Biased Greedy, RBG( $N$ ))

Given a norm  $N$ , define  $w^0(x) = N(x)$  for all  $x$  and  $w^t(x) = \min_y \{w^{t-1}(y) + c^t(y) + N(x - y)\}$ . Generate a random number  $r$  uniformly in  $(-1, 1)$ . For each time step  $t$ , go to the state  $x^t$  which minimizes  $Y^t(x^t) = w^{t-1}(x^t) + rN(x^t)$ .

$$\text{MTS: } s^t = \operatorname{argmin}_x w^t(x) + \|x - s^{t-1}\|$$

**Theorem 7** For a SOCO problem in a one-dimensional normed space  $\|\cdot\|$ , running RBG( $N$ ) with a one-dimensional norm having  $N(1) = \theta\|1\|$  as input (where  $\theta \geq 1$ ) attains an  $\alpha$ -unfair competitive ratio  $CR_1^\alpha$  of  $(1 + \theta) / \min\{\theta, \alpha\}$  and a regret  $R'_0$  of  $O(\max\{T/\theta, \theta\})$ .

**RBG( $\|\bullet\|$ ) is 2-competitive, has linear regret**

**RBG( $N$ ) encourages its actions to change less**



# Drawbacks of RBG(N)

- Metrics are still unfair
- Can't guarantee two kinds of performance in one problem setting  
 $c^t(x^t), c^t(x^{t+1})$
- How to simultaneously guarantee good  $CR_0^\alpha(A)$  and  $R_0^\alpha(A)$ , or competitive difference?

# Lessons learned

- Negative results are wonderful
- New metrics matter if meaningful. ( $\alpha$ -unfair competitive ratio, competitive difference)
- Classical algorithms (or variants) may have good performance in other metrics (WFA)

# Open questions

- What if input is i.i.d.?
- Does their algorithm work well in the combinatorial version of MTS?