# Analysis of Adaptive Streaming for CDN/P2P Live Video Systems

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#### Introduction

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 Highlight: Build the model with existing models as building blocks.

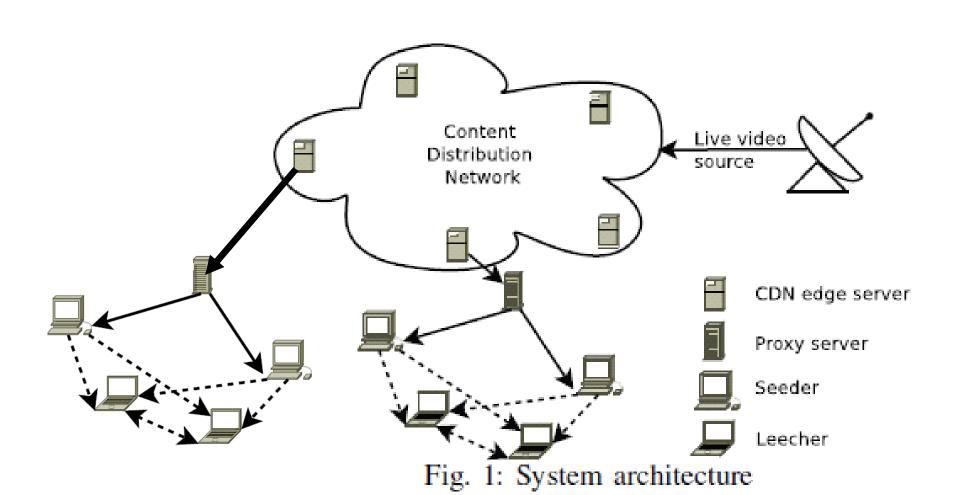
### Significance

- Existing researches have explored
  - Adaptive streaming systems (mainly layered streaming)
  - Hybrid CND/P2P streaming systems
- respectively, but not yet explored the combination of the two.

#### Definition

- Adaptive streaming (HTTP v.s. layered):
  - A video can be streamed at multiple qualities
  - Users can switch among different qualities
  - Download rate is assumed to be variable among clients but constant over time.
- Hybrid CDN/P2P
  - Users may receive data either from the server or other peers streaming the same quality of the same video

### System Architecture



#### Questions

- How to find a way to switch the operation of the system between the CDN and P2P modes
- How to find the best bitrate adaption strategy
- Is a hybrid adaptive system better than a classic CDN adaptive system? How much better will it be? (can not answer without a quantitative model)

# Approach: add dimensions one by one

- Hybrid system model with single rate
  - Unconstrained churnless system (fluid model)
  - Unconstrained system with churn (M/G/∞ queueing model)
  - Constrained churnless system
  - Constrained system with churn
- Adaptive hybrid system model
  - ..... (similar as the previous part)
- CDN adaptive model

# Single rate, unconstrained, churnless

- Relationship between
  - Streaming rate (or capacity of proxy server)
  - # of seeders and leechers
  - Bw. of seeders and of leechers

$$r \le \frac{n_s u_s + n_l u_l}{n_l}$$

# Single rate, unconstrained, with churn

- The key is to find the probablity in terms of peer population
- M/G/∞

$$P(\text{support bitrate } r) = P(r \le \frac{n_s u_s + N u_l}{N})$$

$$= P(N \le \frac{n_s u_s}{r - u_l}) = F(\frac{n_s u_s}{r - u_l})$$

 Relationship between streaming rate, # of seeders and leechers, bw. of seeders and of leechers

$$n_s \ge \frac{(\phi_{1-\alpha}\sqrt{\rho} + \rho)(r - u_l)}{u_s}$$

• Approximation:

$$\sqrt{\hat{\rho}} = a + b\hat{\rho}$$

### Single rate, constrained, churnless

Avg. download bw. of leechers

$$d = \sum_{x} E[d|\text{leecher is connected to } x \text{ seeders}] \times Pr\{x\}$$

$$= \sum_{x} \left(\frac{xu_s}{S_{in}} + \frac{(Y_{out} - x)\eta u_l}{Y_{in}}\right) \times Pr\{x\}$$

$$= \frac{Y_{out}\eta u_l}{Y_{in}} + \left(\frac{u_s}{S_{in}} - \frac{\eta u_l}{Y_{in}}\right) \sum_{x} xPr\{x\}$$

$$Y_{out} = (n_s S_{in} + n_l Y_{in})/n_l$$

$$(4)$$

• Therefore,

$$d = \frac{n_s u_s + \eta n_l u_l}{n_l}$$

Compare with the results for churnless system

$$r \le \frac{n_s u_s + n_l u_l}{n_l}$$

### Adaptive, unconstrained, churnless

- Proxy can offer multiple bitrates.
- Upon arrival, a client requests a bitrate and then the proxy assigns it a bitrate according to the optimization results.
  - The difference between the two is called "client dissatisfaction"
- The objective is to minimize total client dissatisfaction over all clients.

### Adaptive, unconstrained, churnless

Mininze peers' dissatisfaction under capacity constraints

$$\min \sum_{i=1}^{R} \sum_{j=i}^{R} x_{ij} n_{l_i} (r_i - r_j)$$
 (7)

subject to: 
$$\sum_{j=i}^{R} x_{ij} = 1$$
,  $0 \le x_{ij} \le 1$  for  $i = 1, \ldots, R$ 

$$n_{s_i} u_s \ge \left(n_{l_i} x_{ii} + \sum_{k=1}^{i-1} n_{l_k} x_{ki} - n_{s_i}\right) (r_i - u_l)$$
 (8)

$$r \le \frac{n_s u_s + n_l u_l}{n_l} \qquad \sum_{i=1}^R n_{s_i} r_i \le C_{proxy} \tag{9}$$

#### Observation:

 Minimizing client dissatisfaction is equivalent to miximizing inter-client fairness

$$\sum_{i=1}^{R} \sum_{j=i}^{R} x_{ij} n_{l_i} r_i (1 - \frac{r_j}{r_i}) = \sum_{i=1}^{R} n_{l_i} r_i - \sum_{i=1}^{R} r_i \left( \sum_{j=i}^{R} n_{l_i} x_{ij} \frac{r_j}{r_i} \right)$$

### Feasibility

 Only if the system can support the lowest bitrate for all clients,

$$\frac{C_{proxy}}{r_R} \ge \frac{r_R - u_l}{u_s} \sum_{i=1}^R n_{l_i}$$

# Adaptive, unconstrained, with churn

 Relationship between streaming rate, # of seeders and leechers, bw. of seedrs and leechers

$$n_{s_i}u_s \ge (\phi_{1-\alpha/2}\sqrt{\hat{\rho}_i} + \hat{\rho}_i)(r_i - \eta u_l - \epsilon)$$

where  $\hat{\rho}$  is defined as in the previous section.

### **CDN** adaptive

 Mininze peers' dissatisfaction under capacity constraints.

$$\min \sum_{i=1}^{R} \sum_{j=i}^{R} x_{ij} n_i (r_i - r_j)$$

lacktriangle

subject to: 
$$\sum_{j=i}^{R} x_{ij} = 1$$
,  $0 \le x_{ij} \le 1$  for  $i = 1, \dots, R$ 

$$\sum_{i=1}^{R} r_i \left( n_i x_{ii} + \sum_{k=1}^{i-1} n_k x_{ki} \right) \le C_e \tag{12}$$

#### Validation

- Validate single-rate model
  - compare results gained from fluid model and simulation. The accuracy of the whole model relies largely on the fluid model
- Validate the CDN adaptive model
- Fix bandwidth capacity and compare QoS
- Fix QoS and compare bandwidth savings

# Numerical results of single-rate model

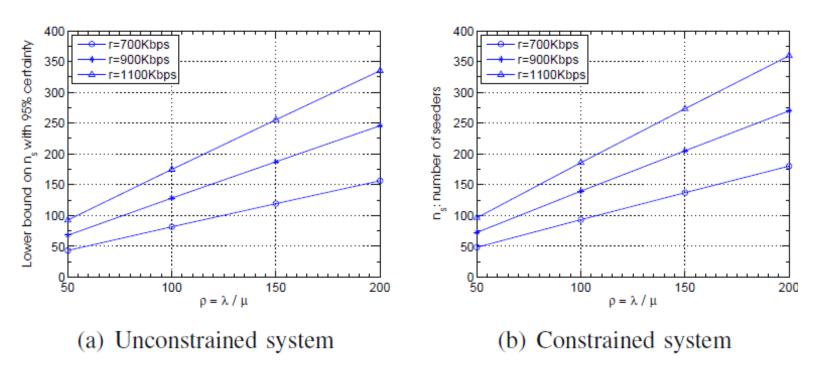


Fig. 2:  $n_s$  vs  $\rho$  for different video bitrates for systems with churn,  $\alpha = 0.05$ 

### Validate single-rate model

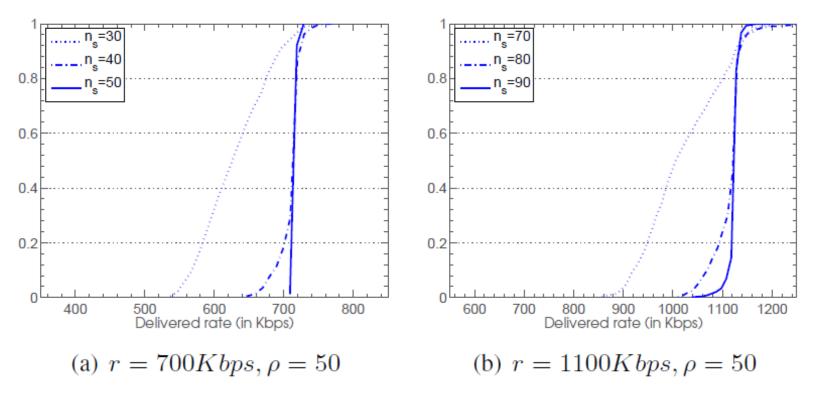


Fig. 3: CDF of average delivered rate for unconstrained system with churn

### Validate single-rate model

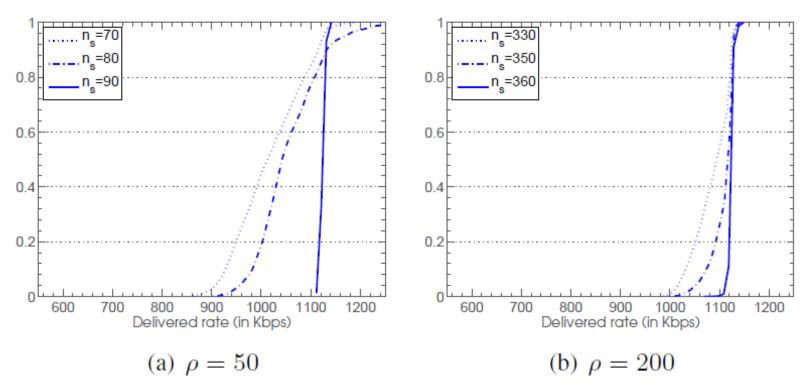


Fig. 4: CDF of average delivered rate for constrained system with churn, r = 1100Kbps

### Validate the CDN adaptive model

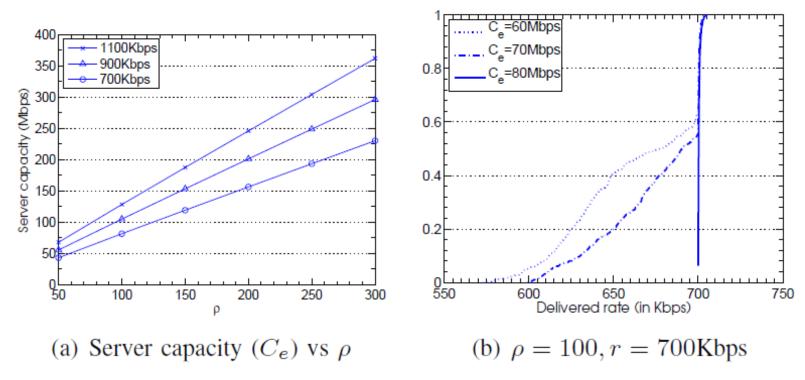
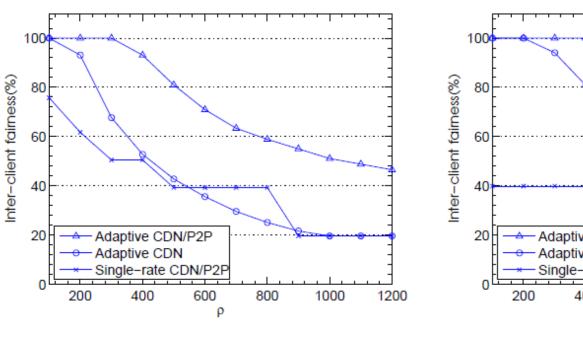
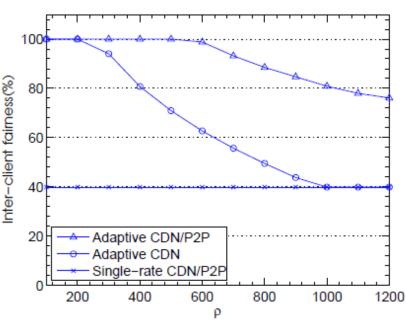


Fig. 5: CDN system with churn

# Fix bandwidth capacity and compare QoS

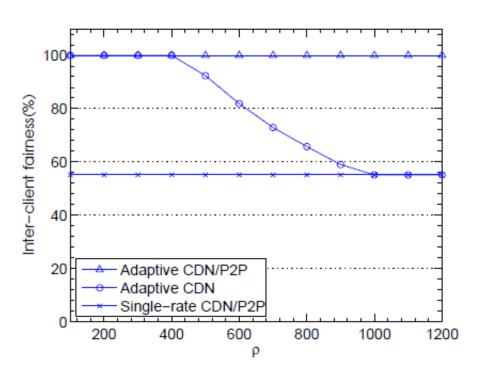


(a) High bitrate bias



(b) Uniform distribution over bitrates

## Fix bandwidth capacity and compare QoS



(c) Low bitrate bias

# Fix QoS and compare bandwidth saving

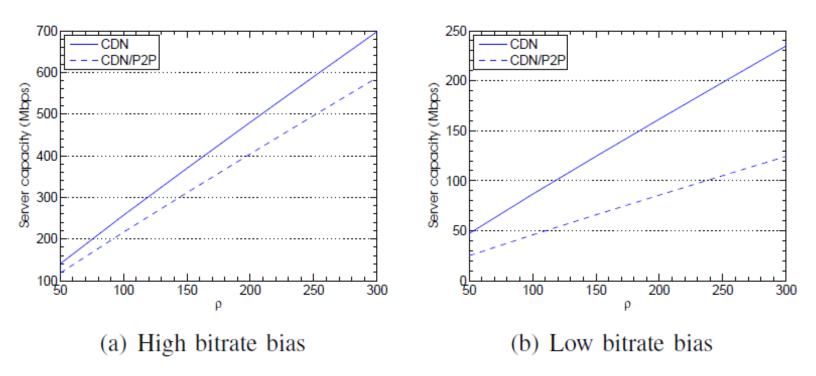


Fig. 7: Required server capacity for CDN/P2P and CDN systems with churn

Q & A

Thank you!