## Modeling of the P2P service migration problem 1

We suppose there are M videos, and N ISPs. There are one on-premise server and one cloud node in each ISP.

Notation definition:

 $C_s^j$ : storage capacity of the on-premise server at the j-th ISP

 $C_u^j$ : upload bandwidth capacity of the on-premise server at the j-th ISP

 $h_j$ : charging rate for storage on the cloud at the j-th ISP

 $k_i$ : charging rate for upload bandwidth on the cloud at the j-th ISP

 $s_m$ : storage of m - th video

 $x_m^j = \{0,1\}, m=1,...,M$ :  $x_m^j = 1$  if the placement of the m-th video is on the on-premise server at the j-th ISP;  $x_m^j=0$  otherwise;

 $y_m^j = \{0,1\}, m=1,...,M$ :  $y_m^j = 1$  if the placement of the m-th video is on the cloud at the j - th ISP;  $y_m^j = 0$  otherwise;

 $r_m^j$ : request rate of the m-th video from the j-th ISP, i.e., the bandwidth demand is  $s_m r_m^j$ .

 $R_{ji}^m$ : percentage of requests from j for video m is routed to on-premise server i  $T_{ji}^m$ : percentage of requests from j for video m is routed to cloud i

## 1.1 Optimization of the problem without Lyapunov optimization

 $min\sum_{m=1}^M\sum_{j=1}^N\sum_{i=1}^N(s_mr_m^jT_{ji}k+s_mh)y_m^j-\alpha\sum_{m=1}^M\sum_{j=1}^Ns_mr_m^j(T_{jj}+R_{jj})$  (maximize local traffic, i.e., minimize delay) subject to:

$$y_m^j = \{0, 1\}, \forall j = 1, ..., N, \forall m = 1, ...M$$
  
 $x_m^j = \{0, 1\}, \forall j = 1, ..., N, \forall m = 1, ...M$ 

$$\begin{aligned} x_m^j &= \{0,1\}, \forall j=1,...,N, \forall m=1,...M \\ \sum_{i=1}^N (R_{ji}^m + T_{ji}^m) &= 1, \forall j=1,...N, \forall m=1,...,M \end{aligned}$$

$$0 \leq R_{ji}^m \leq x_m^i, \forall j=1,..,N, \forall i=1,...,N, \forall m=1,...,N$$

$$0 \leq T_{ji}^{m} \leq y_{m}^{i}, \forall j = 1, ..., N, \forall i = 1, ..., N, \forall m = 1, ..., N$$

$$\sum_{m=1}^{M} s_m x_m^j \le C_s^j, \forall j$$
 (on-premise server's storage constraint)

 $\sum_{m=1}^{M} s_m x_m^j \leq C_s^j, \forall j \text{ (on-premise server's storage constraint)}$   $\sum_{m=1}^{M} \sum_{j=1}^{N} s_m r_m^j R_{ji}^m \leq C_u^i, \forall i=1,...,N \text{ (on-premise server's upload bandwidth}$ constraint)

known values:  $C_s^j$ ,  $C_u^j$ ,  $h_j$ ,  $k_j$ ,  $s_m$ ,  $r_m^j$ optimization variables:  $x_m^j, y_m^j, R_{ji}^m, T_{ii}^m$ 

## Optimization of the problem with Lyapunov optimization

Request from video m from ISP j is modeled as a queue, whose length in time slot t is  $Q_m^j(t)$ .

The queue update is:  $Q_m^j(t+1) = \max[Q_m^j(t) + r_m^j(t) - \sum_{i=1}^N R_{ji}^m - \sum_{i=1}^N T_{ji}^m, 0]$  Different from the previous sub section,  $R_{ji}^m(t)$  and  $T_{ji}^m(t)$  is not a schedule of fraction of arrival rates for all times alots. Note that of arrival rates for all time slots. Now they are schedule of number of requests (integers) for each time slot.

 $D_{ii}^{s}$  is the delay from source j to on premise server i, and  $D_{ii}^{c}$  is the delay from source

$$D_{ji}^s$$
 is the delay from source  $j$  to on premise server  $i$ , and  $D_{ji}^c$  is the delay from source  $j$  to on cloud node  $i$ .

minimize  $k \overline{\sum_{m=1}^M \sum_{j=1}^N \sum_{i=1}^N (s_m T_{ji}(t))} + \alpha \sum_{m=1}^M \sum_{j=1}^N \overline{s_m} \overline{R_{ji}(t)} + \beta h \sum_{m=1}^M \sum_{j=1}^N (s_m y_m^j) + \gamma \sum_{m=1}^M \sum_{j=1}^N (s_m x_m^j) - \rho \sum_{j=1}^N \sum_{i=1}^N \overline{\sum_{m=1}^M s_m} (T_{ji}^m(t) D_{ji}^c + R_{ji}^m(t) D_{ji}^s)$  subject to:

$$y_m^j = \{0, 1\}, \forall j = 1, ..., N, \forall m = 1, ...M$$

$$x_m^j = \{0, 1\}, \forall j = 1, ..., N, \forall m = 1, ...M$$

$$\begin{aligned} y_m^j &= \{0,1\}, \forall j=1,...,N, \forall m=1,...M \\ x_m^j &= \{0,1\}, \forall j=1,...,N, \forall m=1,...M \\ 0 &\leq R_{ji}^m(t) \leq R_{ji}^m(t) x_m^j, \forall j=1,...,N, \forall i=1,...,N, \forall m=1,...,N, \forall t \\ 0 &\leq T_{ji}^m(t) \leq T_{ji}^m(t) y_m^t, \forall j=1,...,N, \forall i=1,...,N, \forall m=1,...,N, \forall t \end{aligned}$$

$$0 \le T_{ji}^{\tilde{m}}(t) \le T_{ji}^{\tilde{m}}(t)y_m^t, \forall j = 1, ..., N, \forall i = 1, ..., N, \forall m = 1, ..., N, \forall t = 1,$$

$$\sum_{m=1}^{M} s_m x_m^j \leq C_s^j, \forall j$$
 (on-premise server's storage constraint)

 $\sum_{m=1}^{M} s_m x_m^j \leq C_s^j, \forall j \text{ (on-premise server's storage constraint)}$   $\sum_{m=1}^{M} \sum_{j=1}^{N} s_m R_{ji}^m(t) \leq C_u^i, \forall i=1,...,N, \forall t \text{ (on-premise server's upload band-premise server)}$ width constraint)

Queues 
$$Q_m^j(t)$$
 is stable,  $\forall m,j,$  i.e.,  $\overline{r_m^j(t)} \leq \overline{\sum_{i=1}^N R_{ji}^m + \sum_{i=1}^N T_{ji}^m}$  Note:

known values:  $C_s^j, C_u^j, h_j, k_j, s_m$ , optimization variables:  $x_m^j, y_m^j, R_{ji}^m(t), T_{ji}^m(t)$