Axiomatic Foundations of Physics: Problems and Ideas

Yuri Ximenes Martins¹

Department of Mathematics, UFMG, Brazil Math-Phys-Cat Group

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https://math-phys.group/~yxmartins

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- The problem of axiomatizing Mathematics was explicitly considered in his *Hilbert's Program* which is about finding a formalist approach to Mathematics through a consistent axiomatic language;
- Similarly, one could consider an Axiomatization Program of Physics as the search for a formalist approach to Physics through a consistent axiomatic language.

Part I: The Program

1. Foundations of Mathematics

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- 2. Formalism

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- 1. what are "mathematical objects"?
- 2. are they "real"? Are they "abstract"?
- 3. what it means to "know" in mathematics?
- 4. when is a mathematical sentence true? In other words, what is "mathematical truth"?
- 5. what is the role of axioms?
- 6. what is their nature?

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Phil. Found. = Ontology + Epistemology + Logic + \dots

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- mathematical epistemology reduces to logic;

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- none requirement on mathematical **ontology**;
- mathematical **epistemology** reduces to an axiomatic logic.

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- mathematical **ontology** reduces to pure intuition;
- mathematical **epistemology** reduces to pure intuition;
- mathematical **logic** reduces to pure intuition.

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Mathematics is fruit of our minds...

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 - Fregue-Hilbert debate;
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- In the Brouwer view, intuitionistic logics are allowed to be axiomatic, but they avoid some axioms like Law of Excluded Middle and Axiom of Choice.
- On the other hand, these axioms are considered fundamental in Hilbert's formalism.

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- intuitionistic logicism: is logicism for intuitionistic logic, i.e., it is the claim that mathematics reduces to intuitionistic logic
- intuitionistic formalism: is formalism for axiomatic intuitionistic logic, i.e., it is the claim that mathematics can be formalized in an axiomatic intuitionistic logic.

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- a class T of terms (or things);
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- a family S of relations in $C \sqcup E$ defining the *structure* of S.

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- Curiously, Bunge does not even define what is a "morphism" between systems, although there is a natural notion of that.

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- ... we need "higher systems", so that "higher (or homotopical) systemics" (see Section 2.2 of the thesis).
- Thus, instead of fixing an *a priori* formalization to "area of study", we will first work *naively*.

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- thus, we can think of \mathcal{O} as a "class of systems" of a specific type.

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Phil. Found. = Ontology + Epistemology + Logic + \dots

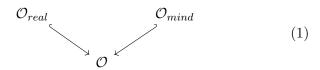
 $\begin{array}{c} \text{Phil. Found.} = \text{Ontology} + \text{Epistemology} + \text{Logic} + \textcolor{red}{\text{Phil.}} \\ \textcolor{blue}{\text{Mind}} \end{array}$

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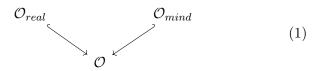
Definition. Let \mathcal{O} be an area of study.

• An ontological theory for \mathcal{O} is a subclass $\mathcal{O}_{real} \subset \mathcal{O}$ of real systems.



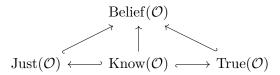
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- A philosophy of mind for \mathcal{O} is a subclass $\mathcal{O}_{mind} \subset \mathcal{O}$ of mental systems.

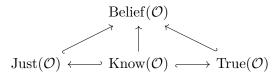


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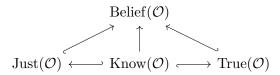
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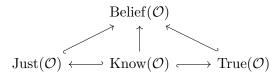
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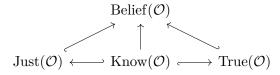
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 - a subclass $Know(\mathcal{O}) \subset True(\mathcal{O}) \cap Just(\mathcal{O})$, describing knowledge.



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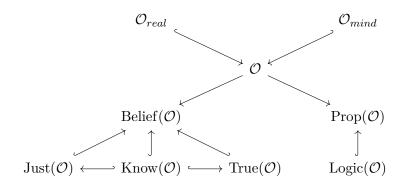


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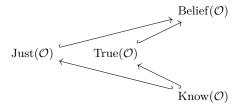
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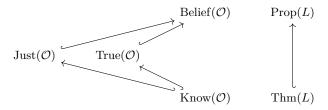
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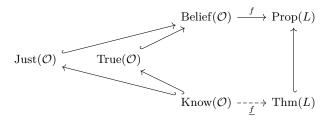
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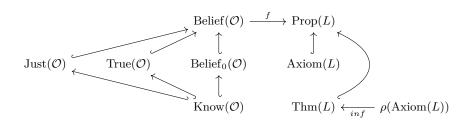
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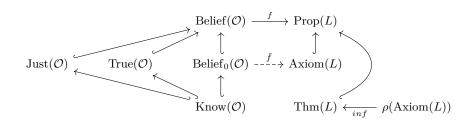
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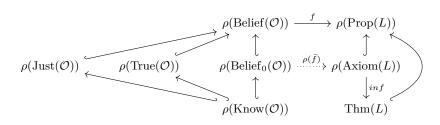
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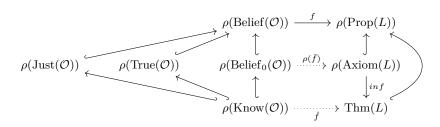
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Hilbert's Program

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Plain

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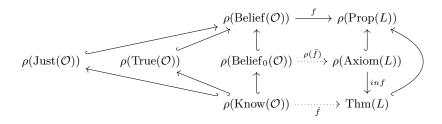
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Thus, given an area of study \mathcal{O} we want to

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We presented three strategies, all of them based on the following additional epistemological requirement.

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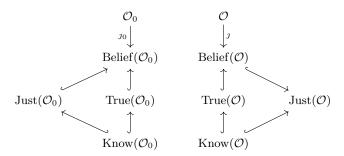
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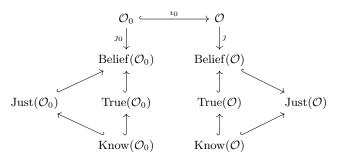
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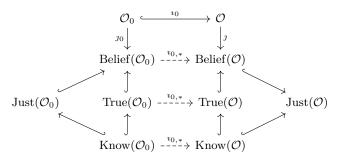
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Strategies

Given an area of study \mathcal{O} we suggested three strategies, corresponding to the following cases:

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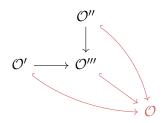
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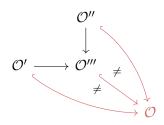
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- 3 when has a span $\mathcal{O}' \to \mathcal{O}''' \leftarrow \mathcal{O}''$ of subclasses $\mathcal{O}', \mathcal{O}'', \mathcal{O}''' \subset \mathcal{O}$.



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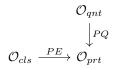
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The strategy of each approach is the following:

1 use of induction on the subclasses \mathcal{O}'_i of $\mathcal{O}' \subset \mathcal{O}$;

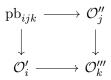
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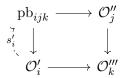
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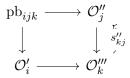
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In the case of $\mathcal{O} = \text{Physics}$, the three strategies above can be interpreted as follows:

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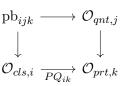
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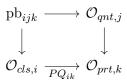
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- ...however it tells nothing about philosophical quantization or axiomatic quantization relative to other axiomatic languages.

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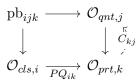
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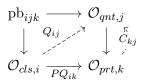
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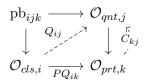
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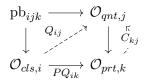
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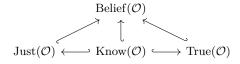
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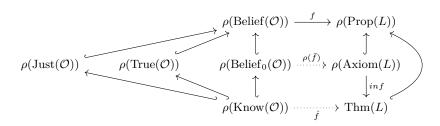
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• Let **H** be a category. An object $X \in \mathcal{H}$ is *injective* if every monomorphism $f: X \to Y$ is a split monomorphism.

Theorem.

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 - $\operatorname{Ep}(\mathcal{O}_0) \in \mathbf{H}$ is an injective object;
 - the pair $(\text{Ep}(\mathcal{O}_0), L_0)$ is a a formalist approach \mathcal{O}_0 internal to \mathbf{H} ,

- Let \mathcal{O} be an area of study internal to a category \mathbf{H} with finite limits and let \mathcal{O}' be a distinguished subclass \mathcal{O}' . Suppose given:
 - an axiomatic language $L_0 \in \mathbf{Axiom_H}$ internal to \mathbf{H} ;
 - a coherent epistemological theory $\mathrm{Ep}(\mathcal{O})$ internal to \mathbf{H} .
- In this case, if there is some $\mathcal{O}_0 \subset \mathcal{O}'$ such that:
 - $\operatorname{Ep}(\mathcal{O}_0) \in \mathbf{H}$ is an injective object;
 - the pair $(\text{Ep}(\mathcal{O}_0), L_0)$ is a a formalist approach \mathcal{O}_0 internal to \mathbf{H} ,
- then every abstraction process internal to \mathbf{H} defines an "universal" axiomatic language L such that $(\mathrm{Ep}(\mathcal{O}',L)$ is a formalist approach for \mathcal{O}' internal to \mathbf{H} .

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• Let \mathcal{O} an area of study internal to ZFC or NBG. If there is some $\mathcal{O}_0 \subset \mathcal{O}$ which admits a formalist approach, then every abstraction process extends this formalist approach from \mathcal{O}_0 to \mathcal{O} .

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- Thus, if every branch of Physics could be internalized in ZFC/NBG, then the choice of an abstraction process would provide a formalist approach to the entire Physics
- But gauge theories, supergravity and string theory are examples of classes of Physical systems in which "homotopies" play a fundamental role...
- ... which forces us to look for languages internal to more abstract categories than ZFC/NBG.

That is all...

A minha vó caiu e quebrou o fêmur Caiu o meu avô e perdeu os dentes.

Espero que a chuva que cai lá fora lave a culpa da ausência e leve embora o desatino que assola minh'alma fina.