

Examination Process Algebra (2IMF10)

28 June 2022, 13:30 – 16:30

The number between parentheses in front of a problem indicates how many points you score with a correct answer to it. A partially correct answer is sometimes awarded with a fraction of those points. The grade for this examination will be determined by dividing the total number of scored points by 2.

1. A word $\sigma \in A^*$ is a trace of a state s in a transition-system space over the set of labels A if, and only if, there exists a state s' such that $s \xrightarrow{\sigma}^* s'$. We denote by $tr(s)$ the set of all traces of s . Two states are trace equivalent (notation: $s =_T t$) if they have the same sets of traces, i.e., if $tr(s) = tr(t)$.

- (3) (a) Give an equation that is valid in $\mathbb{P}(\text{MPT}(A))/_{=_T}$ but not in $\mathbb{P}(\text{MPT}(A))/_{=_\leftrightarrow}$. Support your answer with two proofs: one that the equation is valid in $\mathbb{P}(\text{MPT}(A))/_{=_T}$ and one that the equation is not valid in $\mathbb{P}(\text{MPT}(A))/_{=_\leftrightarrow}$.
- (2) (b) Is the process theory $\text{MPT}(A)$ a ground-complete axiomatisation of $\mathbb{P}(\text{MPT}(A))/_{=_T}$? (Motivate your answer.)

- (1) 2. (a) Prove that $\text{BCP}(A, \emptyset) \vdash x \parallel y = y \parallel x$.

Let E be the recursive specification over $\text{BCP}(A, \emptyset)$ consisting of the following equation:

$$X = 1 + a.(X \parallel b.1) \ .$$

(Throughout this exam, feel free to abbreviate $\mu X.E$ by simply writing X .)

- (4) (b) Define an infinite guarded recursive specification F over $\text{BSP}(A)$ (i.e., not using parallel composition) defining a recursion variable Y_0 such that

$$\text{BCP}_{\text{rec}}(A, \emptyset) + \text{RSP} \vdash \mu X.E = \mu Y_0.F$$

and prove that your answer is correct.

- (3) (c) Does there also exist a finite guarded recursive specification G over $\text{BSP}(A)$ defining a recursion variable Z such that

$$\text{BCP}_{\text{rec}}(A, \emptyset) + \text{RSP} \vdash \mu X.E = \mu Z.G \ ?$$

Support your answer with a proof.

3. Let E be the recursive specification over $\text{TSP}(A)$ consisting of the following equation:

$$X = X \cdot a.1 + 1$$

(Throughout this exam, feel free to abbreviate $\mu X.E$ by simply writing X .)

- (2) (a) Find two distinct closed $\text{TSP}_{\text{rec}}(A)$ -terms p and q such that $\mu X.E \xrightarrow{a} p$ and $\mu X.E \xrightarrow{a} q$, and prove the correctness of your answers by giving formal derivations of the transitions in accordance with the term deduction system (i.e., the structural operational semantics) for $\text{TSP}_{\text{rec}}(A)$.
- (1) (b) Sketch (a representative part of) the transition system associated with $\mu X.E$ by the term deduction system for $\text{TSP}_{\text{rec}}(A)$.

- (2) (c) Define a closed $\text{TSP}_{rec}(A)$ -term r that denotes a solution for E in $\mathbb{P}(\text{TSP}_{rec}(A))/\underline{\leftrightarrow}$ but is not bisimilar to $\mu X.E$ (i.e., $r \not\sim \mu X.E$).
- (2) (d) Let F be the recursive specification consisting of the following equations:

$$\begin{aligned} Y_0 &= \tau.Y_1 + 1 \\ Y_{n+1} &= \tau.Y_{n+2} + \tau.Y_n + a^{n+1}.1 \quad (n \in \mathbb{N}) \end{aligned}$$

(Here $a^n.1$ is inductively defined by $a^0.1 \equiv 1$ and $a^{n+1}.1 \equiv a.a^n.1$ for all $n \in \mathbb{N}$.)

Prove that $\mu Y_0.F \underline{\leftrightarrow}_b \mu X.E$.

- (2) 4. Let A be a set of actions that contains at least three distinct actions a , b and c . Prove that there exists a communication function γ on A and a set $I \subseteq A$ such that

$$\text{TCP}_\tau(A, \gamma) \not\models \tau_I(x \parallel y) = \tau_I(x) \parallel \tau_I(y) .$$