

## CSE213L: DSA Lab

### LNMIIT, Rupa Ki Nangal

### Training Set 09

A code implementing ADT interface for 2-3 trees is provided. The code is large but is not difficult for the best students in a DSA class to understand. To support the reading of the code, ADT implementation has been divided into three files: `tree23Display.c`, `tree23Insert.c` and `tree23Delete.c`. File `tree23.h` lists the interface functions and other declarations.

#### File `tree23.h`

```
#include <stdio.h>
#include <assert.h>
#include <stdlib.h>

struct node23 *toRoot;

struct node23 {
    int data[2];
    struct node23 *subT[3];
    int countKeys;
};
typedef struct node23 *node23p;

int search(int val);
void insert(int val);
void rid(int val);
void display();
```

First please familiarise yourself with the display format used for 2-3 trees. This format is like the nesting and indenting we use to indent and print C programs. A brace-pair encloses information about a 2-3 node and its subtrees. The first line within a brace-pair reveals the number of subtrees of the node. The key values and the subtree details are included in the brace-pair and are indented for a clearer reading.

It is suggested that you refresh your classroom learnings related to the 2-3 trees and B-tree before reading this document any further.

A tree constructed from a test run of the program is provided here. Function `main()` code can be found at the end of this document.

## Sample display of created 2-3 tree (two columns):

```

    }
{ count of keys = 1                                }
<12                                                >12
    { count of keys = 1                                { count of keys = 1
    <7                                                    <30
        { count of keys = 2                                { count of keys = 1
        <3                                                    <20
            { NULL }                                          { NULL }
        <5 && >3                                              >20
            { NULL }                                          { NULL }
        >5                                                    }
            { NULL }                                          >30
    }                                                        { count of keys = 1
    >7                                                    <37
        { count of keys = 2                                { NULL }
        <9                                                    >37
            { NULL }                                          { NULL }
        <10 && >9                                              }
            { NULL }
        >10
            { NULL }
    }
    }

```

## Training Set 09: Task 01

As demonstration of your understanding of the code provided to you, augment the program provide a new ADT interface function `int search(int key)`. To demonstrate the function implemented by you, please change function `main()` as follows:

1. Define two functions: `int nextAction()` and `int nextValue()`. Both these functions use standard C function `rand()` to return a random value. Function `nextAction()` returns int values specified by three macros with probabilities as shown in the parentheses: `SEARCH` (5%), `INSERT` (70%), and `DELETE` (25%). Function `nextValue()` returns a randomly generated integer value between 3001 and 5000.
2. Add an ADT interface function to compute height of the tree. Feel free to include additional member in `struct node23` to track the tree height.
3. The idea is to create a 2-3 tree with about 100 values for demonstration and more importantly as a preparation for the remaining tasks in the training set.

## Training Set 09: Task 02

Alter the data structure declarations to make the 2-3 tree, a B-tree up to 7-way branching. This can be done by making necessary changes to the array sizes in `struct node23`.

Modify implementations of B-tree ADT interface functions related to `search()`, `display()` and `insert()`.

Interface function `rid()` could immediately return without any real change in the B-tree.

What is the height of the tree with following number of entries in the B-tree?

1. 100
2. 1000
3. 10000
4. 20000

## Training Set 09: Task 03

It should already be obvious to the trainee that this task requires them to complete the implementation of function `rid()`.

File `main.c` defines function `main()`:

```
#include "tree23.h"

int main(void) {
    insert(10); insert(5); insert(7); insert(12);
    insert(3); insert(20); insert(9); insert(30);
    insert(54); insert(36); insert(63);
    insert(72); insert(90); insert(81); insert(37);
    rid(37); rid(36);
    insert(18); insert(27); insert(45);
    insert(39); insert(99); insert(42);
    rid(63);
    insert(37); insert(47);
    rid(99); rid(54); rid(42); rid(27);
    rid(81);
    rid(72); rid(39); rid(45);
    rid(90);
    rid(47); rid(18);
    display();

    return 0;
}
```

The code from files implementing 2-3 tree is placed in the appendices to this training document.

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# Appendices

## File tree23Display.c

```
#include "tree23.h"

static void display2(node23p node, int offset);
static void printSpaces(int offset);

void display() {
    if (toRoot == NULL) {
        printf("NULL TREE\n");
        return;
    }
    display2(toRoot, 0);
}

static void printSpaces(int offset) {
    int i;
    for (i=0; i<offset; i++)
        printf(" ");
}

void display2(node23p node, int offset) {
    int kys;

    if (node == NULL) {
        printSpaces(offset);
        printf("{ NULL }\n");
        return;
    }

    printSpaces(offset);
    printf("{ count of keys = %d\n", node->countKeys);
    for (kys = 0; kys <= node->countKeys; kys++)
        switch(kys) {
            case 0:
                printSpaces(offset);
                printf(" <%d\n", node->data[0]);
                display2(node->subT[0], offset+5);
                break;
            case 1:
                if (kys < node->countKeys) {
                    printSpaces(offset);
                    printf(" <%d && >%d\n", node->data[1],
                        node->data[0]);
                    display2(node->subT[1], offset+5);
                    break;
                } else {
                    printSpaces(offset);
                    printf(" >%d\n", node->data[0]);
                    display2(node->subT[1], offset+5);
                    break;
                }
            case 2:
                printSpaces(offset);
                printf(" >%d\n", node->data[1]);
                display2(node->subT[2], offset+5);
                break;
        }
    printSpaces(offset);
    printf("}\n");
}
```

## File tree23Insert.c

```
#include "tree23.h"

struct node23 *toRoot = NULL;

/* Calls to insert2() and remove2() will use
   these variables to tell their parents if
   the parent node needs to accommodate a new
   key after the call */
struct node23 *newRight = NULL;
int newKey = 0;

static void insert2(node23p node, int data);
static void splitLeaf(node23p node, int data);
static void insertInternal(node23p node, int data);

static node23p mkNode(int val) {
    node23p tmp = malloc(sizeof(struct node23));
    assert(tmp != NULL);
    tmp->data[0] = val;
    tmp->data[1] = 0;
    tmp->countKeys = 1;
    tmp->subT[0] = tmp->subT[1] =
        tmp->subT[2] = NULL;
    return tmp;
}

void insert(int data) {
    node23p oldRoot = toRoot;

    if (toRoot == NULL) {
        toRoot = mkNode(data);
        return;
    }
    newRight = NULL;
    insert2(toRoot, data);
    /* check if tree height needs increase */
    if (newRight != NULL) {
        toRoot = mkNode(newKey);
        toRoot->subT[0] = oldRoot;
        toRoot->subT[1] = newRight;
    }
}

static void insert2(node23p node, int data) {
    assert(node != NULL);
    int kys, temp;
    // Primarily to insert in a LEAF
    if (node->subT[0] == NULL) { // This is a leaf
        if (node->data[0] == data)
            return; // Duplicate data
        if (node->countKeys == 2) { // Full leaf node
            if (node->data[1] == data)
                return; // Duplication of data
            // Tell parent to care for new right sibling.
            splitLeaf(node, data);
            return;
        }
        // Add data in this leaf node
        node->data[node->countKeys] = data;
        node->countKeys += 1;
        kys = node->countKeys-1;
        while (kys>0) {
            if (node->data[kys-1] <= node->data[kys])
                return; // Node set up
            // Swap

```

```

        temp = node->data[kys-1];
        node->data[kys-1] = node->data[kys];
        node->data[kys] = temp;
    }
    return;
}
// Insertion into an internal node
assert(node->subT[0] != NULL && node->subT[1] != NULL);
assert(node->countKeys == 1 || node->subT[2] != NULL);
/* When inserting data at an internal node there are
   two phases of importance. So we separately code it */
insertInternal(node, data);
}

static void splitLeaf(node23p node, int data) {
    // Use newRight to tell the parent to update.
    assert(node->countKeys == 2);

    int temp = node->data[1];

    node->countKeys -= 1;
    if (temp < data) {
        newRight = mkNode(data);
        newKey = temp;
        return;
    }

    newRight = mkNode(node->data[1]);
    if (node->data[0] < data) {
        newKey = data;
        return;
    }

    newKey = node->data[0];
    node->data[0] = data;
    return;
}

static insertInternal(node23p node, int data) {
    // Not a leaf node

    int kys;
    int median;
    int who; // Subtree index
    node23p newSibling;

    assert(node != NULL);
    assert(node->subT[0] != NULL);
    newRight = NULL; // Perhaps superfluous
    /* If duplicate data do not insert */
    if (node->data[0] == data) return;
    if (node->countKeys == 2 && node->data[1] == data) return;

    if (node->data[0] > data) {
        who = 0;
        insert2(node->subT[0], data);
    } else if (node->countKeys == 2 && node->data[1] < data) {
        who = 2;
        insert2(node->subT[2], data);
    } else {
        who = 1;
        insert2(node->subT[1], data);
    }
    /* On return determine if this node is obliged to accommodate
       a new key? */
    if (newRight == NULL) // All done

```

```

        return;
    if (node->countKeys == 1) { // There is space for key 2
        node->countKeys += 1;
        switch (who) {
            case 0:
                node->subT[2] = node->subT[1];
                node->data[1] = node->data[0];
                node->data[0] = newKey;
                node->subT[1] = newRight;
                newRight = NULL;
                return;
            case 1:
                node->subT[2] = newRight;
                newRight = NULL;
                node->data[1] = newKey;
                return;
        }
    }

    switch (who) {
        case 2: // newKey is largest key
            newSibling = mkNode(newKey);
            newSibling->subT[1] = newRight;
            newSibling->subT[0] = node->subT[2];
            node->subT[2] = NULL;
            node->countKeys -= 1;
            newKey = node->data[1];
            newRight = newSibling;
            return;
        case 1: // Right key in node is largest
            newSibling = mkNode(node->data[1]);
            newSibling->subT[1] = node->subT[2];
            newSibling->subT[0] = newRight;
            node->subT[2] = NULL;
            node->countKeys -= 1;
            newRight = newSibling;
            return;
        case 0: // newKey is smallest
            newSibling = mkNode(node->data[1]);
            newSibling->subT[1] = node->subT[2];
            newSibling->subT[0] = node->subT[1];
            node->subT[2] = NULL;
            median = node->data[0]; // Median key
            node->data[0] = newKey;
            node->subT[1] = newRight;
            node->countKeys -= 1;
            newKey = median;
            newRight = newSibling;
            return;
    }
}

```

## File tree23Delete.c

```
#include "tree23.h"

struct node23 *toRoot = NULL;

/* Calls to insert2() and remove2() will use
   these variables to tell their parents if
   the parent node needs to accommodate a new
   key after the call */
struct node23 *newRight = NULL;
int newKey = 0;

static void insert2(node23p node, int data);
static void splitLeaf(node23p node, int data);
static insertInternal(node23p node, int data);

static node23p mkNode(int val) {
    node23p tmp = malloc(sizeof(struct node23));
    assert(tmp != NULL);
    tmp->data[0] = val;
    tmp->data[1] = 0;
    tmp->countKeys = 1;
    tmp->subT[0] = tmp->subT[1] =
        tmp->subT[2] = NULL;
    return tmp;
}

void insert(int data) {
    node23p oldRoot = toRoot;

    if (toRoot == NULL) {
        toRoot = mkNode(data);
        return;
    }
    newRight = NULL;
    insert2(toRoot, data);
    /* check if tree height needs increase */
    if (newRight != NULL) {
        toRoot = mkNode(newKey);
        toRoot->subT[0] = oldRoot;
        toRoot->subT[1] = newRight;
    }
}

static void insert2(node23p node, int data) {
    assert(node != NULL);
    int kys, temp;
    // Primarily to insert in a LEAF
    if (node->subT[0] == NULL) { // This is a leaf
        if (node->data[0] == data)
            return; // Duplicate data
        if (node->countKeys == 2) { // Full leaf node
            if (node->data[1] == data)
                return; // Duplication of data
            // Tell parent to care for new right sibling.
            splitLeaf(node, data);
            return;
        }
        // Add data in this leaf node
        node->data[node->countKeys] = data;
        node->countKeys += 1;
        kys = node->countKeys-1;
        while (kys>0) {
            if (node->data[kys-1] <= node->data[kys])
                return; // Node set up
            // Swap

```



```

        temp = node->data[kys-1];
        node->data[kys-1] = node->data[kys];
        node->data[kys] = temp;
    }
    return;
}
// Insertion into an internal node
assert(node->subT[0] != NULL && node->subT[1] != NULL);
assert(node->countKeys == 1 || node->subT[2] != NULL);
/* When inserting data at an internal node there are
   two phases of importance. So we separately code it */
insertInternal(node, data);
}

static void splitLeaf(node23p node, int data) {
    // Use newRight to tell the parent to update.
    assert(node->countKeys == 2);

    int temp = node->data[1];

    node->countKeys -= 1;
    if (temp < data) {
        newRight = mkNode(data);
        newKey = temp;
        return;
    }

    newRight = mkNode(node->data[1]);
    if (node->data[0] < data) {
        newKey = data;
        return;
    }

    newKey = node->data[0];
    node->data[0] = data;
    return;
}

static insertInternal(node23p node, int data) {
    // Not a leaf node

    int kys;
    int median;
    int who; // Subtree index
    node23p newSibling;

    assert(node != NULL);
    assert(node->subT[0] != NULL);
    newRight = NULL; // Perhaps superfluous
    /* If duplicate data do not insert */
    if (node->data[0] == data) return;
    if (node->countKeys == 2 && node->data[1] == data) return;

    if (node->data[0] > data) {
        who = 0;
        insert2(node->subT[0], data);
    } else if (node->countKeys == 2 && node->data[1] < data) {
        who = 2;
        insert2(node->subT[2], data);
    } else {
        who = 1;
        insert2(node->subT[1], data);
    }
    /* On return determine if this node is obliged to accommodate
       a new key? */
    if (newRight == NULL) // All done
        return;
    if (node->countKeys == 1) { // There is space for key 2

```

```

node->countKeys += 1;
switch (who) {
    case 0:
        node->subT[2] = node->subT[1];
        node->data[1] = node->data[0];
        node->data[0] = newKey;
        node->subT[1] = newRight;
        newRight = NULL;
        return;
    case 1:
        node->subT[2] = newRight;
        newRight = NULL;
        node->data[1] = newKey;
        return;
}

switch (who) {
    case 2: // newKey is largest key
        newSibling = mkNode(newKey);
        newSibling->subT[1] = newRight;
        newSibling->subT[0] = node->subT[2];
        node->subT[2] = NULL;
        node->countKeys -= 1;
        newKey = node->data[1];
        newRight = newSibling;
        return;
    case 1: // Right key in node is largest
        newSibling = mkNode(node->data[1]);
        newSibling->subT[1] = node->subT[2];
        newSibling->subT[0] = newRight;
        node->subT[2] = NULL;
        node->countKeys -= 1;
        newRight = newSibling;
        return;
    case 0: // newKey is smallest
        newSibling = mkNode(node->data[1]);
        newSibling->subT[1] = node->subT[2];
        newSibling->subT[0] = node->subT[1];
        node->subT[2] = NULL;
        median = node->data[0]; // Median key
        node->data[0] = newKey;
        node->subT[1] = newRight;
        node->countKeys -= 1;
        newKey = median;
        newRight = newSibling;
        return;
}
}

```