

# algorithm template

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December 22, 2025

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# 数论

## 素性检测

```
1 #include<vector>
2 namespace PrimeTest {
3     long long mul(long long a, long long b, long long mod){
4         return (__int128) a * b % mod;
5     }
6
7     long long Pow(long long a, long long b, long long mod){
8         //mod <= 10^18.
9         long long res = 1;
10        while(b){
11            if (b&1) res = mul(res, a, mod);
12            b >>= 1;
13            a = mul(a, a, mod);
14        }
15        return res;
16    }
17
18    std::vector<long long> pr = {2, 3, 5, 7, 11, 13, 17, 19, 23, 29, 31,
19      ↵ 37};
20
21    bool rabin_test(long long a, long long n, long long s, long long d){
22        long long u = Pow(a, d, n);
23        if (u == 1 or u == n - 1) return false;
24
25        for(long long i = 1; i < s; i++){
26            u = mul(u, u, n);
27            if (u == n - 1) return false;
28        }
29        return true;
30    }
31
32    bool rabin_miller(long long n){
33        if (n < 2) return false;
34        if (n % 2 == 0) return n==2;
35        long long res = 1;
36        long s = 0, d = n-1;
37        while(d%2==0) {
38            s++;
39            d>>=1;
40        }
```

```
41     for(long long i = 0;i<pr.size();i++){
42         if (n%pr[i] == 0) {
43             return n == pr[i];
44         }
45         if (rabin_test(pr[i], n, s, d)){
46             return false;
47         }
48     }
49     return true;
50 }
51 }
```

# 树

## 树的直径

```
1 #include <vector>
2 #include <tuple>
3 namespace TreeDiameter {
4     /*
5      * 无向正权树的最大直径，限制：
6      * 1. 直径需要 $\leq LONG\_LONG\_MAX$ 
7      * 2. 单颗树，而非森林
8      */
9     using namespace std;
10    using Graph = vector<vector<pair<int, long long>>>; // 起点对应的
11    // 边(终点 & 权值)
12    /*
13     * Input: 树，起始点
14     * Output: 离起始点最大的距离，对应的点
15     * 复杂度:  $O(\text{边数})$ 
16     */
17    pair<long long, int> dfs(const vector<vector<pair<int, long long>>>
18    // &g, int cur, int par = -1) {
19        pair<long long, int> ret(0, cur);
20        for (auto e : g[cur]) {
21            if (e.first == par) continue;
22            auto cost = dfs(g, e.first, cur);
23            cost.first += e.second;
24            ret = max(ret, cost);
25        }
26        return ret;
27    /*
28     * Input: 树
29     * Output: 直径起点，直径终点，直径长度
30     */
31    tuple<int, int, long long> tree_diameter(const
32        vector<vector<pair<int, long long>>> &g) {
33        auto u = dfs(g, 0, -1).second;
34        long long dist;
35        int v;
36        tie(dist, v) = dfs(g, u, -1);
37        return make_tuple(u, v, dist);
38    }
```

```

39  /*
40   * 会搜索出一条从cur到goal的路径，结果会放在path里面
41   * Input: 树
42   * Output: 路径
43   * 复杂度:  $O(\text{边数})$ 
44   */
45   void path_restoration(const vector<vector<pair<int, long long>>> &g,
46     vector<int> &path, int cur, int par, int &goal) {
47     path.push_back(cur);
48     if (cur == goal) {
49       goal = -1;
50       return;
51     }
52     for (auto e : g[cur]) {
53       int nxt = e.first;
54       if (nxt == par) continue;
55       path_restoration(g, path, nxt, cur, goal);
56       if (goal == -1) return;
57     }
58     if (goal == -1) {
59       return;
60     }
61     path.pop_back();
62   }
63 }
64 }
```

## 最近公共祖先

```

1 #include <vector>
2
3 namespace LCA {
4 #define V vector
5   /*
6    * 最近公共祖先，限制：
7    * 1. root必须为0 / 1
8    * 2. 总复杂度  $O(q \log(n) + n \log(n))$  // q次查询，一共n个节点，建树过
9    * 程  $n \log(n)$ ，每次查询  $\log(n)$ 
10   */
11   using namespace std;
12   class Tree {
13     private:
14       int n_;
```

```

14     int root_;
15     int lg;
16     V<int> depth;
17     V<V<int>> father;
18     V<V<int>> son;
19     /*
20      * 从跟节点dfs, 来构建depth数组
21      * 复杂度: O(n)
22      */
23     void dfs(int now, int pre = -1, int dep = 1) {
24         depth[now] = dep;
25         for(auto s : son[now]) {
26             dfs(s, now, dep+1);
27         }
28     }
29     /*
30      * 构建祖先关系, 倍增构建
31      * 复杂度: O(n*log(n))
32      */
33     void build_father() {
34         for(int i = 1; i < lg; i++) {
35             for(int j = root_; j < n_ + (root_==1); j++) {
36                 father[i][j] = (father[i-1][j] == -1) ? -1 :
37                     father[i-1][father[i-1][j]];
38             }
39         }
40     }
41     public:
42     Tree(int root, int n) {
43         root_ = root;
44         n_ = n;
45         lg = 1;
46         while((1<<lg) < n) lg++;
47         depth.resize(n+(root==1));
48         father = V<V<int>>(lg, V<int>(n+(root==1), -1));
49         son = V<V<int>>(n+(root==1), V<int>(0));
50     }
51     /*
52      * 增加一个父子关系, now的父亲是pre
53      * 复杂度: O(1)
54      */
55     void add(int pre, int now) {
56         father[0][now] = pre;
57         son[pre].push_back(now);
58     }

```

```

59
60     /*
61      * 完整建树，需要在add完所有父子关系才可以调用
62      * 复杂度：  $O(n \log(n))$ 
63      */
64     void build() {
65         this->dfs(root_);
66         this->build_father();
67     }
68
69     /*
70      * 查询u和v的最近公共祖先
71      * 复杂度：  $O(\log(n))$ 
72      */
73     int query_lca(int u, int v) {
74         if(depth[u] > depth[v]) {
75             swap(u, v);
76         }
77         int depth_diff = depth[v] - depth[u];
78         for(int i = 0; i < lg; i++) {
79             if(depth_diff & (1<<i)) {
80                 v = father[i][v];
81             }
82         }
83         if(u == v) {
84             return u;
85         }
86         for(int i = lg - 1; i >= 0; i--) {
87             if(father[i][u] != father[i][v]) {
88                 u = father[i][u];
89                 v = father[i][v];
90             }
91         }
92         return father[0][u];
93     }
94 };
95
96
97

```