# Review Notes: Midterm 1

CSCI2270: Data Structures

## C++ Fundamentals

1. Command-line arguments: Remember to typecast necessary arguments (stof(), stoi())



2. A structure in C++ is a way to bundle data together. For example, the student information can be maintained using a struct which may contain their student ID, full name, age, GPA, etc.

```
struct student {
     unsigned int student_id; string name; int age; float gpa;
};
```

A class is an extension of a struct which also combines data with associated functions and introduces membership criteria (private v/s public). A struct is like a class with all members being public.

```
class student {
    unsigned int student_id; string name; int age; float gpa;
    void print_student_info();
    void update_gpa(float updated_gpa);
};
```

- 3. Pointers are variables that store (or point to) a memory address. References (&) give memory address of a variable. e.g.: int x = 10; int \*ptrX = &x; student\* stu = new student;
  - a. Use new, delete keywords to allocate memory and store address in a pointer
  - b. Dereferencing Operators (\*, ->): \*ptrX == 10; stu->name = "John Doe";
- 4. A function in C++ is a subroutine that can be called by other code routines or subroutines. A function may take arguments and may return a value (both optional). e.g.

```
void reset_value(int c, int d) { c = 7; d = 8; }
void reset_pointers(int * c, int * d) { *c = 7; *d = 8; }
void reset_reference(int & c, int & d) { c = 7; d = 8; }
int main() {
   int a = 3; int b = 4;
   reset(a, b);
   reset_pointer(&a, &b);
   reset_reference(a, b); }
```

- a. **Pass by value**: When the **reset** function is called from main, the value of a in main will be copied over to the argument c, and the value of b in main will be copied over to argument d.
- b. <u>Pass by pointers</u>: When the *reset\_pointers* function is called from main, the address of **a** in main will be copied over to the argument **c**, and the address of **b** in main will be copied over to argument **d**. Both parameters **c** and **d** are pointers to int, and contain the addresses of **a** and **b**. Dereferencing **c** and **d** will change the values of the original arguments **a** and **b**.
- c. <u>Pass by reference</u>: When the *reset\_reference* function is called from main, parameters **c** and **d** will act as references to the original variables **a** and **b** in main. Updating **c** and **d** will change the values of the original arguments **a** and **b**.

## Arrays

- 1. **Basics:** A contiguous chunk of memory, indexed with an integer.
- 2. Allocation: Static (int array[100]) or dynamic (int \* array = new int[100])
- 3. Accessing notation (array/pointer): a[i] == \*(a+i) == \*(i+a) == i[a]; // true
  - a. Pointer accessing the first element of the array: int \* ptr = &a[0];
  - b. Pointer accessing the i<sup>th</sup> element of the array: int \* ptr = &a[i];
  - c. Pointer arithmetic: ptr++ will access the next element, ptr-- will access the previous.

Don't access the array out of bounds! (< 0 or > array size)!

4. **Deallocation:** You can only manually delete a dynamic array: delete [] array

## **Linked List**

- Basics: A linked list can only be accessed using its first node, pointed to by a pointer head. Each node contains data and a pointer next. next points to the next node in the list; for the last node next = nullptr or NULL. Never modify head unless inserting or deleting at head.
- 2. **Traversing the linked list and searching for an element:** Initialize a pointer *temp* to *head*, and hop *temp* through the linked list (*temp* = *temp*->*next*) using a loop. <u>Hint</u>: If searching based on index, a for loop is a good idea (or a while loop with a counter). If searching based on the node value, a while loop is easier. (However, with the right design, either loop can be used for both the cases)

```
while(pres) { pres = pres->next; }
```

3. **Deleting a node:** Use two pointers *prev* and *pres* to traverse the linked list such that *prev* is one node behind *pres*, until *pres* points to the node to be deleted. Then *prev*->next = *pres*->next, and then delete *pres*. Recall how you can do the two pointer traversal:

```
Node * pres = head;
while(pres) {
    prev = pres;
    pres = pres->next;
}
```

4. **Inserting a node after a node prev:** Create a new node and manipulate the **next** pointers.

```
Node * temp = new Node;
temp->next = prev->next;
prev->next = temp;
```

Important: Always handle corner cases for the above functions!

## **Stack and Queue**

- 1. Stack: Last-In First-Out (LIFO), First-In Last-Out (FILO)
- 2. Queue: First-In First-Out (FIFO), Last-In Last-Out (LILO)
- 3. Linked List Implementation:

	Linked List	Stack (LL)	Queue (LL)
Members	head	head	head, tail
Insertion	Anywhere	Push: At head	Enqueue: At tail
Deletion	Anywhere	Pop: At head	Dequeue: At head

## **Complexity**

- 1. Time complexity: number of times a set of operations is repeated, parameterized by input size n
- 2. Space complexity: additional space required, parameterized by input size n
- 3. Big-Oh notation: Keep the highest power, ignore constants. E.g.:  $O(n^3 + 100n^2 + 1000) = O(n^3)$

$$O(1) < O(log(n)) < O(n) < O(nlog(n)) < O(n^k) < O(2^n) < O(n!)$$

- 4. Loops: Literally repetitive sections of code.
  - a. O(n): for (int i = 0; i < n; i++) { ... }
  - b. O(mn): for (int i = 0; i < m; i++) { for (int j = 0; j < n; j++) { ... } }

## **Practice**

## **Array**

Recall array doubling, except here we are going to use an extend\_factor as well as a twist in the copying mechanism. We will call this function extend array, and it has the following prototype:

```
int* extend array(int* A, int &array size, int extend factor);
```

- 1. Create a new dynamic array B of size = (array size \* extend factor)
- 2. Then, copy all the elements from **A** to **B** in an interleaving manner, i.e., copy all odd-index elements contiguously to B followed by all even-index elements.
- 3. Modify array size to the new size
- 4. Finally, return the pointer **B** and (in main) store this returned pointer in **A**.

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#### 5. Questions

a. What is the algorithmic complexity of this function?

1

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0

b. If extend array is a void function, how will you modify the implementation?

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c. Complete using dereference operator (\*) and pointer arithmetic instead of [] to index the array.

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### **Linked List**

Implement a class called LinkedList. The class should maintain a single data member (also private), head which is a pointer to a Node. Create separate header and source files.

Use the following definition for a Linked List node: struct Node { int id; string name; float gpa; };

Implement the following subroutines for this class:

- 1. LinkedList(): Initialize head to NULL
- 2. ~LinkedList(): Delete the entire list (don't just delete head)
- 3. Node\* CreateNode(int id, string name, float gpa): Dynamically allocate a new node with the provided arguments and return a pointer to it
- 4. void PrintNode (Node\* n): Print passed node n's data members in a single line
- 5. void PrintList(): Traverse the list and print each node's details in sequence using PrintNode
- 6. Node\* Tail(): Find and return a pointer to the tail of the list; return NULL if list is empty
- 7. Node\* Search (string name): Find and return the first node (starting at head) in the list which consists of the same name data member as the passed argument.
- 8. Node\* Search(int ix): Find and return the node at index = ix.
- 9. int Index (Node\* n): Find and return the index of the passed node n. head has index 0.
- 10. void Filter(float low, float high): Traverse the list & print nodes whose gpa ∈ [low, high]
- 11. Node\* CreateLoop(string name): Similar to A4, create a loop in the list, return tail pointer.
- 12. void Insert (Node\* p, int ix): Insert node p at index ix. If the list is empty, insert as head.
- 13. void Delete (int ix): Delete node at index ix.
- 14. int Length (): Compute and return the number of nodes in the list

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