

Chapter 1: Numbers and Expressions

“All the mathematical sciences are founded on relations between physical laws and laws of numbers, so that the aim of exact science is to reduce the problems of nature to the determination of quantities by operations with numbers”

James Clerk Maxwell 1856

Topics:

- ▶ numbers like 2, 3.75, π and $\sqrt{19}$
- ▶ evaluation of expressions like

$$\frac{-b \pm \sqrt{b^2 - 4ac}}{2a}$$

- ▶ representation and approximation of numbers and expressions
- ▶ computational errors, can they be avoided or at least controlled?

1.1 Numbers

Numbers for Computations

- ▶ integers, rational, real and complex numbers

$$\mathbb{Z} \subset \mathbb{Q} \subset \mathbb{R} \subset \mathbb{C}$$

- ▶ evaluate expressions
- ▶ for solution of equations
 - ▶ linear systems of equations $\mathbb{Q}, \mathbb{R}, \mathbb{C}$
 - ▶ polynomial equations \mathbb{C}
- ▶ for continuous functions: \mathbb{R} and \mathbb{C}

Definition: Ring

- ▶ A ring R is a set with two binary operations $+$ and $*$ with the following properties
- ▶ $(R, +)$ is an abelian group which satisfies, for all $a, b, c \in R$:
 - ▶ $a + (b + c) = (a + b) + c$, *associative law*
 - ▶ $a + b = b + a$, *commutative law*
 - ▶ there exists $0 \in R$ such that $a + 0 = a$
 - ▶ there exists $-a$ such that $a + (-a) = 0$
- ▶ $(R, *)$ is a monoid where for all $a, b, c \in R$:
 - ▶ $a * (b * c) = (a * b) * c$
 - ▶ there exists $1 \in \mathbb{R}$ such that $1 * a = a$
- ▶ distributive law

$$(a + b) * c = a * c + b * c, \quad a * (b + c) = a * b + a * c$$

Rings in Computation

- ▶ all the number sets considered are rings including \mathbb{Z} , \mathbb{Q} , \mathbb{R} and \mathbb{C}
- ▶ the sets of functions considered are rings including
 - ▶ continuous functions
 - ▶ polynomials
- ▶ set of square n by n matrices with elements from a ring is a ring
- ▶ the arithmetic laws lead to efficient expression evaluation
- ▶ the distributive law is the basis for fast algorithms like fast matrix multiplication, the FFT but also for machine learning and dynamic programming
 - ▶ note that the number of multiplications is 2 in $a * b + a * c$ but 1 in $a * (b + c)$

\mathbb{Q} and \mathbb{R}

- ▶ \mathbb{Q} is a countable subset of \mathbb{R}
- ▶ \mathbb{R} is used for theory but (subsets of) \mathbb{Q} used for actual computations
- ▶ \mathbb{Q} is dense in \mathbb{R} , i.e. $\forall x \in \mathbb{R}, \epsilon > 0, \exists u \in \mathbb{Q} :$

$$|x - u| \leq \epsilon$$

Decimal and binary fractions

- ▶ for computations we consider subsets of \mathbb{Q} for a fixed positive $B \in \mathbb{Z}$

$$\mathbb{Q}_B = \{p / B^k \mid p, k \in \mathbb{Z}, k \geq 0\}$$

- ▶ here we only consider $B = 10$ or $B = 2$:
 - ▶ decimal fractions \mathbb{Q}_{10} used for manual computations
 - ▶ binary fractions \mathbb{Q}_2 implemented in computer hardware
- ▶ we use the decimal point, e.g.,

$$44.78 = \frac{4478}{10^2}$$

Fractions

$$\mathbb{Z} = \mathbb{Q}_1 \subset \mathbb{Q}_2 \subset \mathbb{Q}_{10} \subset \mathbb{Q}$$

- ▶ \mathbb{Q}_2 (and \mathbb{Q}_{10}) is a dense subset of \mathbb{Q} and thus of \mathbb{R} . In particular, each real number can be written as an infinite decimal or binary fraction.
- ▶ The sets of fractions \mathbb{Q}_B are all rings and contain the integers

Proposition $\mathbb{Q}_2 \subset \mathbb{Q}_{10}$

Proof.

- ▶ $x \in \mathbb{Q}_2$, thus

$$x = \frac{p}{2^k}$$

for some integers p and $k \geq 0$.

- ▶ consequently

$$x = \frac{5^k p}{10^k} \in \mathbb{Q}_{10}$$

1.2 Representation of Numbers

Simple representation of Integers

- ▶ small integers
 - ▶ counts: $|$, $||$, $|||$, $||||$, \dots , each number is the cardinality of a set
 - ▶ digits: $0, 1, 2, \dots, 9$, each number has a symbol
 - ▶ roman numbers: $I, II, III, IV, \dots, XXXII, \dots$
- ▶ numbers need to be represented in order to do arithmetics
- ▶ all computers (including us) are finite

implementing your own numbers in Python using strings --

```
x = '||||'
```

```
y = '||'
```

```
z = x + y    # concatenation is addition
```

```
print("OUTPUT:")
```

```
print('sum x + y = {}'.format(z))
```

```
n = len(z)    # conversion to ordinary integers
```

```
print('sum in decimal number system x+y = {}'.format(n))
```

```
u = 13*'|'    # conversion to our system
```

```
print('conversion of 13: {}'.format(u))
```

OUTPUT:

```
sum x + y = |||||
```

```
sum in decimal number system x+y = 6
```

```
13 = 6 + 6 + 1  ||||| ||||| ||
```

Representation of Integers

- ▶ computer and human number representations are similar, and of the form

$$n = \pm \sum_{k=0}^t n_k B^k$$

B : *basis* (humans: $B = 10$, computers: $B = 2$)

The Polynomials which represent integers

- ▶ polynomial of degree n

$$p(x) = c_0 + c_1x + \cdots + c_nx^n$$

- ▶ set P of all polynomials is an infinite dimensional linear (vector) space with basis $1, x, x^2, \dots$
- ▶ is also a ring with multiplication defined by

$$p * q(x) = p(x)q(x)$$

- ▶ a ring which is also a vector space is called an *algebra*
- ▶ for representing the positive integer $n = n_0 + n_1B + \cdots + n_tB^t$ choose the polynomial

$$p(x) = n_0 + n_1x + \cdots + n_tx^t$$

and one has

$$n = p(B)$$

- ▶ example $n = 739$ and $B = 10$:

Representation of Rationals

- ▶ rationals $x \in \mathbb{Q}$ are represented as pairs of integers (p, q) and written as

$$x = \frac{p}{q}$$

- ▶ uniqueness is achieved by choosing the $\gcd(p, q) = 1$ (use Euclid's algorithm)
- ▶ as rational numbers are ratios of integers, and each integer is represented by a polynomial and a basis B , each rational number $x \in \mathbb{Q}$ is represented by a rational function

$$r(x) = \frac{p(x)}{q(x)}$$

and

$$x = r(B)$$

Rational numbers in Python

```
from fractions import Fraction
```

```
x = Fraction(16, -10)
```

```
print("OUTPUT:")
```

```
print(x)
```

```
y = Fraction('3/7')
```

```
print(y)
```

```
z = Fraction('1.2341')
```

```
print(z)
```

OUTPUT:

-8/5

3/7

12341/10000

Representation of decimal and binary (and other) fractions

standard integer format

- ▶ any $q \in Q_B$ is of the form

$$q = \frac{n}{B^k}$$

for some integers n and k

- ▶ example – approximation of $1/3$:

$$\frac{333}{1000}$$

- ▶ uses the rational function

$$r(x) = \frac{n_0 + n_1x + \cdots + n_tx^t}{x^k}$$

scientific format

- ▶ any $q \in Q_B$ is of the form

$$q = (n_t + n_{t-1}B^{-1} + \cdots + n_0B^{-t})B^e$$

- ▶ example – approximation of $1/3$:

0.333

- ▶ uses the rational function

$$f(x) = (n_t + n_{t-1}x^{-1} + \cdots + n_0x^{-t})x^e$$

where $e = t - k$

Representation of Real Numbers

- ▶ real numbers $x \in \mathbb{R}$ are represented as (potentially infinite) power series in the basis

$$x = \pm \sum_{j=-\infty}^t n_j B^j$$

again, the basis $B = 10$ is used in human computation and $B = 2$ is used by computers

- ▶ the digits $n_j \in \{0, \dots, B - 1\}$
- ▶ real numbers need to be approximated

Representation of Complex Numbers

- ▶ complex numbers $z \in \mathbb{C}$ are represented as pairs of reals

$$z = x + iy$$

- ▶ addition like vectors
- ▶ complex multiplication
- ▶ conjugate complex
- ▶ imaginary unit i

Complex numbers in Python

```
z = 4.0 + 5.0j
```

```
print("OUTPUT:")
```

```
print("Re(z) = {zr:g}, Im(z) = {zi:g}"  
      .format(zr=z.real,zi=z.imag))
```

```
I = 1j      # sqrt(-1)
```

```
print("I**2 = {}".format(I**2))
```

OUTPUT:

Re(z) = 4, Im(z) = 5

I**2 = (-1+0j)

Floating Point Numbers – approximations of real numbers

The set of floating point numbers with t digits to base B is

$$\mathbb{F}_B(t) = \{x = \pm \sum_{j=1}^t c_j B^{-j+e} \mid e \in \mathbb{Z}, c_j \in \mathbb{Z}_B, c_1 \neq 0\} \cup \{0\}$$

where $\mathbb{Z}_B = \{0, \dots, B-1\}$.

$$\mathbb{F}_B(t) \subset \mathbb{Q}_B \subset \mathbb{Q} \subset \mathbb{R}$$

- ▶ $\mathbb{F}_B(t)$ is *not*
 - ▶ a ring
 - ▶ dense in \mathbb{Q}
- ▶ motivation: subset of $\mathbb{F}_B(t)$ with $e = e_{\min}, \dots, e_{\max}$ is computationally feasible
- ▶ challenge: floating point arithmetic has to be (re-)defined

Floating point numbers in Python

```
x = 0.7  
print("OUTPUT:")  
print("computer approximation \nx= {0:2.53g}".format(x))
```

OUTPUT:

computer approximation

x= 0.699999999999999955910790149937383830547332763671875

The IEEE standard 754 floating point numbers

- ▶ $B = 2$ and $t = 53$, each number is stored in 64 bit
- ▶ special numbers are: 0 , $\pm\infty$, some non-normalised numbers, NaNs
- ▶ the exponents $e = -1022, \dots, 1023$
- ▶ the standard also specifies details about the arithmetic

Representation of floating point numbers

- ▶ representation of any floating point number

$$x = \pm 0.d_1 d_2 \dots d_t \cdot B^e$$

B : base, $d_j \in \{0, \dots, B-1\}$: digits, e : exponent, number of digits t

- ▶ written as a sum

$$x = sB^e \sum_{j=1}^t d_j B^{-j}$$

$s = +1, -1$ sign

- ▶ normalised $d_1 \neq 0$
- ▶ $\mathbb{F}_{\leq}(B, t)$ denotes floating point numbers with t digits in base B (here we allow any $e \in \mathbb{Z}$, in practice it is a finite range, see notes)

IEEE 754 standard on representation

- ▶ most commonly used system today: IEEE double precision

number system	base B	number of digits t	exponent range
IEEE double precision	2	53	$[-1022, 1023]$
IEEE single precision	2	24	$[-126, 127]$

- ▶ Note: there are many other formats, both decimal and binary

Effect of finite exponent

There is a smallest (strictly positive) floating point number

- ▶ normalised: $x_{\min} = B^{e_{\min}-1}$
- ▶ denormalised: $B^{e_{\min}-t}$

There is a largest (positive) floating point number *

$$x_{\max} = B^{e_{\max}} - B^{e_{\max}-t}$$

Errors occur when computations exceed these limits

- ▶ *underflow* occurs for $0 < x < x_{\min}/2$
- ▶ *overflow* occurs when numbers exceed x_{\max} We ignore these types of errors in the remainder by allowing $e \in \mathbb{Z}$

Questions