

* (Linear) Least Squares: 对称正定矩阵

Given a symmetric PSD matrix A, we want to minimize

$$f(x) = \frac{1}{2} \langle Ax, x \rangle - \langle b, x \rangle$$

$$\nabla f(x) = Ax - b$$

$$\nabla^2 f(x) = A \succeq 0 \Rightarrow f \text{ is convex}$$

if $\lambda_{\max}(A)$ is the max eval of A, then $f(x)$ is $\lambda_{\max}(A)$ -smooth

if A^+ is the Moore-Penrose pseudo inverse of A, then:

$$\forall x: \|\nabla f(x)\|_2 \geq \|\nabla f(A^+ b)\|_2 \quad \|Ax-b\|_2 \geq \|A A^+ b - b\|_2$$

→ If the system $Ax=b$ is solvable, then

$$x^* = A^+ b \in \arg \min_{x \in \mathbb{R}^d} f(x) \quad \nabla^2 f \text{ is constant.}$$

$$f(y) = f(x) + \langle \nabla f(x), y-x \rangle + \frac{1}{2} \underbrace{\langle \nabla^2 f(x+\gamma(y-x)), y-x \rangle}_{A}, \quad \gamma \in (0,1).$$

$$= f(x) + \langle \nabla f(x), y-x \rangle + \frac{1}{2} \langle A(y-x), y-x \rangle$$

→ Use this when computing the exact L.s. step size in HJB.

* Other forms of linear least square problems:

$$\tilde{f}(x) = \frac{1}{2} \|Mx - c\|_2^2 = \frac{1}{2} \langle M^T M x, x \rangle - \langle M^T c, x \rangle + \frac{1}{2} \|c\|_2^2$$

$$A = M^T M, \quad b = M^T c$$

$$\arg \min_x \tilde{f}(x) = \arg \min_x \left\{ \underbrace{\frac{1}{2} \langle Ax, x \rangle - \langle b, x \rangle}_{f(x)} \right\}$$

$$\tilde{f}(x) - \tilde{f}(x^*) = f(x) - f(x^*)$$

共轭梯度

* Method of Conjugate Gradients: 对称正定.

* Here, we will take A to be symmetric & PD.

* Consider methods of the form: $x_k = x_0 - \sum_{i=0}^{k-1} h_{i,k} \nabla f(x_i)$, where $h_{i,k} \in \mathbb{R}$

$$(*) \quad x_k = x_0 - \sum_{i=0}^{k-1} h_{i,k} \nabla f(x_i), \text{ where } h_{i,k} \in \mathbb{R}$$

Both GD and AGD take the form $(*)$.

+ As A symm. & PD, it is invertible. $x^* = A^{-1}b = \underset{x}{\operatorname{argmin}} f(x)$.

x_k defined by $(*)$ is from $x_0 + \text{Lin} \{ \nabla f(x_0), \dots, \nabla f(x_{k-1}) \}$.

$$\nabla f(x_0) = Ax_0 - b = A(x_0 - x^*), \text{ as } b = Ax^*$$

$$\nabla f(x_1) = Ax_1 - b$$

$$= A(\underline{x_0} - h_{0,1} \nabla f(x_0)) - \underline{Ax^*}$$

$$= A(\underline{x_0} - x^*) - h_{0,1} A^2 (\underline{x_0} - x^*) \in \text{Lin} \{ A(x_0 - x^*), A^2(x_0 - x^*) \}$$

Suppose $x_k \in x_0 + \text{Lin} \{ A(x_0 - x^*), A^2(x_0 - x^*), \dots, A^k(x_0 - x^*) \}$

Claim $x_{k+1} \in x_0 + \text{Lin} \{ A(x_0 - x^*), \dots, A^{k+1}(x_0 - x^*) \}$

$$x_{k+1} = x_0 - \sum_{i=0}^k h_{i,k+1} \nabla f(x_i) = x_0 - \underbrace{\sum_{i=0}^{k-1} h_{i,k+1} \nabla f(x_i)}_{\in x_0 + \text{Lin} \{ A(x_0 - x^*), \dots, A^k(x_0 - x^*) \}} - h_{k,k+1} \nabla f(x_k)$$

$$\nabla f(x_k) = A(x_k - x^*)$$

$$= A(x_0 + \sum_{i=1}^k \alpha_i A^i (x_0 - x^*) - x^*)$$

$$= A(x_0 - x^*) + \sum_{i=1}^k \alpha_i A^{i+1} (x_0 - x^*)$$

$$\in \text{Lin} \{ A(x_0 - x^*), \dots, A^{k+1}(x_0 - x^*) \}$$

K_k is krylov subspace

$\mathcal{K}_k = \text{Lin} \{ A(x_0 - x^*), \dots, A^k(x_0 - x^*) \}$ - Krylov subspace of order k

- Method of Conjugate Gradients:

$$(CG) \quad x_k^{out} = \underset{x \in x_0 + \mathcal{K}_k}{\operatorname{argmin}} f(x)$$

* Lemma (1.3.1 in Nesterov's book) For any $k \geq 1$, we have $\mathcal{K}_k = \text{Lin} \{ \nabla f(x_0^{out}), \dots, \nabla f(x_{k-1}^{out}) \}$

Proof: By induction on k .

Base case : $k=1$

$$\nabla f(x_0) = A(x_0 - x^*) \Rightarrow \mathcal{K}_1 = \text{Lin} \{ A(x_0 - x^*) \} = \text{Lin} \{ \nabla f(x_0) \}.$$

Suppose the lemma holds for some $k \geq 1$.

Any point $x_k \in x_0 + \mathcal{K}_k$ can be expressed as :

$$x_k = x_0 + \sum_{i=1}^k \beta_{i,k} A^i (x_0 - x^*)$$

$$\nabla f(x_k) = A(x_0 - x^*) + \sum_{i=1}^k \beta_{i,k} A^{i+1} (x_0 - x^*)$$

$$= A(x_0 - x^*) + \underbrace{\sum_{i=1}^{k-1} \beta_{i,k} A^{i+1} (x_0 - x^*)}_{\in \mathcal{K}_k} + \beta_{k,k} A^{k+1} (x_0 - x^*)$$

$$\in \mathcal{K}_k$$

$$\mathcal{K}_{k+1} = \text{Lin} \{ \mathcal{K}_k \cup A^{k+1} (x_0 - x^*) \} = \text{Lin} \{ \mathcal{K}_k \cup \nabla f(x_k) \}.$$

$\Rightarrow CG$ outputs x_k^{out} s.t. $f(x_k^{out}) - f(x^*) \leq \epsilon$ in at most

$$\mathcal{O} \left(\min \left\{ \sqrt{\frac{\epsilon}{L}} \|x_0 - x^*\|_2, \sqrt{\frac{\epsilon}{m}} \log \left(\frac{L \|x_0 - x^*\|_2^2}{\epsilon} \right) \right\} \right)$$

$$L = \lambda_{\max}(A), \quad m = \lambda_{\min}(A).$$

* Lemma (1.3.2 in Nes'18 book) If x_k^{out} is generated by CG, then $\forall i < k \quad \langle \nabla f(x_k^{\text{out}}), \nabla f(x_i^{\text{out}}) \rangle = 0$.

Proof: Let $k > i$. Define:

$$\bar{f}(h_k) = f\left(x_0 - \sum_{i=0}^{k-1} h_{i,k} \nabla f(x_i^{\text{out}})\right)$$

$x_k^{\text{out}} \in x_0 + \mathbb{R}^d$

$$h_k = \underset{h \in \mathbb{R}^d}{\operatorname{argmin}} \bar{f}(h) \quad \text{then } x_k^{\text{out}} = x_0 - \sum_{i=0}^{k-1} h_{i,k} \nabla f(x_i^{\text{out}})$$

$$= \underset{x \in \mathbb{R}^d}{\operatorname{argmin}} f(x)$$

$$\frac{\partial \bar{f}(h_k)}{\partial h_{i,k}} = 0 = \langle \nabla f(x_k^{\text{out}}), -\nabla f(x_i^{\text{out}}) \rangle. \quad \square$$

* Corollary CG finds $x^* = \underset{x \in \mathbb{R}^d}{\operatorname{argmin}} f(x)$ in at most d iterations.

* Corollary. $\forall p \in \mathbb{R}^d, \quad \langle \nabla f(x_k^{\text{out}}), p \rangle = 0$.