

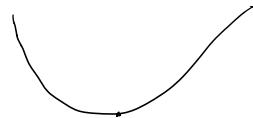
* Minima of convex functions:

(P) $\min_{x \in X} f(x)$: f is convex, X is convex, closed, and nonempty

* Thm 2.6 Consider (P). We have the following:

(a) Any local solution to (P) is also a global solution.

(b) The set of global solutions to (P) is convex.



Proof:

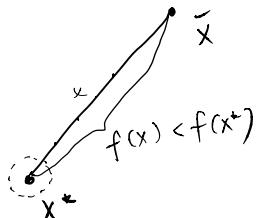
(a) Suppose f.p.o.c. that x^* is a local but not a global solution. Then $\exists \bar{x} \in X$, s.t. $f(\bar{x}) < f(x^*)$.

As X is convex, $\forall \alpha \in (0, 1)$:

$$(1-\alpha)x^* + \alpha\bar{x} \in X$$

As f is convex, $\forall \alpha \in (0, 1)$:

$$f((1-\alpha)x^* + \alpha\bar{x}) \leq (1-\alpha)f(x^*) + \alpha f(\bar{x}) < f(x^*)$$



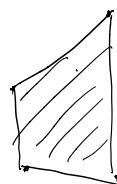
\Rightarrow Every neighborhood of x^* must include a point $(1-\alpha)x^* + \alpha\bar{x}$ for some $\alpha > 0$ that will have a strictly lower function value. $\Rightarrow x^*$ cannot be a local solution.

(b) Let $x^*, \bar{x} \in X$ be any two global solutions.

X is convex $\Rightarrow \forall \alpha \in (0, 1)$: $(1-\alpha)x^* + \alpha\bar{x} \in X$.

f is convex $\Rightarrow \forall \alpha \in (0, 1)$:

$$\begin{aligned} f((1-\alpha)x^* + \alpha\bar{x}) &\leq (1-\alpha)f(x^*) + \alpha f(\bar{x}) \\ &= f(x^*) = f(\bar{x}) \end{aligned}$$



$$\Rightarrow f((1-\alpha)x^* + \alpha\bar{x}) = f(x^*)$$

$\Rightarrow (1-\alpha)x^* + \alpha\bar{x}$ is a global solution.

\Rightarrow the set of global solutions must be convex.

□

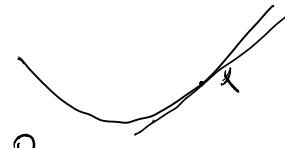
* Thm.

(a) Let f be cont. by diff. 'able. f is convex if and only if

$$\forall x, y : f(y) \geq f(x) + \langle \nabla f(x), y - x \rangle .$$

(b) Let f be twice cont. by diff. 'able.

f is convex if and only if $\forall x : \nabla^2 f(x) \geq 0$.

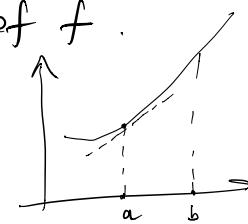


* Thm 2.7 Let f be cont. by diff. 'able and convex.

If $\nabla f(x^*) = 0$, then x^* is a global min of f .

Pf: Use Part (a) of the Thm above:

$$\forall x : f(x) \geq f(x^*) + \langle \nabla f(x^*), x - x^* \rangle .$$



(for constrained setups, we would use $\langle \nabla f(x^*), x - x^* \rangle \geq 0, \forall x$)

* Strongly convex functions:

* Def. Given $m > 0$, we say that $f: \mathbb{R}^d \rightarrow \bar{\mathbb{R}}$ is m -strongly convex (or strongly convex w) modulus m), if $\forall x, y \in \mathbb{R}^d$:

$$f((1-\alpha)x + \alpha y) \leq (1-\alpha)f(x) + \alpha f(y) - \frac{m}{2}(1-\alpha)\alpha \|y - x\|^2 .$$

* Ex: 1) When f is cont. by diff. 'able, equivalently

$$\forall x, y : f(y) \geq f(x) + \langle \nabla f(x), y - x \rangle + \frac{m}{2} \|y - x\|^2 .$$

2) When f is twice cont. by diff. 'able, equivalently:

$$\forall x : \nabla^2 f(x) \geq m I .$$

* Thm 2.8. Suppose that $f: \mathbb{R}^d \rightarrow \bar{\mathbb{R}}$ is cont. by diff. 'able and m -strongly convex for some $m > 0$. If $\nabla f(x^*) = 0$, then x^* is the unique global min of f .

Proof: From Ex 1):

$$\forall x : f(x) \geq f(x^*) + \langle \nabla f(x^*), x - x^* \rangle + \underbrace{\frac{m}{2} \|x - x^*\|^2}_{>0 \text{ unless } x = x^*} .$$

■

* Growth of sequences:

$$\{a_k\}_{k \geq 1}, \{b_k\}_{k \geq 1}, \forall k: a_k, b_k \geq 0. \quad a_k \leq 10 b_k \\ a_k = O(b_k)$$

* "Big-Oh" notation:

$$a_k = O(b_k) \Leftrightarrow (\exists M > 0) (\exists K < \infty) (\forall k \geq K): a_k \leq M b_k.$$

$$(E.g., k = O(\frac{1}{10} k^2), k = O(\frac{1}{10!} k))$$

* If $a_k = O(b_k)$ and $b_k = O(a_k)$, we write $a_k = \Theta(b_k)$.

* "Little-Oh" notation:

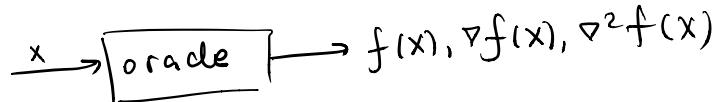
$$a_k = o(b_k) \Leftrightarrow \lim_{k \rightarrow \infty} \frac{a_k}{b_k} = 0.$$

* Algorithmic setup:

1) first-order oracle model:



2) second-order oracle model:



* All algorithms we consider in this class are iterative:

- start w/ some x_0 , get oracle answers for x_0 , choose x_1 ,
- at iteration k , get oracle answers for x_k , choose x_{k+1}

* Basic Descent Methods:

* Assumptions for this part:

(A1) f is L -smooth for some $L < \infty$ (thus also cont.'ly diff.able)

(A2) $X = \mathbb{R}^d$, i.e., the problem is unconstrained

Note: for now, and until explicitly stated otherwise, we are not assuming that f is convex.

* Def. $p \in \mathbb{R}^d$ is a descent direction for f at x if $f(x+tp) < f(x)$ for all suff. small $t > 0$.

* Prop 3.2. If f is cont. /ly diff. /able (in a neighborhood of x), then any p s.t. $\langle \nabla f(x), p \rangle < 0$ is a descent direction.

Proof: TT + continuity of ∇f : $y = x + tp$

$f(x+tp) = f(x) + t \langle \nabla f(x+ptp), p \rangle$ for some $p \in [0,1]$.
We know that $\langle \nabla f(x), p \rangle < 0$. As ∇f is continuous,
for all suff. small $t > 0$:

$$t \langle \nabla f(x+ptp), p \rangle < 0$$

$$\Rightarrow f(x+tp) < f(x)$$

* What would be a good descent direction?

- could try to move in the direction of $-\nabla f(x)$

- justification:

Look at all p w/ $\|p\|_2 = 1$. Then:

$$\inf \langle \nabla f(x), p \rangle = -\|\nabla f(x)\|_2 \text{ attained for } p = -\frac{\nabla f(x)}{\|\nabla f(x)\|_2}$$

$$\|p\|_2 = 1$$

