

* Our standard optimization problem:

$$\min_{x \in X} f(x) \quad (\text{P})$$

↓ ↓
vector feasible set

objective function

$$\max_{x \in X} f(x) \Leftrightarrow \min_{x \in X} -f(x)$$

$$*\text{ the value of } (\text{P}): \text{val}(\text{P}) = \inf_{x \in X} f(x).$$

- * To give (P) a meaning, we need to specify:
 - vector space, feasible set, objective function
 - what it means to "solve" (P)

Γ Q. Can we even hope to solve an arbitrary opt. problem?

Ex. Can you come up with an example of positive integers x, y, z s.t.

$$x^2 + y^2 = z^2 \leftarrow (\text{Pythagorean triples})$$

$$(3, 4, 5); (5, 12, 13); (8, 15, 17) \dots$$

$$\text{How about } x^3 + y^3 = z^3 ? \leftarrow$$

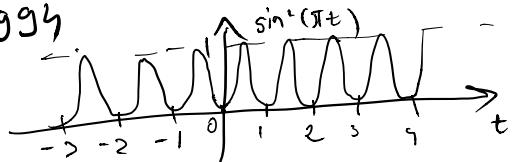
* Fermat's conjecture [Fermat's Last Theorem] (1637)

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 For any $n \geq 3$, $x^n + y^n = z^n$ has no solutions over positive integers.

→ Proved by Andrew Wiles in 1994

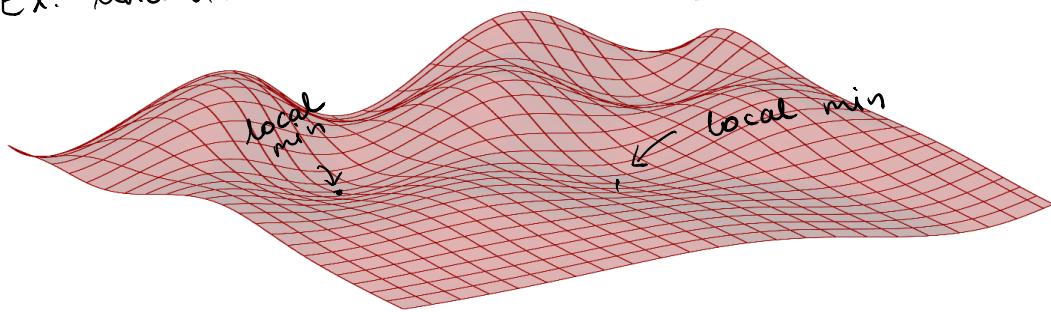
Consider:

$$(\text{PF}) \left\{ \begin{array}{l} \min_{x,y,z \geq 1} (x^n + y^n - z^n)^2 \\ \text{s.t. } x \geq 1, y \geq 1, z \geq 1, n \geq 3 \\ \sin^2(\pi x) + \sin^2(\pi y) + \sin^2(\pi z) = 0 \end{array} \right.$$



If you could certify whether $\text{val}(\text{PF}) \neq 0$, you would have found a proof for Fermat's conjecture.

Ex. Unconstrained optimization, many local minima:



"Arbitrary optimization problems are hopeless, we always need some structure"

* Specifying the optimization problem:

$$\min_{x \in X} f(x) \quad (\text{P}) \leftarrow$$

① Vector space (where the optimization variables and the feasible set "live")

$(\mathbb{R}^d, \|\cdot\|)$: normed vector space; "primal space"

tells us that x is a vector in \mathbb{R}^d tells us how to measure distances in \mathbb{R}^d

$$x = \begin{bmatrix} x_1 \\ x_2 \\ \vdots \\ x_d \end{bmatrix}$$

* Most often, we will take $\|x\| = \|x\|_2 = \left(\sum_{i=1}^d |x_i|^2 \right)^{1/2} \leftarrow$
(Euclidean norm)

* We might sometimes also consider l_p norms!

$$\|x\|_p = \left(\sum_{i=1}^d |x_i|^p \right)^{1/p}, \quad p \geq 1$$

$$\boxed{\|x\|_1 = \sum_i |x_i|},$$

$$\boxed{\|x\|_\infty = \max_{1 \leq i \leq d} |x_i|}.$$

* We will use $\langle \cdot, \cdot \rangle$ to denote inner products.

standard inner product:

$$\langle x, y \rangle = x^T y = \sum_{i=1}^d x_i y_i$$

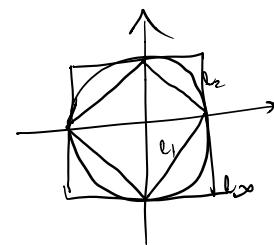
! When we work with $(\mathbb{R}^d, \|\cdot\|_p)$, view $\langle y, x \rangle$ as the value of a linear function y at x . So, if we are measuring x using the $\|\cdot\|_p$, we should measure the length of y using $\|\cdot\|_{p^*}$, where $\frac{1}{p} + \frac{1}{p^*} = 1$.

dual norm: $\|z\|_* = \sup_{\substack{\|x\| \leq 1}} \langle z, x \rangle$

$$\Rightarrow \forall z, x : \boxed{\langle z, x \rangle \leq \|z\|_* \cdot \|x\|}.$$

pf: Fix any two vectors x, z . Assume $x \neq 0, z \neq 0$, o.w. trivial. Define $\hat{x} = \frac{x}{\|x\|}$.

$$\|z\|_* \geq \langle z, \hat{x} \rangle = \frac{\langle z, x \rangle}{\|x\|}.$$



② Feasible set:

- specifies what solution points we are allowed to output

$x \subseteq \mathbb{R}^d$. If $x = \mathbb{R}^d$, we say that (P) is unconstrained.

O.w., we say that (P) is constrained.

x can be specified:

- as an abstract geometric body (a ball, a box, a polyhedron)

- via functional constraints: $g_i(x) \leq 0, i=1, 2, \dots, m$,

$$h_i(x) = 0, i=1, \dots, p$$

$$f_i(x) \geq c \quad g_i(x) = c - f_i(x)$$

E.g., $x = \underbrace{B_2(0, 1)}_{\text{unit Euclidean ball}}$

$$x = \{x \in \mathbb{R}^d : \|x\|_2 \leq 1\}$$

- * In this class, we will always assume that X is closed and convex.

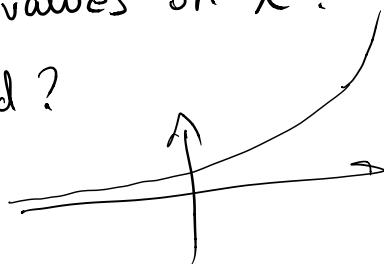
* Heine-Borel Thm: If X is closed and bounded, then it is compact [if $X \subseteq \bigcup_{\alpha \in A} U_\alpha$ for some family of open sets $\{U_\alpha\}$, then \exists a finite subfamily $\{U_{\alpha_i}\}_{i=1}^n$ s.t. $X \subseteq \bigcup_{1 \leq i \leq n} U_{\alpha_i}$]

* Weierstrass Extreme Value Theorem: If X is compact and f is a function that is defined and continuous on X , then f attains its extreme values on X .

* What if X is not bounded?

Consider $f(x) = e^x$

$$\inf_{x \in \mathbb{R}} f(x) = 0.$$

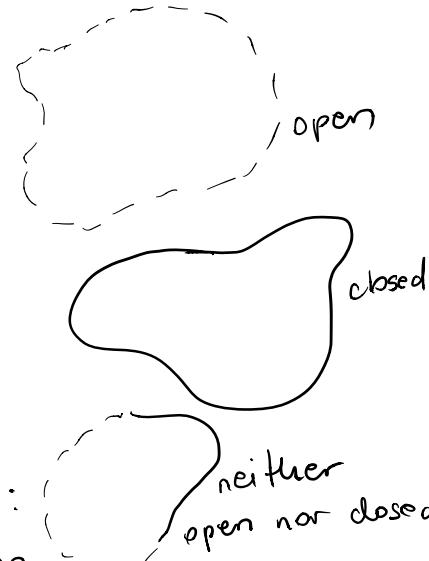
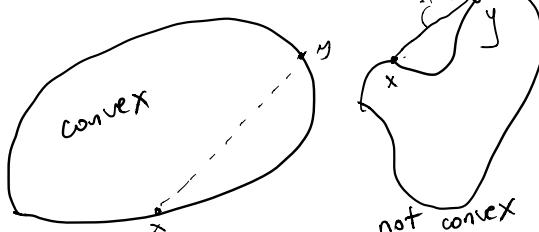


* When we work with unconstrained problems, we will normally assume that f is bounded below.

* Convex sets:

* Def. A set $X \subseteq \mathbb{R}^d$ is convex if

$$(\forall x, y \in X) (\forall \alpha \in (0, 1)): (1-\alpha)x + \alpha y \in X. \\ x + \alpha(y-x).$$



③ Objective function:

- "cost", "loss"
- Extended real valued functions: $f: \overline{\mathbb{Q}} \rightarrow \mathbb{R} \cup \{-\infty, +\infty\} \subseteq \bar{\mathbb{R}}$.
- We will define f on all of \mathbb{R}^d by assigning if value $+\infty$ at each point $x \in \mathbb{R}^d \setminus \mathcal{D}$.
- Effective domain: $\text{dom}(f) = \{x \in \mathbb{R}^d : f(x) < \infty\}$.
- "nonlinear opt" $\stackrel{?}{=}$ "continuous opt"
- * Lower semi continuous functions:
- * Def. A function $f: \mathbb{R}^d \rightarrow \bar{\mathbb{R}}$ is said to be lower semi continuous (l.s.c.) at $x \in \mathbb{R}^d$ if $\liminf_{y \xrightarrow{y \rightarrow x} x} f(y) \geq f(x)$.

f is l.s.c. on \mathbb{R}^d if it is l.s.c. at all $x \in \mathbb{R}^d$.

Ex. indicator of a closed set is l.s.c.

$$I_X(x) = \begin{cases} 0, & x \in X \\ \infty, & x \notin X \end{cases}$$

$$\min_{x \in X} f(x) \equiv \min_{x \in \mathbb{R}^d} \{f(x) + I_X(x)\}$$

* Unless we are abstracting away constraints, the least we will assume about f is that it is continuous.

* Def. $f: \mathbb{R}^d \rightarrow \bar{\mathbb{R}}$ is said to be:

i) Lipschitz-cont. on $X \subseteq \mathbb{R}^d$ if $\exists M < \infty$

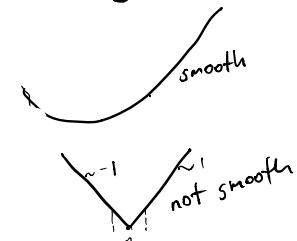
$$\forall x, y \in X : |f(x) - f(y)| \leq M \|x - y\|.$$

l.s.c.

2) Smooth on $X \subseteq \mathbb{R}^d$ if f 's gradients are Lipschitz-cont.'s, i.e., $\exists L < \infty$ s.t. $\forall x, y \in X$:

$$\|\nabla f(x) - \nabla f(y)\|_* \leq L \|x - y\|.$$

$$\nabla f(x) = \begin{bmatrix} \frac{\partial f}{\partial x_1} \\ \vdots \\ \frac{\partial f}{\partial x_d} \end{bmatrix}$$



* Def. $f : \mathbb{R}^d \rightarrow \bar{\mathbb{R}}$ is convex if $\forall x, y \in \mathbb{R}^d, \forall \alpha \in (0, 1)$:

$$f((1-\alpha)x + \alpha y) \leq (1-\alpha)f(x) + \alpha f(y).$$

* Ex. Function that is differentiable on its domain but not smooth:

$$f(x) = \frac{1}{x}$$

$\text{dom}(f) = \mathbb{R}_{++}$

