Homework 1: Answers

1.

1)Total number of normalized floating-point number:

There are β -1 possible choices for the leading digit d_0 , β choices for p-1 fractions, and U + L - 1 possible values for the exponent. And because the number could be zero, 1 should be added. There is a positive and negative sign in front of values which leads double of the total numbers.

Therefore, the total numbers of normalized floating-point number is

$$2(\beta-1)\beta^{p-1}(U+L-1)+1$$

2) Smallest positive normalized number:

For the smallest positive normalized number, the leading digit is 1 and the remaining digits of the mantissa are 0, and the exponent is smallest possible value

$$UFL = \beta^L$$

3) Largest floating-point number:

Largest floating-point number has $\beta - 1$ as the value for each digit of the mantissa and the largest possible value for the exponent.

$$OFL = (\beta - 1)(1 + \frac{1}{\beta} + \dots + \frac{1}{\beta^{p-1}})\beta^{U} = \beta^{U+1} \left(1 - \frac{1}{\beta^{p}}\right)$$

2 Exercise 1.4

a) Absolute value = $\sin(x + h) - \sin(x)$

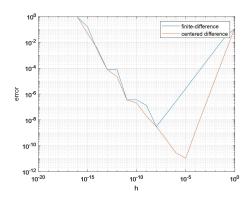
b) Relative value = $\frac{\sin(x+h) - \sin(x)}{\sin(x)}$

c) Condition number =
$$\frac{\frac{\sin(x+h) - \sin(x)}{\sin(x)}}{\frac{x+h-x}{x}} = \frac{|x\cos(x)|}{\sin(x)} = \frac{x}{\tan(x)}$$

d) When condition number is much larger or smaller than 1, i.e. x is near any integer multiple of π except 0, the problem is highly sensitive.

3 Programming

Code: refer to Q3 in zip file



As the figure shows, there is a minimum value error = 10^{-8} for the magnitude of the error when $h=10^{-8}$, which corresponds to the rule in Example 1.3. While for centered difference

approximation, there is a minimum value for the magnitude of the error when $h = 10^{-5}$.

4.

stopping criterion:

(1) $x^n/n! < epsilon$

(2)
$$[(x^n/n!) / (1+...+x^n(n-1)/(n-1)!)] < epsilon$$

Both (1) and (2) is right, it does not make contributions to the series sum. The answer (2) is more reasonable. More details can be found in Q4.mlx

5.

$$x = (x_1, ..., x_n)^T$$
 $||x||_{\infty} = \max_{0 \le i \le n} |x_i| = |x_j|$

For
$$||x||_2^2 = \sum_{i=1}^n |x_i|^2 \le (\sum_{i=1}^n |x_i|)^2 = ||x||_1^2$$

So: $||x||_2 \le ||x||_1$

For:
$$||x||_1^2 = (\sum_{i=1}^n |x_i|)^2 \le n \sum_{i=1}^n |x_i|^2 = n * ||x||_2^2$$

So: $||x||_1 \le \sqrt{n} ||x||_2$

So:
$$||x||_2 \le ||x||_1 \le \sqrt{n} ||x||_2$$

(2)

For:
$$||x||_{\infty}^2 = |x_j|^2 \le \sum_{i=1}^n |x_i|^2 = ||x||_2^2$$

So:
$$||x||_{\infty} \le ||x||_2$$

For :
$$||x||_2^2 = \sum_{i=1}^n |x_i|^2 \le \sum_{j=1}^n |x_j|^2 \le n |x_j|^2 = n||x||_{\infty}^2$$

So:
$$||x||_2 \le \sqrt{n} ||x||_{\infty}$$

So:
$$||x||_{\infty} \le ||x||_2 \le \sqrt{n} ||x||_{\infty}$$

6.

(a)

$$\mathbf{A} = \begin{bmatrix} 1 & 1 & 0 \\ 1 & 2 & 1 \\ 1 & 3 & 2 \end{bmatrix} \xrightarrow{c2-c1} \begin{bmatrix} 1 & 1 & 0 \\ 0 & 1 & 1 \\ 0 & 2 & 2 \end{bmatrix} \xrightarrow{c3-2*c2} \begin{bmatrix} 1 & 1 & 0 \\ 0 & 1 & 1 \\ 0 & 0 & 0 \end{bmatrix}$$

So: A is singular

(b)

$$[A|b] = \begin{bmatrix} 1 & 1 & 0 & 2 \\ 1 & 2 & 1 & 4 \\ 1 & 3 & 2 & 6 \end{bmatrix} \xrightarrow{c2-c1} \begin{bmatrix} 1 & 1 & 0 & 2 \\ 0 & 1 & 1 & 2 \\ 0 & 2 & 2 & 4 \end{bmatrix} \xrightarrow{c3-2*c2} \begin{bmatrix} 1 & 1 & 0 & 2 \\ 0 & 1 & 1 & 2 \\ 0 & 0 & 0 & 0 \end{bmatrix}$$

So: when $b = \begin{bmatrix} 2 & 4 & 6 \end{bmatrix}^T$, there are infinite solutions.

(c)
$$||A||_1 = \max_{j} \sum_{i=1}^{n} |a_{ij}| = 6$$

$$||A||_{\infty} = \max_{i} \sum_{j=1}^{n} |a_{ij}| = 6$$

(d)

For A is singular, so the condition number of A is infinite.

7.8.

More details can be found in Q7.mlx and Q8.mlx