

# Cálculo de Programas

## *Algebra of Programming*

Lic. Ciências da Computação (2º ano)  
Lic./Mest.Int. em Engenharia Informática (3º ano)  
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1. Considere o diagrama

*Consider the following diagram*

$$\begin{array}{ccc} & \xrightarrow{\text{assocr}} & \\ (A \times B) \times C & \cong & A \times (B \times C) \\ & \xleftarrow{\text{assocl}} & \end{array}$$

onde  $\text{assocl} = \langle id \times \pi_1, \pi_2 \cdot \pi_2 \rangle$ . Apresente justificações para o cálculo que se segue em que se resolve em ordem a  $\text{assocr}$  a equação  $\text{assocl} \cdot \text{assocr} = id$ :

*where  $\text{assocl} = \langle id \times \pi_1, \pi_2 \cdot \pi_2 \rangle$ . The reasoning below solves the equation  $\text{assocl} \cdot \text{assocr} = id$  for variable  $\text{assocr}$ . Fill in justifications for each step in the reasoning:*

$$\begin{aligned} & \text{assocl} \cdot \text{assocr} = id \\ \equiv & \quad \{ \dots\dots\dots \} \\ & \left\{ \begin{array}{l} (id \times \pi_1) \cdot \text{assocr} = \pi_1 \\ \pi_2 \cdot \pi_2 \cdot \text{assocr} = \pi_2 \end{array} \right. \\ \equiv & \quad \{ \dots\dots\dots \} \\ & \left\{ \begin{array}{l} \left\{ \begin{array}{l} \pi_1 \cdot \text{assocr} = \pi_1 \cdot \pi_1 \\ \pi_1 \cdot \pi_2 \cdot \text{assocr} = \pi_2 \cdot \pi_1 \end{array} \right. \\ \pi_2 \cdot \pi_2 \cdot \text{assocr} = \pi_2 \end{array} \right. \\ \equiv & \quad \{ \dots\dots\dots \} \\ & \left\{ \begin{array}{l} \pi_1 \cdot \text{assocr} = \pi_1 \cdot \pi_1 \\ \left\{ \begin{array}{l} \pi_1 \cdot \pi_2 \cdot \text{assocr} = \pi_2 \cdot \pi_1 \\ \pi_2 \cdot \pi_2 \cdot \text{assocr} = \pi_2 \end{array} \right. \end{array} \right. \\ \equiv & \quad \{ \dots\dots\dots \} \\ & \left\{ \begin{array}{l} \pi_1 \cdot \text{assocr} = \pi_1 \cdot \pi_1 \\ \pi_2 \cdot \text{assocr} = \langle \pi_2 \cdot \pi_1, \pi_2 \rangle \end{array} \right. \\ \equiv & \quad \{ \dots\dots\dots \} \\ & \text{assocr} = \langle \pi_1 \cdot \pi_1, \pi_2 \times id \rangle \end{aligned} \tag{F1}$$

2. (a) Codifique (F1) directamente em Haskell e verifique o comportamento dessa função no GHCi; (b) De seguida, converta — por igualdade extensional — (F1) para notação Haskell *pointwise* que não recorra a nenhum combinador nem projecção e verifique no GHCi que as duas versões dão os mesmos resultados.

(a) Encode (F1) directly in Haskell and check its behavior in GHCi; (b) Then convert — by extensional equality — (F1) to pointwise Haskell code dispensing with combinators or projections, and let GHCi check that both versions give the same results.

3. Recorde a propriedade universal do combinador  $[f, g]$ ,

Recall the universal property of the  $[f, g]$  combinator,

$$k = [f, g] \equiv \begin{cases} k \cdot i_1 = f \\ k \cdot i_2 = g \end{cases}$$

Demonstre a igualdade

Prove the equality

$$[\underline{k}, \underline{k}] = \underline{k} \quad (\text{F2})$$

recorrendo à propriedade universal acima e a uma lei que qualquer função constante  $\underline{k}$  satisfaz. (Ver no formulário.)

using the universal property given above and a law that any constant function  $\underline{k}$  satisfies. (Check the reference sheet.)

4. Os isomorfismos

Both isomorphisms of

$$A \times (B + C) \begin{array}{c} \xrightarrow{\text{distr}} \\ \cong \\ \xleftarrow{\text{undistr}} \end{array} A \times B + A \times C$$

estudados na aula teórica estão codificados na biblioteca *Cp.hs*. Supondo  $A = \text{String}$ ,  $B = \mathbb{B}$  e  $C = \mathbb{Z}$ , (a) aplique no GHCi `undistr`, alternativamente, aos pares `("CP", TRUE)` ou `("LEI", 1)`; (b) verifique que  $(\text{distr} \cdot \text{undistr}) x = x$  para essas (e quaisquer outras) situações que possa testar.

studied in the theory class are encoded in the library *Cp.hs*. Assuming that  $A = \text{String}$ ,  $B = \mathbb{B}$  and  $C = \mathbb{Z}$ , (a) apply in GHCi `undistr`, alternatively, to the pairs `("CP", TRUE)` or `("LEI", 1)`; (b) check that  $(\text{distr} \cdot \text{undistr}) x = x$  for these (and any other) situations you can test.

5. Recorde a função

Recall

$$\alpha = [\langle \text{FALSE}, \text{id} \rangle, \langle \text{TRUE}, \text{id} \rangle]$$

da ficha anterior. Mostre, usando a propriedade *universal-+*, que  $\alpha$  se pode escrever em Haskell da forma seguinte:

from the previous exercise sheet. Show, using the *+universal* law, that  $\alpha$  can be written in pointwise Haskell as follows:

$$\begin{aligned} \alpha (i_1 a) &= (\text{FALSE}, a) \\ \alpha (i_2 a) &= (\text{TRUE}, a) \end{aligned}$$

Codifique  $\alpha$  e teste-a no GHCi, onde  $i_1$  (resp.  $i_2$ ) se escreve `Left` (resp. `Right`).

Encode  $\alpha$  and test it on GHCi.

6. Recorra às leis dos coprodutos para mostrar que a definição que conhece da função factorial,

*Show by coproduct laws that the usual definition of the factorial function,*

$$\begin{aligned} fac\ 0 &= 1 \\ fac\ (n + 1) &= (n + 1) * fac\ n \end{aligned}$$

é equivalente à equação seguinte

*is equivalent the following equation,*

$$fac \cdot [\underline{0}, succ] = [\underline{1}, mul] \cdot \langle succ, fac \rangle$$

onde

*where*

$$\begin{aligned} succ\ n &= n + 1 \\ mul\ (a, b) &= a * b \end{aligned}$$

7. A função  $in = [\underline{0}, succ]$  da questão anterior exprime, para  $succ\ n = n + 1$ , a forma como os números naturais ( $\mathbb{N}_0$ ) são gerados a partir do número 0, de acordo com o diagrama seguinte:

*Function  $in = [\underline{0}, succ]$  in the previous question (where  $succ\ n = n + 1$ ) expresses the way natural numbers ( $\mathbb{N}_0$ ) are generated from the number 0, according to the diagram below:*

$$\begin{array}{ccccc} 1 & \xrightarrow{i_1} & 1 + \mathbb{N}_0 & \xleftarrow{i_2} & \mathbb{N}_0 \\ & \searrow \underline{0} & \downarrow in = [\underline{0}, succ] & \swarrow succ & \\ & & \mathbb{N}_0 & & \end{array} \quad (F3)$$

Sabendo que o tipo 1 coincide com o tipo () em Haskell e é habitado por um único elemento, também designado por (), calcule a inversa de  $in$ ,

*Knowing that type 1 matches type () in Haskell and is inhabited by a single element, also denoted by (), find the inverse of  $in$ ,*

$$\begin{cases} out\ 0 = i_1\ () \\ out\ (n + 1) = i_2\ n \end{cases} \quad (F4)$$

resolvendo em ordem a  $out$  a equação

*by solving the equation*

$$out \cdot in = id \quad (F5)$$

e introduzindo variáveis.

*for  $out$  and adding variables.*

8. Verifique no GHCi que a seguinte função

*On GHCi check that function*

$$fac = [\underline{1}, mul] \cdot (id + \langle succ, fac \rangle) \cdot out$$

a que corresponde o diagrama

*captured by diagram*

$$\begin{array}{ccc} \mathbb{N}_0 & \xrightarrow{out} & 1 + \mathbb{N}_0 \\ fac \downarrow & & \downarrow id + \langle succ, fac \rangle \\ \mathbb{N}_0 & \xleftarrow{[\underline{1}, mul]} & 1 + \mathbb{N}_0 \times \mathbb{N}_0 \end{array}$$

calcula o factorial da sua entrada, assumindo out (F4) e mul  $(a, b) = a * b$  já definidas.

*computes de factorial of its input, assuming out (F4) and mul  $(a, b) = a * b$  already defined.*

9. **Questão prática** — Este problema não irá ser abordado em sala de aula. Os alunos devem tentar resolvê-lo em casa e, querendo, publicarem a sua solução no canal **#geral** do Slack, com vista à sua discussão com colegas. Os requisitos do problema são dados abaixo. **NB:** usa-se a notação  $X^*$  para designar o tipo  $[X]$  em Haskell.

**Open assignment** — This assignment will not be addressed in class. Students should try to solve it at home and, wishing so, publish their solution in the **#geral** Slack channel, so that it can be discussed among colleagues. The requirements of the problem are given below. **NB:** notation  $X^*$  is used to denote the type  $[X]$  in Haskell.

**Problem requirements:**

*The automatic generation of bibliographies in the  $\text{\LaTeX}$  text preparation system is based bibliographic databases from which the following information can be extracted:*

$$\text{Bib} = (\text{Key} \times \text{Aut}^*)^*$$

*It associates authors (Aut) to citation keys (Key).*

*Whenever  $\text{\LaTeX}$  processes a text document, it compiles all occurrences of citation keys in an auxiliary file*

$$\text{Aux} = (\text{Pag} \times \text{Key}^*)^*$$

*associating pages (Pag) to the citation keys that occur in them.*

*An **author index** is an appendix to a text (e.g. book) indicating, in alphabetical order, the names of authors mentioned and the ordered list of pages where their works are cited, for example:*

*Arbib, M. A. – 10, 11*

*Bird, R. – 28*

*Horowitz, E. – 2, 3, 15, 16, 19*

*Hudak, P. – 11, 12, 29*

*Jones, C. B. – 3, 7, 28*

*Manes, E. G. – 10, 11*

*Sahni, S. – 2, 3, 15, 16, 19*

*Spivey, J.M. – 3, 7*

*Wadler, P. – 2, 3*

*The above structure can be represented by the type*

$$\text{Ind} = (\text{Aut} \times \text{Pag}^*)^*$$

*listing authors (Aut) and the respective pages where they are mentioned (Pag).*

*Write a Haskell function  $\text{mkInd} : \text{Bib} \times \text{Aux} \rightarrow \text{Ind}$  that generates author indices (Ind) from Bib and Aux.*

**Important:** Structure your solution across the  $f \cdot g$ ,  $\langle f, g \rangle$  and  $f \times g$  combinators that can be found in library *Cp.hs*. Use **diagrams** to plan your proposed solution, which should avoid re-inventing functions over lists already available in the Haskell standard libraries.

□