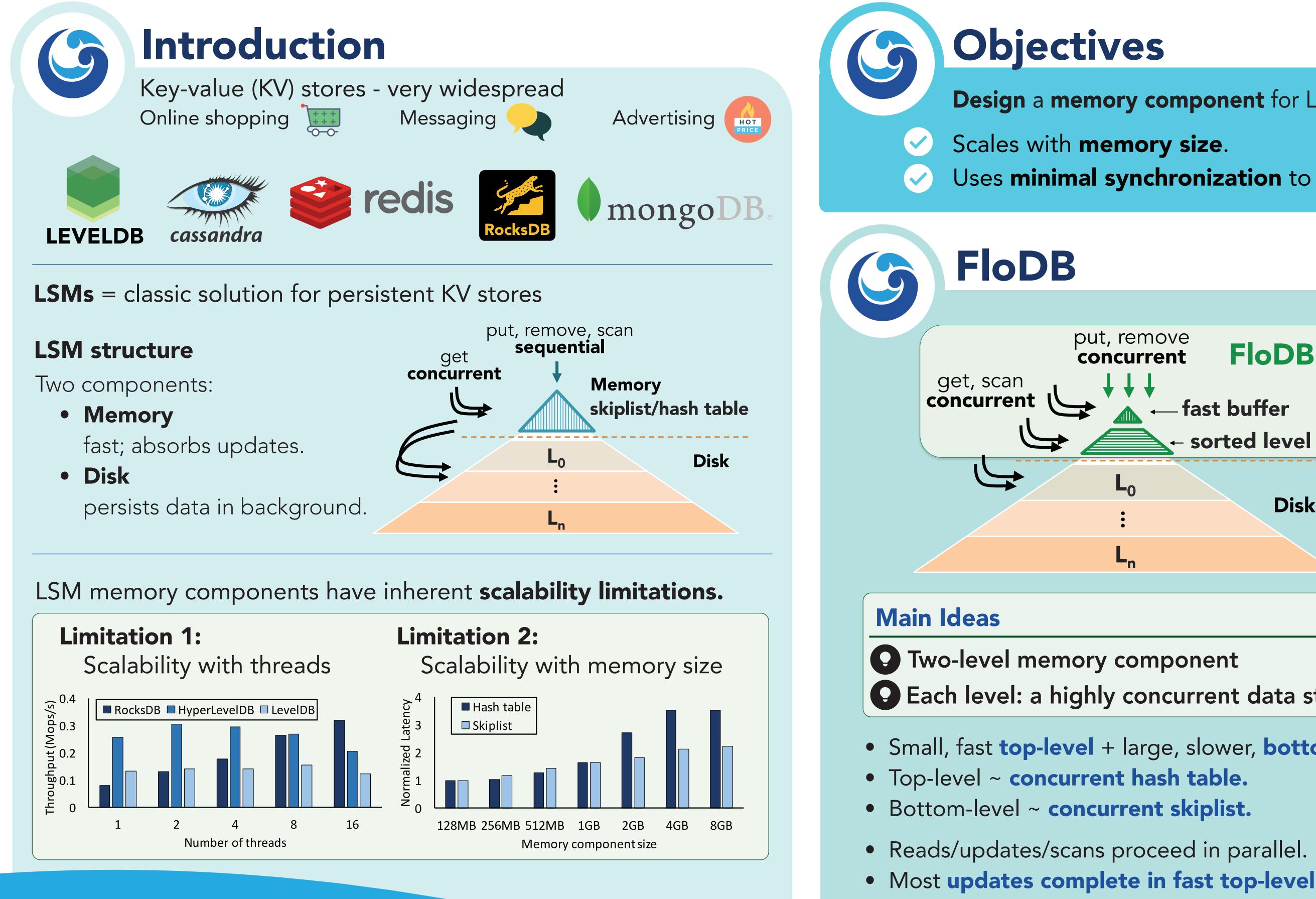
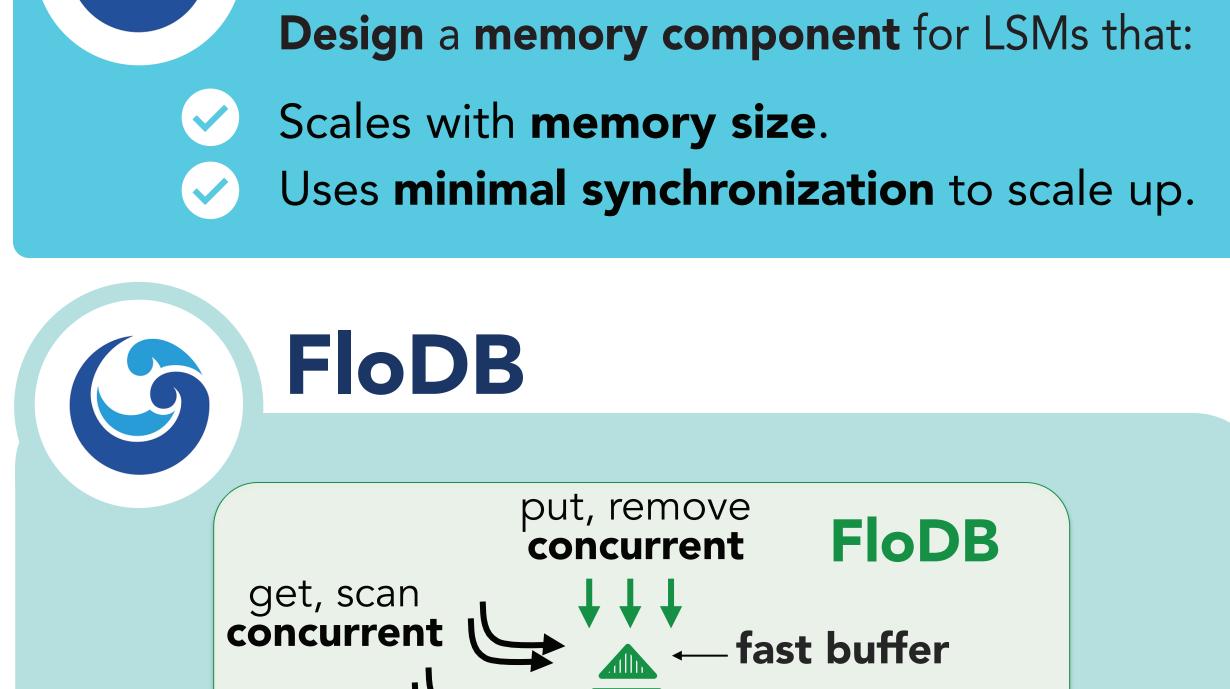
## FloDB: Unlocking Memory in Persistent Key-Value Stores

Oana Balmau Rachid Guerraoui Vasileios Trigonakis Igor Zablotchi





- Two-level memory component
- Each level: a highly concurrent data structure

Disk

- Small, fast top-level + large, slower, bottom-level.

- Most updates complete in fast top-level.
- Can increase memory size w/o hurting performance.
- Data **flows**: top-level  $\rightarrow$  bottom-level  $\rightarrow$  disk.

## Results **Read-only** Write-only Mixed ■ FloDB ■ FloDB ■ FloDB ■ RocksDB ■ RocksDB ■ RocksDB Throughput (Mops/s) 0 0 0 0 0 0 0 Throughput (Mops/s) ■ RocksDB/cLSM ■ RocksDB/cLSM ■ RocksDB/cLSM Average persistence □ HyperLevelDB throughput ■ LevelDB ■ LevelDB ■ LevelDB 16 32 64 128 16 16 Number of threads Number of threads Number of threads Skewed workload (98%–2%) and varying memory size ■ FloDB put (Mops/s) ■ RocksDB ■ RocksDB/cLSM Through ■ LevelDB 128MB 256MB 512MB 1GB 2GB 8GB 16GB 32GB 64GB 128GB 192GB 4GB Memory component size FloDB outperforms state-of-the-art LSMs in write-intensive scenarios and scales with memory size.

Find out more in our paper:

O. Balmau, R. Guerraoui, V. Trigonakis, and I. Zablotchi. FloDB: Unlocking Memory in Persistent Key-Value Stores. In Proceedings of the 12th European Conference on Computer Systems (EuroSys '17) Check out our website: lpd.epfl.ch/site/flodb

