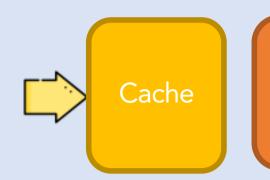
Log-Free Concurrent Data Structures

Tudor David¹, Aleksandar Dragojević², Rachid Guerraoui³, Igor Zablotchi³ ¹IBM Research Zurich, ²MSR Cambridge, ³EPFL

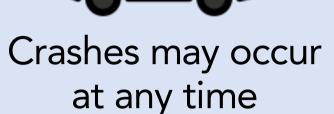
Introduction

Designing fast data structures for NVRAM: challenging









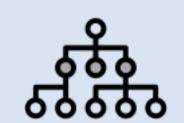


Stores may not reach NVRAM in program order



Need expensive instructions for ordering & persistence

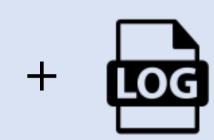
Previous approaches: transactions & logging



Writes might not

reach NVRAM

immediately



1. Log intention

2. Apply to data structure

3. Log completion



easy to use

high overhead (need to wait for log)

Our objective

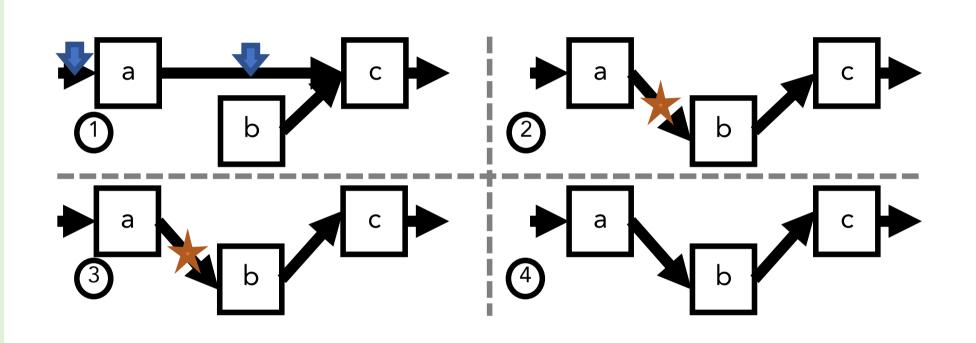
© Design durable concurrent data structures that do not perform any logging in the common case.

Our insights

- Description Lock-free algorithms are good candidates: they are always in a consistent (recoverable) state.
- Pointer marking can ensure store atomicity w/o logging
- Batching write-backs can improve throughput
- © Locality in allocation & deallocation can be exploited to reduce or eliminate logging from memory management

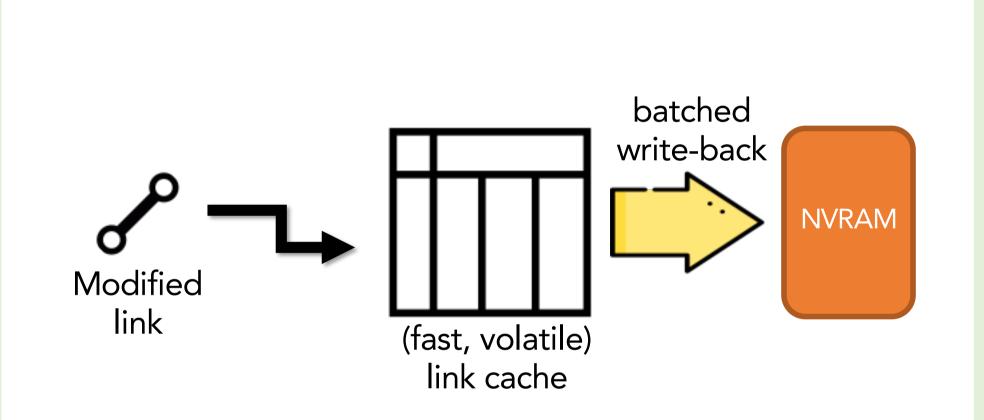
Our Techniques

Link-and-Persist



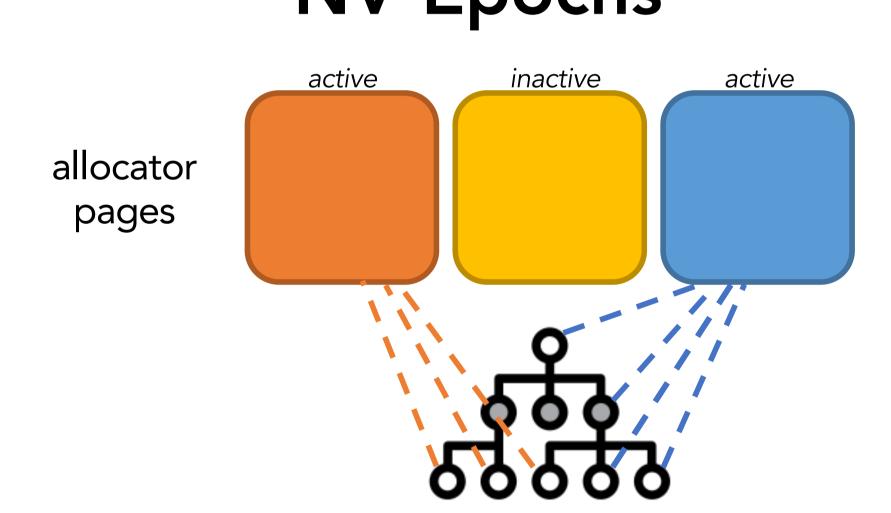
- 1. Prepare new node & persist dependent links
- 2. Link new node & add mark (★)
- 3. Make the new link persistent
- 4. Remove mark

Link Cache



^Q Cache modified data structure links and write them to NVRAM in batches \rightarrow better performance.

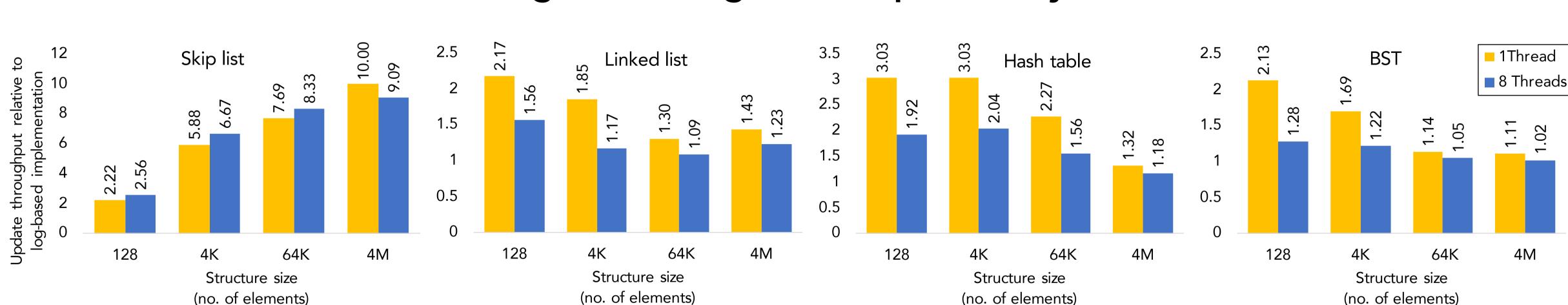
NV-Epochs



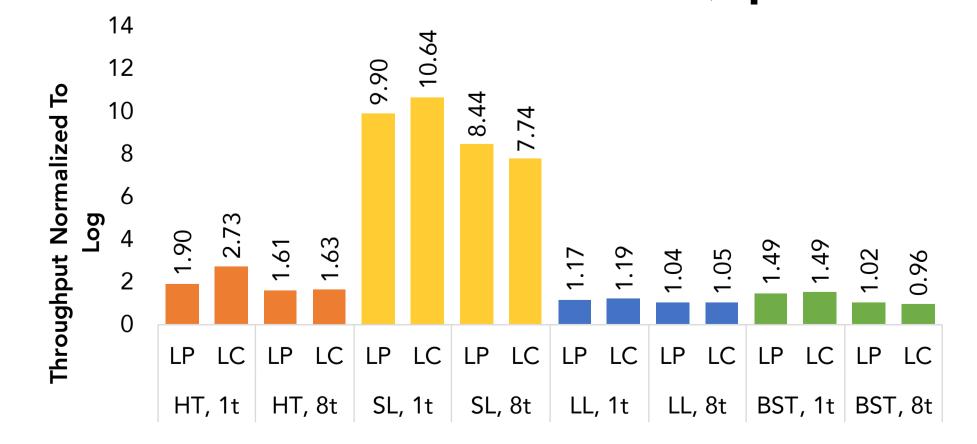
- Instead of logging each allocation/deallocation, keep track of active pages
- When there is locality, no need to log at all

Results

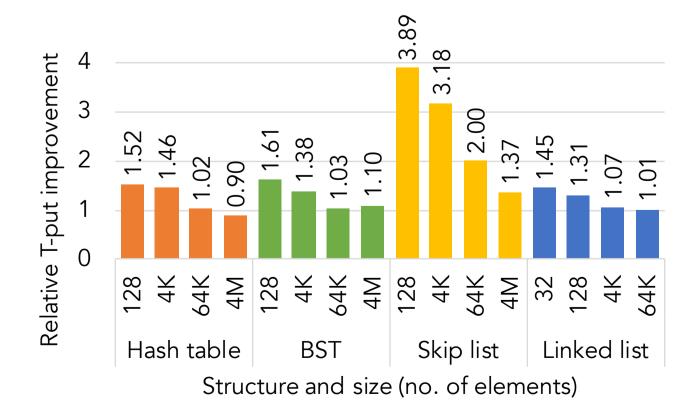
Log-free vs. log-based (update-only)



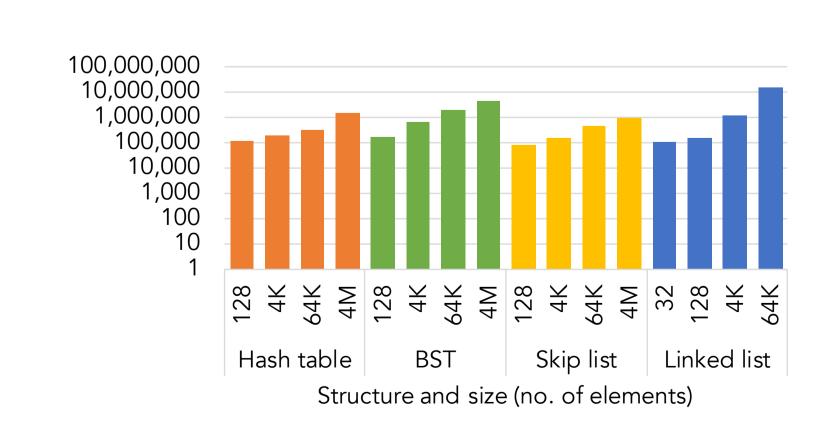
Link-and-Persist & Link Cache (update-only)



NV-Epochs (update-only)



Recovery time (ns)



Find out more in our paper & on our website:

T. David, A. Dragojević, R. Guerraoui, I. Zablotchi. Log-Free Concurrent Data Structures. 2018 USENIX Annual Technical Conference. https://lpd.epfl.ch/site/nvram

