

Analysis, Big O and Growth of Functions

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1 Book Keeping

- Reading posted
- Lab 1 available

2 Analysis of Algorithms

Problem: a general description of input parameters and the properties that an optimal solution should have

Instance: a specific example of a problem with all parameters specified

- Example: Given a weighted graph, find the cheapest Hamiltonian Cycle (TSP)
- A "problem" can have many instances

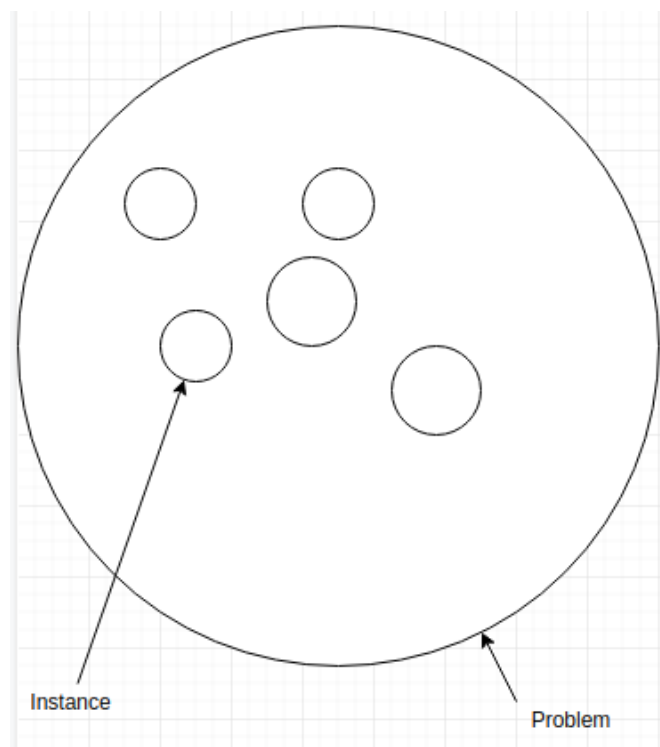


Figure 1: instance_problem

- An algorithm solves all instances of problem

- Many algorithms, what is most efficient?
- What is efficient?
 - Memory
 - Time
 - CPU cycles
 - Disk Space
 - I/O bandwidth
 - Power
- Efficiency usually defined as using smallest time
- Index runtimes by instance size
- "Instance Size" not always well defined - can have multiple params (edges, nodes)

3 Example: Insertion Sort

INSERTION-SORT(<i>A</i>)	<i>cost</i>	<i>times</i>
1 for $j \leftarrow 2$ to $\text{length}[A]$	c_1	n
2 do $\text{key} \leftarrow A[j]$	c_2	$n - 1$
3 ▷ Insert $A[j]$ into the sorted sequence $A[1..j-1]$.	0	$n - 1$
4 $i \leftarrow j - 1$	c_4	$n - 1$
5 while $i > 0$ and $A[i] > \text{key}$	c_5	$\sum_{j=2}^n t_j$
6 do $A[i+1] \leftarrow A[i]$	c_6	$\sum_{j=2}^n (t_j - 1)$
7 $i \leftarrow i - 1$	c_7	$\sum_{j=2}^n (t_j - 1)$
8 $A[i+1] \leftarrow \text{key}$	c_8	$n - 1$

Figure 2: Cor p.24

- Best case: already sorted. $T(j) = 1, T(n) = an + b \rightarrow$ linear
- Worst case: reverse sorted: $T(j) = j, T(n) = \frac{n(n+1)}{2} \approx an^2 + bn + c \rightarrow$ quadratic

Time Complexity Function: The largest amount of time for an algorithm needed to solve the problem for a given instance size.

- Even Time-Complexity function considered too complicated for daily use
- Asymptotic notation used instead

4 Asymptotic Notation

For a given function $g(n)$, $O(g(n)) = f(n)$ there exist positive constants k and n_0 such that $f(n) \leq K g(n)$ for all $n \geq n_0$

Less formally: $O(g(n))$ is the set of functions that are asymptotically less than $g(n)$ for large n .

Example

I claim that $f(n) = an^2 + bn + c = O(n^2)$. If so, then there should exist positive constants k and n_0 such that

$$an^2 + bn + c \leq kn^2$$
$$a + b/n + \frac{c}{n^2} \leq k$$

$$k = a + 1$$

n_0 is intersection