

Homework #2

Solve the following problems.

1. Define the following terms:
 - (a) the complexity class NPC , and
 - (b) a reduction from problem P_1 to problem P_2 .
2. Argue that each of the following problems is or is not in NP (note that this is a tricky problem). A tour is a Hamiltonian Cycle.
 - (a) Given a weighted graph $G = (V, E)$ and an integer k , is there at least one tour with cost $\leq k$?
 - (b) Given a weighted graph $G = (V, E)$ and an integer k , is the number of tours with cost $\leq k$ equal to one million?
 - (c) Given a weighted graph $G = (V, E)$ and an integer k , is the number of tours with cost $\leq k$ greater than or equal to one million?
 - (d) Given a weighted graph $G = (V, E)$ and an integer k , does the one millionth-shortest tour have cost $\leq k$? (a tour t is the n -shortest tour if the number of tours whose cost is less than the cost of t is $n - 1$. Assume that no two different tours have the same cost.)
 - (e) Given a weighted graph $G = (V, E)$ and an integer k , does the one millionth-shortest tour have cost $> k$? Assume that no two different tours have the same cost.
3. Prove that $HC \leq TSP$.
4. Prove that $SP \leq SS$.
5. Prove that for decision problems A_1 and A_2 such that $A_1 \leq A_2$, if $A_1 \notin P$ then $A_2 \notin P$.
6. Prove that for decision problems A_1 and A_2 , if $A_1 \leq A_2$, $A_1 \in NPC$, and $A_2 \in NP$, then $A_2 \in NPC$.
7. Prove that for decision problems A_1 and A_2 , if $A_1, A_2 \in NPC$, then $A_1 \leq A_2$ and $A_2 \leq A_1$.
8. Prove that for decision problems A_1 and A_2 , if $A_1, A_2 \in NPC$ and $A_1 \in P$ then $A_2 \in P$.
9. Why is it considered unlikely that any NP-complete problem is tractable? Use the definition of NPC to give a precise argument.
10. Prove the following or give a counter example. Given decision problems $\Pi(D, Y)$ and $\Pi'(D', Y')$ such that Π' is a subproblem of Π , if Π' is NP-complete, then Π is NP-complete. Assume $\Pi \in NP$.