## **Assignment 2**

## 1.1 Are either $\lceil \lg n \rceil!$ or $\lceil \lg \lg n \rceil!$ polynomially bounded?

Polynomially bounded means  $f_n = O(n^k)$  for some constant k (e.g., whether  $f_n \le c \cdot n^k$  for constants c and k as n approaches  $\infty$ ). For the first function  $\lceil \lg n \rceil!$ , without loss of generality, assume  $n = 2^a$  (where  $a \in \mathbb{N}$ ).

$$\lceil \lg n \rceil! \le c \cdot n^k$$
$$\lg(2^a)! \le c \cdot (2^a)^k$$
$$a! \le c \cdot 2^{ak}$$

The statement  $a! \le c \cdot 2^{ak}$  is a contradiction, as the factorial function a! is not exponentially bounded. Therefore,  $\lceil \lg n \rceil !$  is not polynomially bounded (via proof by contradiction). For the second function  $\lceil \lg \lg n \rceil !$ , without loss of generality, assume  $n = 2^{2^a}$  (where  $a \in \mathbb{N}$ ).

$$\begin{split} \lceil \lg \lg n \rceil! &\leq c \cdot n^k \\ \lg \lg \left( 2^{2^a} \right)! &\leq c \cdot \left( 2^{2^a} \right)^k \\ a! &\leq c \cdot 2^{k \cdot 2^a} \\ 1 \cdot 2 \cdot 3 \cdots a &\leq c \cdot \left( 2^{2k} \cdot 2^{4k} \cdot 2^{8k} \cdots 2^{2^a \cdot k} \right) \end{split}$$

The statement  $1 \cdot 2 \cdot 3 \cdots a \leq c \cdot (2^{2k} \cdot 2^{4k} \cdot 2^{8k} \cdots 2^{2^{a_k}})$  is obviously true. Therefore  $\lceil \lg \lg n \rceil!$  is polynomially bounded (via direct proof).

- 1.2 Use induction to prove  $F_i = \frac{\phi^i \hat{\phi}^i}{\sqrt{5}}$ ; where  $F_i = F_{i-2} + F_{i-1}$ , and  $\phi$  is the golden ratio  $\frac{1+\sqrt{5}}{2}$ .
- 1.3 Show that  $k \lg k = \Theta(n)$  implies  $k = \Theta(\frac{n}{n \ln n})$ .
- 1.4 Are either  $2^{n+1}$  or  $2^{2n}$  big-O of  $2^{n}$ ?
- 1.5 For each pair of functions (A,B), indicate whether A is  $O,o,\Omega,\omega$ , or  $\Theta$  of B. Assume  $k\geq 1,\ \epsilon>0,\ c>1$  are constants.

	$\boldsymbol{A}$	B	0	o	Ω	ω	Θ
a.	$\lg^k n$	$n^{\epsilon}$	yes	yes			
b.	$n^k$	$c^n$	yes	yes			
c.	$\sqrt{n}$	$n^{\sin n}$					
d.	$2^n$	$2^{n/2}$			yes	yes	
e.	$n^{\lg c}$	$c^{\lg n}$	yes		yes		yes
f.	$\lg(n!)$	$\lg(n^n)$	yes		yes		yes

1.6 Order the following functions such that  $f_1 = \Omega(f_2), f_2 = \Omega(f_3), ..., f_{29} = \Omega(f_{30})$ , and partition them into equivalence classes such that each function is big- $\Theta$  of each other.

$$2^{2^{n+1}} = \Omega(2^{2^n}),$$

$$2^{2^n} = \Omega((n+1)!),$$

$$(n+1)! = \Omega(n!),$$

$$n! = \Omega(e^n),$$

$$e^n = \Omega(n \cdot 2^n),$$

$$n \cdot 2^n = \Omega(2^n),$$

$$2^n = \Omega(\left(\frac{3}{2}\right)^n),$$

$$\left(\frac{3}{2}\right)^n = \Omega(\log \log n),$$

$$n^{\lg \lg n} = \Omega((\lg n)^{\lg n}),$$

$$(\lg n)^{\lg n} = \Omega((\lg n)!),$$

$$(\lg n)! = \Omega(N^3),$$

$$N^3 = \Omega(n^2),$$

$$n^2 = \Omega(4^{\lg n}),$$

$$4^{\lg n} = \Omega(\lg(n!)),$$

$$\lg(n!) = \Omega(n \lg n),$$

$$n \lg n = \Omega(2^{\lg n}),$$

$$2^{\lg n} = \Omega(n),$$

$$n = \Omega(\left(\sqrt{2}\right)^{\lg n}),$$

$$2^{\lg n} = \Omega(\ln n),$$

$$\ln n = \Omega(2^{\log n}),$$

$$2^{\log n} = \Omega(\ln n),$$

$$\ln n = \Omega(\sqrt{\lg n}),$$

$$2^{\log n} = \Omega(\ln n),$$

$$\ln n = \Omega(\sqrt{\lg n}),$$

$$2^{\lg^n} = \Omega(\ln n),$$

$$\ln n = \Omega(2^{\lg^n}),$$

$$2^{\lg^n} = \Omega(\log \ln n),$$

$$\log (\log n) = \Omega(\log (\log n)),$$

$$\log (\log n) = \Omega(\log (\log n)),$$

$$\log (\log n) = \Omega(\log (\log n)),$$

$$n^{\frac{1}{\lg n}}$$