low local 3 4 ROP local 2 E RBP old RBP 1et addr 4 ROP org 2 4 ROP local 2 local 2 local 2 Callee local (local (E RBP E RBP OH RBP OU RBP OH RBP OH RBP Ret addr Ret addr Ret addr Ret addr ary 1 arg 2 (d) org 2 GOAL: Slot in ms Idea: just slot in in He old an dors, new args & ignore Keeping some retaddr fle extra reserved space & saved RBP args at RBP + 8 (1+1) Note: can use more locals, so reset RSP to RBP before jumping and first thing - reserve space - we get to scribble over old locals * or some cycle detection thing ... por not worth layementation note if any of the values we are slotting in as new ags are flemselves coming from argumn's (specifically, ones we are overwhim), we have to push -> pop into place * TODO: how to deck this safety condition?

Compiling a decl (basic)	Muchion	Compiling a push any n	Anction	call (basic)
START. push RBP mov RBP, sub RSP, END: mov RSP, R pop RBP ret	RSP 8 + N C (pority -aljusted deepest stack for		8* M & G	ileer off ass tatget pushed
	, RSP HERST H	Cali	il call)	
Returning (bax mov RSP, pop RBP ret	RBP mov	bar body	estoe und tack space before taking fail Call,	
breaks w/	lambdas 60 hime elecch	e ur don't	Coo wh	attly Call

Scenario: - F (nomal) calls G, pulling on N ags low (no reg.s CC) - G (tail) calls H, w/ greater arity (M) - if we shifted everthing up how when we return, Forly adds N to RSP, so (M-N) ags are still there Idea 1: pad all Anction calls I max called arity - G would have to be called w/ M arg slots, the bottom (M-N) of which are padding until H slots things in there 4 ROP local 2 pragmatics: lots of stack space usage local ((always max possible, along any path of calls) E RBP OH RBP implementation: would basically have to Ret addr construct a directed graph of anchon calls ary - find all SCCs -> max will SCC -> build DAG of SCCs - Could be hard to compute w/o mutation (should be) Idea 2: Change He calling convention - we need to know smed RSP, RBP, and return address - evoything else can be calculated from those - both parametes and locals will be above RBP he wont be able to use call Mayere could use 3 calling comertions: - one for extragal calls - one for internal calls - one for intenal tail calls what if we have a Metion being called in both tail & non-tail Problem: whose do we find args?

Callee responsibility: - load in locals low - Stash He old RBP - return to ret address - callee-sure reg.s. - RSP, RBP Caller responsibility. - load in orgs - load in Net address ← ROP ← RSP local 2 - caller-save registers loca (2 local 1 local 1 ERBP OH RBP are (Ret addr arg 2 6 RBP ary I 1. Load new ares onto stack old RBP arg 2 Ret addr - Here may use ago or locals, and may overnite ass solely (ven) or locals -> crayers le that does this fix? Compiling a function call (basic) push return address reposely hard to do [Jurino 1002]]

push org_M Strick looks like rsp - now, when He Callee returns, He entire stack ars 1 (no longer have push ag_2 (locals AND ons) og Z To worm about missing the bottom ← RBP get paired down old RBP push an-1 M-Ni orgs) jump foo ret sub RSP, 8+ N damn tail call... mov RSP, RBP ANS: just reset to RBP + anity pop RBP of next fuetion ret