Java and CII & Course Wrap-Up

CSE 351 Spring 2019

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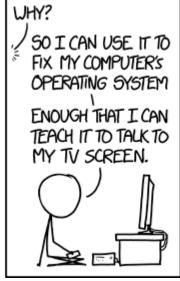
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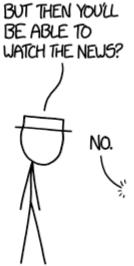
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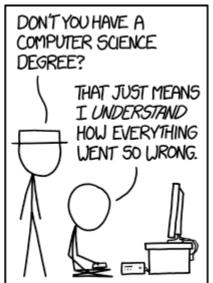
Jack Skalitzky

Sam Wolfson









https://xkcd.com/1760/

Administrivia

- Lab 5, due TONIGHT, Friday (6/7)
 - Memory Allocation
 - Recommended that you watch the Lab 5 helper videos
 - Sunday 6/9 is last day Lab 5 may be submitted (if one late day is used)

- ❖ Final Exam: Wed, 6/12, 12:30-2:20 pm in KNE 130
 - Review session Tuesday June 11, 3-6pm in ECE 105
 - Check course calendar for office hours for next week
- Course evaluations now open, please fill out!

Today

- Finish Java & C
- End-to-end Review
 - What happens after you write your source code?
 - How code becomes a program
 - How your computer executes your code
- Review of high-level concepts & course themes
 - More useful for "5 years from now" than "next week's final"

C: The Low-Level High-Level Language

- C is a "hands-off" language that "exposes" more of hardware (especially memory)
 - Weakly-typed language that stresses data as bits
 - Anything can be represented with a number!
 - Unconstrained pointers can hold address of anything
 - And no bounds checking buffer overflow possible!
 - Efficient by leaving everything up to the programmer

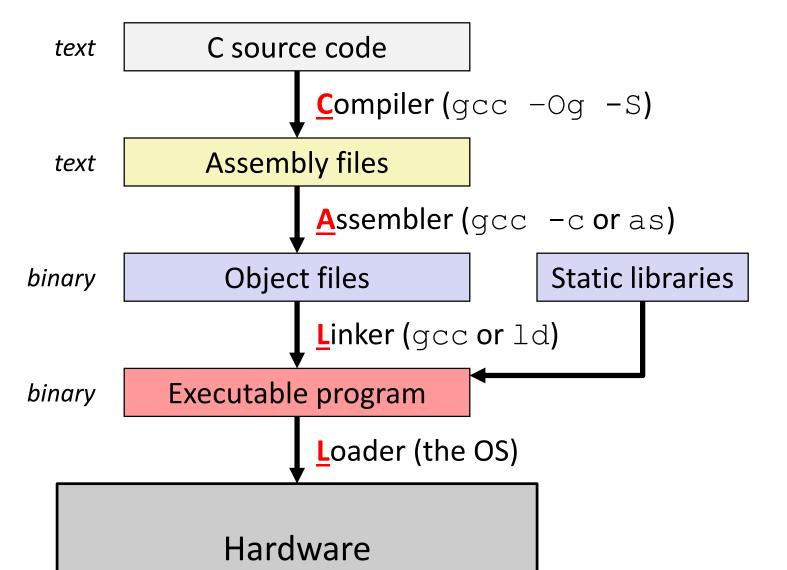
C Data Types

- C Primitive types
 - Fixed sizes and alignments
 - Characters (char), Integers (short, int, long), Floating Point (float, double)
- C Data Structures
 - Arrays contiguous chunks of memory
 - Multidimensional arrays = still one continuous chunk, but row-major
 - Multi-level arrays = array of pointers to other arrays
 - Structs structured group of variables
 - Struct fields are ordered according to declaration order
 - Internal fragmentation: space between members to satisfy member alignment requirements (aligned for each primitive element)
 - External fragmentation: space after last member to satisfy overall struct alignment requirement (largest primitive member)

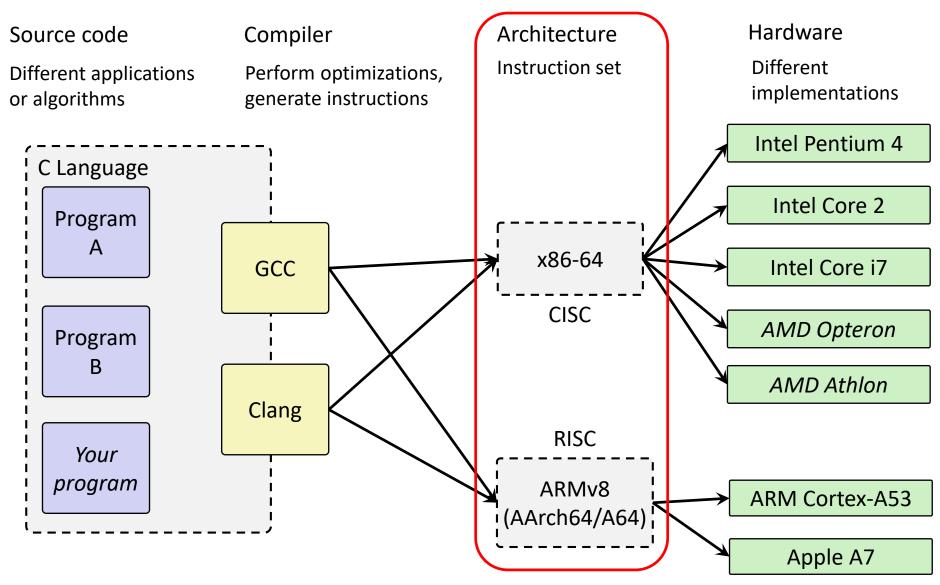
C and Memory

- Using C allowed us to examine how we store and access data in memory
 - Endianness (only applies to memory)
 - Is the first byte (lowest address) the least significant (little endian) or most significant (big endian) of your data?
 - Array indices and struct fields result in calculating proper addresses to access
- Consequences of your code:
 - Affects performance (locality)
 - Affects security
- But to understand these effects better, we had to dive deeper...

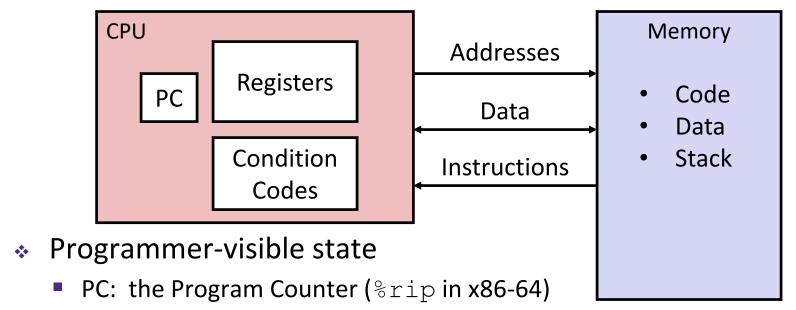
How Code Becomes a Program



Instruction Set Architecture



Assembly Programmer's View

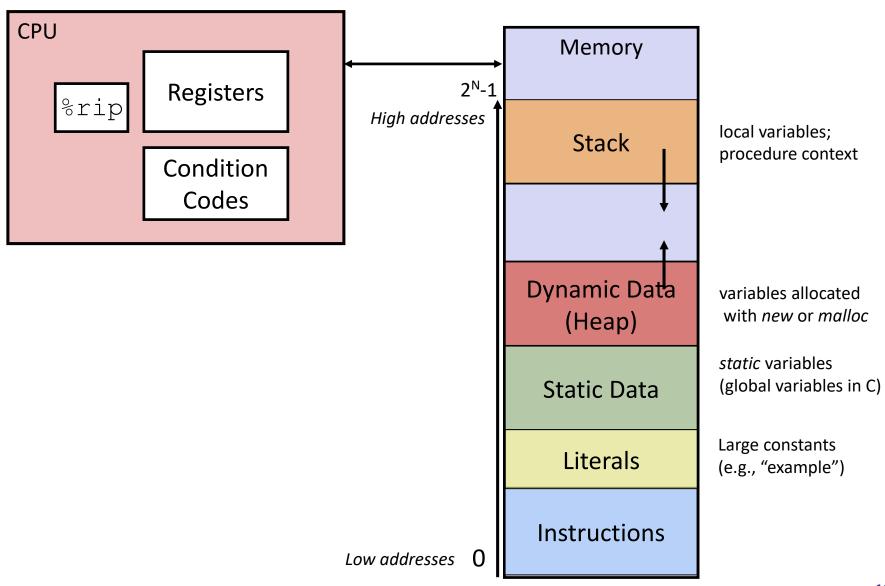


- Address of next instruction
- Named registers
 - Together in "register file"
 - Heavily used program data
- Condition codes
 - Store status information about most recent arithmetic operation
 - Used for conditional branching

Memory

- Byte-addressable array
- Huge virtual address space
- Private, all to yourself...

Program's View



Program's View

Instructions

- Data movement
 - mov, movz, movz
 - push, pop
- Arithmetic
 - add, sub, imul
- Control flow
 - · cmp, test
 - jmp, je, jgt, ...
 - call, ret

Operand types

- Literal: \$8
- Register: %rdi, %al
- Memory: D(Rb,Ri,S) = D+Rb+Ri*S
 - lea: not a memory access! Low addresses

Memory $2^{N}-1$ High addresses Stack Dynamic Data (Heap) Static Data Literals

Instructions

local variables; procedure context

variables allocated with *new* or *malloc*

static variables (global variables in C)

(e.g., "example")

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Program's View

- **Procedures**
 - **Essential abstraction**
 - Recursion...
- Stack discipline
 - Stack frame per call
 - Local variables
- Calling convention
 - How to pass arguments
 - Diane's Silk Dress Costs \$89
 - How to return data
 - Return address
 - Caller-saved / callee-saved registers

Memory $2^{N}-1$ High addresses local variables; Stack procedure context Dynamic Data variables allocated with *new* or *malloc* (Heap) static variables (global variables in C) Static Data Large constants Literals (e.g., "example") **Instructions**

High addresses

 $2^{N}-1$

Program's View



- Variable size
- Variable lifetime

Allocator

- Balance throughput and memory utilization
- Data structures to keep track of free blocks

Garbage collection

- Must always free memory
- Garbage collectors help by finding anything reachable
- Failing to free results in memory leaks

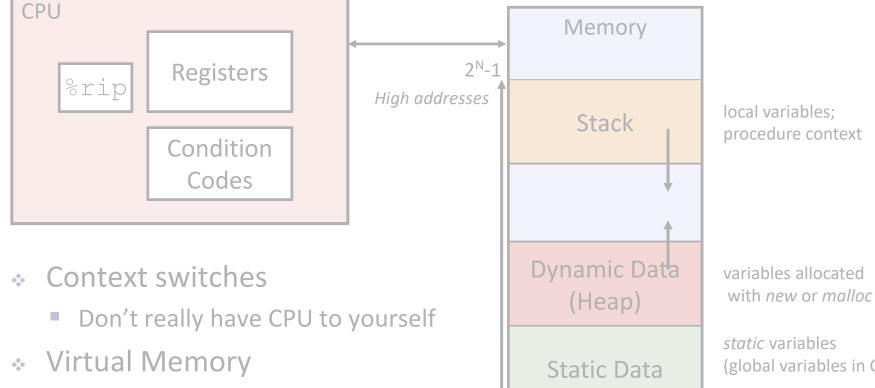
Low addresses (

local variables; Stack procedure context Dynamic Data variables allocated with *new* or *malloc* (Heap) static variables (global variables in C) Static Data Large constants Literals (e.g., "example") **Instructions**

Memory

But remember... it's all an illusion! (3)





static variables (global variables in C)

Large constants (e.g., "example")

- Don't really have 2⁶⁴ bytes of
 - memory all to yourself
 - Allows for *indirection* (remap physical pages, sharing...)

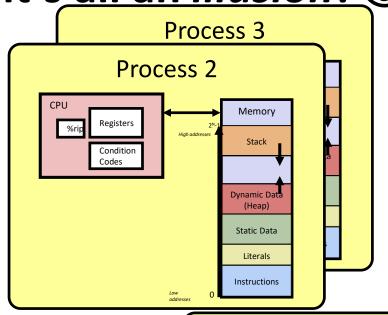
Low addresses

Instructions

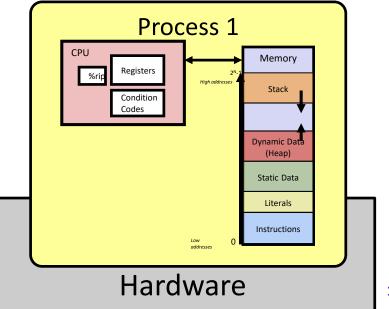
Literals

But remember... it's all an illusion!

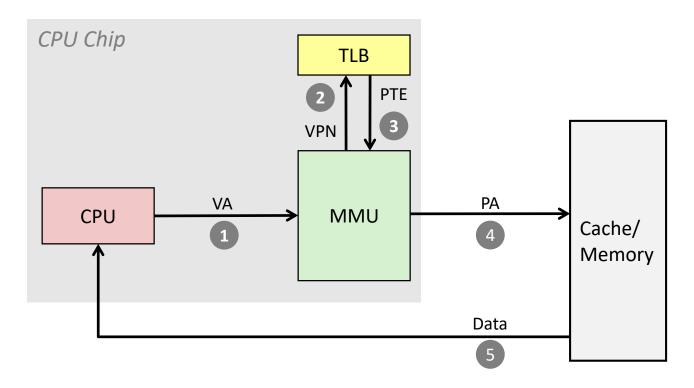
VIRTUAL ADDRESSES EVERYWHERE!



- * fork
 - Creates copy of the process
- execv
 - Replace with new program
- wait
 - Wait for child to die (to reap it and prevent zombies)



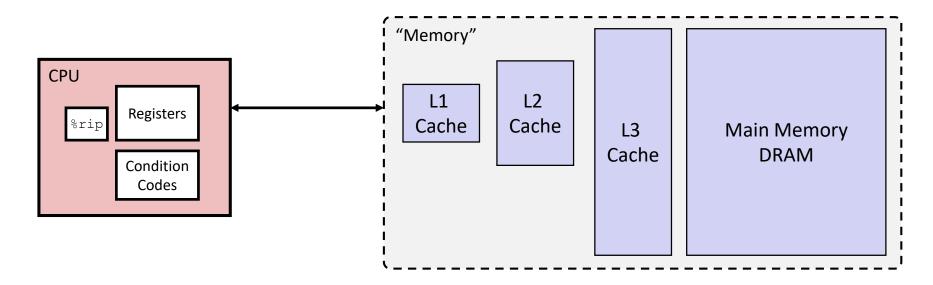
Virtual Memory



Address Translation

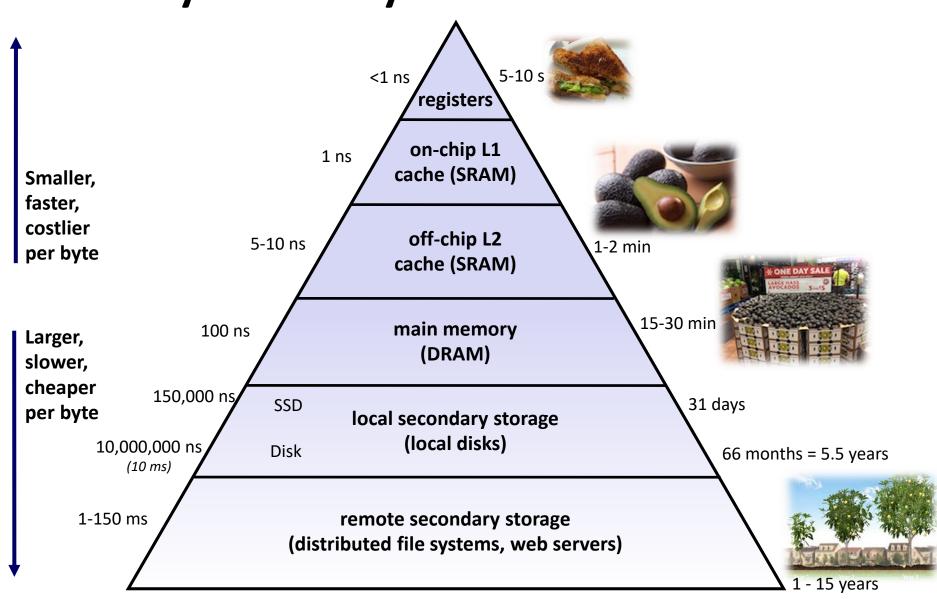
- Every memory access must first be converted from virtual to physical
- Indirection: just change the address mapping when switching processes
- Luckily, TLB (and page size) makes it pretty fast

But Memory is Also a Lie! 🐷



- Illusion of one flat array of bytes
 - But caches invisibly make accesses to physical addresses faster!
- Caches
 - Associativity tradeoff with miss rate and access time
 - Block size tradeoff with spatial and temporal locality
 - Cache size tradeoff with miss rate and cost

Memory Hierarchy



Review of Course Themes

- Review course goals
 - They should make much more sense now!

Big Theme: Abstractions and Interfaces

- Computing is about abstractions
 - (but we can't forget reality)
- What are the abstractions that we use?
- What do you need to know about them?
 - When do they break down and you have to peek under the hood?
 - What bugs can they cause and how do you find them?
- * How does the hardware relate to the software?
 - Become a better programmer and begin to understand the important concepts that have evolved in building ever more complex computer systems

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Little Theme 1: Representation

- All digital systems represent everything as 0s and 1s
 - The 0 and 1 are really two different voltage ranges in the wires
 - Or magnetic positions on a disc, or hole depths on a DVD, or even DNA...
- "Everything" includes:
 - Numbers integers and floating point
 - Characters the building blocks of strings
 - Instructions the directives to the CPU that make up a program
 - Pointers addresses of data objects stored away in memory
- Encodings are stored throughout a computer system
 - In registers, caches, memories, disks, etc.
- They all need addresses (a way to locate)
 - Find a new place to put a new item
 - Reclaim the place in memory when data no longer needed

Little Theme 2: Translation

- There is a big gap between how we think about programs and data and the 0s and 1s of computers
 - Need languages to describe what we mean
 - These languages need to be translated one level at a time
- We know Java as a programming language
 - Have to work our way down to the 0s and 1s of computers
 - Try not to lose anything in translation!
 - We encountered C language, assembly language, and machine code (for the x86 family of CPU architectures)

Little Theme 3: Control Flow

- How do computers orchestrate everything they are doing?
- Within one program:
 - How do we implement if/else, loops, switches?
 - What do we have to keep track of when we call a procedure, and then another, and then another, and so on?
 - How do we know what to do upon "return"?
- Across programs and operating systems:
 - Multiple user programs
 - Operating system has to orchestrate them all
 - Each gets a share of computing cycles
 - They may need to share system resources (memory, I/O, disks)
 - Yielding and taking control of the processor
 - Voluntary or "by force"?

Course Perspective

- CSE351 will make you a better programmer
 - Purpose is to show how software really works
 - Understanding the underlying system makes you more effective
 - Better debugging
 - Better basis for evaluating performance
 - How multiple activities work in concert (e.g., OS and user programs)
 - Not just a course for hardware enthusiasts!
 - What every CSE major needs to know (plus many more details)
 - See many patterns that come up over and over in computing (like caching and indirection)
 - "Stuff everybody learns and uses and forgets not knowing"
- CSE351 presents a world-view that will empower you
 - The intellectual and software tools to understand the trillions+ of 1s and 0s that are "flying around" when your program runs

Courses: What's Next?

- Staying near the hardware/software interface:
 - **EE271/CSE369:** Digital Design basic hardware design using FPGAs
 - **EE/CSE474:** Embedded Systems software design for microcontrollers
- Systems software
 - CSE341: Programming Languages (or CSE413 for non-majors)
 - CSE332: Data Structures and Parallelism (or CSE373 for non-majors)
 - CSE333: Systems Programming building well-structured systems in C/C++ (or CSE374 for non-majors)
- Looking ahead
 - CSE401: Compilers (pre-reqs: 332) (or CSE413 for non-majors)
 - CSE451: Operating Systems (pre-reqs: 332, 333)
 - CSE461: Networks (pre-reqs: 332, 333)

Thanks for a great quarter!

Huge thanks to your awesome TAs!











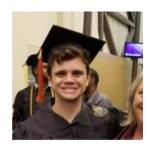












- Don't be a stranger!
 - I'll likely be teaching this course again next year