The Traveling Salesman Problem

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Problem Statement

- **Input** : A complete undirected graph with non-negative edge costs.
- **Output** : A minimum cost tour i.e. a cycle that visits all the vertices exactly once.

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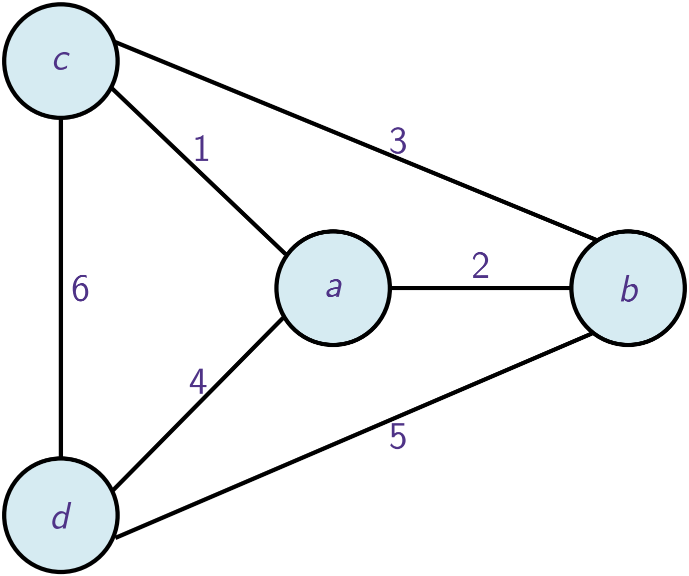
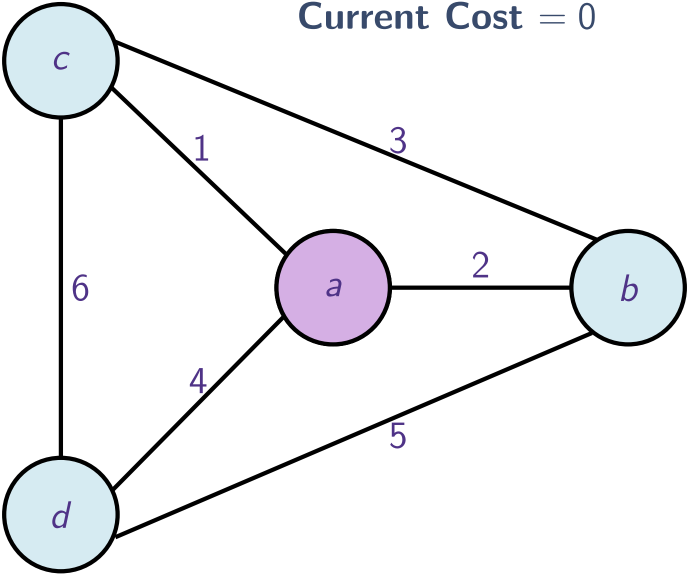
An Example



An Example

Current Cost $= 0$



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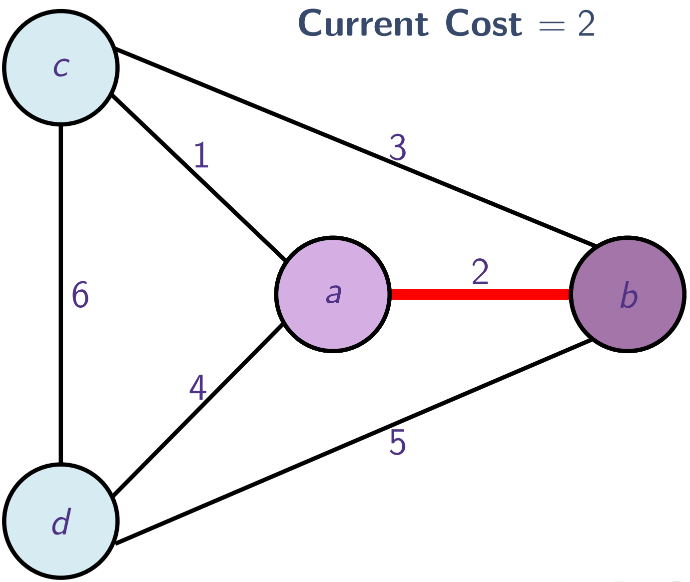

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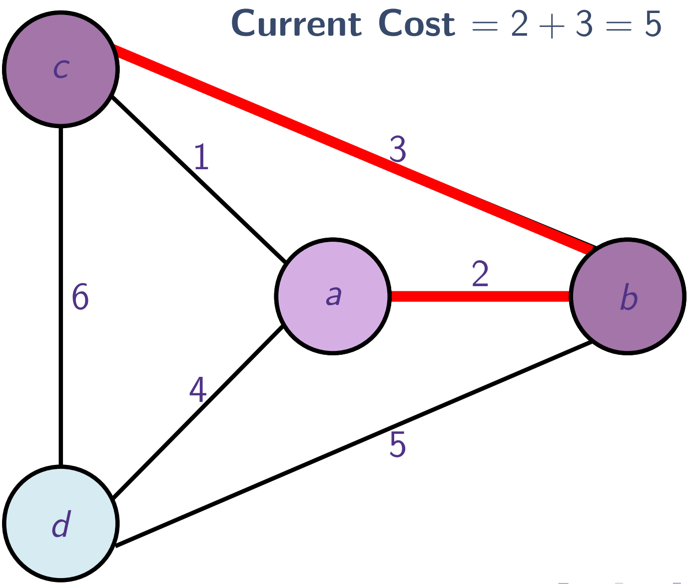

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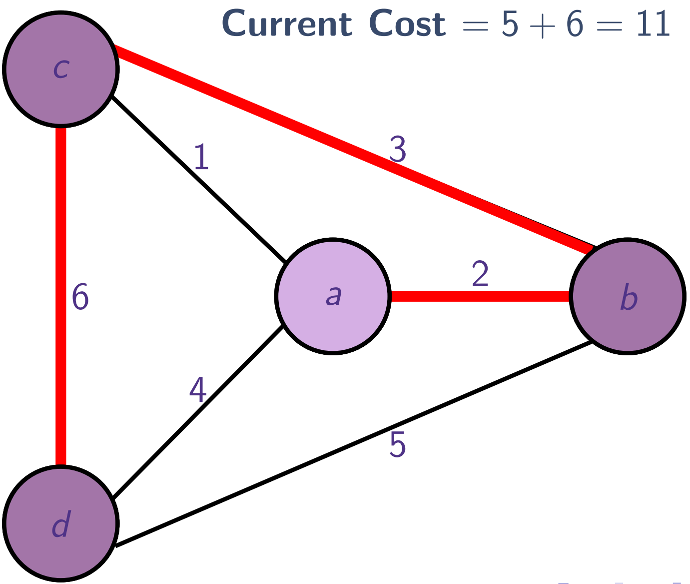

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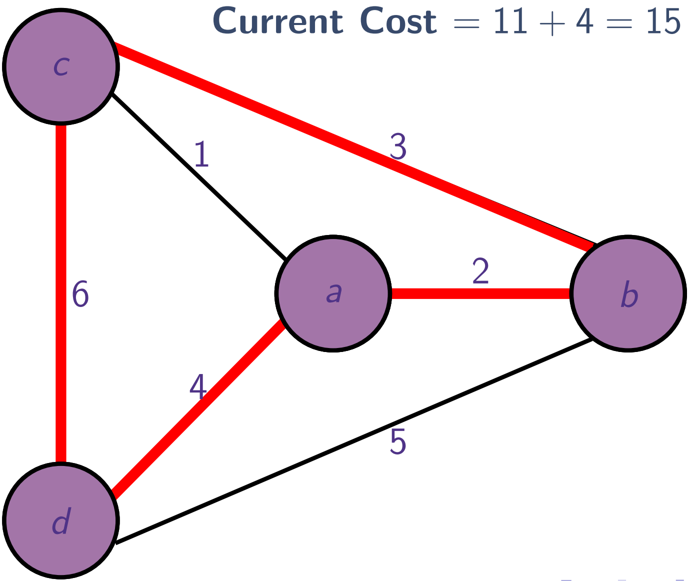
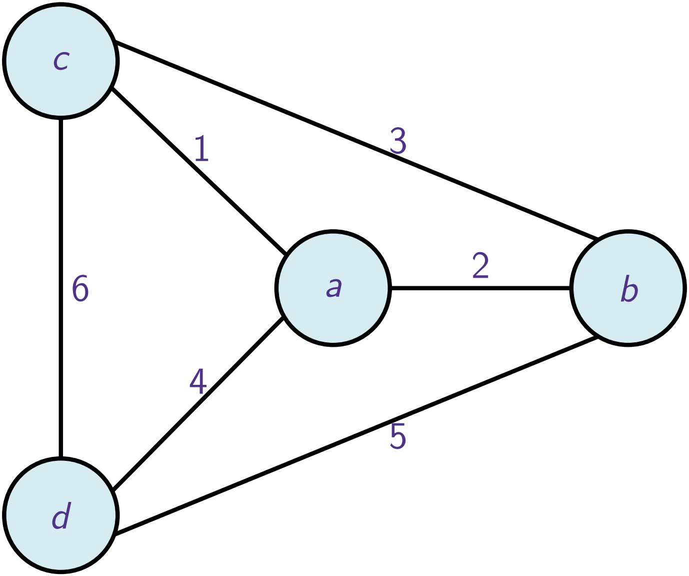


An Example (Continued)



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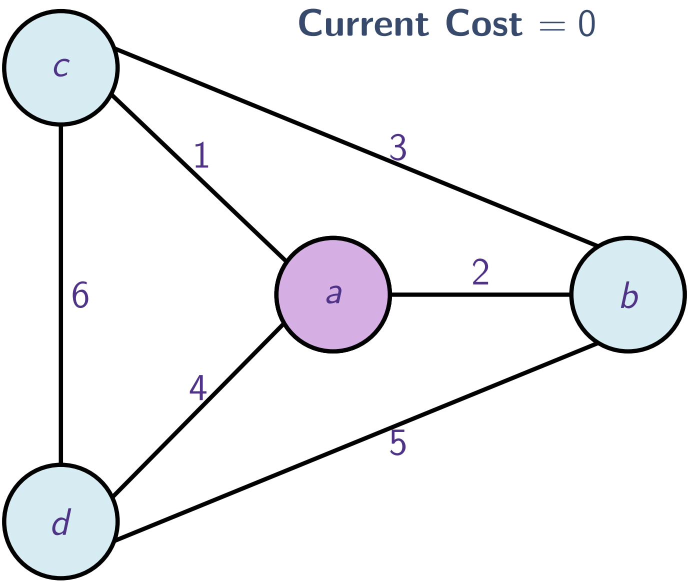
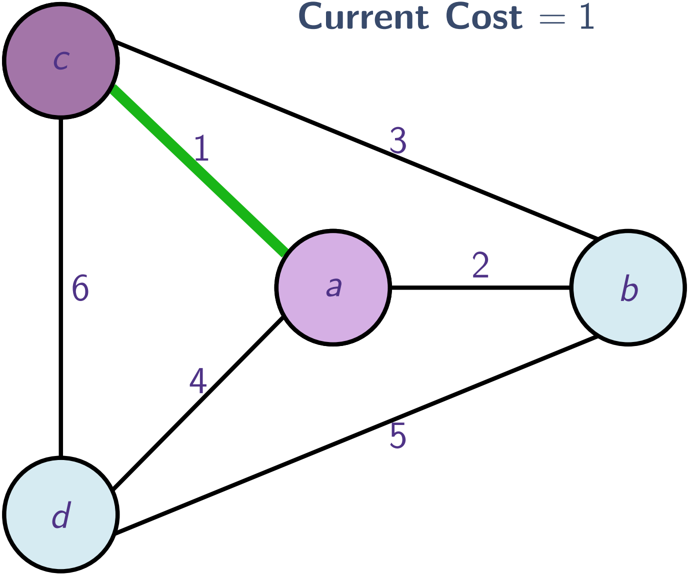


An Example (Continued)

Current Cost $= 1$



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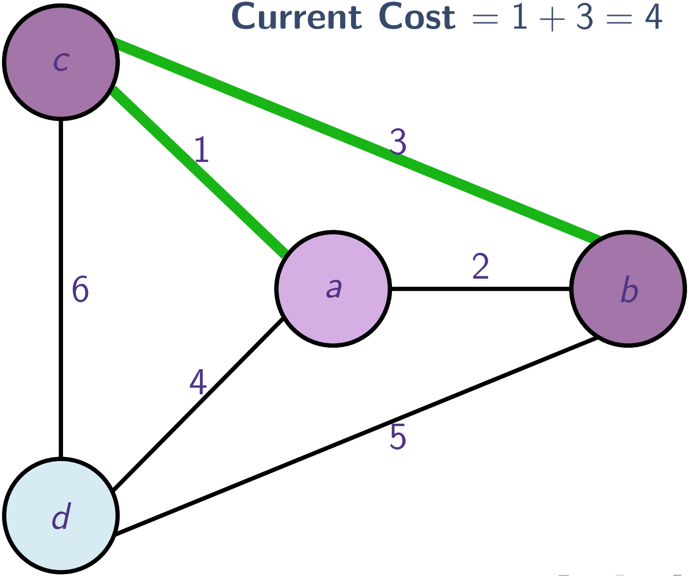
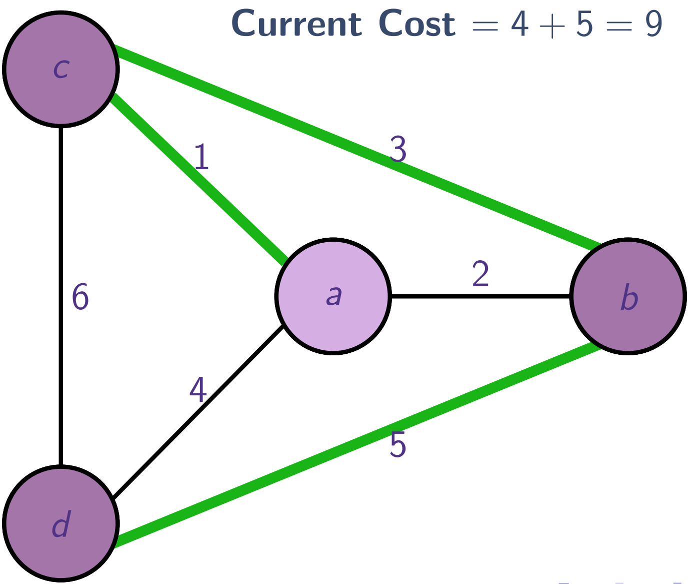


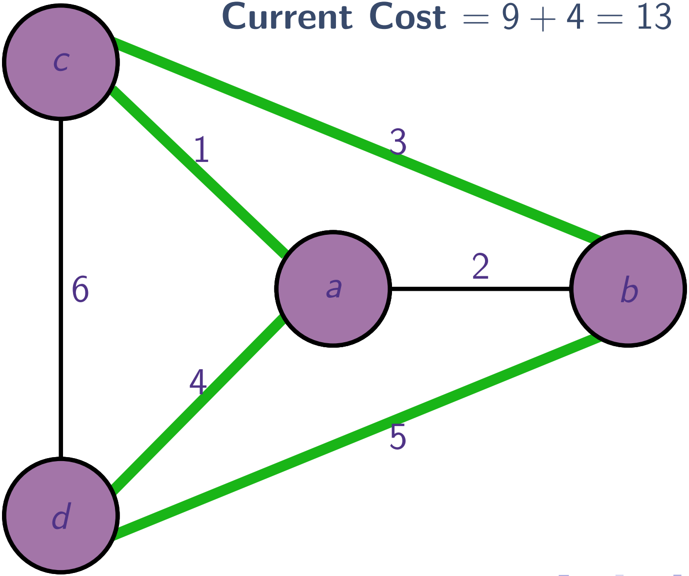
Current Cost $= 1 + 3 = 4$

An Example (Continued)



An Example (Continued)

Current Cost $= 9 + 4 = 13$



Question of the Day

How do we find an optimal tour?

The Brute Force Algorithm

The Brute Force approach:

- Look at all possible tours in the graph.
- Compute their costs.
- Pick the minimum from them.

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Brute Force Algorithm Complexity

However,

- In a graph on n vertices, there are $(n-1)!$ TSP tours.
- Computing the cost of a tour takes linear time.
- **Brute-Force Algorithm running time:**
  \# of tours $\times$ cost of computing one tour $= (n-1)! \times O(n) = O(n!)$

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An Efficient Algorithm?

A Question We Should be Asking Everyday

Can we do better?

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We can. But a polynomial time algorithm doesn't seem likely.

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- Edmonds' Conjecture equivalent to $P \neq NP$.
- *The Traveling Salesman Problem* is NP-Complete!

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Hard to even approximate

There is a catch.

Theorem

Unless $P = NP$, there does not exist a polynomial time α- approximation algorithm for the Traveling Salesman Problem.

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Metric TSP

Edge costs satisfy the triangle inequality i.e. the shortest path between vertices = the one-hop path between them.

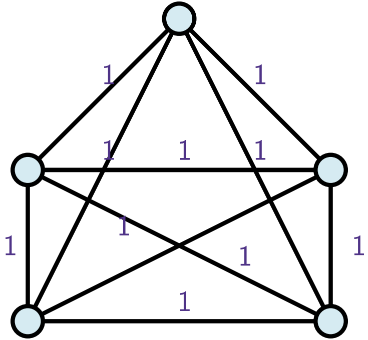

Figure: A Metric TSP instance.

Approximation Algorithms for Metric TSP

- Still NP-Complete!
- But there are good approximation algorithms.
  - **The MST Heuristic** (a 2-approximation algorithm)
  - **Christofides's Algorithm (1976)** (a $\frac{3}{2}$-approximation algorithm)

To Summarize

- The *Traveling Salesman Problem* is interesting.
- *The Traveling Salesman Problem* is hard!
- Approximation algorithms for NP-Complete Problems are still an active area of research.

Acknowledgements I

Tim Roughgarden.
Stanford CS261 Lecture Notes.
2016.

Jack Edmonds.
Paths, trees, and flowers.
Can. J. Math. 17: 449467, 1965.