

Strong 1-dimensional Clobber Solver

TheSolvers

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Outline

- Correctness
- Abstract Code Flow
- Database Structure
- Improvements
 - Dominated Moves
 - Game Values
 - Rules
 - Statistics
- Results
- Next Step

Correctness

Bugs

- Generating illegal moves
- Database generation
- Non null terminated strings
- 0-0 instead of None

Correctness

Fixing bugs

- We also conducted a pipeline for testing our solver based on the reference solver that we have.
- We recognized the source of our few wrong results and debugged our code to fix them.
- Finally, having a correct verified solver.

Abstract Code Flow

Simplify board()

lookup TT, if found → return the result

generating subgames()

for each subgame:

 outcome, game value, dominated moves ← db.get(subgame)

if the game can be solved using just the outcomes → return the result

else if the game has a value using sum of subgames → return the result

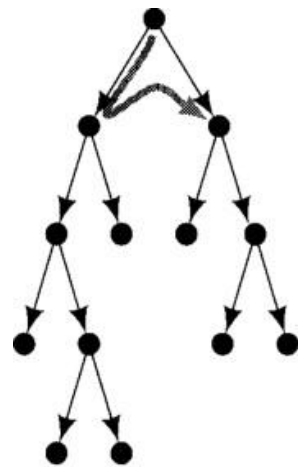
else:

solve the game with the help of dominated moves, based on TT, and ID search.

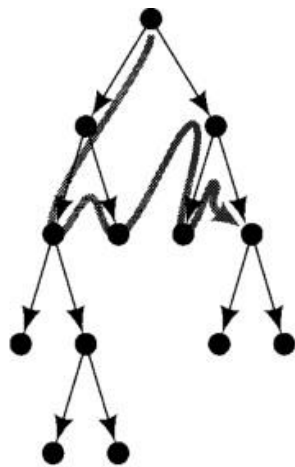
The Search Algorithm

Iterative Deepening:

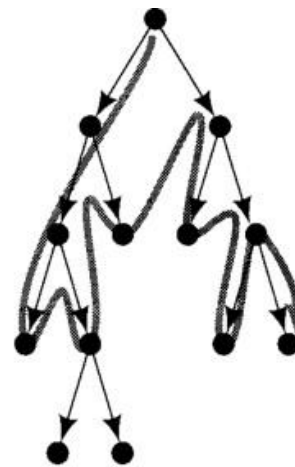
We use the iterative deepening search, with a limit on the number of searches that are allowed to reach terminal nodes (nodes at the current maximum depth, or actual terminal nodes).



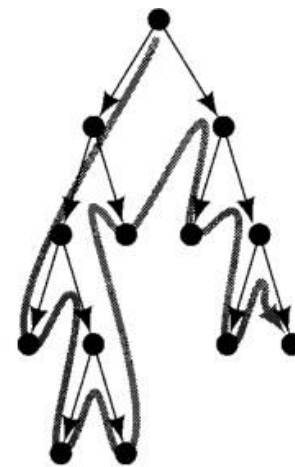
Depth bound = 1



Depth bound = 2



Depth bound = 3



Depth bound = 4

Database Structure

		outcome	player1 dominated moves	player2 dominated moves	game value	link to the equivalent node
entry:	WBBBBBBB	OC_B	1001...1	0100..11	05000	4

- Creating a database for up to 20 tiles with outcome and game values while computing dominated moves for up 12 to tiles in 30m 12.221s
- Computing a database up to 20 tiles without computing dominated moves and outcome takes 52m 47.186s.

Improvements - dominated moves

Removing dominated moves

Steps of determining if a move is dominated

- generates the list of moves for a player at a state
- compare all pairs of moves (i, j) and their respective resulting games I, J
- check if a single player wins I - J regardless of who goes first
- If this happens -> one of moves i and j are marked as dominated
- save the dominated move in the database

For now, we store dominated moves for all connected boards from sizes 1 through 12.

Improvements - Game Values

Game Values

We extended our database to be able to save some game values.

- First of all, saved some initial values (rule 3-5) in our db.
- Saved the number of ups, downs and star values.
 - (two digits for up, two for down, and one for star)

$$WBBB = WB^3 = \uparrow\uparrow^* = 2001$$

Improvements - Game Values

For the remaining boards, if the board value is not set:

- Find the game value for left and right
 - Play all the left and right options for one level
- Return the max/min value for L/R
 - Comparing left values (or right values) using rules in fig1 and fig2
- Set the final value based on the rules (1) using L and R values.

$$\uparrow\uparrow > \uparrow + \uparrow^2 > \uparrow + \uparrow^3 > \uparrow > \uparrow^2 > \uparrow^3 > 0 > \downarrow^3 > \downarrow^2 > \downarrow > \downarrow + \downarrow^3 > \downarrow + \downarrow^2 > \downarrow\downarrow.$$

fig 1

$$\uparrow\uparrow * > \uparrow * > * > \downarrow * > \downarrow\downarrow *.$$

fig 2

Game value Rules

<p>1</p> $n \cdot \uparrow = \{0 \mid (n - 1) \cdot \uparrow^*\} \quad \text{if } n \geq 1$ $n \cdot \uparrow^* = \begin{cases} \{0 \mid (n - 1) \cdot \uparrow\} & \text{if } n > 1 \\ \{0, * \mid 0\} & \text{if } n = 1 \end{cases}$	<p>2</p> $n \cdot \downarrow = \{(n - 1) \cdot \downarrow^* \mid 0\}$ $n \cdot \downarrow^* = \begin{cases} \{(n - 1) \cdot \downarrow \mid 0\} & \text{if } n > 1 \\ \{0 \mid 0, *\} & \text{if } n = 1 \end{cases}$
$(BBW)^n = \left\lfloor \frac{n+1}{2} \right\rfloor \cdot \uparrow$ <p>3</p>	$B^m W^n = 0, \text{ for } m, n \geq 2$ <p>4</p>
$WB^n = (n - 1) \cdot \uparrow + n \cdot *$ <p>5</p>	

Game Value Statistics

Number of exact game values within the $\uparrow\uparrow$ and $\downarrow\downarrow$ bound

*	48	\downarrow	32
0	74	$\uparrow\uparrow$	2
\uparrow	32	$\downarrow\downarrow$	2

Examples

example of some initial patterns:

- $\uparrow\uparrow = (\text{BBW})^4, (\text{BBW})^3$
- $\uparrow = (\text{BBW})^2, (\text{BBW}), \text{WBB},$
- $* = \text{BW}, \text{WB}$
- $\downarrow\downarrow = (\text{WWB})^4, (\text{WWB})^3$
- $\downarrow = (\text{WWB})^2, (\text{WWB}), \text{BWW},$
- $\uparrow\uparrow* = \text{WBBB}$
- $\downarrow\downarrow* = \text{BWWW}$
- $0 = 34 (\text{B}^m\text{W}^n, \text{W}^m\text{B}^n, m,n \geq 2)$

Other Improvements

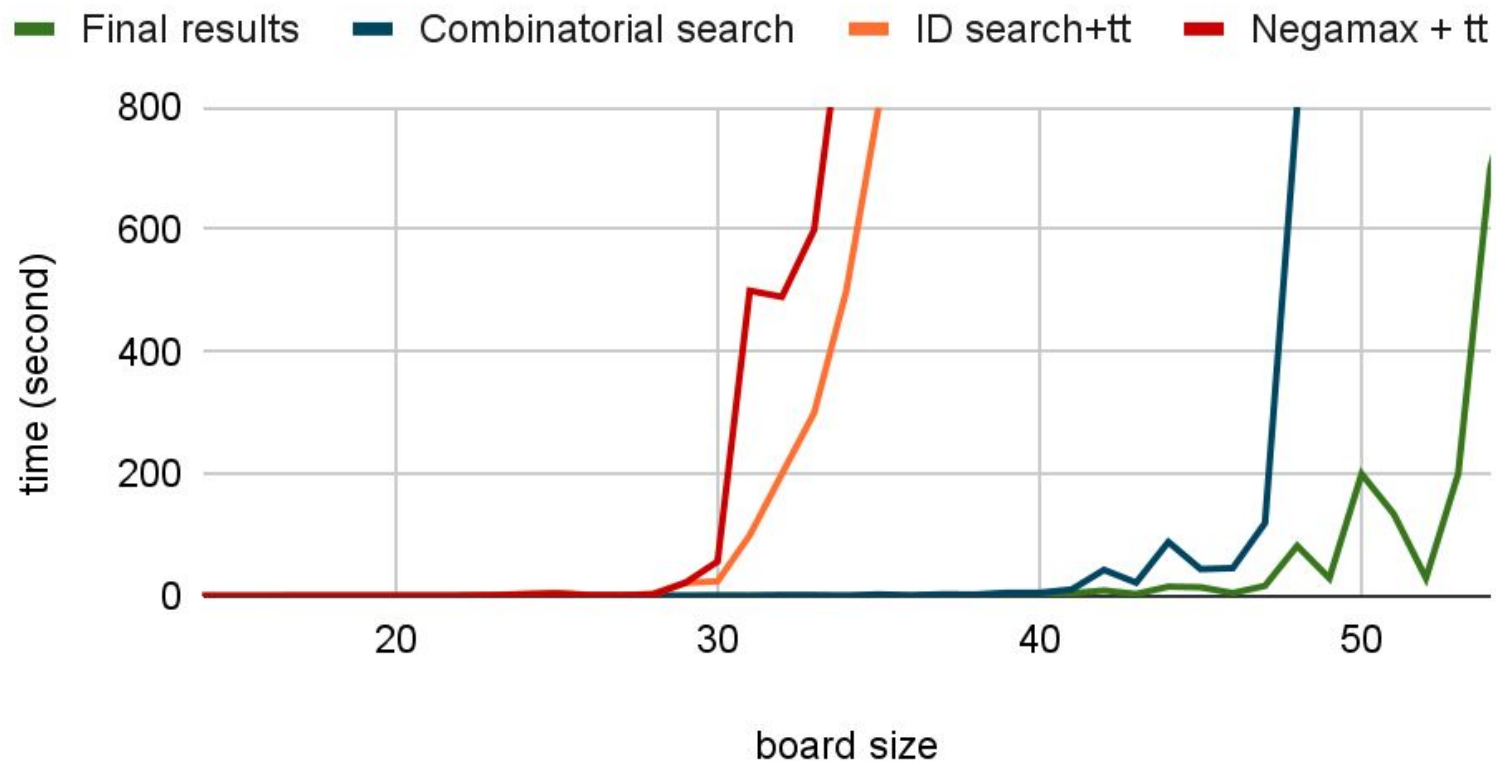
We implemented the rule “One N game with several L/R subgames is the first-player win for Left/Right.”

- The correctness of this improvement is also tested.
- It speeded up the solver a bit (not noticeable), and the number of nodes that were searched is decreased.

Trying alpha-beta pruning on heuristic values → did not improve our solver with our current heuristic.

Results - runtime

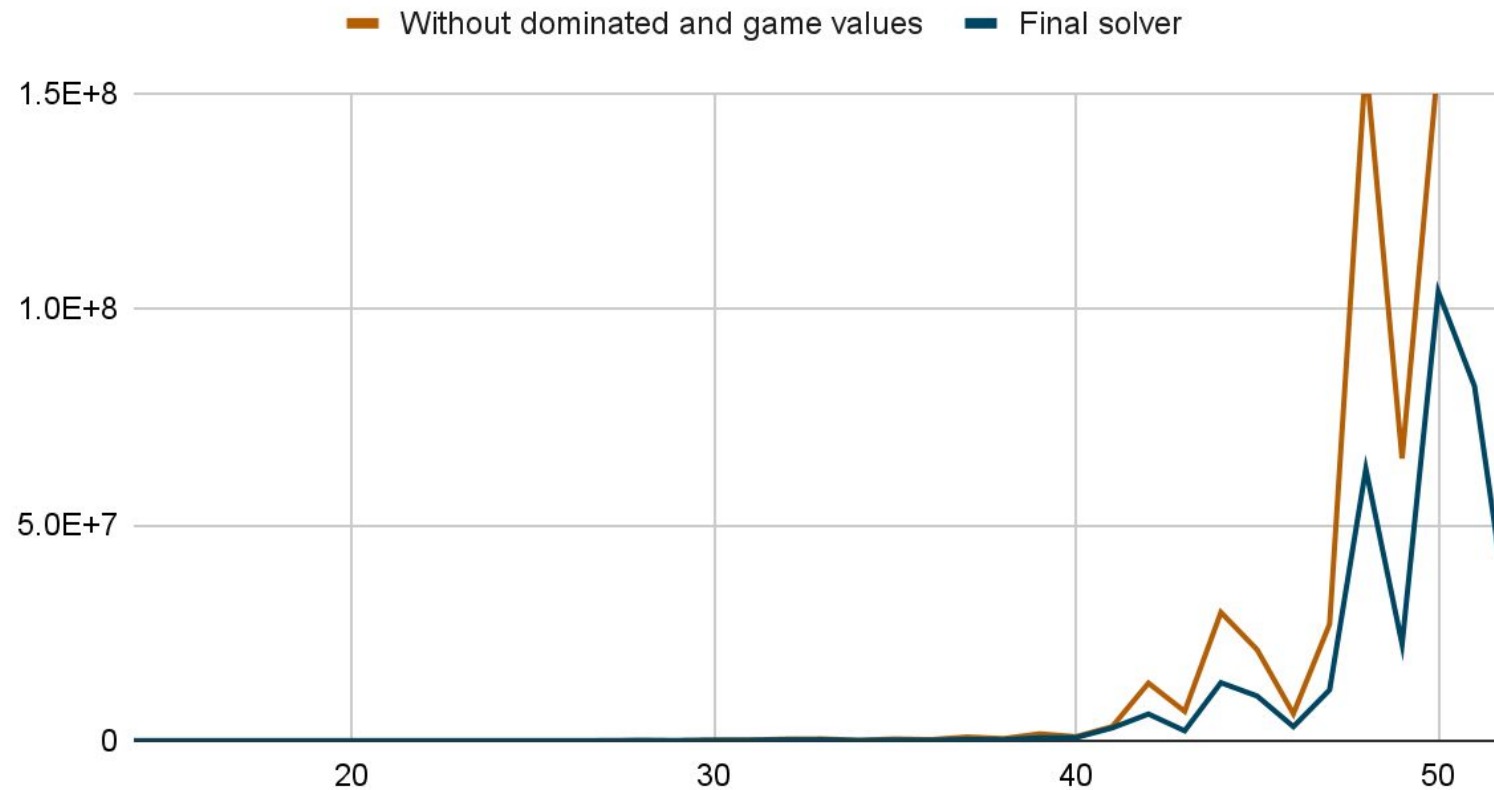
Comparing four solvers



The runtime result of consecutive B and W stones of different length

Node counts vs board length

Node counts



Result of consecutive B and W stones of different length

Overall runtime on 500 random test cases

Overall runtime for solving 500
randomly generated boards:

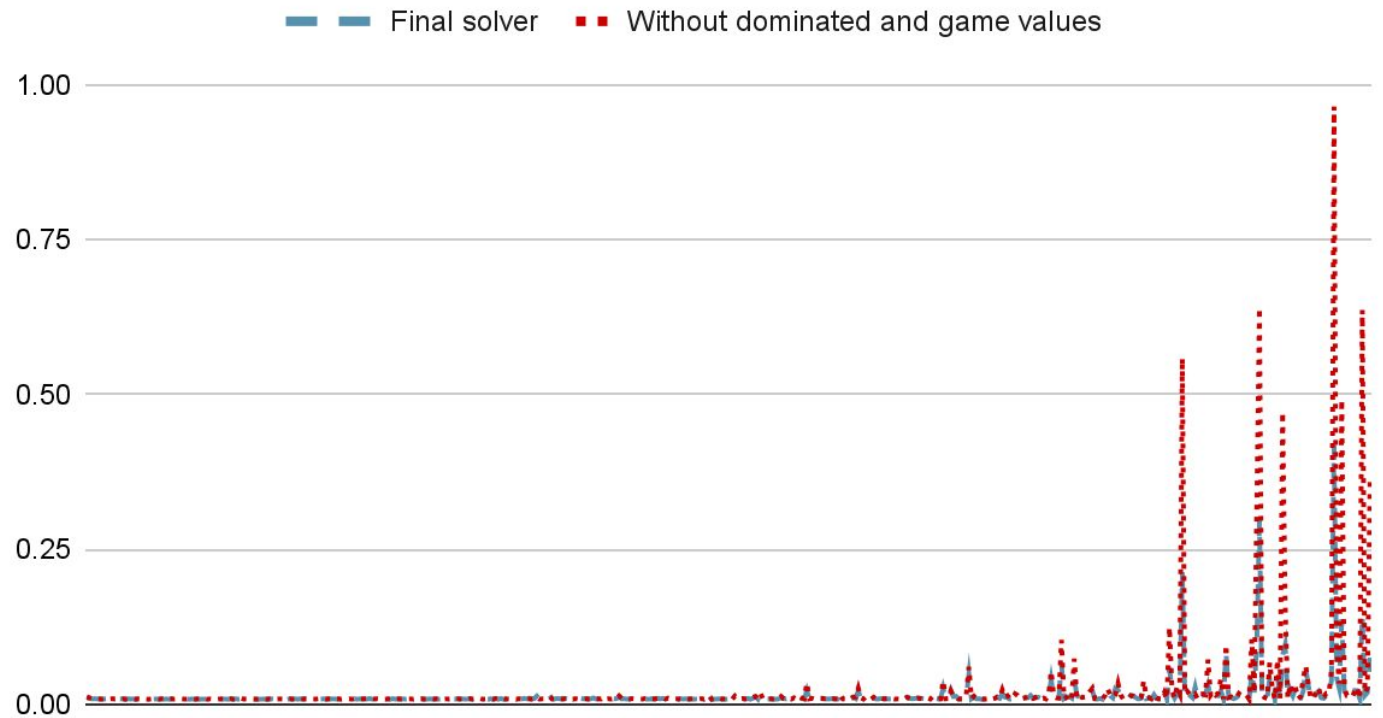
Final results 7.18s

Previous best
combinatorial solver:

15.93s

[10-52 tiles]

Runtime for 500 randomly generated boards



Next Steps

- Simplifying subgames
- Lower/upper bounds of subgames
- Improve subgame deletion heuristic



Demo

References

- [1] J Fernando Hernandez, Using Left and Right Stop Information for Solving Clobber.
- [2] Michael Albert et al., Introduction to Clobber.
- [3] Claessen. “Combinatorial Game Theory in Clobber.” Maastricht University. 2011.
- [4] Kristopher De Asis, A CGT-informed clobber player.
- [5] Max Blankestijn, Stops, All-Small & Infinitesimals, Lessons In Play.