

EXHAUSTIVE SEARCH ALGORITHMS

Brute Force Approach

Brute force is a **straightforward approach** to solving a problem, directly based on the **problem statement** and **definitions** of the concepts involved.

Exhaustive Search

Exhaustive search is a brute-force approach to a problem involving search for an **object with a special property**, usually among **combinatorial objects** such as **subsets** of a set, **permutations**, or **combinations** of sequences.

Exhaustive Search

Exhaustive search is a brute-force approach to a problem involving search for an **object with a special property**, usually among **combinatorial objects** such as **subsets** of a set, **permutations**, or **combinations** of sequences.

Method:

- generate a list of all potential solutions to the problem in a systematic manner,
- **evaluate potential solutions one by one**, disqualifying infeasible ones and, for an optimization problem, **keeping track of the best one** found so far,
- when search ends, **announce the solution(s) found**.

Activity (in group)

- **Design your original algorithm**
 - TSP / Knapsack / Job Assignment Problem
 - Show how you can improve better than exhaustive approach
- Analyze: running time
- Write a pseudocode/program
- Share

[https://padlet.com/
sem2_2018_2019/dsa2](https://padlet.com/sem2_2018_2019/dsa2)



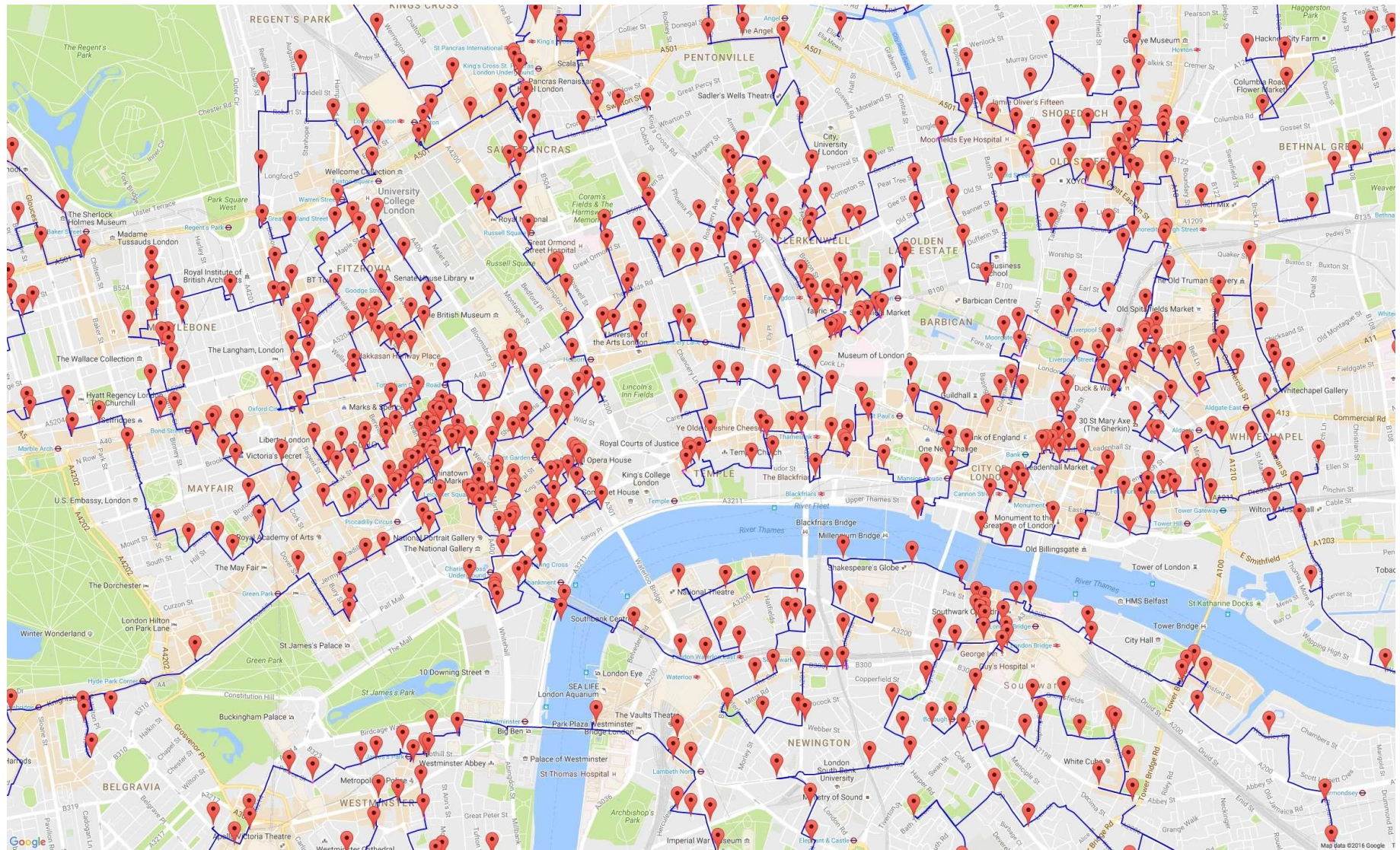
Traveling Salesman Problem

Problem: Given n cities with known distances between each pair, find the **shortest tour** that passes through **all the cities exactly once** before returning to the starting city.

Traveling Salesman Problem



Traveling Salesman Problem



Traveling Salesman Problem

Problem: Given n cities with known distances between each pair, find the **shortest tour** that passes through all the cities exactly once before returning to the starting city.

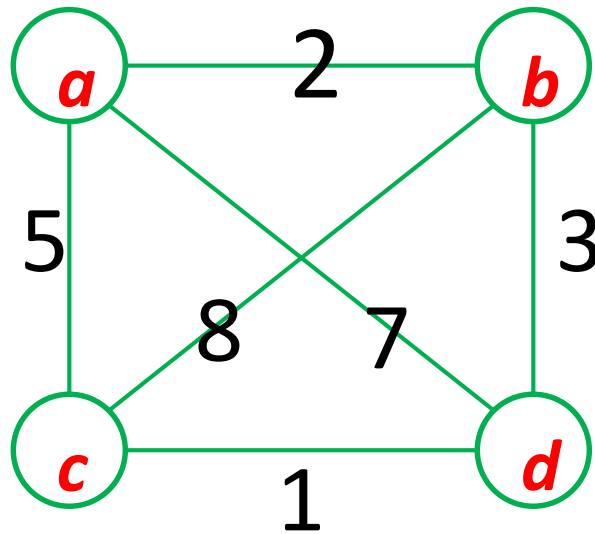
Alternatively: Find the shortest **Hamiltonian cycle** (a cycle that passes through all vertices (v_1, v_2, \dots, v_n) exactly once) in a weighted connected graph:

$$v'_1, \underbrace{v'_2, v'_3, \dots, v'_n}, v'_1$$

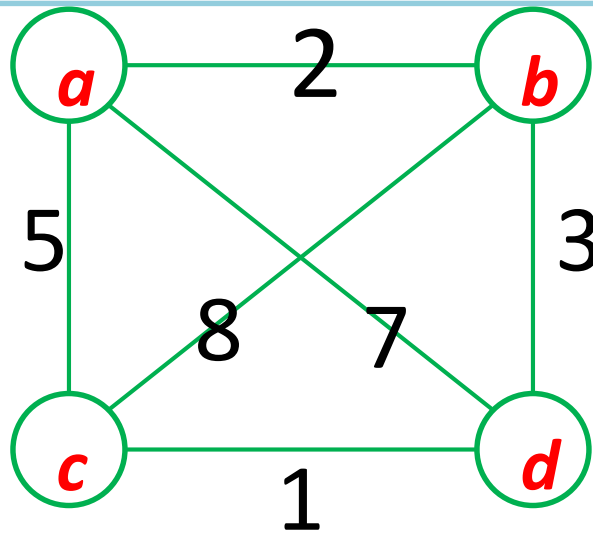
distinct vertices

Traveling Salesman Problem

Example:



Example:



Traveling Salesman Problem

- | | |
|--|--------------------------|
| 1. $a \rightarrow b \rightarrow c \rightarrow d \rightarrow a$ | $d = 2 + 8 + 1 + 7 = 18$ |
| 2. $a \rightarrow b \rightarrow d \rightarrow c \rightarrow a$ | $d = 2 + 3 + 1 + 5 = 11$ |
| 3. $a \rightarrow c \rightarrow b \rightarrow d \rightarrow a$ | $d = 5 + 8 + 3 + 7 = 23$ |
| 4. $a \rightarrow c \rightarrow d \rightarrow b \rightarrow a$ | $d = 5 + 1 + 3 + 2 = 11$ |
| 5. $a \rightarrow d \rightarrow b \rightarrow c \rightarrow a$ | $d = 7 + 3 + 8 + 5 = 23$ |
| 6. $a \rightarrow d \rightarrow c \rightarrow b \rightarrow a$ | $d = 7 + 1 + 8 + 2 = 18$ |

Traveling Salesman Problem

The efficiency:

$$v'_1, \underbrace{v'_2, v'_3, \dots, v'_n}_{\text{distinct vertices}}, v'_1$$

distinct vertices

- The number of all permutations: $(n - 1)!$
- The half tour differs only by their directions: $\frac{1}{2}(n - 1)!$
- $O(n!)$

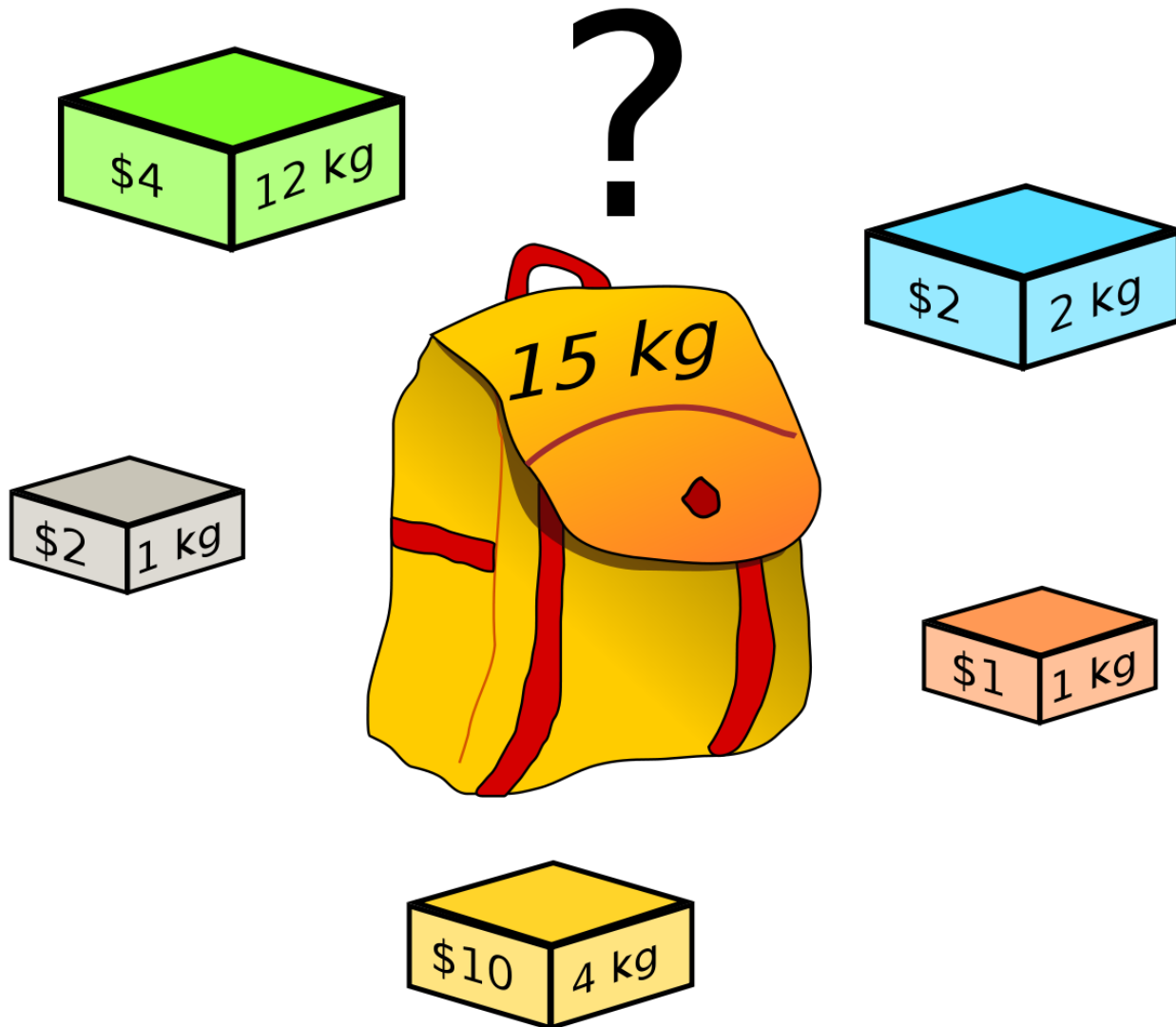
Knapsack Problem

Given n items i_1, i_2, \dots, i_n of known:

- **weights:** w_1, w_2, \dots, w_n
- **values:** v_1, v_2, \dots, v_n
- a **knapsack** of capacity W

Problem: Find the **most valuable subset** of the items i_1, i_2, \dots, i_n that **fit into the knapsack**.

Knapsack Problem



Knapsack Problem

Example: Knapsack capacity $W = 16$

ITEM	WEIGHT	VALUE
1	2	\$20
2	5	\$30
3	10	\$50
4	5	\$10

The exhaustive search approach:

Knapsack Problem

Example: Knapsack capacity $W = 16$

ITEM	WEIGHT	VALUE
1	2	\$20
2	5	\$30
3	10	\$50
4	5	\$10

The exhaustive search approach: (1) generate **all subsets** of the set of n items; (2) compute the **total weight** of each subset; (3) find a **subset of largest value** among them.

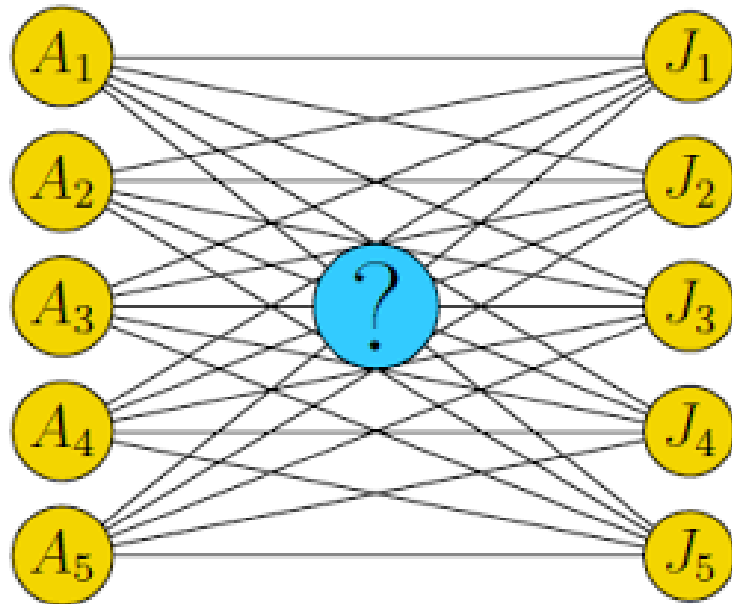
- The **total # of all subsets**: 2^n **Efficiency**: $\Omega(2^n)$

SUBSETS	TOTAL WEIGHT	TOTAL VALUE
{1}	2	\$20
{2}	5	\$30
{3}	10	\$50
{4}	5	\$10
{1,2}	7	\$50
{1,3}	12	\$70
{1,4}	7	\$30
{2,3}	15	\$80
{2,4}	10	\$40
{3,4}	15	\$60
{1,2,3}	17	not feasible
{1,2,4}	12	\$60
{1,3,4}	17	not feasible
{2,3,4}	20	not feasible
{1,2,3,4}	22	not feasible

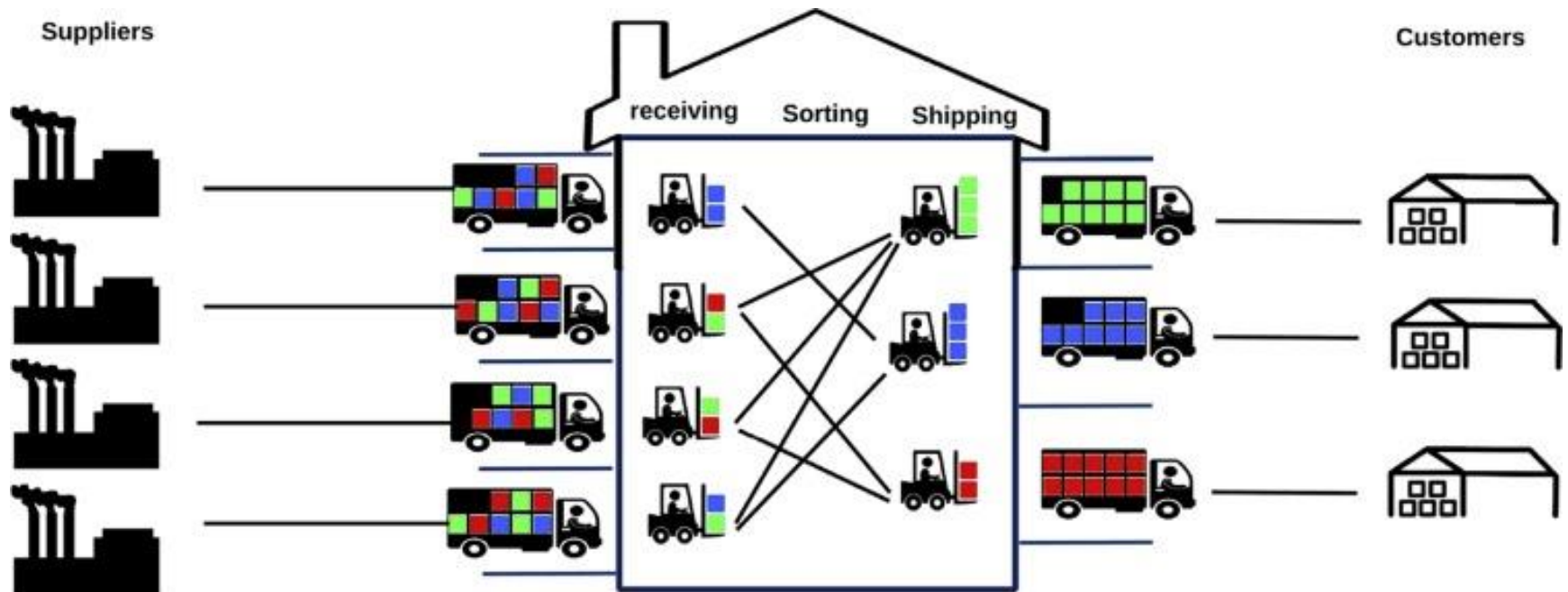
Assignment Problem

There are n people who need to be assigned to n jobs, one person per job. The cost of assigning person i to job j is $C[i, j]$. **Find an assignment that minimizes the total cost.**

Assignment Problem



Assignment Problem



Assignment Problem

There are n people who need to be assigned to n jobs, one person per job. The cost of assigning person i to job j is $C[i, j]$. **Find an assignment that minimizes the total cost.**

	Job 1	Job 2	Job 3	Job 4
Person 1	9	7	2	8
Person 2	6	4	3	7
Person 3	5	8	1	8
Person 4	7	6	9	4

Assignment Problem

Solution: Select one element in each row so that all selected elements are in different columns and the total sum of the selected elements is the smallest possible.

- | | P1 | P2 | P3 | P4 | |
|----|----|----|----|----|------------------------------------|
| 1. | 9 | 4 | 1 | 4 | $\text{cost} = 9 + 4 + 1 + 4 = 18$ |
| 2. | 9 | 4 | 8 | 9 | $\text{cost} = 9 + 4 + 8 + 9 = 30$ |
| 3. | 9 | 3 | 8 | 4 | $\text{cost} = 9 + 3 + 8 + 4 = 24$ |
| 4. | 9 | 3 | 8 | 6 | $\text{cost} = 9 + 3 + 8 + 6 = 26$ |
| 5. | 9 | 7 | 1 | 6 | $\text{cost} = 9 + 7 + 1 + 6 = 23$ |

Efficiency: $O(n!)$

Activity (in group)

- **Design your original algorithm**
 - TSP / Knapsack / Job Assignment Problem
 - Show how you can improve better than exhaustive approach
- Analyze: running time
- Write a pseudocode/program
- Share

[https://padlet.com/
sem2_2018_2019/dsa2](https://padlet.com/sem2_2018_2019/dsa2)

