

**Algorithms with order  
 $\log_2 n$**

# Analysis of Nonrecursive Algos

**Example 1:** The **number of binary digits** in the binary representation of a positive integer.

**Algorithm** Binary( $n$ )

// Input: A positive decimal integer  $n$

//Output: The number of binary digits in the bin. repr.

count  $\leftarrow 1$

**while**  $n > 1$  **do**

    count  $\leftarrow$  count + 1

$n \leftarrow \lfloor n/2 \rfloor$

**return** count

# Analysis of Nonrecursive Algos

- **The basic operation:** the number of times of the **comparisons** executed (which is larger than the number of repetitions of the loop's body by exactly one)
- **The input size:** the loop variable  $n$  takes only a few values between its lower and upper limits:

$$n, \left\lfloor \frac{n}{2} \right\rfloor, \left\lfloor \frac{n}{4} \right\rfloor, \left\lfloor \frac{n}{8} \right\rfloor, \dots, 1$$

- **The total number:** In each loop repetition,  $n$  is about halved, thus,

$$\text{the total \#} \approx \log_2 n$$

# Analysis of Recursive Algos

**Example 2:** Find the **number of binary digits** in the binary representation of a positive integer.

**Algorithm** BinRec( $n$ )

// Input: A positive decimal integer  $n$

//Output: The number of binary digits in the bin. repr.

**if**  $n = 1$  **return** 1

**else return** BinRec( $\lfloor n/2 \rfloor$ ) + 1

# Analysis of Recursive Algos

**Example 2:** Find the **number of binary digits** in the binary representation of a positive integer.

**Algorithm** BinRec(n)

// Input: A positive decimal integer n

//Output: The number of binary digits in the bin. repr.

**if** n = 1 **return** 1

**else return** BinRec( $\lfloor n/2 \rfloor$ ) + 1

Let  $A(n)$  be the **total number of additions**.

$$A(n) = A(\lfloor n/2 \rfloor) + 1, n > 1, \quad A(1) = 0$$

# Analysis of Recursive Algos

For the simplicity we choose  $n = 2^k$  (the smoothness rule)

$$\begin{aligned} A(2^k) &= A\left(\frac{2^k}{2}\right) + 1 = A(2^{k-1}) + 1 \\ &= (A(2^{k-2}) + 1) + 1 = A(2^{k-2}) + 2 \\ &= (A(2^{k-3}) + 1) + 2 = A(2^{k-3}) + 3 \\ &= \dots \\ &= A(2^{k-i}) + i \\ &= \dots \\ &= A(2^{k-k}) + k = A(1) + k = k = \log_2 n \end{aligned}$$