

CS2030S

Finals Cheatsheet for 23/24 S2
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Types

Subtypes $S <: T$ or $T >: S$ (S is subtype of T)

1. **Reflexive** $S <: S$
2. **Transitive** $S <: T \wedge T <: U \Rightarrow S <: U$
3. **Antisymmetry** $S <: T \wedge T <: S \Rightarrow S = T$

Subtyping Relationship

byte <: short <: int <: long <: float <: double
char <: int

Widening and Narrowing Type Conversion

Widening Variable of type T can hold a value of type S if $S <: T$

Narrowing Variable of type S cannot hold a value of type T if $S <: T$

Variance of Types

Covariance $S <: T \Rightarrow C(S) <: C(T)$

Contravariance $S <: T \Rightarrow C(T) <: C(S)$

Invariant Neither

OOP Principles

Encapsulation

Information Hiding Private instance fields

Abstraction

Do not show actual implementation of methods.

Inheritance

Ability to reuse code of existing super classes.
Models **is-a** relationship.

Polymorphism

Using same method signatures in subclasses to determine behaviour for specific subclasses.

Tell Don't Ask

The client should not be doing computation on the object's behalf.

Method Signature and Descriptor

Method Signature method name, number of parameters, type of each parameter and order of parameters

Method Descriptor method signature + return type

Liskov Substitution Principle

Let $\phi(x)$ be a property provable about objects x of type T .

Then $\phi(y)$ should be true for objects y of type S where $S <: T$

Replacing **super** with **child** should not change result.

Overriding & Overloading

Overriding

Same method signature: method name, number of parameters, type of parameters and order

Must use `@Override` annotation

Return type must be a subtype of the overridden method's return type.

Overloading

Same method name in the same class

Different method signature (cannot be different return type with same method signature)

Interfaces and Abstract Classes

Interface

Can implement multiple

Only abstract methods

Can type-cast regardless of subtype

Abstract

Can only extend one

Abstract & non-abstract
Cannot

Wrapper Classes

Autoboxing Convert primitive to wrapper

Unboxing Convert wrapper to primitive

Allowed Unboxing & Widening

Not Allowed Autoboxing & Widening

Dynamic Binding

1. Determine compile-time type of target
2. Check all accessible methods (including inherited ones)
3. Most specific one callable
4. Determine run-time type of target
5. Determine method called.

CTT | Methods Callable | Most Specific | RTT

Wildcards

PECS Producer Extends, Consumer Super

Upper-Bounded Wildcards

$A < ? \text{ extends } S >$ **Covariant.**

Lower-Bounded Wildcards

$A < ? \text{ super } S >$ **Contravariant.**

Raw Type

$A<?>$ Complex type of a specific but unknown type.
 $A<Object>$ Complex type of Object instances with type checking.

A Complex type of Object instances without type checking.

Raw types **throw unchecked warnings.**

Type Inference

Constraints

1. Target Type
2. Argument Type
3. Type Parameter Bound

Functional Interface

`@FunctionalInterface`
interface `Producer<T> {T produce();}`

Exceptions

Checked Exceptions

Exception programmer has **no control** over

All checked exceptions $E <: \text{Exception}$

If a checked exception, must put **throws** keyword

Unchecked Exceptions

Exception programmer has **control** over

All unchecked exceptions <: RuntimeException
Does not require **throws** keyword

Exceptions can only compile if:

The order of **catch** must be so that all exceptions are accessible (for checked exceptions)

Immutability

Advantages

1. Ease of understanding
2. Enabling safe sharing of objects
3. Enabling safe sharing of internals
4. Enabling safe concurrent execution

Nested Classes

Inner class access to outer class

Static nested inner class (only static vars)

Local class variable capture, effectively final

Anonymous class a local class without a name

Monads and Functors

Monads

Left-identity law:

Monad.**of**(x).**flatMap**(y -> f(y)) = f(x)

Right-identity law:

monad.**flatMap**(y -> Monad.**of**(y)) = monad

Associative law:

```
monad.flatMap(x -> f(x))
    .flatMap(x -> g(x)) =
monad.flatMap(x -> f(x))
    .flatMap(x -> g(x))
```

Functors

Identity

functor.**map**(x -> x) = functor

Composition:

```
functor.map(x -> f(x)).map(x -> g(x))
= functor.map(x -> g(f(x)))
```

Streams

```
allMatch(Predicate<? super T> predicate)
anyMatch(Predicate<? super T> predicate)
noneMatch(Predicate<? super T> predicate)
takeWhile(Predicate<? super T> predicate)
```

// Bounded

```
count()
distinct()
reduce(BinaryOperator<T> accumulator)
toArray()
```

// Mutation

```
flatMap(Function<? super T,
    ? extends Stream<? extends R>> mapper)
map(Function<? super T, ? extends R> mapper)
filter(Predicate<? super T> predicate)
```

// Parallel

```
parallel()
sequential()
unordered()
```

When using reduce,

1. combiner.**apply**(identity, i) = i
2. Combiner and accumulator must be associative
3. accumulator.**apply**(u, t) = combiner.**apply**(u, accumulator.**apply**(identity, t))

Parallel

Concurrency Divides computation into subtasks

Parallelism Subtasks are running at same time

Asynchronous

CompletableFuture

// instantiation

```
runAsync(Runnable runnable)
supplyAsync(Supplier<U> supplier)
completedFuture(U value)
```

// mapping

```
thenApply(Function<? super T,
    ? extends U> fn)
```

// map

```
thenCompose(Function<? super T,
    ? extends CompletionStage<U>> fn)
// flatMap
```

```
thenCombine(CompletionStage<? extends U> other,
    BiFunction<? super T, ? super U,
    ? extends V> fn)
// combine
```

Fork-Join-Pull

Order

```
f1.fork();
f2.fork();
f2.join();
f1.join();
```

Fork adds task to the **front** of the queue

Join blocks computation until completed

Task Stealing idle threads will take from the **back** of the queue