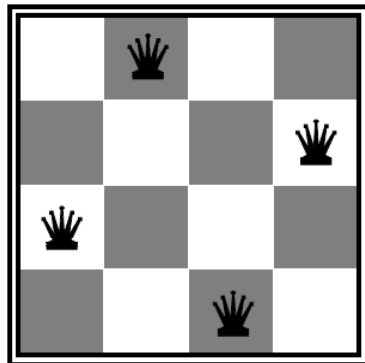
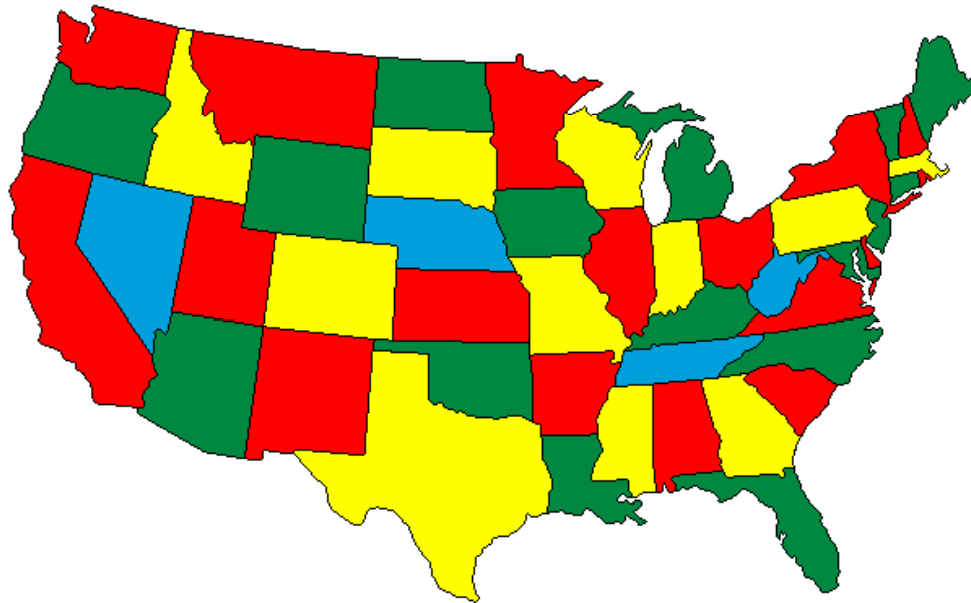


Constraint Satisfaction Problems



8			4		6			7
						4		
	1					6	5	
5		9		3		7	8	
				7				
	4	8		2		1		3
	5	2					9	
		1						
3			9		2			5

Constraint satisfaction problems (CSPs)

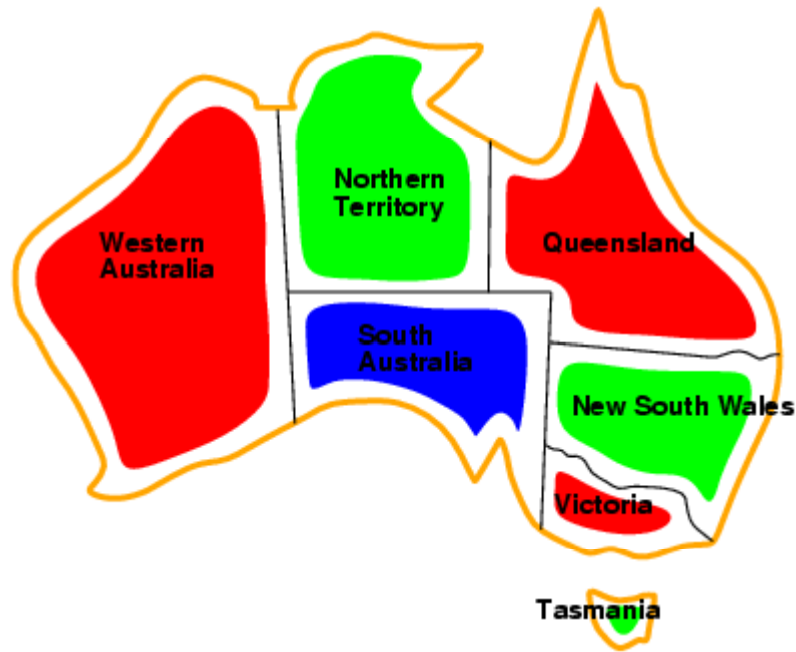
- Standard search problem:
 - **State** is a “black box” – any data structure that supports successor function, heuristic function, and goal test
- Constraint satisfaction problem:
 - **State** is defined by **variables** X_i with **values** from **domain** D_i
 - **Goal test** is a set of **constraints** specifying allowable combinations of values for subsets of variables
- A simple example of a formal representation language
- Allows useful **general-purpose** algorithms with more power than standard search algorithms

Example: Map Coloring



- **Variables:** WA, NT, Q, NSW, V, SA, T
- **Domains:** {red, green, blue}
- **Constraints:** adjacent regions must have different colors
e.g., $WA \neq NT$, or $(WA, NT) \in \{(\text{red}, \text{green}), (\text{red}, \text{blue}), (\text{green}, \text{red}), (\text{green}, \text{blue}), (\text{blue}, \text{red}), (\text{blue}, \text{green})\}$

Example: Map Coloring



- **Solutions** are *complete* and *consistent* assignments, e.g., WA = red, NT = green, Q = red, NSW = green, V = red, SA = blue, T = green

Example: N-Queens

- **Variables:** X_{ij}
- **Domains:** $\{0, 1\}$
- **Constraints:**

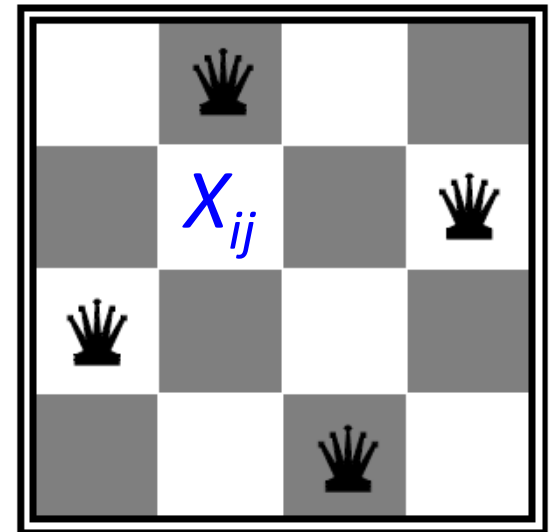
$$\sum_{i,j} X_{ij} = N$$

$$(X_{ij}, X_{ik}) \in \{(0, 0), (0, 1), (1, 0)\}$$

$$(X_{ij}, X_{kj}) \in \{(0, 0), (0, 1), (1, 0)\}$$

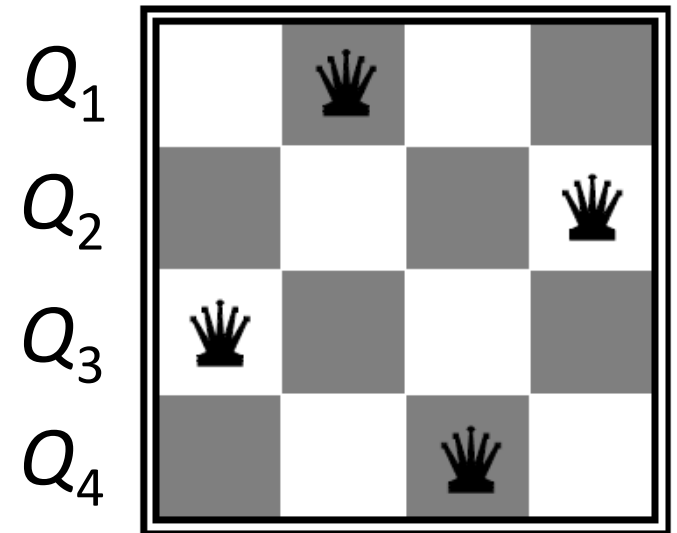
$$(X_{ij}, X_{i+k, j+k}) \in \{(0, 0), (0, 1), (1, 0)\}$$

$$(X_{ij}, X_{i+k, j-k}) \in \{(0, 0), (0, 1), (1, 0)\}$$



N-Queens: Alternative formulation

- **Variables:** Q_i
- **Domains:** $\{1, \dots, N\}$
- **Constraints:**
 $\forall i, j$ non-threatening (Q_i, Q_j)



Example: Cryptarithmic

- **Variables:** T, W, O, F, U, R

X_1, X_2

- **Domains:** $\{0, 1, 2, \dots, 9\}$

- **Constraints:**

$\text{Alldiff}(T, W, O, F, U, R)$

$$O + O = R + 10 * X_1$$

$$W + W + X_1 = U + 10 * X_2$$

$$T + T + X_2 = O + 10 * F$$

$$T \neq 0, F \neq 0$$

$$\begin{array}{r}
 X_2 \ X_1 \\
 T \ W \ O \\
 + \ T \ W \ O \\
 \hline
 F \ O \ U \ R
 \end{array}$$

Example: Sudoku

- **Variables:** X_{ij}
- **Domains:** $\{1, 2, \dots, 9\}$
- **Constraints:**
 $\text{Alldiff}(X_{ij} \text{ in the same } unit)$

					8			4
	8	4		1	6			
			5			1		
1		3	8			9		
6		8		X_{ij}		4		3
		2			9	5		1
		7			2			
			7	8		2	6	
2			3					

Real-world CSPs

- Assignment problems
 - e.g., who teaches what class
- Timetable problems
 - e.g., which class is offered when and where?
- Transportation scheduling
- Factory scheduling
- More examples of CSPs: <http://www.csplib.org/>

Standard search formulation (incremental)

- **States:**
 - Values assigned so far
- **Initial state:**
 - The empty assignment { }
- **Successor function:**
 - Choose any unassigned variable and assign to it a value that does not violate any constraints
 - Fail if no legal assignments
- **Goal test:**
 - The current assignment is complete and satisfies all constraints

Standard search formulation (incremental)

- What is the depth of any solution?
 - n (with n variables assigned)
 - This is the good news (why?)
- Given that there are m possible values for any variable, how many paths are there in the search tree?
 - $n! \cdot m^n$
 - This is the bad news
- How can we reduce the branching factor?

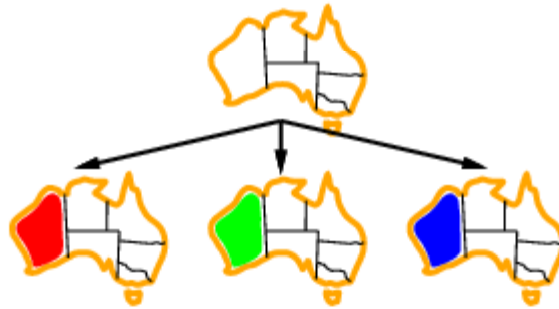
Backtracking search

- In CSP's, variable assignments are **commutative**
 - For example, $[WA = \text{red then } NT = \text{green}]$ is the same as $[NT = \text{green then } WA = \text{red}]$
- We only need to consider assignments to a single variable at each level (i.e., we fix the order of assignments)
 - Then there are only m^n leaves
- Depth-first search for CSPs with single-variable assignments is called **backtracking search**

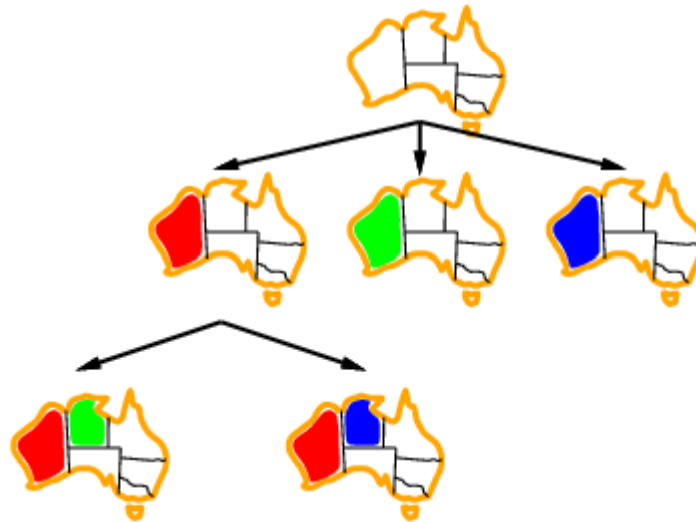
Example



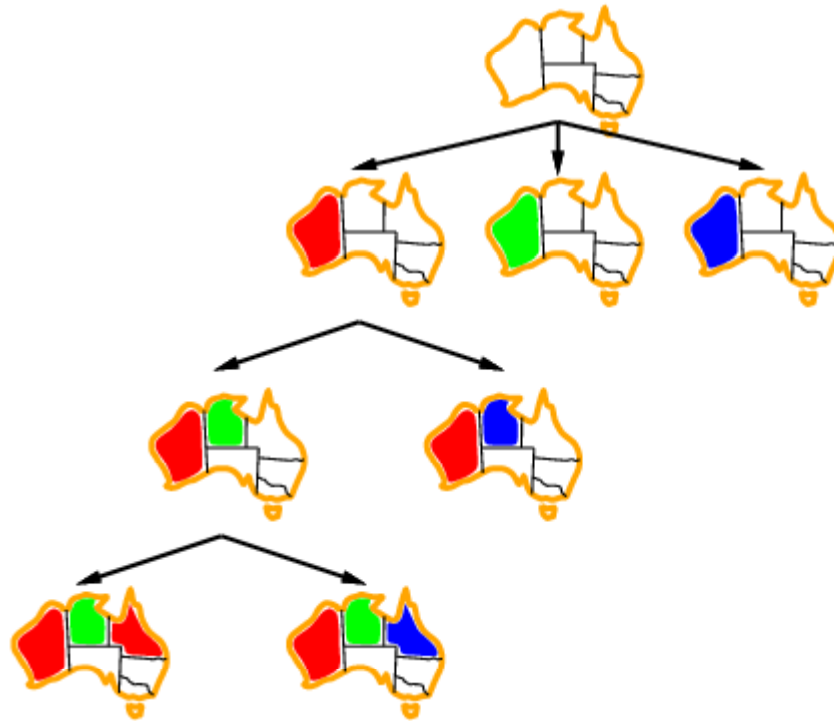
Example



Example



Example



Backtracking search algorithm

```
function RECURSIVE-BACKTRACKING(assignment, csp)  
  if assignment is complete then return assignment  
  var ← SELECT-UNASSIGNED-VARIABLE(VARIABLES[csp], assignment, csp)  
  for each value in ORDER-DOMAIN-VALUES(var, assignment, csp)  
    if value is consistent with assignment given CONSTRAINTS[csp]  
      add {var = value} to assignment  
      result ← RECURSIVE-BACKTRACKING(assignment, csp)  
      if result ≠ failure then return result  
      remove {var = value} from assignment  
  return failure
```

- Improving backtracking efficiency:
 - Which variable should be assigned next?
 - In what order should its values be tried?
 - Can we detect inevitable failure early?