-Create a new tbl with foreign key and unique constraint CREATE TABLE 'person tbl' ('id' INT AUTO INCREMENT PRIMARY KEY, 'first name' VARCHAR(255) NOT NULL, 'last name' VARCHAR(255) NOT NULL, 'dob' DATE. CONSTRAINT full name UNIOUE (first name,last name). FOREIGN KEY('sid') REFERENCES 'student tbl' ('id') ENGINE=InnoDB: --Add new entry INSERT INTO client (first_name,last_name,dob,sid) VALUES ('Sara', 'Smith', '1970-01-02'); --Add new entry with foreign key INSERT INTO project (cid.name.notes) VALUES ((SELECT id FROM client WHERE first name='Sara' AND last name='Smith'), 'diamond'. 'Should be done by Jan 2017'); --find film with max length and min rental duration separate tbl SELECT film id FROM film WHERE length = (SELECT MAX(length) FROM film AS max len) rental duration=(SELECT MIN(rental duration) FROM film AS min rent); -- the number of times ed chases acts in each category SELECT cate.name, COUNT(act.actor id) FROM category cate LEFT JOIN film category AS fc ON cate.category id=fc.category id LEFT JOIN film AS f ON fc.film id=f.film id LEFT JOIN film actor AS fa ON f.film id=fa.film id LEFT JOIN actor AS act ON (fa.actor id=act.actor id) AND (act.first name='ED' AND act.last name='CHASE') GROUP BY cate.category id ORDER BY cate.name -- length of time each actor has starred in sci-fi movies SELECT act.first name, act.last name, sci fi.sum FROM actor act LEFT JOIN (SELECT act.actor id AS actor id, SUM(f.length) AS sum FROM actor act INNER JOIN film actor fa ON act.actor id=fa.actor id INNER JOIN film f ON fa.film id=f.film id INNER JOIN film category fc ON f.film id=fc.film id INNER JOIN category cate ON fc.category id=cate.category id WHERE cate.name='Sci-Fi' GROUP BY act actor id AS sci fi ON act.actor id=sci fi.actor id

```
-- find the actors who never starred in sci-fi films
SELECT act.first name, act.last name FROM actor act
WHERE act.actor id NOT IN
  (SELECT act actor id FROM actor act
    INNER JOIN film actor fa ON act.actor id=fa.actor id
    INNER JOIN film f ON fa.film id=f.film id
    INNER JOIN film category fc ON f.film id=fc.film id
    INNER JOIN category cate ON fc.category id=cate.category id
WHERE cate.name='Sci-Fi')
GROUP BY act.actor id;
-- joining two different queries into one
(SELECT f.title FROM film f
  INNER JOIN film actor fa ON f.film id=fa1.film id
  INNER JOIN actor act ON fa1.actor id=act1.actor id
  INNER JOIN film actor fal ON f.film id=fa.film id
  INNER JOIN actor act1 ON fa.actor id=act.actor id
WHERE (act.first_name='WARREN')
  AND (act1.first_name='KIRSTEN')
ORDER BY f.title DESC;
-- Other SOL
SELECT * FROM foo tbl
DELETE FROM foo tbl WHERE col=...;
UPDATE uni student
SET firstname='Kevin', lastname='Turkington' WHERE id=1
-- Attribute Types
INT, BIGINT, FLOAT, DECIMAL, VARCHAR(255), TEXT,
DATE
-- Aggregate Functions
Works on multiple rows, Group by is used to group rows into sets
AVG(/* col or select */). SUM(/* col or select */). COUNT()
-- Joins
INNER JOIN: intersection of two tables
UNION (full outer): combines two tables
LEFT JOIN: all records from the MAIN and matched from the
ioining tables.
RIGHT JOIN: all records from the JOINING and matched from the
MAIN tables.
JOIN: all from JOINING TO MAIN
-- Lavers of Abstraction
External: front end user data must make sence
Conceptual: how data is organized for the database desiginer
Internal: Mappings to actual data on physical storage media
-- Helpful MvSOL hints
```

- -Delete can violate Referential integrity (foreign keys)
- -Insert can violate Referential and domain integrity (data length or
- -Cascade delete for a foreign key + delete all corresponding data will also be deleted

-- Schemas

exist for entities and relationships and are composed of attributes and constraints

--Integrity constraints

- -Domain: restricted domain of an attribute (INT, VARCHAR, ETC)
- -Key: requires that entries in a column combo must be unque
- -Referential: requires that an attribute be present in another table (foreign key)
- -Semantic: rules about the system outside of the database
- -NOT NULL: requires a value to be specified
- -Entity: primary key cannot be null

-- Database normalization

- -reduction of redundant data
- -functional dependencies if x determine y, y is functionally dependent on x. x-> y
- 1NF: ex. tuples fixed by dividing into seperate columns
- 2NF: Every non primary attribute must not be partially dependant on another col. Ex. course name can be turned into a lookup table.
- 3NF: non primary attribute can be transitively dependant on any key.

Ex. role priority on a person and role on a person is redundant. Fixed by creating a lookup table.

- BCNF: like 3NF but requires no attribute be transitively dependent on any key. Ex. toppings table has meat(beef) and cheese(cheddar) those can be its own table (meat table, cheese table).

-- Relational Algebra

 $\sigma_{\text{col}=\text{val}}$: Where $\Pi_{\text{col}1 \text{ col}2 \text{ col}3}$: Select $\blacktriangleright \blacktriangleleft_{\text{(col}=\text{col or val)}}$: Union

X- cross product AxB

- ∩- Intersection: must have identical domains
- U Union: will combine duplicate rows, must have identical domains
- (-) Difference: A-B, must have identical domains
- σ selection, filter rows
- Π projection, filter columns

Examples:

(Model ► (Model.model id=Vehicle.fk model id) Vehicle)

 $\Pi_{\text{make.make name, model.model name}}$ (

σ_{year=1976} (Make • (Make.make_id=Vehicle.fk_make_id)</sub>

Model ► (Model.model_id=Vehicle.fk model id) Vehicle)))

 $\sigma_{\text{foo a} < 100} (\text{foo } \triangleright \P_{\text{foo b=bar c}} \text{Barr})$