

Homework 11

Problem 1. Use the Lovasz local lemma to show that if

$$4 \binom{k}{2} \binom{n}{k-2} 2^{1-\binom{k}{2}} \leq 1$$

then it is possible to color the edges of K_n with two colors so that it has no monochromatic (i.e. single color) K_k subgraph.

Solution. E_i : the i -th K_k is monochromatic. $Pr(E_i) = 2^{1-\binom{k}{2}}$. Consider the dependency graph, for any different E_i and E_j , they are adjacent if the corresponding K_k share at least one edge. Thus the degree of the dependency graph is bounded by $\binom{k}{2} \binom{n}{k-2}$.

According to the Lovasz local lemma, it is possible that none of the E_i happens under the given inequality. \square

Problem 2. What is the expected number of trees with k vertices in $G \in \mathcal{G}(n, p)$?

Solution. By Cayley's formula and the linearity of expectation, it is $\binom{n}{k} k^{k-2} p^{k-1}$. \square

Problem 3. suppose that $p = \frac{d}{n}$ where $d = o(n^{1/3})$. Show that $G \in \mathcal{G}(n, p)$ almost surely has no copies of K_4 .

Solution. The probability that there exists some 4-vertices sub-graph which induces a K_4 is upper bounded by $p = \binom{n}{4} p^6$. When $d = o(n^{1/3})$, $\lim_{n \rightarrow \infty} p = 0$. \square

Problem 4. Show that if almost all $G \in \mathcal{G}(n, p)$ have a graph property \mathcal{P}_1 and almost all $G \in \mathcal{G}(n, p)$ have a graph property \mathcal{P}_2 , then almost all $G \in \mathcal{G}(n, p)$ have both properties.

Solution. The portion of the graphs have both properties equals 1 minus the portion of the graphs which does not have property \mathcal{P}_1 or \mathcal{P}_2 . However the portion of the graph does not have property \mathcal{P}_1 or \mathcal{P}_2 is bounded by the sum of the portion of the graphs does not have property \mathcal{P}_1 and the the portion of the graphs does not have property \mathcal{P}_2 , which both tend to 0 as n approaches ∞ . The claim in the question then follows. \square

Problem 5. Consider $\mathbf{G}(n, p)$ with $p = \frac{1}{3n}$.

Use the second moment method to show that with high probability there exists a simple path of length 10.

Solution. $I_\ell = \begin{cases} 1 & \ell \text{ is a length 10 simple path} \\ 0 & \text{otherwise} \end{cases}$

X is the number of simple path, then $X = \sum_\ell I_\ell$.

The expectation would be $E(X) = \frac{1}{2}(n)_{11} \times p^{10}$. Thus

$$(E(X))^2 = \frac{1}{4} [(n)_{11}]^2 \times p^{20} \quad (\star)$$

(Note (\star) is about $\Theta(n^2)$.)

Then to calculate $E(X^2)$.

$$E(X^2) = E\left[\left(\sum_\ell I_\ell\right)^2\right] = E\left[\sum_\ell I_\ell \sum_{\ell'} I_{\ell'}\right] = \sum_{\ell, \ell'} E(I_\ell I_{\ell'}) \quad (\Delta)$$

ℓ and ℓ' will be independent to each other unless they have common vertices or edges. We use $k = |\ell \cap \ell'|$ to stand for the number of common edges between ℓ and ℓ' , and $s = |\ell \cap \ell'|$ to stand for the number of common vertices used by ℓ and ℓ' . Obviously $(0 \leq k \leq 10) \wedge (0 \leq s \leq 11) \wedge (k \geq 1 \rightarrow s \geq k + 1)$. (Δ) can be divided into the following subcases:

1. $k = 0$

(a) $s = 0$: $\sum_{\ell, \ell'} E(I_\ell I_{\ell'}) = \frac{1}{8}(n)_{22} \times p^{20} \leq (E(X))^2$;

(b) $1 \leq s \leq 11$:

for each s , $\sum_{\ell, \ell'} E(I_\ell I_{\ell'}) = c \cdot (n)_{22-s} \times (p)^{20} = o((E(X))^2)$, where c is a constant number.

2. $1 \leq k \leq 10$ ($2 \leq s \leq 11$)

The general formula of each of these cases (constant many) would be

$$\begin{aligned} \sum_{\ell, \ell'} E(I_\ell I_{\ell'}) &= d \cdot (n)_{22-s} \times p^{20-k} \\ &\leq d \cdot (n)_{22-s} \times p^{20-(s-1)} \\ &= d \cdot (n)_{22-s} \times p^{21-s} \\ &= o((E(X))^2) \end{aligned}$$

Combining the above results we get that $Var(X) = E(X^2) - (E(X))^2 = \mathbf{o}((E(X))^2)$.
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