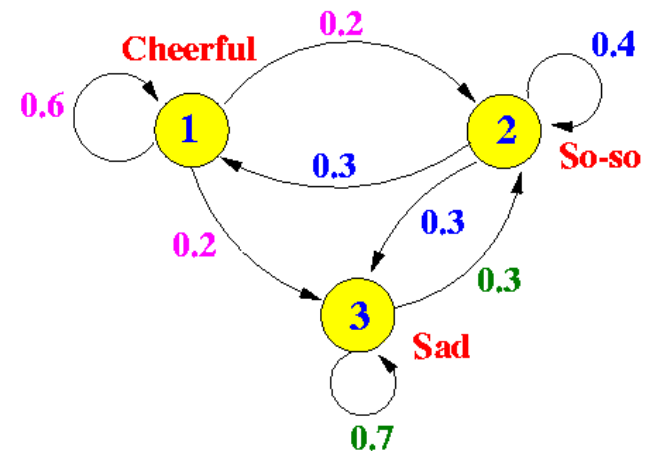


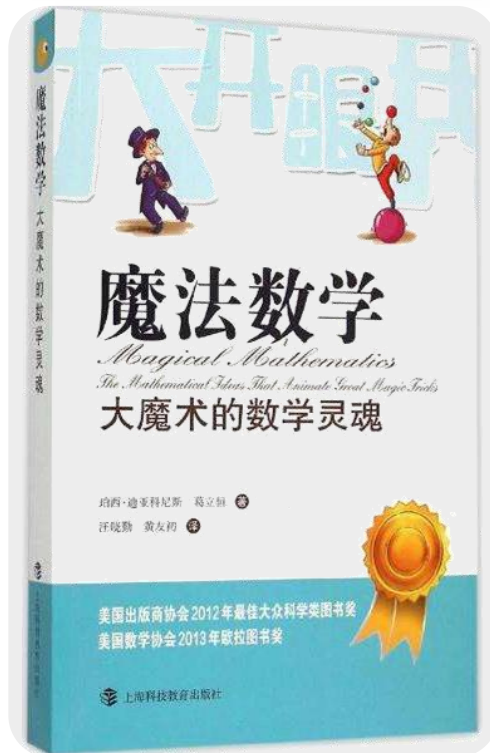
# Random Walks and Markov Chains

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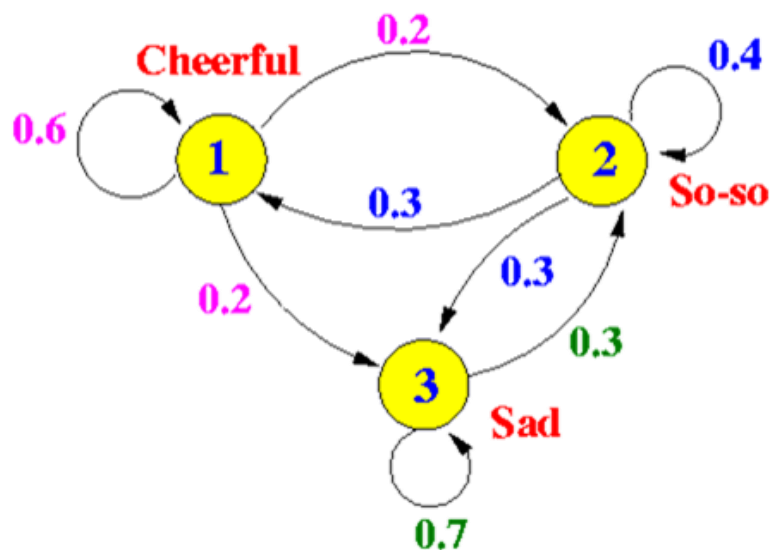


- **戴康尼斯** (Persi Diaconis, 1945年1月31日 - ) : 美国数学家、统计学家。斯坦福大学的数学与统计学教授。
- 他解决了一些随机性的问题，包括掷币和洗牌。1992年，他和David Bayer证明完美的洗牌至少要洗七次。他又和说明从高处跌下的猫为何总能以脚着地的Richard Montgomery合作，证明了掷币哪面向上，物理因素比运气重要得多。
- 自14岁，他便跟随一个叫Dai Vernon的魔术师行走江湖。后来在赌场，他尝试研究防止他和其他魔术师被骗的方法。他18岁时买了一本An Introduction to Probability and Its Applications，但因为不懂微积分而看不明。24岁，他在City College of New York上数学课。其间他在《科学美国人》杂志投稿，介绍了他的两个纸牌戏法。Martin Gardner认为那两个戏法十分精彩，注意到他的才华，为他写了一封推荐信。当时，哈佛大学的统计学家Fred Mosteller正在研究魔术，因此决定让Diaconis成为他的研究生。



# Random walk

- **Random walk.** on a directed graph, a sequence of vertices generated from a start vertex by probabilistically selecting an incident edge, traveling the edge to a new vertex, and repeat the process.



# Random walk

**Probability distribution.**  $p = [p_1, p_2, \dots, p_n]$ , where  $\sum_{i=1}^n p_i = 1$

**Starting.**  $p = p(0) = [p_1(0), p_2(0), \dots, p_n(0)]$ ,  $\sum_{i=1}^n p_i(0) = 1$   
and  $p_x$  is the probability of starting at  $x$ .

The probability of being at vertex  $x$  at time  $t + 1$ :

$$p_x(t + 1) = \sum_{(y,x) \in E} p_y(t) \cdot \Pr(y \rightarrow x)$$

**Transition Matrix  $P$ :**  $P_{ij}$  is the probability of the walk at vertex  $i$  selecting the edge to vertex  $j$ .

$$p(t) \cdot P = p(t + 1)$$

# Random walk

**Fundamental property.** in the limit, the long-term average property of being at a particular vertex is *independent of the start vertex*, or an initial probability distribution over vertices (provided the underlying graph is strongly connected) – the *stationary probabilities*.

# Markov chain

- A finite set  $S$  of **states**
- **Transition probability**: For  $x, y \in S$ ,  $p_{xy}$  is the probability going from state  $x$  to  $y$ .
- $\sum_y p_{xy} = 1$

Markov chain  $\leftarrow$  Random graph

- ① A vertex  $\leftarrow$  a state
- ②  $p_{xy} \leftarrow$  weighted edge from  $x$  to  $y$ .

# Markov chain

Markov chain  $\leftarrow$  Random graph

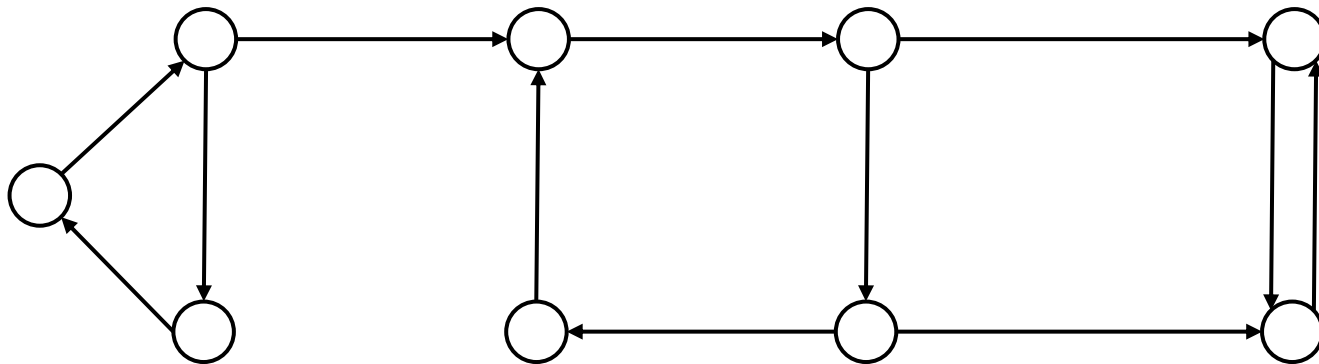
- ① A vertex  $\leftarrow$  a state
- ②  $p_{xy} \leftarrow$  weighted edge from  $x$  to  $y$ .

**Connected Markov chain (irreducible):** if the underlying directed graph is strongly connected.

**Transition probability matrix  $P$ :**  $P_{xy}$  is the probability of changing from state  $x$  to  $y$ .

# Markov chain

**Persistent state (recurrence).** If the state ever be reached, the random process will return to it with probability 1.



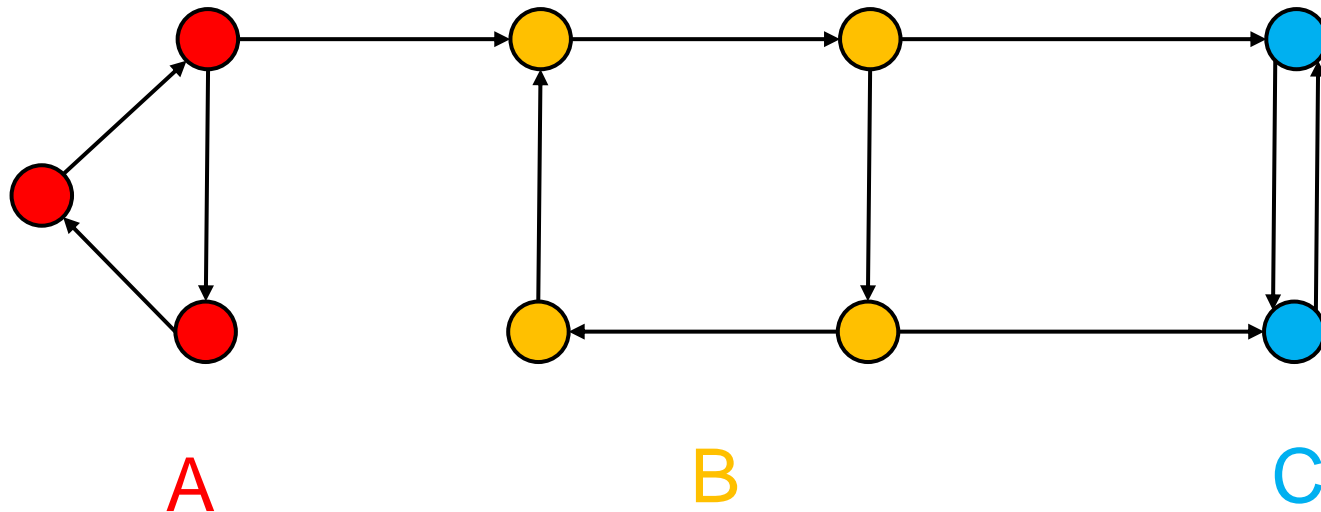
$T(S)$ : a random variable shows the number of steps the chain first goes back to state  $S$ .

Persistent:  $\Pr(T(S) < \infty) = 1$ , then  $S$  is persistent.



# Markov chain

**Persistent state (recurrence).** If the state ever be reached, the random process will return to it with probability 1.



# Markov chain

**Aperiodic.** If the greatest common divisor of the lengths of directed cycles is one.

Random walk	Markov Chain
Graph	Stochastic process
Vertex	State
Strongly connected	Persistent/recurrence
Aperiodic	Aperiodic
Strongly connected +Aperiodic	Ergodic
Undirected graph	Time reversible

We will assume strong connectness by default.

# Stationary distribution

$p(t)$  : the probability distribution after  $t$  steps of a random walk.

Long-term average probability distribution:

$$a(t) = \frac{1}{t} (p(0) + p(1) + \cdots + p(t-1))$$

**Fundamental theorem of Markov chains:**

For a connected MC,  $a(t)$  converges to a limit probability  $x$  which satisfies  $x \cdot P = x$ .

# Fundamental Theorem

**Lemma 1:** Let  $P$  be the transition probability matrix for a connected Markov chain. The  $n \times (n + 1)$  matrix  $A = [P - I, \mathbf{1}]$  obtained by augmenting the matrix  $P - I$  with an additional column of ones has rank  $n$ .

**Fundamental Theorem of Markov Chains:** For a connected Markov chain there is a unique vector  $\pi$  satisfying  $\pi \cdot P = \pi$ . Moreover, for any starting distribution,  $\lim_{t \rightarrow \infty} a(t)$  exists and equals  $\pi$ .

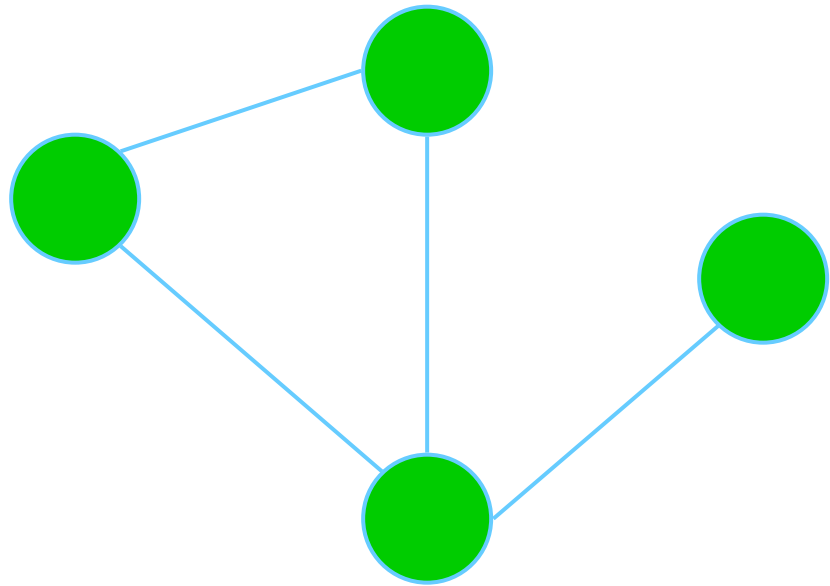
**Lemma 2 (Time reversible):** For a random walk on a strongly connected graph with probabilities on the edge, if the vector  $\pi$  satisfies  $\pi_x p_{xy} = \pi_y p_{yx}$  for all  $x$  and  $y$  and  $\sum_x \pi_x = 1$ , then  $\pi$  is the stationary distribution of the walk.

# Application

$$G = (V, E)$$

$$|V| = n, |E| = m$$

$$\deg(v_i) = d_i$$



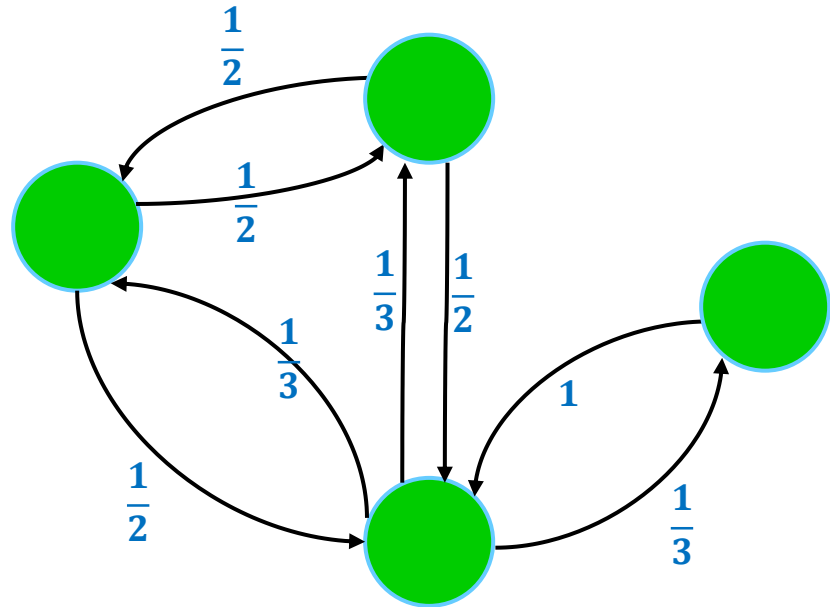
# Application

$$G = (V, E)$$

$$|V| = n, |E| = m$$

$$\deg(v_i) = d_i$$

$$P_{i \rightarrow j} = \begin{cases} \frac{1}{\deg_i} & \text{if } (i, j) \in E \\ 0 & \text{otherwise.} \end{cases}$$



Then the stationary distribution is:

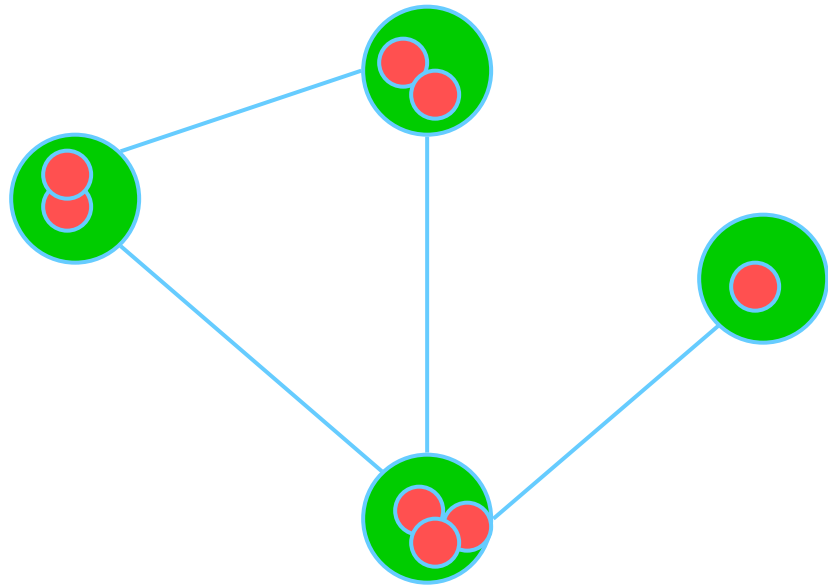
# Application

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$$P_{i \rightarrow j} \begin{cases} \frac{1}{\deg_i} & \text{if } (i, j) \in E \\ 0 & \text{otherwise.} \end{cases}$$



Then the stationary distribution is:  $\pi = \left[ \frac{d_1}{2m}, \frac{d_2}{2m}, \dots, \frac{d_n}{2m} \right]$

# Markov Chain Monte Carlo

**MCMC.** A technique for sampling a multivariate probability distribution  $p(\mathbf{x})$ , where  $\mathbf{x} = (x_1, x_2, \dots, x_d)$ .

**Application.** to estimate the expected value of a function  $f(\mathbf{x})$

$$E(f) = \sum_{\mathbf{x}} f(\mathbf{x}) \cdot p(\mathbf{x})$$



# Markov Chain Monte Carlo

**Application.** to estimate the expected value of a function  $f(x)$

$$E(f) = \sum_x f(x) \cdot p(x)$$

Realization:

- ① Draw a set of samples. Each sample  $x$  is selected with probability  $p(x)$ .
- ② Averaging  $f$  over these samples.

# Markov Chain Monte Carlo

Sample according to  $p(\mathbf{x})$ . Design a MC whose states correspond to the value space of  $\mathbf{x}$  and whose stationary probability distribution is  $p(\mathbf{x})$ .

Recall:

- ✓  $p(t)$  is the row vector of probabilities of the random walk being at each state at time  $t$ .
- ✓  $\mathbf{a}(t) = \frac{1}{t} (p(0) + p(1) + \cdots + p(t-1))$

$$E(r) = \sum_i f_i \left( \frac{1}{t} \sum_{j=1}^t \Pr(\text{walk is in state } i \text{ at time } j) \right) = \sum_i f_i a_i(t)$$

# Markov Chain Monte Carlo

Sample according to  $p(\mathbf{x})$ . Design a MC whose states correspond to the value space of  $\mathbf{x}$  and whose *stationary probability distribution* is  $p(\mathbf{x})$ .

$$E(r) = \sum_i f_i \left( \frac{1}{t} \sum_{j=1}^t \Pr(\text{walk is in state } i \text{ at time } j) \right) = \sum_i f_i a_i(t)$$

$$\begin{aligned} \left| \sum_i f_i p_i - E(r) \right| &\leq f_{\max} \cdot \sum_i |p_i - a_i(t)| \\ &= f_{\max} \cdot \boxed{\|p - a(t)\|_1} \end{aligned}$$

# Markov Chain Monte Carlo

Sample according to  $p(\mathbf{x})$ . Design a MC whose states correspond to the value space of  $\mathbf{x}$  and whose stationary probability distribution is  $p(\mathbf{x})$ .

Two general approach:

- The Metropolis-Hastings algorithm
- The Gibbs sampling

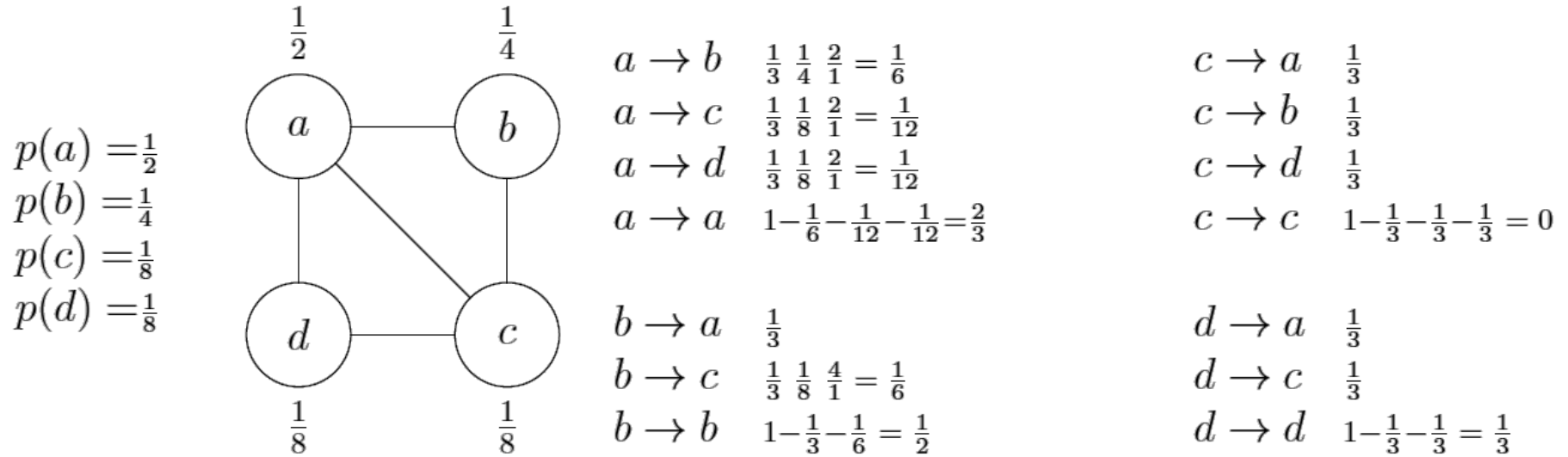
# Metropolis-Hastings Algorithm

**MHA.** A general method to design a Markov chain whose stationary distribution is a given target distribution  $p$ .

Given random graph  $G$ , with  $\Delta(G) = r$ . The transitions of the MC are defined as

$$\begin{cases} p_{ij} = \frac{1}{r} \min\left(1, \frac{p_j}{p_i}\right) \\ p_{ii} = 1 - \sum_{i \neq j} p_{ij} \end{cases}$$

Given random graph  $G$ , with  $\Delta(G) = r$ . The transitions of the MC are defined as  $\begin{cases} p_{ij} = \frac{1}{r} \min\left(1, \frac{p_j}{p_i}\right) \\ p_{ii} = 1 - \sum_{i \neq j} p_{ij} \end{cases}$ .



$$p(a) = p(a)p(a \rightarrow a) + p(b)p(b \rightarrow a) + p(c)p(c \rightarrow a) + p(d)p(d \rightarrow a)$$

$$= \frac{1}{2} \frac{2}{3} + \frac{1}{4} \frac{1}{3} + \frac{1}{8} \frac{1}{3} + \frac{1}{8} \frac{1}{3} = \frac{1}{2}$$

$$p(b) = p(a)p(a \rightarrow b) + p(b)p(b \rightarrow b) + p(c)p(c \rightarrow b)$$

$$= \frac{1}{2} \frac{1}{6} + \frac{1}{4} \frac{1}{2} + \frac{1}{8} \frac{1}{3} = \frac{1}{4}$$

$$p(c) = p(a)p(a \rightarrow c) + p(b)p(b \rightarrow c) + p(c)p(c \rightarrow c) + p(d)p(d \rightarrow c)$$

$$= \frac{1}{2} \frac{1}{12} + \frac{1}{4} \frac{1}{6} + \frac{1}{8} 0 + \frac{1}{8} \frac{1}{3} = \frac{1}{8}$$

$$p(d) = p(a)p(a \rightarrow d) + p(c)p(c \rightarrow d) + p(d)p(d \rightarrow d)$$

$$= \frac{1}{2} \frac{1}{12} + \frac{1}{8} \frac{1}{3} + \frac{1}{8} \frac{1}{3} = \frac{1}{8}$$

# Metropolis-Hastings Algorithm

Given random graph  $G$ , with  $\Delta(G) = r$ . The transitions of the MC are defined as

$$\begin{cases} p_{ij} = \frac{1}{r} \min\left(1, \frac{p_j}{p_i}\right) \\ p_{ii} = 1 - \sum_{i \neq j} p_{ij} \end{cases}$$

## Correctness.

To prove the stationary distribution is indeed the target distribution  $\mathbf{p}$ .

$$p_i p_{ij} = \frac{p_i}{r} \min\left(1, \frac{p_j}{p_i}\right) = \frac{1}{r} \min(p_i, p_j) = \frac{p_j}{r} \min\left(1, \frac{p_i}{p_j}\right) = p_j p_{ji}$$

# Gibbs Sampling

Let  $p(\mathbf{x})$  be the target distribution where  $\mathbf{x} = (x_1, \dots, x_d)$ .  
Now the undirected random graph is a hyper cube:  
there is an edge between  $\mathbf{x}$  and  $\mathbf{y}$  if  $\mathbf{x}$  and  $\mathbf{y}$  differ in only 1 coordinate.

Sampling process: for  $\mathbf{x} = (x_1, \dots, x_d)$

- ① Choose one of the  $x_i$  to update;
- ②  $x_i'$  is chosen based on the marginal probability of  $x_i$

$$p_{xy} = \frac{1}{d} \cdot p(y_i | x_1 \cdots x_{i-1} x_{i+1} \cdots x_d)$$

where  $x_i \neq y_i$  and  $x_j = y_j$  for all  $i \neq j$ ,  
(i.e.,  $x_{i \neq j}$  does not change).

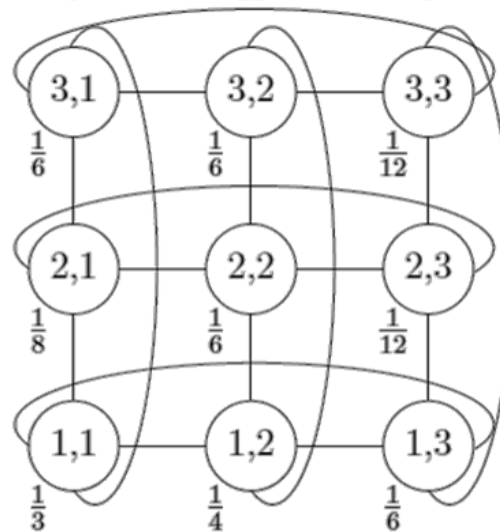


Sampling process: for  $x = (x_1, \dots, x_d)$

① Choose one of the  $x_i$  to update;

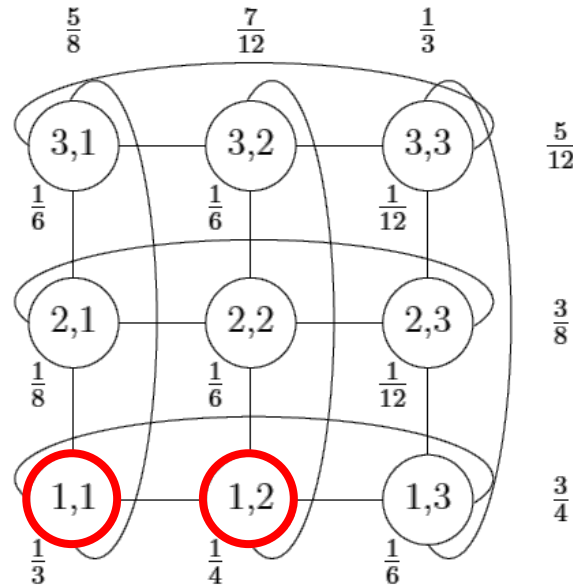
②  $x_i'$  is chosen based on the marginal probability of  $x_i$  (i.e.,  $x_{i \neq j}$  will not change).

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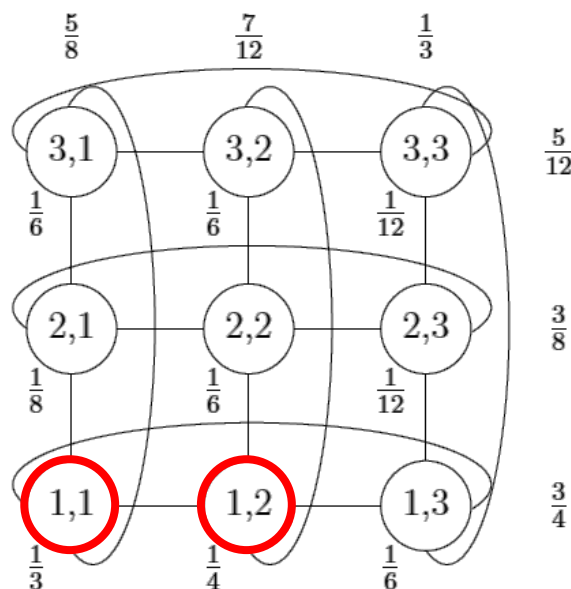


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$$p_{xy} = \frac{1}{d} \cdot p(y_i | x_1 \cdots x_{i-1} x_{i+1} \cdots x_d), \text{ where } x_i \neq y_i \text{ and } x_j = y_j \text{ for all } i \neq j.$$



$$p_{(11)(12)} = \frac{1}{d} p_{12} / (p_{11} + p_{12} + p_{13}) = \frac{1}{2} \left( \frac{1}{4} \right) / \left( \frac{1}{3} \frac{1}{4} \frac{1}{6} \right) = \frac{1}{8} / \frac{9}{12} = \frac{1}{8} \frac{4}{3} = \frac{1}{6}$$

$$\begin{array}{lll} p_{(11)(12)} = \frac{1}{2} \frac{1}{4} \frac{4}{3} = \frac{1}{6} & p_{(12)(11)} = \frac{1}{2} \frac{1}{3} \frac{4}{3} = \frac{2}{9} & p_{(13)(11)} = \frac{1}{2} \frac{1}{3} \frac{4}{3} = \frac{2}{9} \\ p_{(11)(13)} = \frac{1}{2} \frac{1}{6} \frac{4}{3} = \frac{1}{9} & p_{(12)(13)} = \frac{1}{2} \frac{1}{6} \frac{4}{3} = \frac{1}{9} & p_{(13)(12)} = \frac{1}{2} \frac{1}{4} \frac{4}{3} = \frac{1}{6} \\ p_{(11)(21)} = \frac{1}{2} \frac{1}{8} \frac{8}{5} = \frac{1}{10} & p_{(12)(22)} = \frac{1}{2} \frac{1}{6} \frac{12}{7} = \frac{1}{7} & p_{(13)(23)} = \frac{1}{2} \frac{1}{12} \frac{3}{1} = \frac{1}{8} \\ p_{(11)(31)} = \frac{1}{2} \frac{1}{6} \frac{8}{5} = \frac{2}{15} & p_{(12)(32)} = \frac{1}{2} \frac{1}{6} \frac{12}{7} = \frac{1}{7} & p_{(13)(33)} = \frac{1}{2} \frac{1}{12} \frac{3}{1} = \frac{1}{8} \end{array}$$

# Gibbs Sampling

Sampling process: for  $x = (x_1, \dots, x_d)$

- ① Choose one of the  $x_i$  to update;
- ②  $x_i'$  is chosen based on the marginal probability of  $x_i$  (i.e.,  $x_{i \neq j}$  will not change).  $p_{xy} = \frac{1}{d} \cdot p(y_i | x_1 \cdots x_{i-1} x_{i+1} \cdots x_d)$ , where  $x_i \neq y_i$  and  $x_j = y_j$  for all  $i \neq j$ .

**Correctness.** To prove the stationary distribution is indeed the target distribution  $p$ .

$$\begin{aligned} p_{x\mathbf{y}} &= \frac{1}{d} \frac{p(\mathbf{y}_i | x_1 \cdots x_{i-1} x_{i+1} \cdots x_d) p(x_1 \cdots x_{i-1} x_{i+1} \cdots x_d)}{p(x_1 \cdots x_{i-1} x_{i+1} \cdots x_d)} \\ &= \frac{1}{d} \frac{p(x_1 \cdots x_{i-1} \mathbf{y}_i x_{i+1} \cdots x_d)}{p(x_1 \cdots x_{i-1} x_{i+1} \cdots x_d)} = \frac{1}{d} \frac{p(\mathbf{y})}{p(x_1 \cdots x_{i-1} x_{i+1} \cdots x_d)} \end{aligned}$$

$$\text{Similarly } p_{\mathbf{y}x} = \frac{1}{d} \frac{p(x)}{p(x_1 \cdots x_{i-1} x_{i+1} \cdots x_d)}$$

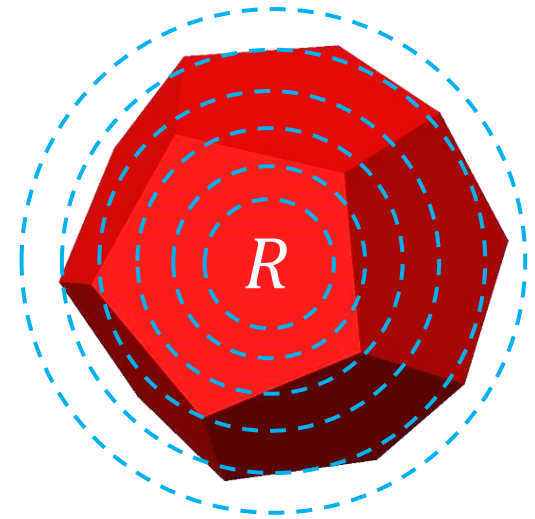
It follows that  $p(\mathbf{x})p_{x\mathbf{y}} = p(\mathbf{y})p_{\mathbf{y}x}$ .

# Areas and Volumes

For general convex sets in  $d$  space, there are no close form formulae for volume.

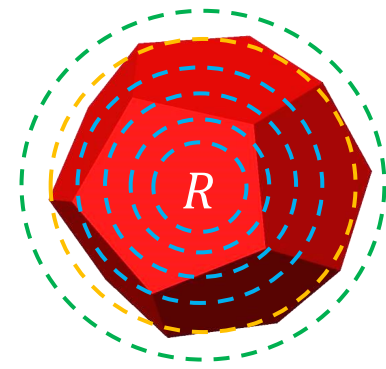
Sequence of concentric spheres:

$$R \supset S_1 \subseteq S_2 \subseteq \dots \subseteq S_k \supset R$$



$$\begin{aligned} Vol(R) &= Vol(S_k \cap R) \\ &= \frac{Vol(S_k \cap R)}{Vol(S_{k-1} \cap R)} \cdot \frac{Vol(S_{k-1} \cap R)}{Vol(S_{k-2} \cap R)} \cdots \frac{Vol(S_2 \cap R)}{Vol(S_1 \cap R)} \cdot Vol(S_1) \end{aligned}$$

# Areas and Volumes



$$Vol(R) = \frac{Vol(\textcolor{teal}{S}_k \cap R)}{Vol(\textcolor{brown}{S}_{k-1} \cap R)} \cdot \frac{Vol(S_{k-1} \cap R)}{Vol(S_{k-2} \cap R)} \cdots \frac{Vol(S_2 \cap R)}{Vol(S_1 \cap R)} \cdot Vol(S_1)$$

$$Radius(S_i) = \left(1 + \frac{1}{d}\right) \cdot Radius(S_{i-1})$$

$$\text{Thus } 1 \leq \frac{Vol(S_i \cap R)}{Vol(S_{i-1} \cap R)} = \left(1 + \frac{1}{d}\right)^d \leq e$$

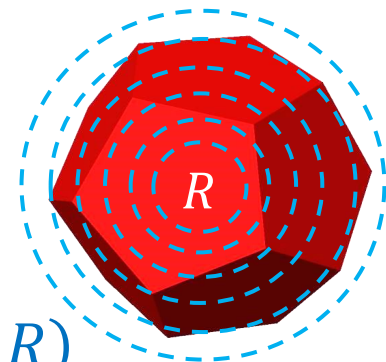
Let  $r = \left(1 + \frac{1}{d}\right)^k$  then the **number of spheres  $k$**  is at most

$$O\left(\log_{1+\frac{1}{d}} r\right) = O(d \ln(r))$$

To estimate the overall volume to error  $1 \pm \epsilon$ :

Estimate each **volume ratio** to a factor of  $1 \pm \frac{\epsilon}{ed \ln(r)}$ .

# Areas and Volumes

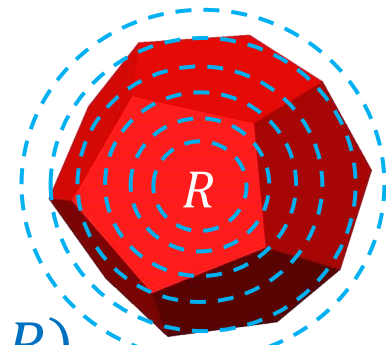


$$\text{Radius}(S_i) = \left(1 + \frac{1}{d}\right) \cdot \text{Radius}(S_{i+1}), \quad 1 \leq \frac{\text{Vol}(S_i \cap R)}{\text{Vol}(S_{i-1} \cap R)} \leq e$$

Estimate the ratio  $\frac{\text{Vol}(S_i \cap R)}{\text{Vol}(S_{i-1} \cap R)}$ :

- ① Selecting points in  $S_i \cap R$  uniformly at random;
- ② Computing the fraction in  $S_{i-1} \cap R$ .

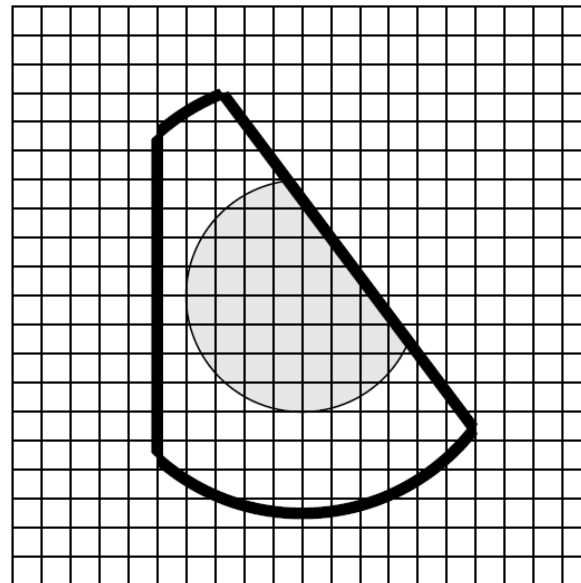
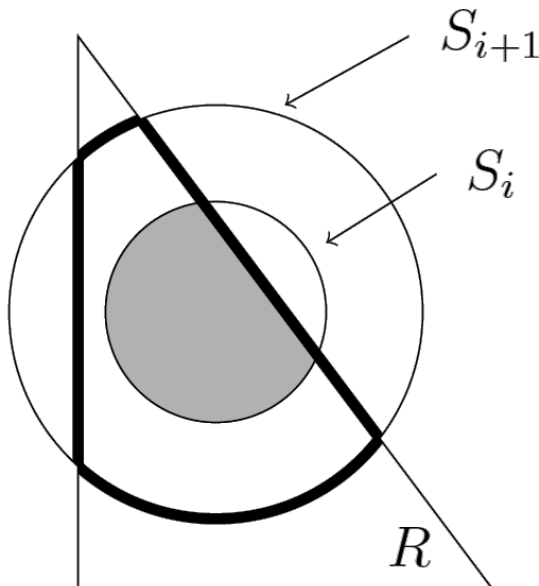
# Areas and Volumes



$$\text{Radius}(S_i) = \left(1 + \frac{1}{d}\right) \cdot \text{Radius}(S_{i+1}), \quad 1 \leq \frac{\text{Vol}(S_i \cap R)}{\text{Vol}(S_{i-1} \cap R)} \leq e$$

Estimate the ratio  $\frac{\text{Vol}(S_i \cap R)}{\text{Vol}(S_{i-1} \cap R)}$ :

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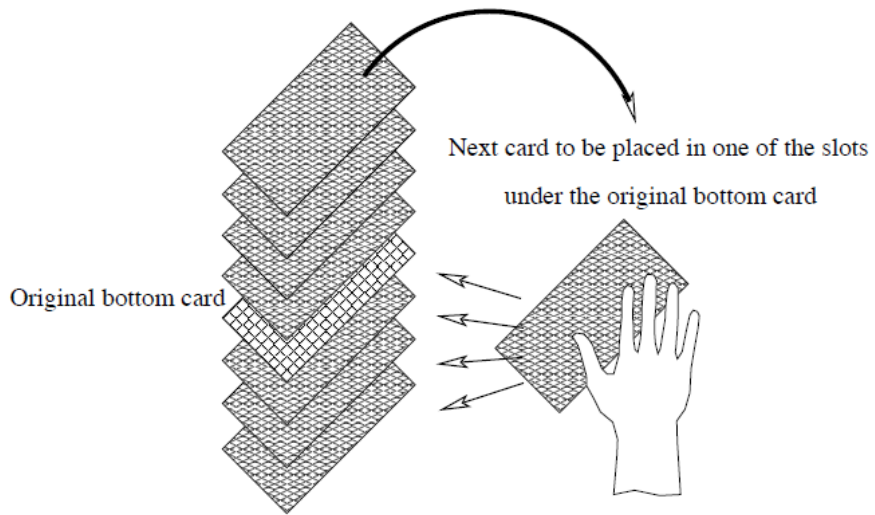


# Some conceptions

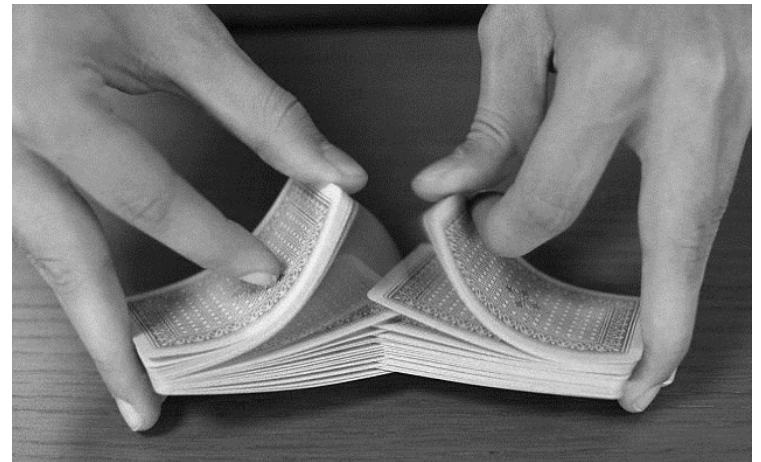
**Mixing time.** Fix  $\epsilon > 0$ . The  $\epsilon$  –mixing time of a MC is the minimum integer  $t$  such that for any starting distribution  $p_0$ , the 1-norm distance between the  $t$  –step running average probability distribution and the stationary distribution is at most  $\epsilon$ .

**Hitting time  $h_{xy}$ .** The expected time of a random walk starting at vertex  $x$  (or a starting probability distribution) to reach vertex  $y$ .

**Cover time.** The expected time of a random walk starting at vertex  $x$  in the graph  $G$  to reach each vertex at least once.

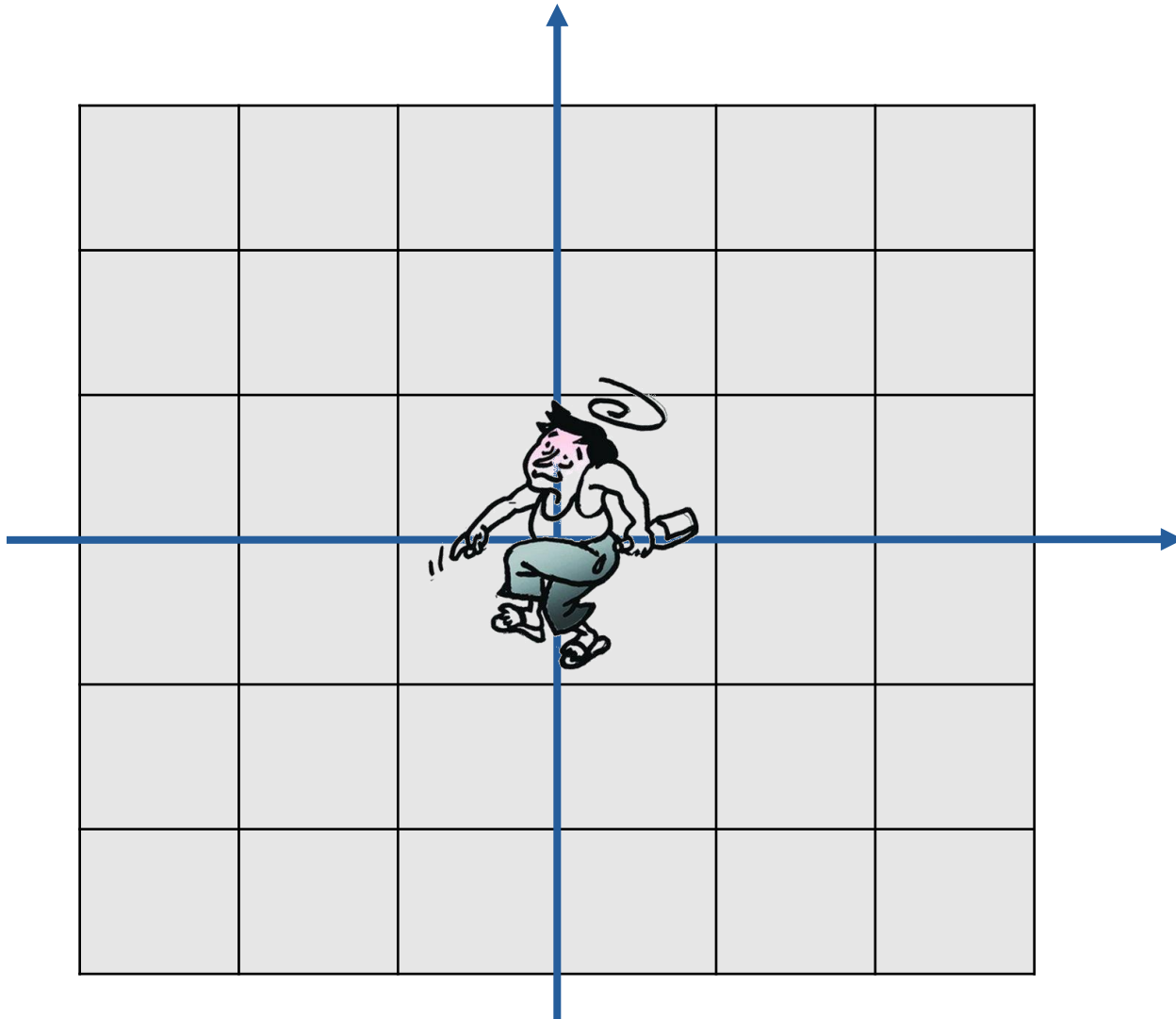


$$O(n \ln n)$$



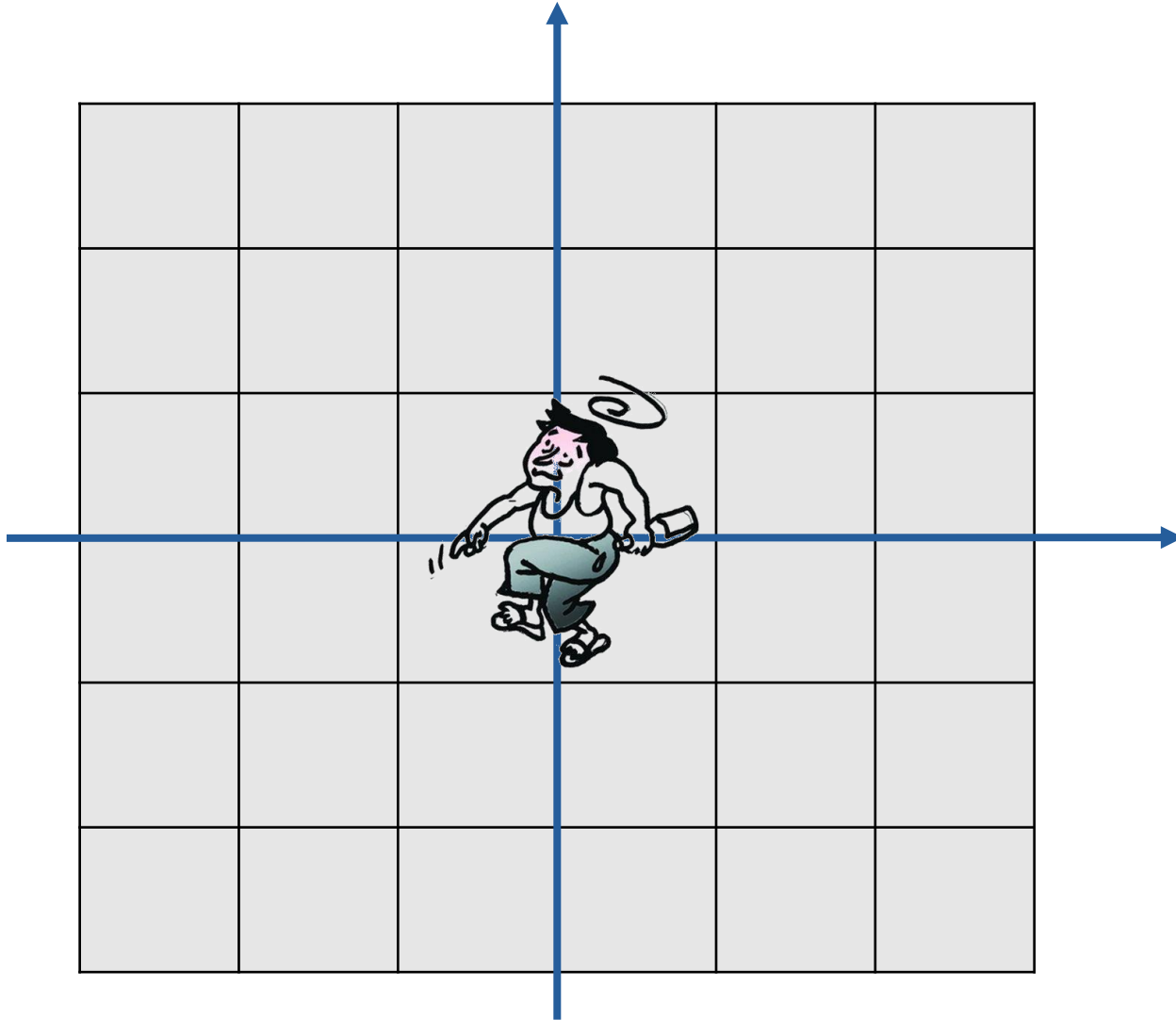
$$O(\ln n)$$

# Random walks in Euclidean Space



George Pólya, 1921

# Random walks in Euclidean Space



# Random walks in Euclidean Space



# Random walks in Euclidean Space



“A drunk person will always find their way home, while a drunk bird may get lost forever.”