## **Virtual Lib**

Let  $p_{\mathsf{fail}}$ , a failure probability and  $z^*$  ( $p_{\mathsf{fail}}$ ), the associated standard score i.e.

$$z^*\left(p_{\mathsf{fail}}
ight) = \sqrt{2} \cdot \mathsf{erf}^{-1}(1-p_{\mathsf{fail}})$$

An error happens when we decode the plaintext and instead of getting the right value, we end up with our value plus or minus a small error i.e.  $\left\lfloor \frac{\Delta \cdot m + e}{\Delta} \right\rceil = m + \left\lfloor \frac{e}{\Delta} \right\rceil$ . The purpose of this document is to find the distribution X of this error.

To satisfy the  $p_{\text{fail}}$  constraint, the optimizer will find parameters such that with e the noise in the ciphertext coming from a normal distribution  $\mathcal{N}(0, \sigma^2)$  and  $\Delta$ , the scaling factor (where the message starts), we have

$$\mathbb{P}\left(|e|<rac{\Delta}{2}
ight)=1-p_{\mathsf{fail}} ext{ i.e. } \mathbb{P}\left(\underbrace{\left\lfloorrac{e}{\Delta}
ight
ceil}_{=F}=0
ight)=1-p_{\mathsf{fail}}$$

by enforcing that  $z^*\left(p_{\mathsf{fail}}\right)\cdot\sigma\leq\Delta/2$  . To simplify the rest of the document, we will assume  $z^*\left(p_{\mathsf{fail}}\right)\cdot\sigma=\Delta/2$ 

and so we have

$$\mathbb{P}\left(F=0
ight)=\mathbb{P}\left(\left|e
ight|<rac{\Delta}{2}
ight)=1-p_{\mathsf{fail}}$$

Now let's find  $\mathbb{P}(F = k), \forall k$ 

$$egin{aligned} \mathbb{P}\left(F=1
ight) &= \mathbb{P}\left(rac{\Delta}{2} < |e| < rac{3 \ \Delta}{2}
ight) \ &= \mathbb{P}\left(|e| < rac{3 \ \Delta}{2}
ight) - \mathbb{P}\left(|e| < rac{\Delta}{2}
ight) \ &= \mathbb{P}\left(|e| < 3 \cdot z^* \left(p_{\mathsf{fail}}
ight) \cdot \sigma
ight) - (1 - p_{\mathsf{fail}}) \end{aligned}$$

We can once again use the normal confidence interval to compute the first probability as we have

$$\mathbb{P}\left(\left|e
ight|<3\cdot z^{st}\left(p_{\mathsf{fail}}
ight)\cdot\sigma
ight)=\mathsf{erf}\left(rac{3\cdot z^{st}\left(p_{\mathsf{fail}}
ight)}{\sqrt{2}}
ight)$$

We can keep going for other values

$$egin{aligned} \mathbb{P}\left(F=2
ight) &= \mathbb{P}\left(rac{3}{2} igwedge \left|e
ight| < rac{5}{2} 
ight) \ &= \mathbb{P}\left(\left|e
ight| < rac{5}{2} rac{\Delta}{2} 
ight) - \mathbb{P}\left(\left|e
ight| < rac{3}{2} rac{\Delta}{2} 
ight) \ &= \mathbb{P}\left(\left|e
ight| < 5 \cdot z^* \left(p_{\mathsf{fail}}
ight) \cdot \sigma
ight) - \mathbb{P}\left(\left|e
ight| < 3 \cdot z^* \left(p_{\mathsf{fail}}
ight) \cdot \sigma
ight) \end{aligned}$$

More generally, we have

$$egin{aligned} orall k > 0, \mathbb{P}\left(F = k
ight) &= \mathbb{P}\left(rac{\Delta}{2} + (k-1)\cdot\Delta < |e| < rac{\Delta}{2} + k\cdot\Delta
ight) \ &= \mathbb{P}\left(|e| < \Delta\cdot(k + rac{1}{2})) - \mathbb{P}\left(\Delta\cdot(k - rac{1}{2}) < |e|
ight) \ &= \mathbb{P}\left(|e| < z^*\left(p_{\mathsf{fail}}
ight)\cdot(2\ k + 1)\cdot\sigma
ight) \ &- \mathbb{P}\left(z^*\left(p_{\mathsf{fail}}
ight)\cdot(2\ k - 1)\cdot\sigma < |e|
ight) \end{aligned}$$

Using the formula below, we can compute the distribution of the errors:

$$\mathbb{P}\left(\left|e
ight|<\left(2k+1
ight)\cdot z^{*}\left(p_{\mathsf{fail}}
ight)\cdot\sigma
ight)=\mathsf{erf}\left(rac{\left(2k+1
ight)\cdot z^{*}\left(p_{\mathsf{fail}}
ight)}{\sqrt{2}}
ight)$$