

# Reduction of Key Sizes on Rainbow-like Multivariate Signature Schemes

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# Outline

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# Context

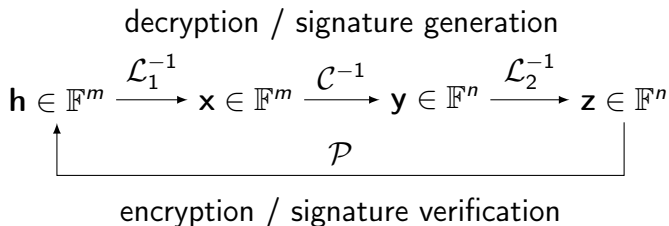
- ▶ Guarantee authenticity of messages sent digitally
- ▶ Security of digital signature schemes is based on problems from number theory
  - ▶ Integer factorization, discrete logarithm
- ▶ There exist quantum algorithms [Sho97] that solve such problems efficiently
- ▶ Post-quantum cryptography aims to create cryptosystems based on problems immune to quantum speed-ups

# Motivation

- ▶ Foreshadowing of quantum computers
- ▶ Several active branches of post-quantum cryptography based on distinct mathematical structures
- ▶ Standardization calls by institutions such as NIST, IRTF and ETSI
- ▶ We focus on cryptosystems that are built upon the difficulty of solving systems of equations

# Multivariate cryptography

- ▶ Cryptography based on systems of multivariate equations over finite fields
- ▶ Bipolar construction, with central trapdoor:

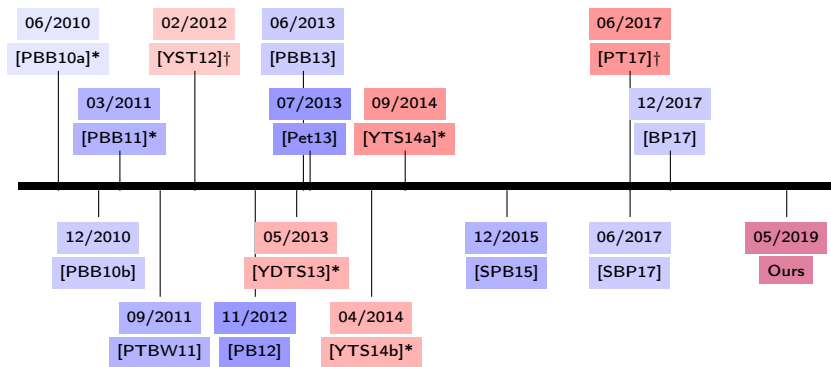


- ▶ Focus on signatures ( $m \leq n$ ): fast operations, small signatures and large keys, compared to current schemes

# Research object

- ▶ Focus on the Rainbow signature scheme [DS05], currently on Round 2 of the NIST standardization process
- ▶ Easy description, good balance between signature and key sizes
- ▶ Keys are one to two orders of magnitude greater than conventional ones (RSA at 512 bytes, ECDSA at 256 bits)
- ▶ Generalized version of Unbalanced Oil and Vinegar [KPG99]

# Related works



Works in blue optimise public keys, while red ones reduce private keys. Asterisks denote reparametrized works and crosses denote broken schemes.

# Hypothesis

- ▶ To the best of our knowledge, works have reduced either private or public keys, but not both
- ▶ Current methods to reduce keys are not compatible between themselves
- ▶ **Can both reductions be achieved simultaneously?**



# Rainbow signature scheme

## Preliminaries

- ▶ Parameters are a finite field  $\mathbb{F}_q$ , integers  $u, n$  such that  $u \leq n$  and  $0 < v_1 < \dots < v_u < v_{u+1} = n$
- ▶ For  $1 \leq \ell \leq u$ , set vinegar variables  $V_\ell = \{1, \dots, v_\ell\}$  and oil variables  $O_\ell = \{v_\ell + 1, \dots, v_{\ell+1}\}$
- ▶ Define vector spaces spanned by quadratic Oil-Vinegar polynomials

$$P_\ell = \sum_{i,j \in V_\ell} \alpha_{ij} x_i x_j + \sum_{i \in V_\ell, j \in O_\ell} \beta_{ij} x_i x_j + \sum_{i \in V_\ell \cup O_\ell} \gamma_i x_i + \delta,$$
$$\alpha_{ij}, \beta_{ij}, \gamma_i, \delta \in \mathbb{F}_q$$

# Rainbow signature scheme

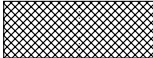



## Key generation

- ▶ Let  $m = n - v_1$  and  $o_\ell = v_{\ell+1} - v_\ell$
- ▶ Randomly pick two affine transformations  $\mathcal{L}_1 : \mathbb{F}_q^m \rightarrow \mathbb{F}_q^m$  and  $\mathcal{L}_2 : \mathbb{F}_q^n \rightarrow \mathbb{F}_q^n$
- ▶ Central map is a function  $\mathcal{C} : \mathbb{F}_q^n \rightarrow \mathbb{F}_q^m$ 
  - ▶ A total of  $o_\ell$  polynomials and respective coefficients are randomly chosen from each  $P_\ell$
- ▶ Private key is the 3-uple  $(\mathcal{L}_1, \mathcal{C}, \mathcal{L}_2)$ , public key is the composition  $\mathcal{P} = \mathcal{L}_1 \circ \mathcal{C} \circ \mathcal{L}_2$

# Rainbow signature scheme

## Inversion of the central map





- ▶ Vinegar variables of a layer are exactly the oil and vinegar variables from the previous layer
- ▶ This enables the inversion of each Oil-Vinegar layer recursively
- ▶ With  $u = 2$ , the initial configuration of  $\mathcal{C}$  is

	$V_1 \times V_1$	$V_1 \times O_1$	$O_1 \times O_1$	$V_1 \times O_2$	$O_1 \times O_2$	$O_2 \times O_2$	$V_1$	$O_1$	$O_2$	$\delta$
$\ell = 1$		0	0	0	0		0	★		
$\ell = 2$		0		★						



# Rainbow signature scheme

## Inversion of the central map

- Randomly choose variables in  $V_1$  and substitute them

	$V_1 \times V_1$	$V_1 \times O_1$	$O_1 \times O_1$	$V_1 \times O_2$	$O_1 \times O_2$	$O_2 \times O_2$	$V_1$	$O_1$	$O_2$	$\delta$
$\ell = 1$	★		0	0	0	0	★		0	★
$\ell = 2$	★					0	★			★

- Solve linear  $o_1$  equations in the first layer to obtain  $V_2$  (if possible), and then solve the remaining  $o_2$  equations

	$V_1 \times V_1$	$V_1 \times O_1$	$O_1 \times O_1$	$V_1 \times O_2$	$O_1 \times O_2$	$O_2 \times O_2$	$V_1$	$O_1$	$O_2$	$\delta$
$\ell = 2$	★	★	★		★	0	★	★		★

# Rainbow signature scheme

## Signature generation

- ▶ Consider a cryptographic hash function  $\mathcal{H}$  and a message  $M$ , and compute the digest  $\mathbf{h} = \mathcal{H}(M)$
- ▶ With possession of the private key, obtain the value  $\mathbf{x} = \mathcal{L}_1^{-1}(\mathbf{h})$
- ▶ Generate the pre-image of  $\mathbf{x}$  under the central map,  $\mathbf{y} = \mathcal{C}^{-1}(\mathbf{x})$ , as per the operations above
- ▶ Compute the final signature  $\mathbf{z} = \mathcal{L}_2^{-1}(\mathbf{y})$

# Rainbow signature scheme

## Signature verification

- ▶ Obtain  $\mathbf{h}$  from the message  $M$
- ▶ With possession of the public key, compute  $\mathbf{h}' = \mathcal{P}(\mathbf{z})$
- ▶ The signature is valid if  $\mathbf{h} = \mathbf{h}'$ , and invalid otherwise

# Our proposal

- ▶ Introduction of structures in the private key may be undesired
- ▶ Recall that, to invert the central map, vinegar variables are chosen randomly every time a preimage is computed
- ▶ If these variables are changed less often, or fixed, they simplify the central map matrix representation
- ▶ Indeed, the central map may be stored in a linearized fashion, and regenerated only occasionally

# Our proposal

- ▶ We propose to fix the  $V_1$  variables throughout the central map
- ▶ A possible implementation of this method is to use a PRNG to regenerate  $\mathcal{C}$  every time it is needed
- ▶ The linear relations described in CyclicRainbow [PBB10b] are another way to obtain the original central map
- ▶ The EUF-CMA variant described in [DCP<sup>+</sup>17] provides a salt that can be modified instead of vinegar variables



# Our proposal

- ▶ Only variables in  $V_1$  are fixed, since  $V_2$  and beyond need the digest to be calculated
- ▶ Strategy is not hindered by current cryptanalytic methods, since the choice of parameters does not change
- ▶ General framework for every Rainbow-like scheme
- ▶ Can be applied on top of variants that reduce the public key, confirming our hypothesis

# Open problems

- ▶ Given multiple signatures created with the same set of vinegar variables, is it possible to unveil information about the private key?
- ▶ Does every signature needs its own set of vinegar variables or is the cost of regenerating the central map amortized?
- ▶ Is it possible to create a constant-time implementation with this strategy?
- ▶ Do there exist parameter sets which optimize the private key size?

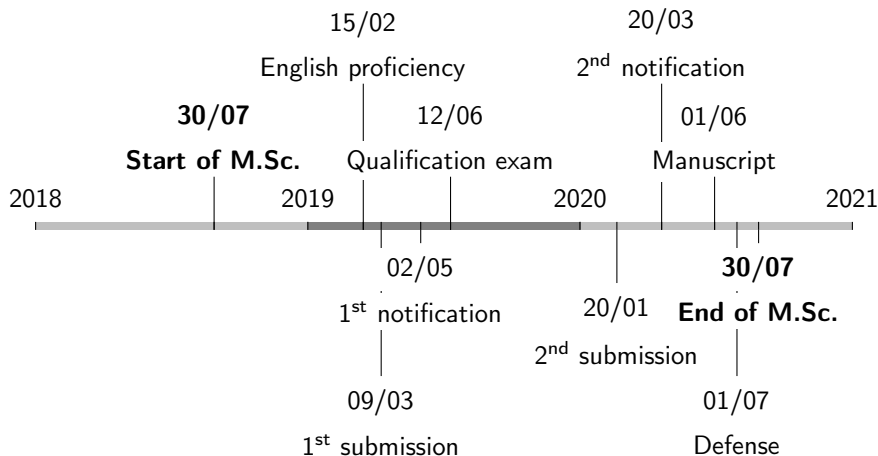
# Preliminary results [ZBC19]

Security	$n$	$m$	$ \mathcal{K}_{Pr} $	$ \mathcal{K}_{Pr}^\eta $	Difference
80	43	26	19208	5914	-69.21%
100	69	43	75440	23193	-69.26%
128	79	43	103704	22110	-78.68%
192	131	68	440638	71773	-83.71%
256	178	93	1086971	164721	-84.85%

# Preliminary results [ZBC19]

Security	Variant	$ \mathcal{K}_{Pr} $	$ \mathcal{K}_{Pr}^\eta $	$ \mathcal{K}_{Pu} $	Difference
80	Classic			25740	-28.76%
	Cyclic	19546	6524	10618	-62.15%
	LRS2			9789	-63.98%
100	Classic			60390	-31.60%
	Cyclic	46131	12474	22246	-67.41%
	LRS2			20662	-68.89%
128	Classic			139320	-32.78%
	Cyclic	105006	24924	48411	-69.98%
	LRS2			45547	-71.16%

# Chronology



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