Explanation of DFSB:

```
function search(csp):
    return backtrack({}, csp)

def backtrack(assignment, csp):
    if assignment is complete:
        return assignment
    else:
        var = Select_Unassigned_Var(csp)
        for each value in Domain_Value(var, csp):
            if value is consistent with assignment given Constraints(csp):
            add{var = value} to assignment
            if (backtrack(assignment, csp) != failure):
                return assignment
            remove{var = value} from assignment
            return failure
```

Plain DFSB loops through every value in the domain for each variable, and it tries to find out a solution that has no conflict with the constraints.

Explanation of DFSB++:

```
function search plus(csp):
  return backtrack_plus({}, csp)
def backtrack_plus(assignment, csp):
  if assignment is complete:
    return assignment
  else:
    var = Select Unassigned MCVar(csp)
    # values are ordered such that LCValue comes first
    for each value in Order_Domain_Value(var, csp):
       if value is consistent with assignment given Constraints(csp):
         AC-3(var, csp)
         add{var = value} to assignment
         if (backtrack_plus(assignment, csp) != failure):
            return assignment
         remove{var = value} from assignment
    return failure
```

DFSB++ is similar to DFSB, except that DFSB++ picks MCVar for var in each step, and it searches the values in an order that from LCValue first and MCValue last. Also, before assigning a consistent value to var and going to the next recursive step, DFSB++ performs AC-3 for the current csp to rule out some values in the domain of the neighbors of current var.

Explanation of Min-Conflicts:

```
function search_minconflicts(csp, max_steps):
    use random method to generate start_state
    return backtrack_plus(csp, max_steps, start_state)

function minconflicts(csp, max_steps, current_state)
    for i=1 to max_steps do
        if current_state is a solution of csp then return current_state
        var <-- a randomly chosen variable from the set of conflicted variables CONFLICTED[csp]
        value <-- the value v for var that minimizes CONFLICTS(var,v,current_state,csp)
        pr = random_number_from_1_to_10
        if pr <= 7:
            set var = value in current_state
        else:
            set var = a random value in the domain in current_state
        return failure
```

Min-Conflicts starts with a random generated state (assignment). In each step, a variable is chosen in random, and the value chosen should minimized the conflicts in the state (assignment). To adjust the algorithm to avoid plateaus or getting stuck in local depressions, this value will be assigned to the variable with a probability of 70%; and a random value will be assigned to the variable with a probability of 30%. If no solution was found after max_steps, return failure. In my implementation, I set max_steps to be 100*N.

Performance Table

	DFSB-	DFSB-	DFSB++-	DFSB++-	MinConflicts-	MinConflicts-
	easy	hard	easy	hard	easy	hard
time	1ms	Not	1ms	1727ms	3ms	172ms
		converging				
Search	8		8	990	113	1051
steps						
Prune			13	15328		
calls						

Difference Observation:

DFSB and DFSB++ have similar performance on the easy case. However, DFSB is not converging for the hard case. The reason for this is that I did not set a max_steps (the search steps allows before giving up). So, while DFSB is searching a result for the hard case, it can go to a very deep level of the states tree, while that could be the wrong trace.

For the Min-Conflicts, it will always return a solution for the easy case. Since 100*N random searches should be enough with high probability to find out a solution with respect to the M constraints and the randomly generated start_state for the easy case. However, for the hard case, the start_state could be totally messed up, and the 100*N search steps might not be enough to find out a solution. So Min-Conflicts might fail with a higher probability for the hard case.