

# Support Vector Machines

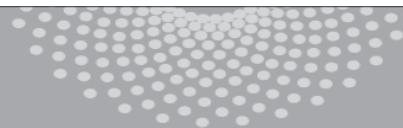
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Machine Learning 10-701



**MACHINE LEARNING** DEPARTMENT

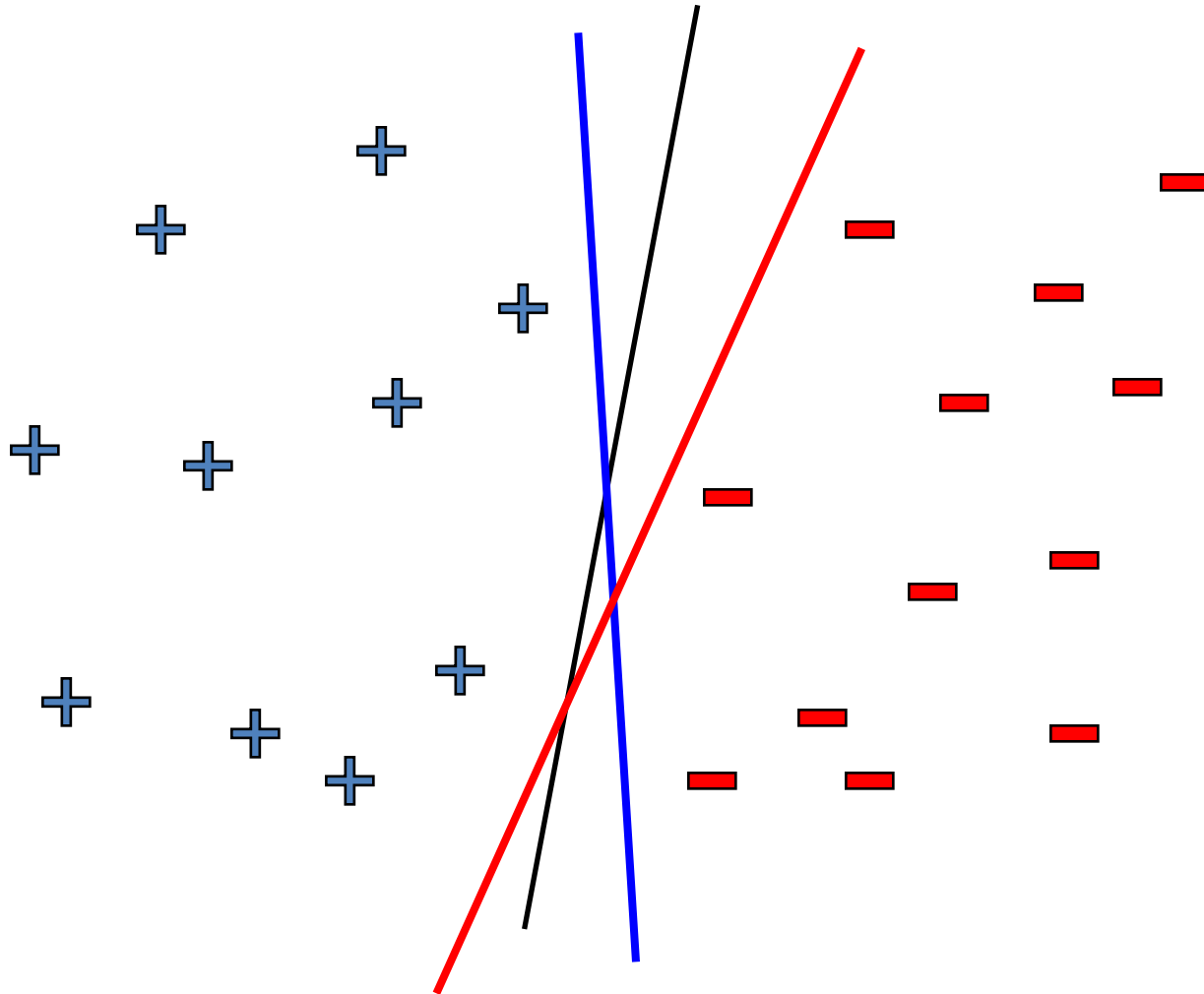


**Carnegie Mellon.**  
School of Computer Science

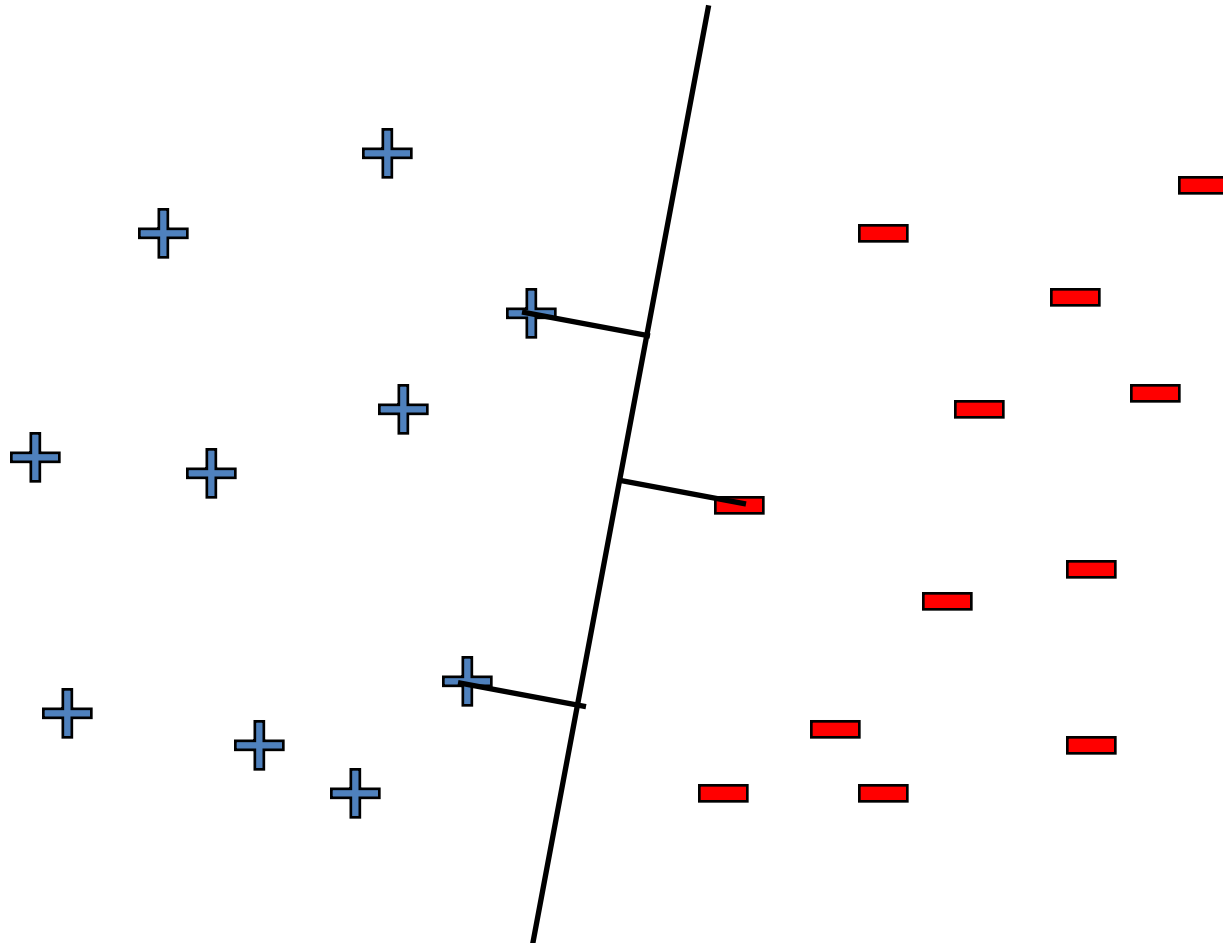
# At Pittsburgh G-20 summit ...



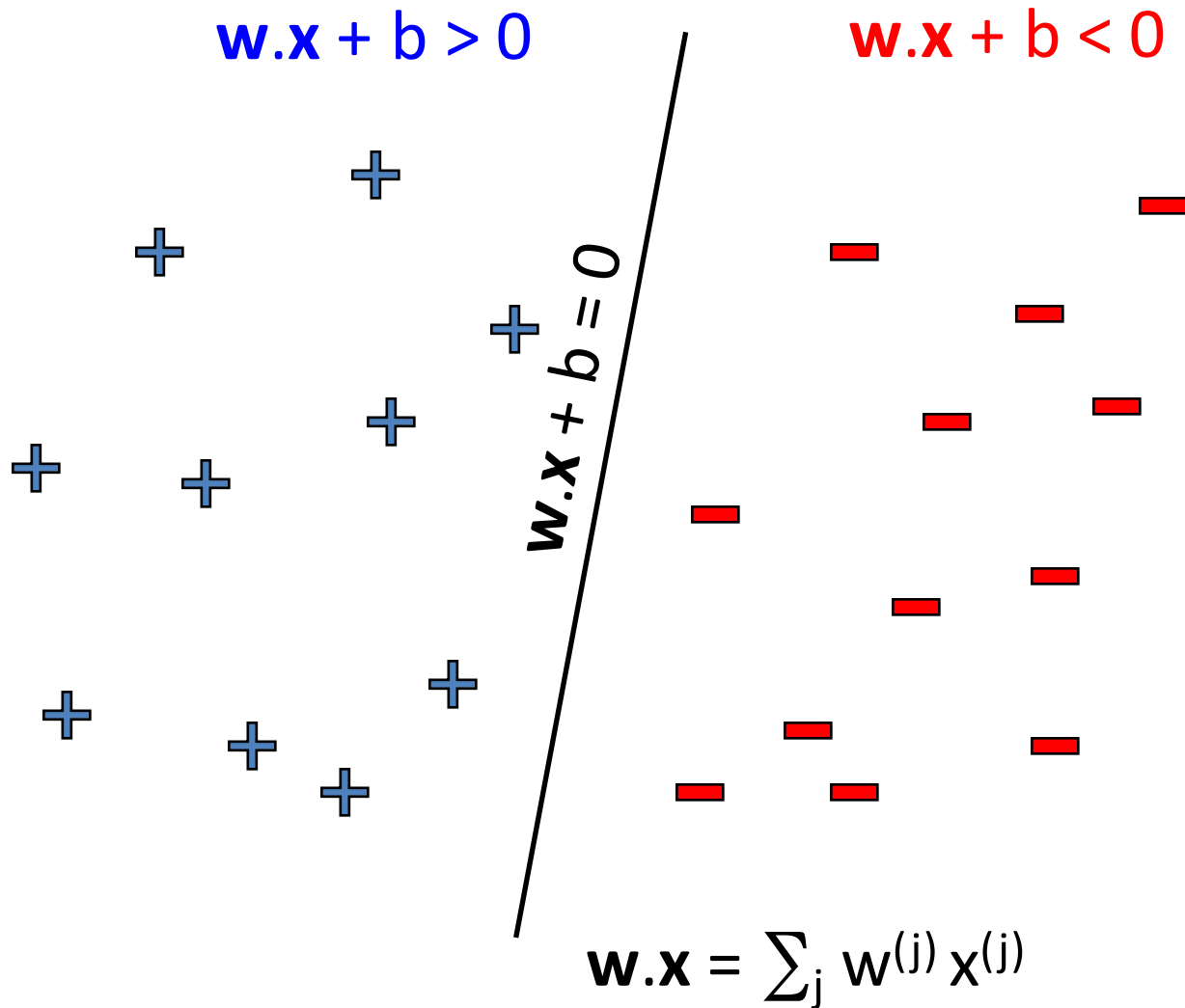
# Linear classifiers – which line is better?



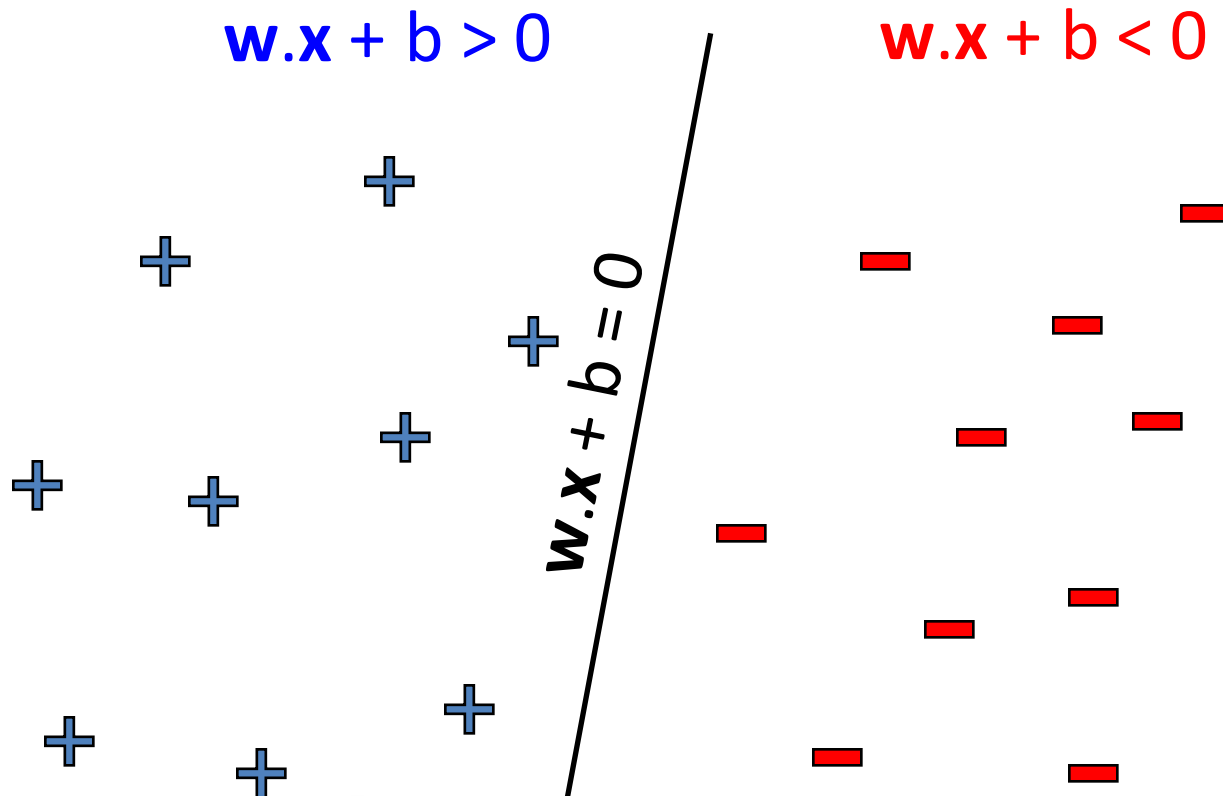
# Pick the one with the largest margin!



# Parameterizing the decision boundary



# Parameterizing the decision boundary



$y_j \in \{-1, +1\}$  — class

$$\text{"confidence"} = (w \cdot x_j + b) y_j$$

# Maximizing the margin

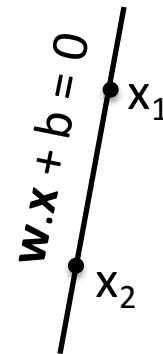
$$\mathbf{w} \cdot \mathbf{x} + b > 0$$

$$\mathbf{w} \cdot \mathbf{x} + b < 0$$

Distance of closest examples from the line/hyperplane

$$\text{margin} = \gamma = 2a / \|\mathbf{w}\|$$

Step 1:  $\mathbf{w}$  is perpendicular to lines since for any  $\mathbf{x}_1, \mathbf{x}_2$  on line  $\mathbf{w} \cdot (\mathbf{x}_1 - \mathbf{x}_2) = 0$



# Maximizing the margin

$$w \cdot x + b > 0$$

$$w \cdot x + b < 0$$

$$\text{margin} = \gamma = 2a / \|w\|$$

Step 1:  $w$  is perpendicular to lines

Step 2: Take a point  $x_-$  on  $w \cdot x + b = -a$  and move to point  $x_+$  that is  $\gamma$  away on line  $w \cdot x + b = a$

$$x_+ = x_- + \gamma w / \|w\|$$

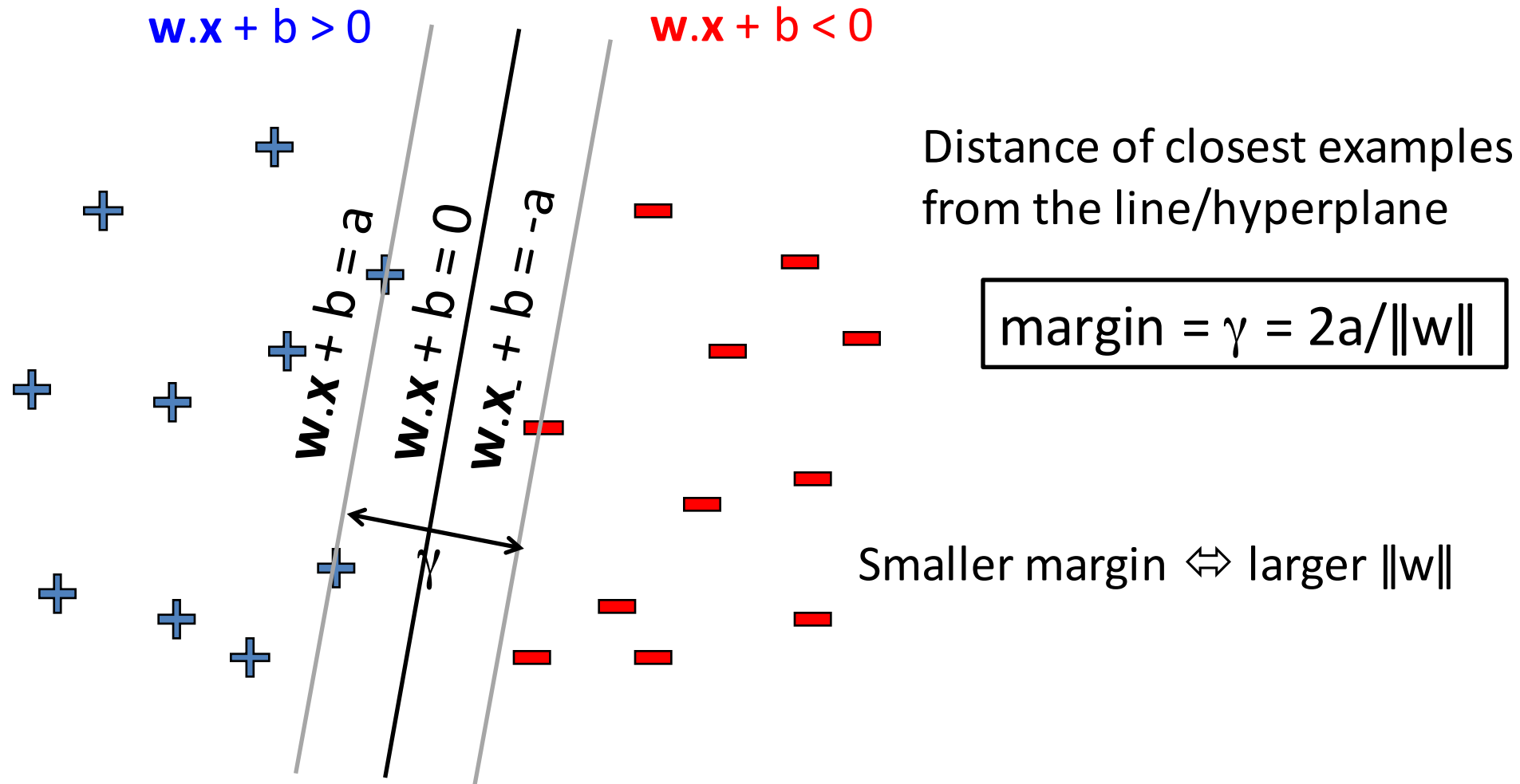
$$w \cdot x_+ = w \cdot x_- + \gamma w \cdot w / \|w\|$$

$$a - b = -a - b + \gamma \|w\|$$

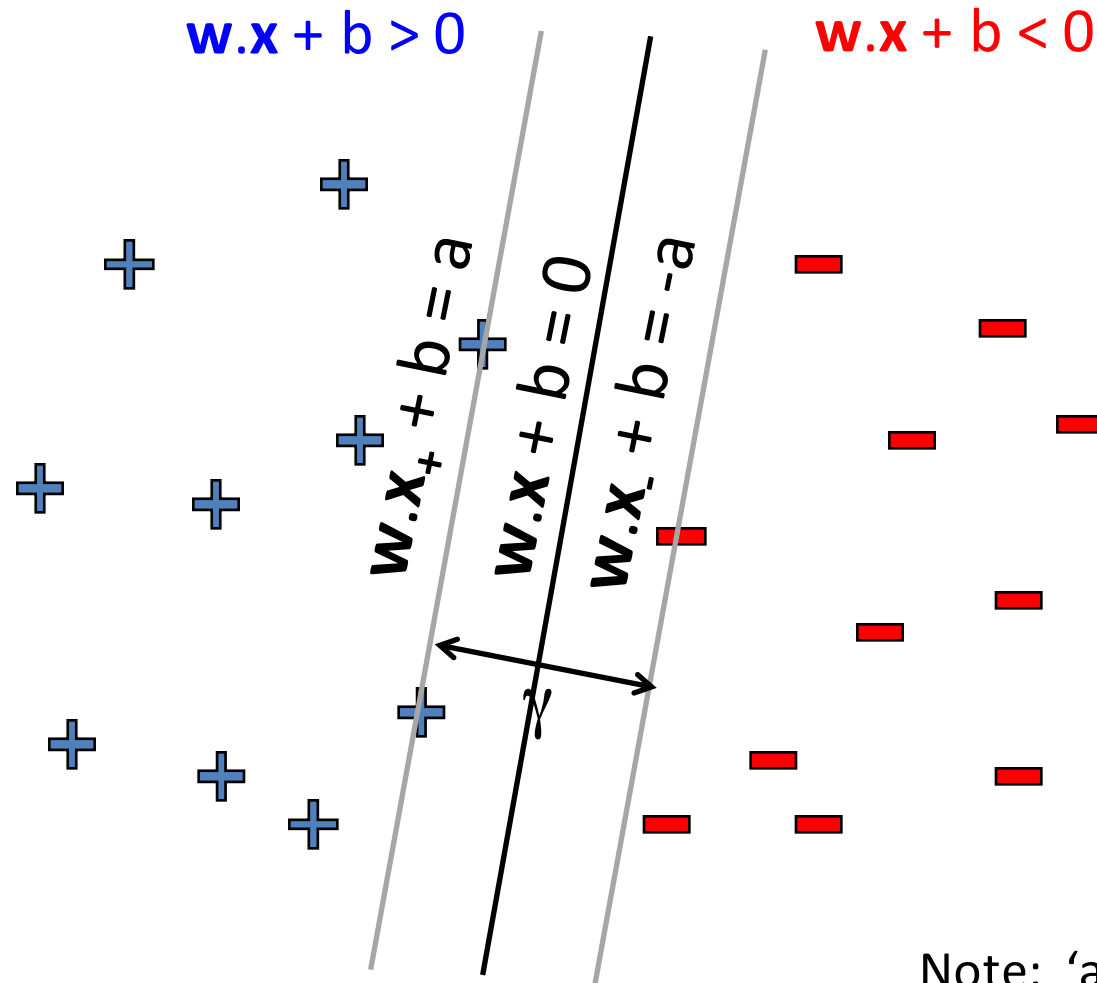
$$2a = \gamma \|w\|$$



# Maximizing the margin



# Maximizing the margin



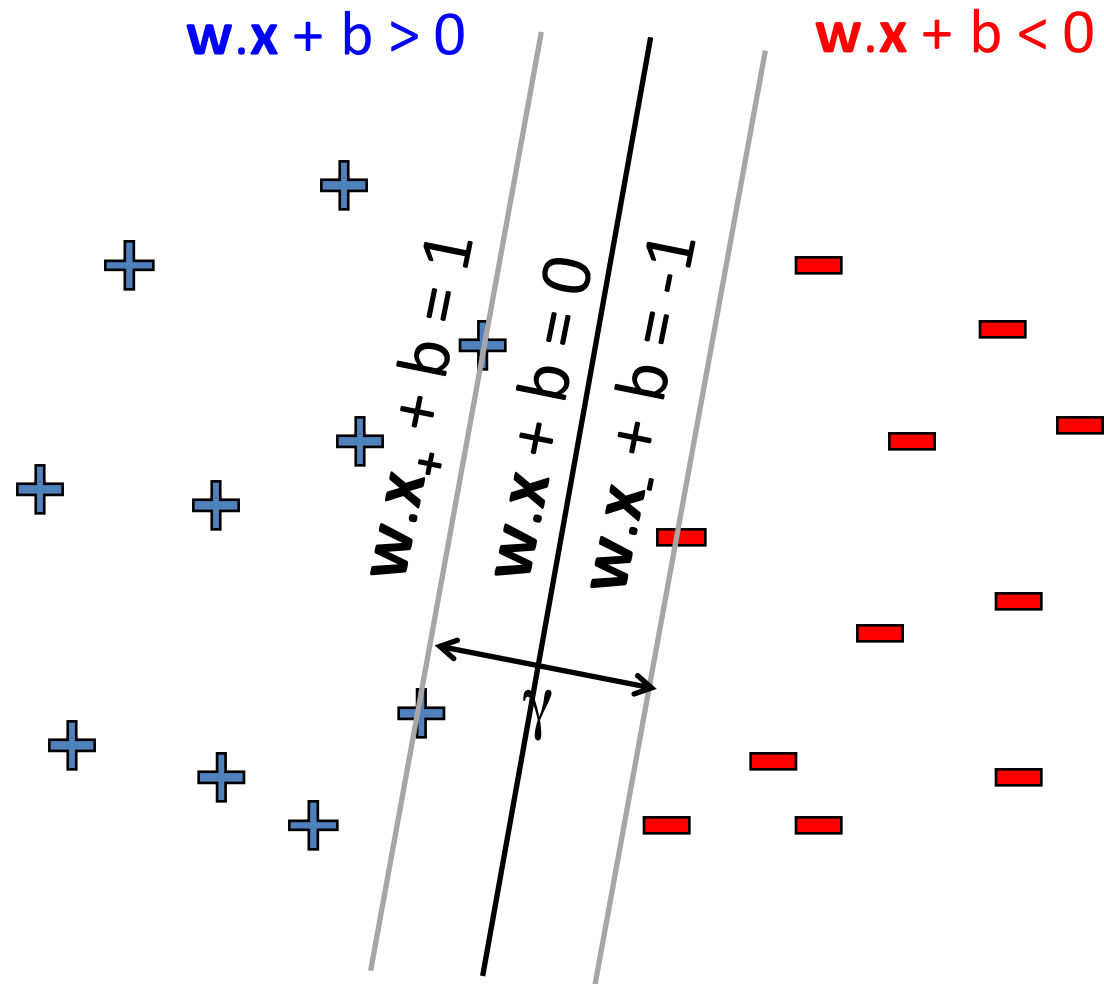
Distance of closest examples from the line/hyperplane

$$\text{margin} = \gamma = 2a / \|w\|$$

$$\begin{aligned} \max_{w,b} \quad & \gamma = 2a / \|w\| \\ \text{s.t.} \quad & (w \cdot x_j + b) y_j \geq a \quad \forall j \end{aligned}$$

Note: 'a' is arbitrary (can normalize equations by a)

# Support Vector Machines



$$\min_{w,b} w \cdot w$$

$$\text{s.t. } (w \cdot x_j + b) y_j \geq 1 \quad \forall j$$

Solve efficiently by quadratic programming (QP)

- Quadratic objective, linear constraints
- Well-studied solution algorithms

# Support Vectors

$$\mathbf{w} \cdot \mathbf{x} + b > 0$$

$$\mathbf{w} \cdot \mathbf{x} + b < 0$$

Linear hyperplane defined by  
“support vectors”

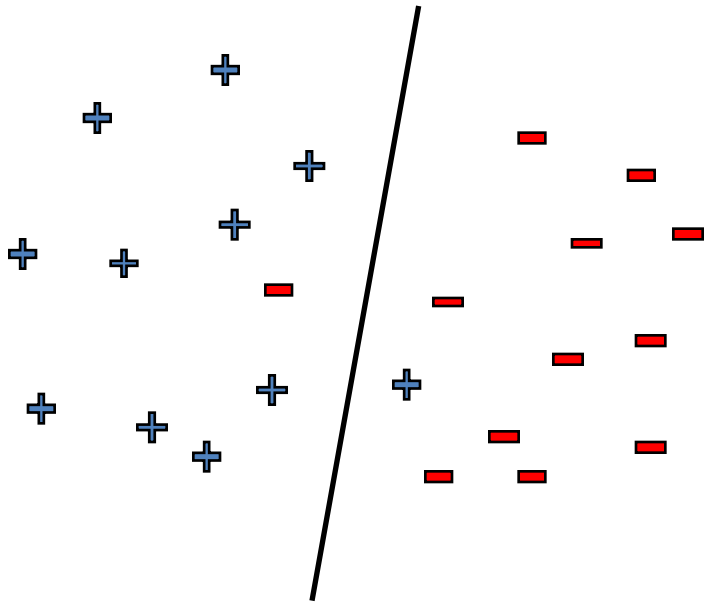
Moving other points a little  
doesn't effect the decision  
boundary

only need to store the  
support vectors to predict  
labels of new points

For support vectors  
 $(\mathbf{w} \cdot \mathbf{x}_j + b) y_j = 1$

# What if data is not linearly separable?

Use features of features  
of features of features....

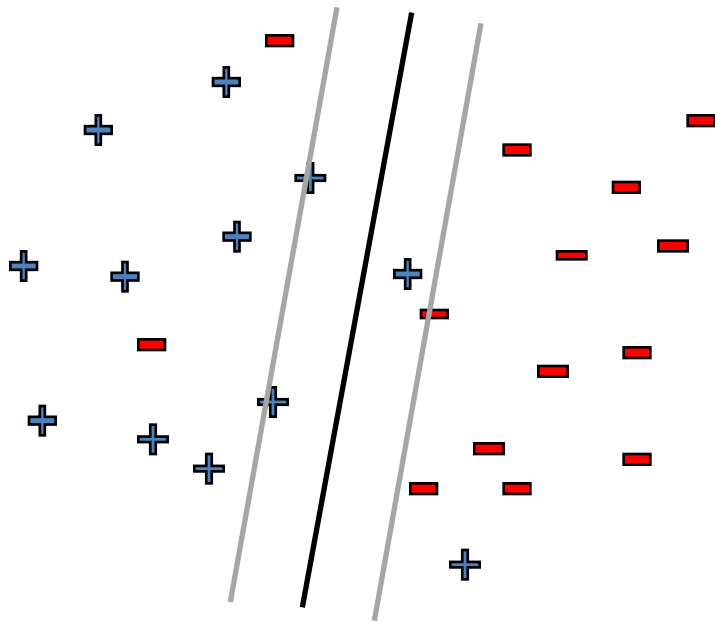


$$x_1^2, x_2^2, x_1x_2, \dots, \exp(x_1)$$

But run risk of overfitting!

# What if data is still not linearly separable?

Allow “error” in classification



Smaller margin  $\Leftrightarrow$  larger  $\|w\|$

$$\min_{w,b} \mathbf{w} \cdot \mathbf{w} + C \# \text{mistakes}$$

$w, b$

$$\text{s.t. } (\mathbf{w} \cdot \mathbf{x}_j + b) y_j \geq 1$$

$$\forall \text{ non-mistakes}$$

Maximize margin and minimize  
# mistakes on training data

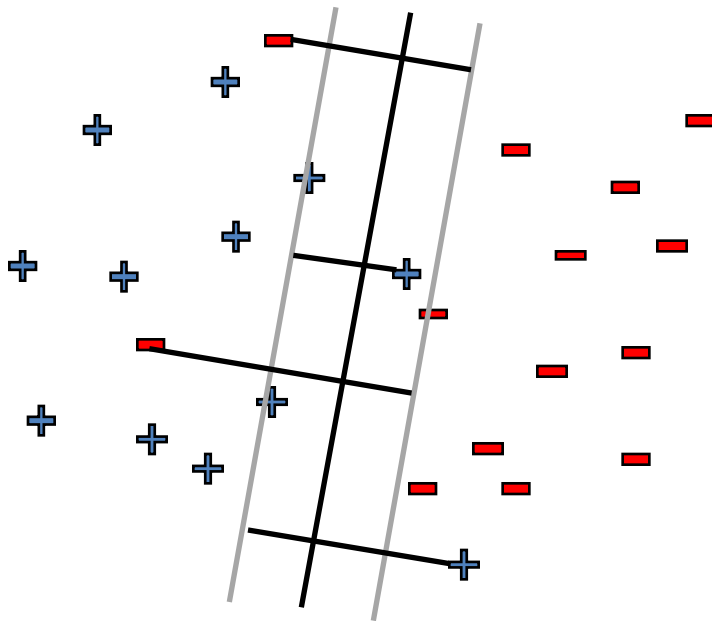
$C$  - tradeoff parameter

Not QP ☹

0/1 loss (doesn't distinguish between  
near miss and bad mistake)

# What if data is still not linearly separable?

Allow “error” in classification



**Soft margin approach**

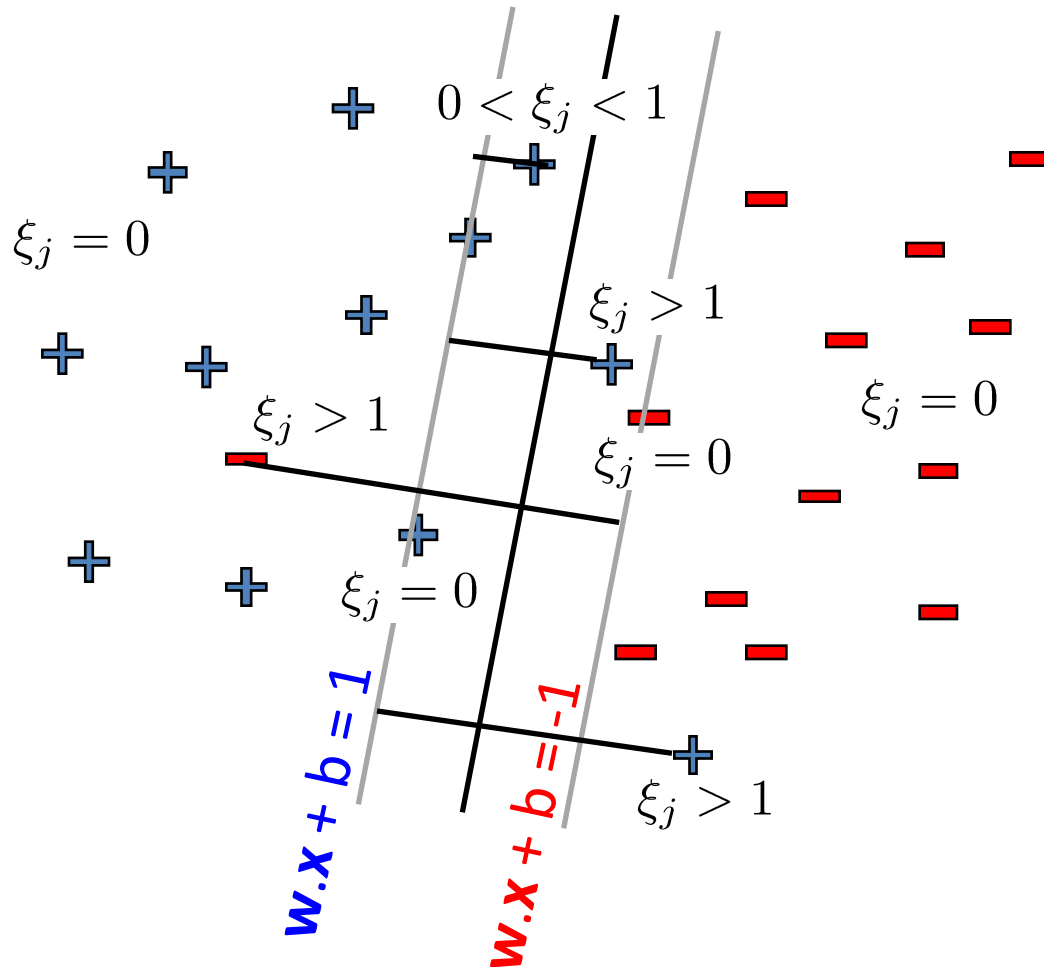
$$\begin{aligned} \min_{\mathbf{w}, b, \{\xi_j\}} \quad & \mathbf{w} \cdot \mathbf{w} + C \sum_j \xi_j \\ \text{s.t.} \quad & (\mathbf{w} \cdot \mathbf{x}_j + b) y_j \geq 1 - \xi_j \quad \forall j \\ & \xi_j \geq 0 \quad \forall j \end{aligned}$$

$\xi_j$  - “slack” variables  
= (>1 if  $x_j$  misclassified)  
pay linear penalty if mistake

$C$  - tradeoff parameter (chosen by cross-validation)

Still QP 😊

# Soft-margin SVM



Soften the constraints:

$$(w \cdot x_j + b) y_j \geq 1 - \xi_j \quad \forall j$$

$$\xi_j \geq 0 \quad \forall j$$

Penalty for misclassifying:

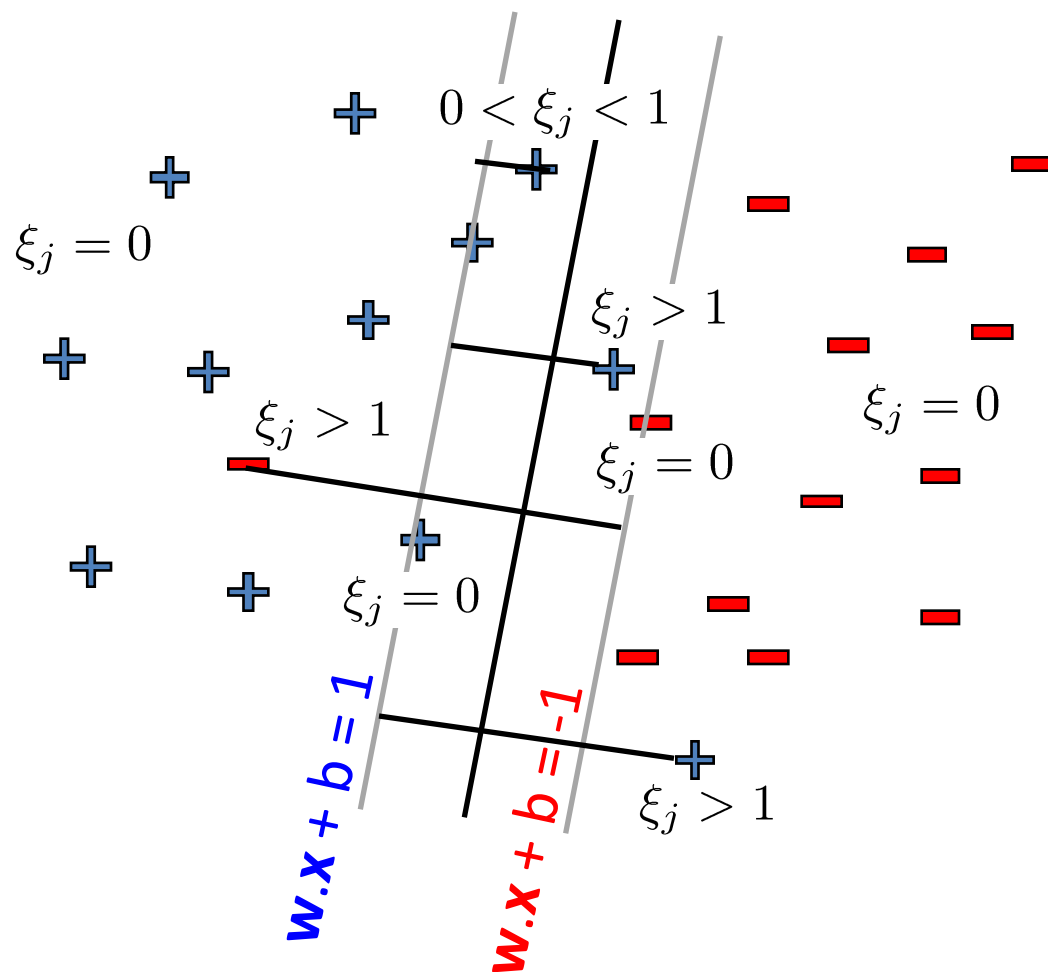
$$C \xi_j$$

How do we recover hard margin SVM?

Set  $C = \infty$



# Slack variables – Hinge loss

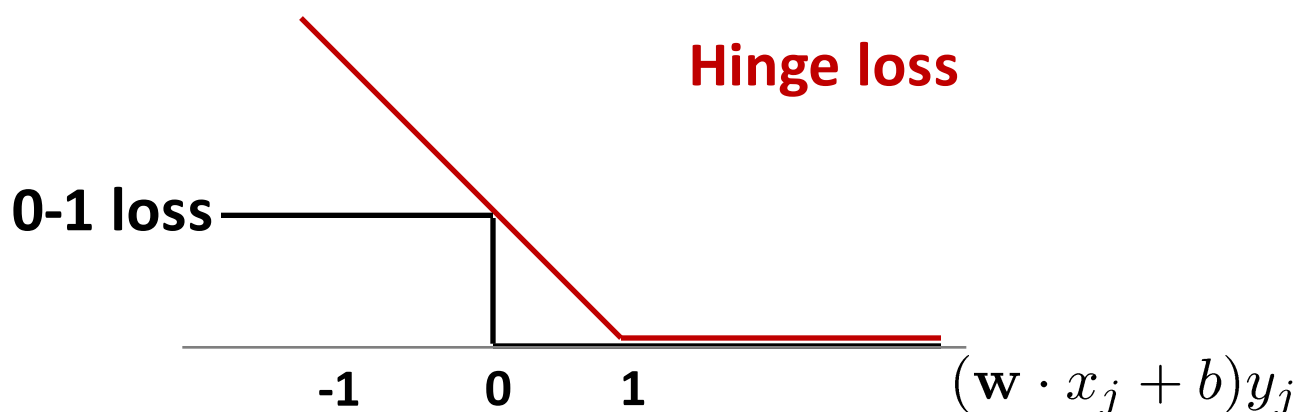


Notice that

$$\xi_j = (1 - (\mathbf{w} \cdot \mathbf{x}_j + b)y_j)_+$$

# Slack variables – Hinge loss

$$\xi_j = (1 - (\mathbf{w} \cdot \mathbf{x}_j + b)y_j)_+$$



$$\min_{\mathbf{w}, b, \{\xi_j\}} \mathbf{w} \cdot \mathbf{w} + C \sum_j \xi_j$$

$$\text{s.t. } (\mathbf{w} \cdot \mathbf{x}_j + b) y_j \geq 1 - \xi_j \quad \forall j$$

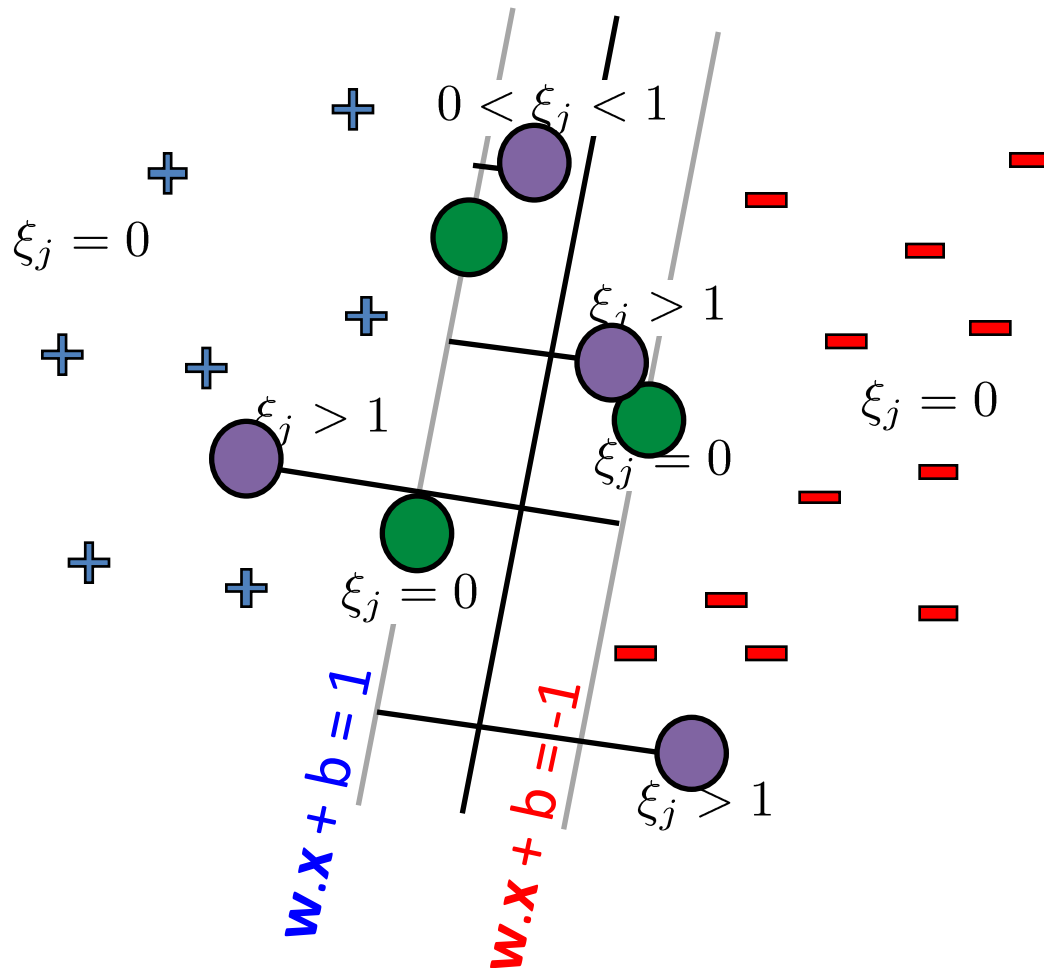
$$\xi_j \geq 0 \quad \forall j$$



Regularized hinge loss

$$\min_{\mathbf{w}, b} \mathbf{w} \cdot \mathbf{w} + C \sum_j (1 - (\mathbf{w} \cdot \mathbf{x}_j + b)y_j)_+$$

# Support Vectors



## Margin support vectors

$\xi_j = 0$ ,  $(w \cdot x_j + b) y_j = 1$   
(don't contribute to objective but enforce constraints on solution)

Correctly classified but on margin

## Non-margin support vectors

$\xi_j > 0$   
(contribute to both objective and constraints)

$1 > \xi_j > 0$  Correctly classified but inside margin

$\xi_j > 1$  Incorrectly classified 19

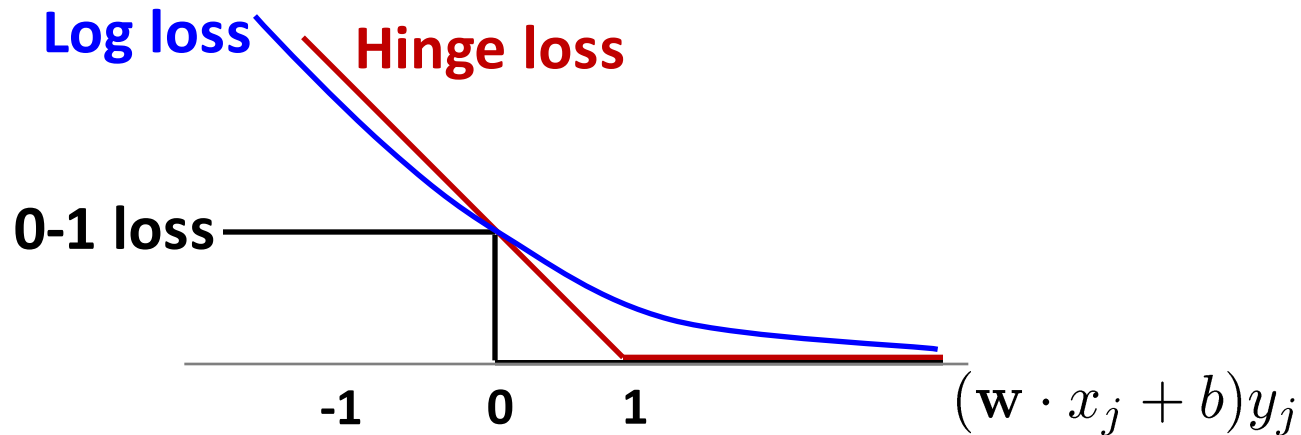
# SVM vs. Logistic Regression

SVM : **Hinge loss**

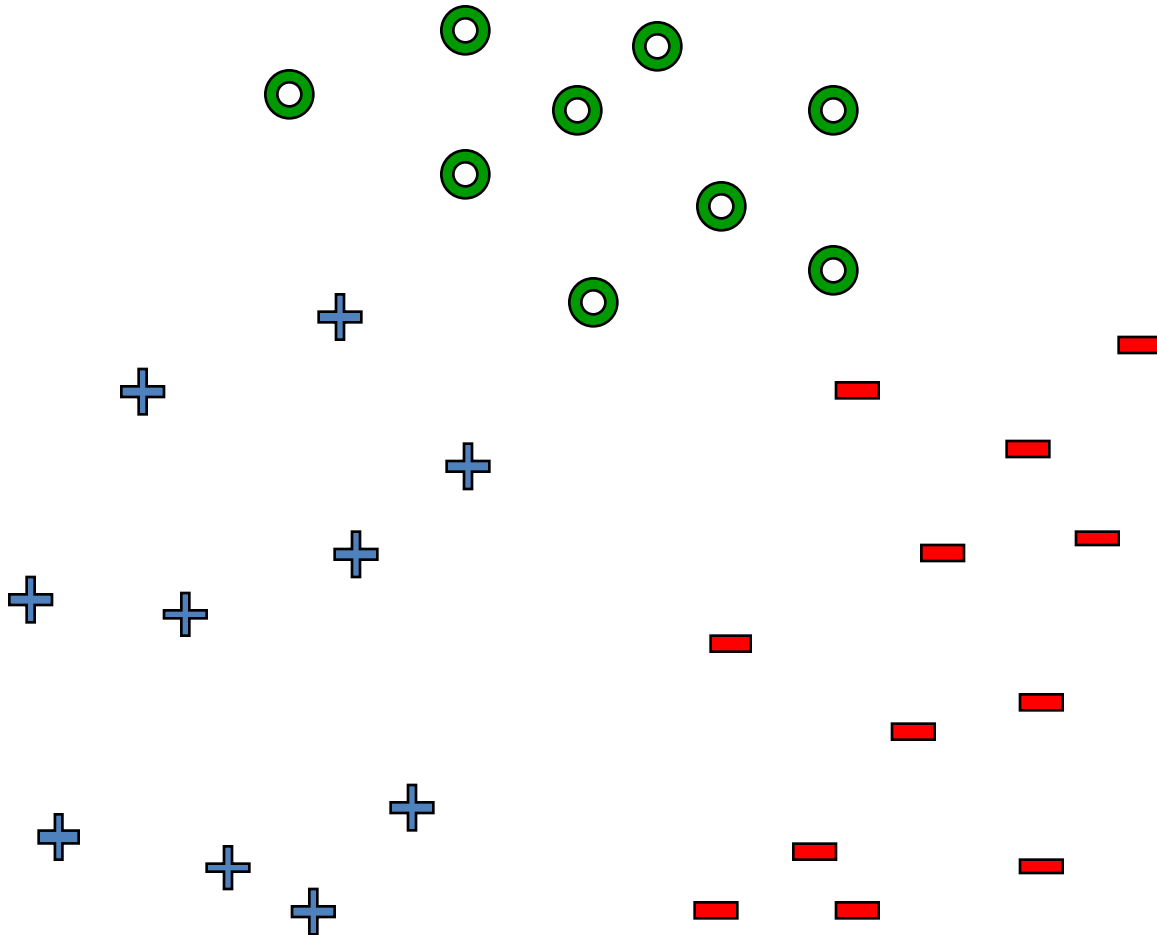
$$\text{loss}(f(x_j), y_j) = (1 - (\mathbf{w} \cdot x_j + b)y_j)_+$$

Logistic Regression : **Log loss** (-ve log conditional likelihood)

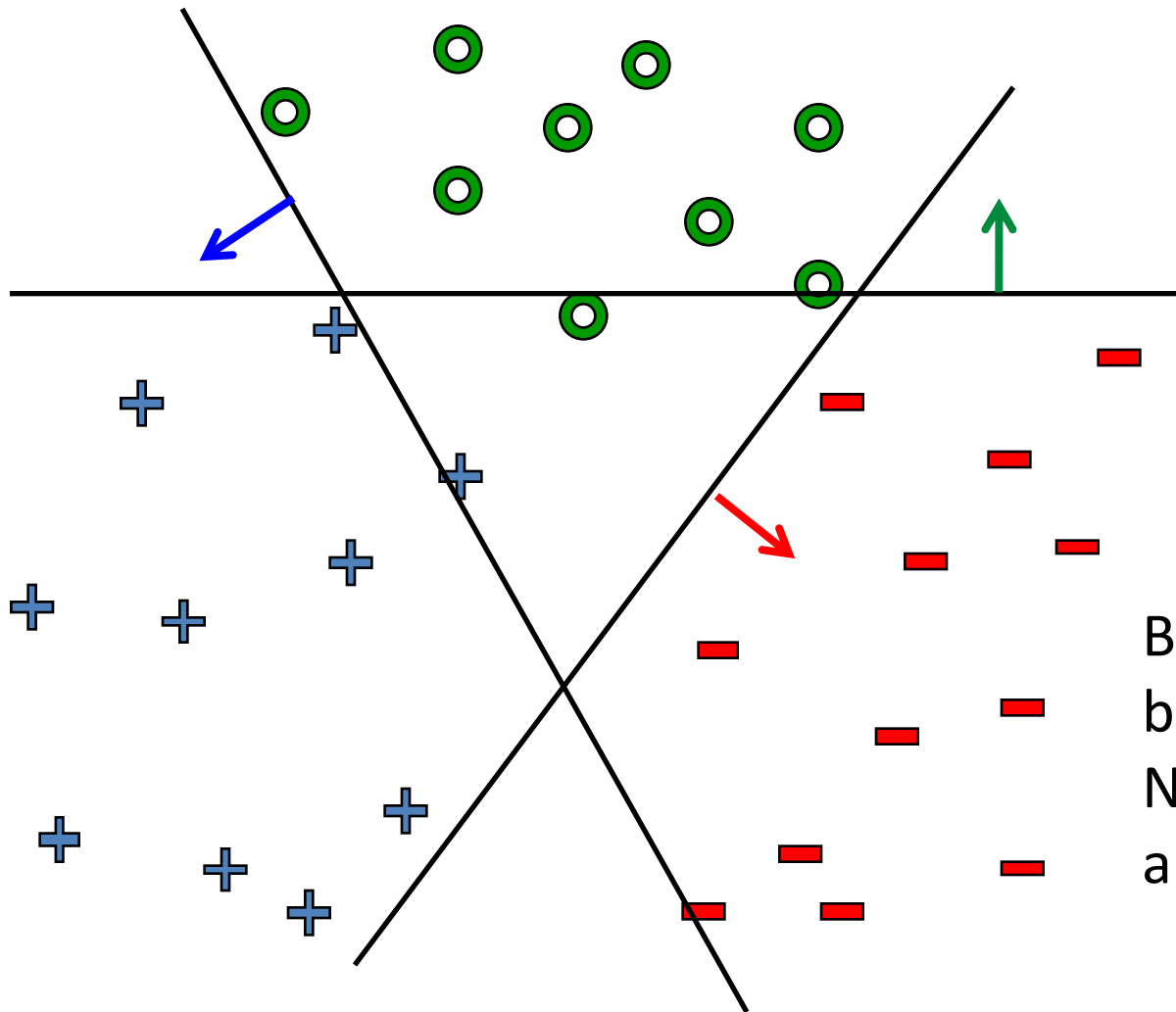
$$\text{loss}(f(x_j), y_j) = -\log P(y_j \mid x_j, \mathbf{w}, b) = \log(1 + e^{-(\mathbf{w} \cdot x_j + b)y_j})$$



# What about multiple classes?



# One vs. rest



Learn 3 classifiers  
separately:

Class  $k$  vs. rest

$$(\mathbf{w}_k, b_k)_{k=1,2,3}$$

$$y = \arg \max_k \mathbf{w}_k \cdot \mathbf{x} + b_k$$

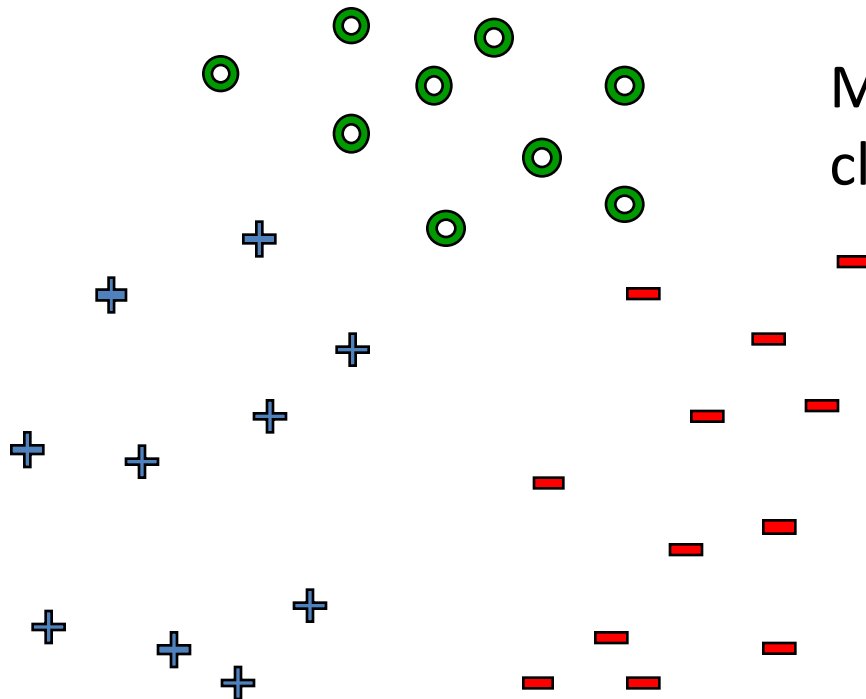
But  $\mathbf{w}_k$ s may not be  
based on the same scale.  
Note:  $(a\mathbf{w}).\mathbf{x} + (ab)$  is also  
a solution

# Learn 1 classifier: Multi-class SVM

Simultaneously learn 3 sets of weights

$$\min_{\{\mathbf{w}^{(y)}\}, \{b^{(y)}\}} \sum_y \mathbf{w}^{(y)} \cdot \mathbf{w}^{(y)}$$

$$\mathbf{w}^{(y_j)} \cdot \mathbf{x}_j + b^{(y_j)} \geq \mathbf{w}^{(y')} \cdot \mathbf{x}_j + b^{(y')} + 1, \quad \forall y' \neq y_j, \quad \forall j$$



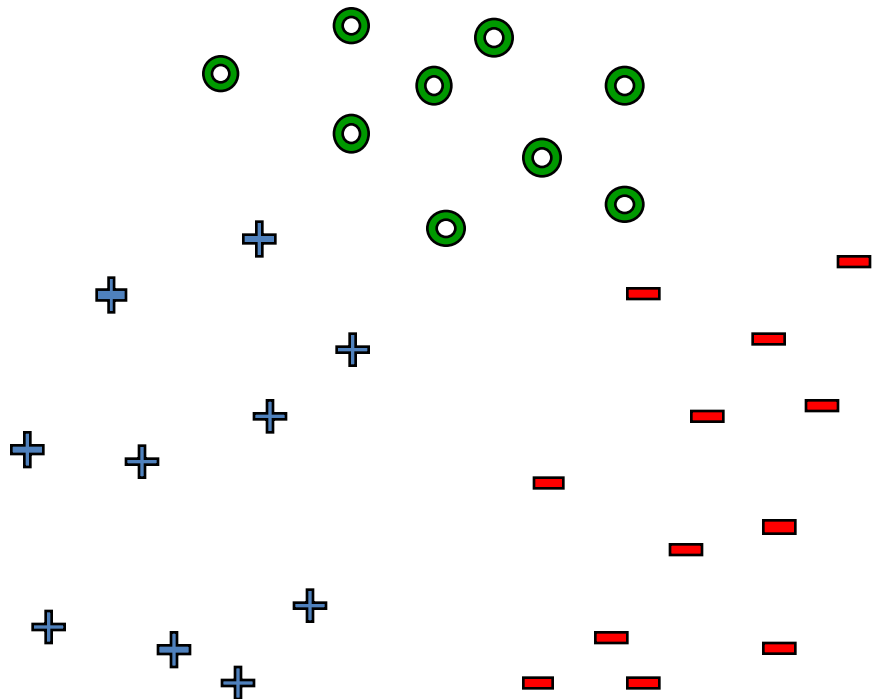
Margin - gap between correct class and nearest other class

$$y = \arg \max_k \mathbf{w}^{(k)} \cdot \mathbf{x} + b^{(k)}$$

# Learn 1 classifier: Multi-class SVM

Simultaneously learn 3 sets of weights

$$\begin{aligned} \text{minimize} \quad & \sum_y \mathbf{w}^{(y)} \cdot \mathbf{w}^{(y)} + C \sum_j \sum_{y \neq y_j} \xi_j^{(y)} \quad \text{over } \{\mathbf{w}^{(y)}\}, \{b^{(y)}\}, \{\xi_j^{(y)}\} \\ & \mathbf{w}^{(y_j)} \cdot \mathbf{x}_j + b^{(y_j)} \geq \mathbf{w}^{(y)} \cdot \mathbf{x}_j + b^{(y)} + 1 - \xi_j^{(y)}, \quad \forall y \neq y_j, \quad \forall j \\ & \xi_j^{(y)} \geq 0, \quad \forall y \neq y_j, \quad \forall j \end{aligned}$$



$$y = \arg \max \mathbf{w}^{(k)} \cdot \mathbf{x} + b^{(k)}$$

Joint optimization:  $\mathbf{w}_k$ s  
have the same scale.

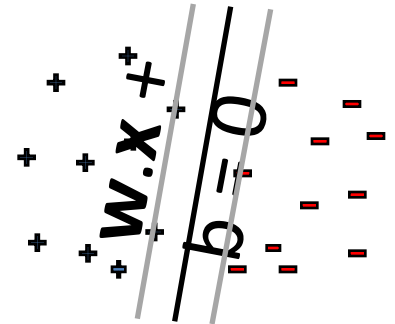


# SVM – linearly separable case

n training points  $(\mathbf{x}_1, \dots, \mathbf{x}_n)$

d features  
vector

$\mathbf{x}_j$  is a d-dimensional



- Primal problem: minimize <sub>$\mathbf{w}, b$</sub>   $\frac{1}{2} \mathbf{w} \cdot \mathbf{w}$   
 $(\mathbf{w} \cdot \mathbf{x}_j + b) y_j \geq 1, \quad \forall j$

**w – weights on features (d-dim problem)**

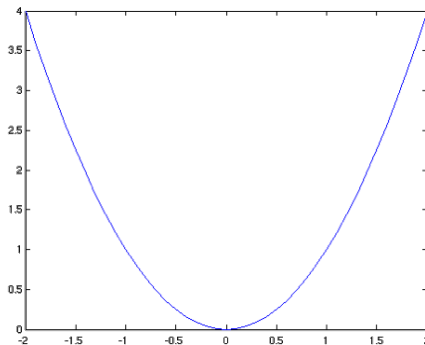
- Convex quadratic program – quadratic objective, linear constraints
- But expensive to solve if d is very large
- Often solved in dual form (n-dim problem)

# Constrained Optimization

$$\begin{array}{ll}\min_x & x^2 \\ \text{s.t.} & x \geq b\end{array}$$

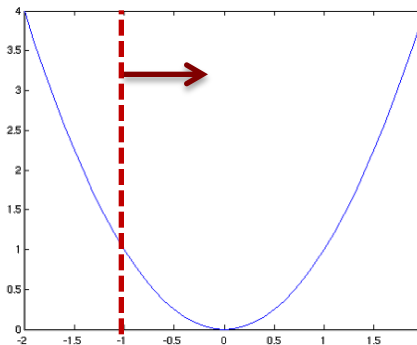
$$x^* = \max(b, 0)$$

$$\min_x x^2$$



$$x^* = 0$$

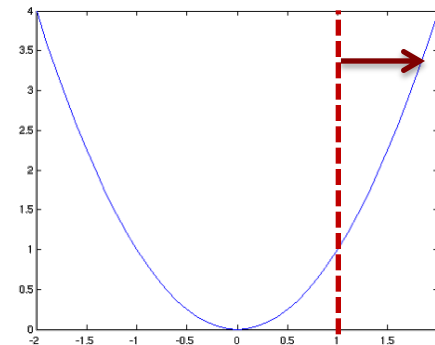
$$\begin{array}{ll}\min_x & x^2 \\ \text{s.t.} & x \geq -1\end{array}$$



$$x^* = 0$$

Constraint inactive

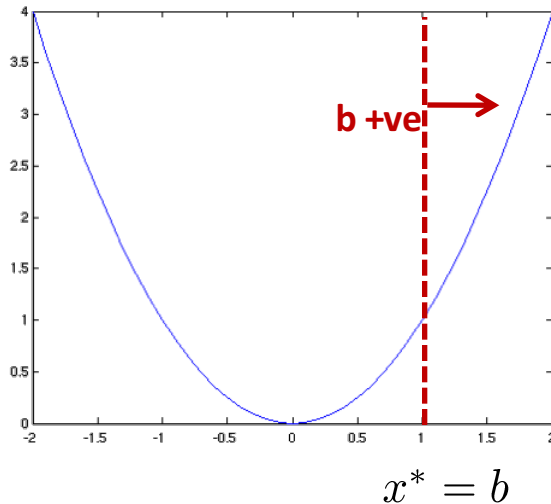
$$\begin{array}{ll}\min_x & x^2 \\ \text{s.t.} & x \geq 1\end{array}$$



$$x^* = 1$$

Constraint active  
and tight

# Constrained Optimization – Dual Problem



$\alpha = 0$  constraint is inactive  
 $\alpha > 0$  constraint is active

Primal problem:

$$\begin{aligned} \min_x \quad & x^2 \\ \text{s.t.} \quad & x \geq b \end{aligned}$$

Moving the constraint to objective function  
Lagrangian:

$$\begin{aligned} L(x, \alpha) &= x^2 - \alpha(x - b) \\ \text{s.t.} \quad & \alpha \geq 0 \end{aligned}$$

Dual problem:

$$\begin{aligned} \max_{\alpha} \quad & d(\alpha) \rightarrow \min_x L(x, \alpha) \\ \text{s.t.} \quad & \alpha \geq 0 \end{aligned}$$

# Connection between Primal and Dual

**Primal problem:**  $p^* = \min_x x^2$   
s.t.  $x \geq b$

**Dual problem:**  $d^* = \max_{\alpha} d(\alpha)$   
s.t.  $\alpha \geq 0$

- **Weak duality:** The dual solution  $d^*$  lower bounds the primal solution  $p^*$  i.e.  $d^* \leq p^*$

To see this, recall  $L(x, \alpha) = x^2 - \alpha(x - b)$

For every feasible  $x$  (i.e.  $x \geq b$ ) and feasible  $\alpha$  (i.e.  $\alpha \geq 0$ ), notice that

$$d(\alpha) = \min_x L(x, \alpha) \leq p^*$$

- **Dual problem (maximization) is always concave even if primal is not convex**

# Connection between Primal and Dual

**Primal problem:**  $p^* = \min_x x^2$   
s.t.  $x \geq b$

**Dual problem:**  $d^* = \max_{\alpha} d(\alpha)$   
s.t.  $\alpha \geq 0$

- **Weak duality:** The dual solution  $d^*$  lower bounds the primal solution  $p^*$  i.e.  $d^* \leq p^*$
- **Strong duality:**  $d^* = p^*$  holds often for many problems of interest e.g. if the primal is a feasible convex objective with linear constraints

# Connection between Primal and Dual

What does strong duality say about  $\alpha^*$  (the  $\alpha$  that achieved optimal value of dual) and  $x^*$  (the  $x$  that achieves optimal value of primal problem)?

Whenever strong duality holds, the following conditions (known as KKT conditions) are true for  $\alpha^*$  and  $x^*$ :

- 1.  $\nabla L(x^*, \alpha^*) = 0$  i.e. Gradient of Lagrangian at  $x^*$  and  $\alpha^*$  is zero.
- 2.  $x^* \geq b$  i.e.  $x^*$  is primal feasible
- 3.  $\alpha^* \geq 0$  i.e.  $\alpha^*$  is dual feasible
- 4.  $\alpha^*(x^* - b) = 0$  (called as complementary slackness)

We use the first one to relate  $x^*$  and  $\alpha^*$ . We use the last one (complimentary slackness) to argue that  $\alpha^* = 0$  if constraint is inactive and  $\alpha^* > 0$  if constraint is active and tight.

# Solving the dual

**Solving:**

$$\begin{array}{ll} \max_{\alpha} \min_x & \overbrace{x^2 - \alpha(x - b)}^{L(x, \alpha)} \\ \text{s.t.} & \alpha \geq 0 \end{array}$$

Optimization over  $x$  is unconstrained.

$$\frac{\partial L}{\partial x} = 2x - \alpha = 0 \Rightarrow x^* = \frac{\alpha}{2}$$

$$\begin{aligned} L(x^*, \alpha) &= \frac{\alpha^2}{4} - \alpha \left( \frac{\alpha}{2} - b \right) \\ &= -\frac{\alpha^2}{4} + b\alpha \end{aligned}$$

Now need to maximize  $L(x^*, \alpha)$  over  $\alpha \geq 0$

Solve unconstrained problem to get  $\alpha'$  and then take  $\max(\alpha', 0)$

$$\frac{\partial}{\partial \alpha} L(x^*, \alpha) = -\frac{\alpha}{2} + b \Rightarrow \alpha' = 2b$$

$$\Rightarrow \alpha^* = \max(2b, 0) \qquad \Rightarrow x^* = \frac{\alpha^*}{2} = \max(b, 0)$$

**$\alpha = 0$  constraint is inactive,  $\alpha > 0$  constraint is active and tight**

# Dual SVM – linearly separable case

n training points, d features  $(\mathbf{x}_1, \dots, \mathbf{x}_n)$  where  $\mathbf{x}_i$  is a d-dimensional vector

- Primal problem: 
$$\begin{aligned} &\text{minimize}_{\mathbf{w}, b} \quad \frac{1}{2} \mathbf{w} \cdot \mathbf{w} \\ &\quad (\mathbf{w} \cdot \mathbf{x}_j + b) y_j \geq 1, \quad \forall j \end{aligned}$$

**w – weights on features (d-dim problem)**

- Dual problem (derivation):

$$\begin{aligned} L(\mathbf{w}, b, \alpha) &= \frac{1}{2} \mathbf{w} \cdot \mathbf{w} - \sum_j \alpha_j [(\mathbf{w} \cdot \mathbf{x}_j + b) y_j - 1] \\ \alpha_j &\geq 0, \quad \forall j \end{aligned}$$

**a – weights on training pts (n-dim problem)**



# Dual SVM – linearly separable case

- Dual problem:

$$\max_{\alpha} \min_{\mathbf{w}, b} L(\mathbf{w}, b, \alpha) = \frac{1}{2} \mathbf{w} \cdot \mathbf{w} - \sum_j \alpha_j \left[ (\mathbf{w} \cdot \mathbf{x}_j + b) y_j - 1 \right]$$
$$\alpha_j \geq 0, \quad \forall j$$

$$\frac{\partial L}{\partial \mathbf{w}} = 0 \quad \Rightarrow \quad \mathbf{w} = \sum_j \alpha_j y_j \mathbf{x}_j$$

$$\frac{\partial L}{\partial b} = 0 \quad \Rightarrow \quad \sum_j \alpha_j y_j = 0$$

If we can solve for  $\alpha$ s (dual problem), then we have a solution for  $\mathbf{w}, b$  (primal problem)

# Dual SVM – linearly separable case

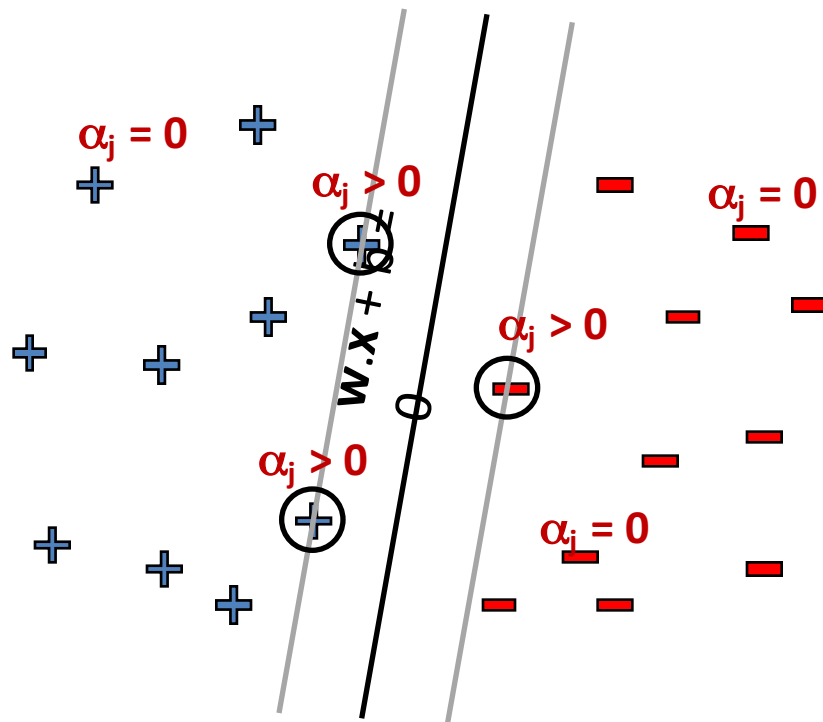
$$\begin{aligned} \text{maximize}_{\alpha} \quad & \sum_i \alpha_i - \frac{1}{2} \sum_{i,j} \alpha_i \alpha_j y_i y_j \mathbf{x}_i \cdot \mathbf{x}_j \\ & \sum_i \alpha_i y_i = 0 \\ & \alpha_i \geq 0 \end{aligned}$$

Dual problem is also QP  
Solution gives  $\alpha_j$ s

$$\mathbf{w} = \sum_i \alpha_i y_i \mathbf{x}_i$$

What about  $b$ ?

# Dual SVM: Sparsity of dual solution



$$w = \sum_j \alpha_j y_j x_j$$

Only few  $\alpha_j$ s can be non-zero : where constraint is active and tight

$$(w \cdot x_j + b) y_j = 1$$

**Support vectors** – training points  $j$  whose  $\alpha_j$ s are non-zero

# Dual SVM – linearly separable case

$$\begin{aligned} \text{maximize}_{\alpha} \quad & \sum_i \alpha_i - \frac{1}{2} \sum_{i,j} \alpha_i \alpha_j y_i y_j \mathbf{x}_i \cdot \mathbf{x}_j \\ & \sum_i \alpha_i y_i = 0 \\ & \alpha_i \geq 0 \end{aligned}$$

Dual problem is also QP  
Solution gives  $\alpha_j$ s

$$\begin{aligned} \mathbf{w} &= \sum_i \alpha_i y_i \mathbf{x}_i \\ b &= y_k - \mathbf{w} \cdot \mathbf{x}_k \\ &\text{for any } k \text{ where } \alpha_k > 0 \end{aligned}$$

Use support vectors with  $\alpha_k > 0$  to  
compute  $b$  since constraint is tight  
 $(\mathbf{w} \cdot \mathbf{x}_k + b) y_k = 1$

# Dual SVM – non-separable case

- Primal problem:

$$\begin{aligned} \text{minimize}_{\mathbf{w}, b, \{\xi_j\}} \quad & \frac{1}{2} \mathbf{w} \cdot \mathbf{w} + C \sum_j \xi_j \\ \text{s.t.} \quad & (\mathbf{w} \cdot \mathbf{x}_j + b) y_j \geq 1 - \xi_j, \quad \forall j \\ & \xi_j \geq 0, \quad \forall j \end{aligned}$$

$$\begin{array}{|c|} \hline \alpha_j \\ \hline \mu_j \\ \hline \end{array}$$

**Lagrange  
Multipliers**

- Dual problem:

$$\begin{aligned} \max_{\alpha, \mu} \quad & \min_{\mathbf{w}, b, \{\xi_j\}} L(\mathbf{w}, b, \xi, \alpha, \mu) \\ \text{s.t.} \quad & \alpha_j \geq 0 \quad \forall j \\ & \mu_j \geq 0 \quad \forall j \end{aligned}$$

# Dual SVM – non-separable case

$$\text{maximize}_{\alpha} \quad \sum_i \alpha_i - \frac{1}{2} \sum_{i,j} \alpha_i \alpha_j y_i y_j \mathbf{x}_i \cdot \mathbf{x}_j$$

$$\sum_i \alpha_i y_i = 0$$

$$C \geq \alpha_i \geq 0$$

comes from  $\frac{\partial L}{\partial \xi} = 0$

Intuition:

If  $C \rightarrow \infty$ , recover hard-margin SVM

Dual problem is also QP  
Solution gives  $\alpha_j$ s

$$\mathbf{w} = \sum_i \alpha_i y_i \mathbf{x}_i$$

$$b = y_k - \mathbf{w} \cdot \mathbf{x}_k$$

for any  $k$  where  $C > \alpha_k > 0$

# So why solve the dual SVM?

- There are some quadratic programming algorithms that can solve the dual faster than the primal, (specially in high dimensions  $d \gg n$ )
- But, more importantly, the “**kernel trick**”!!!  
(later!)