武汉大学计算机学院 2016 - 2017 学年第一学期 2014 级《编译原理》(A) 期末考试参考答案

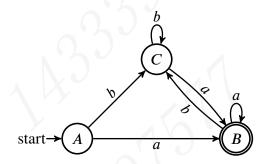
一、 (1)

start
$$\longrightarrow 0$$
 $\stackrel{a}{\longrightarrow} 1$ $\stackrel{b}{\longrightarrow} 4$ $\stackrel{b}{\longrightarrow} 5$ $\stackrel{\varepsilon}{\longrightarrow} 3$ $\stackrel{\varepsilon}{\longrightarrow} 4$ $\stackrel{\varepsilon}{\longrightarrow} 0$ $\stackrel{a}{\longrightarrow} 1$ $\stackrel{a}{\longrightarrow} 1$

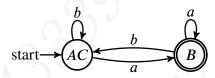
(2)

$$A = \{0,3,4\}, B = \{1,2\}, C = \{0,3,4,5\}.$$

状态转换图为:



(3) 最小 DFA 如下所示:



- (4) 由 a 和 b 且以 a 结尾的字符串.
- (5) (a | b)*a.
- 二、 (1) 语句 "()(())" 的最左推导如下:

$$\begin{array}{cccc} S & \underset{lm}{\Longrightarrow} & SS & \underset{lm}{\Longrightarrow} & ()(S) \\ & \underset{lm}{\Longrightarrow} & (S)S & \underset{lm}{\Longrightarrow} & ()((S)) \\ & \underset{lm}{\Longrightarrow} & ()S & \underset{lm}{\Longrightarrow} & ()(()) \end{array}$$

(2) 消除左递归后的文法如下:

$$\begin{array}{ccc} S & \rightarrow & S' \mid (S) \, S' \\ S' & \rightarrow & S \, S' \mid \varepsilon \end{array}$$

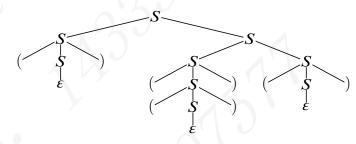
- (3) $\operatorname{First}(S) = \operatorname{First}(S') = \{ \varepsilon, (\}; \\ \operatorname{Follow}(S) = \operatorname{Follow}(S') = \{ \$, (,) \}.$
- (4) LL(1) 分析表如下所示:

	()	\$
S	$S \rightarrow (S)S' \mid S'$	$S \to S$	S o S'
S'	$S' o SS' \mid \varepsilon$	$S' o SS' \mid \varepsilon$	$S' \to SS' \mid \varepsilon$

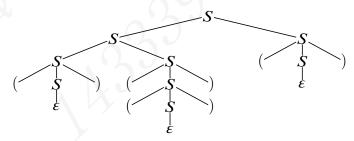
(5) 语句"()()"的分析过程如下所示:

剩余串	分析栈	分析动作		
()()\$	<i>S</i> \$	$S \to (S)S'$		
()()\$	(S)S'\$	match-advance		
)()\$	S)S'\$	$S \to S'$		
)()\$	S')S'\$	$S' o \varepsilon$		
)()\$) <i>S</i> '\$	match-advance		
()\$	S'\$	$S' \rightarrow SS'$		
()\$	<i>SS</i> '\$	$S \to (S)S'$		
()\$	(S)S'S'\$	match-advance		
)\$	S)S'S'\$	$S \rightarrow S'$		
)\$	S')S'S'\$	$S' \to \varepsilon$		
)\$)S'S'\$	match-advance		
\$	S'S'\$	$S' o \varepsilon$		
\$	S'\$	$S' o \varepsilon$		
\$	\$	分析成功		

三、 (1) 语句 "()(())()" 的两颗不同的语法树为: 语法树 1:



语法树 2:



(2) 无二义文法:

$$S \rightarrow S(S) \mid \varepsilon$$

四、 (1) 状态 I_3 的 LR(0) 项目集为

$$\overline{\{S \to S \bullet S, S \to (S \bullet)\}}
= \{S \to S \bullet S, S \to (S \bullet), S \to \bullet SS, S \to \bullet(S), S \to \bullet\}$$

- (2) 识别活前缀的 DFA 在识别正则式 $\underline{(^*S(^*)$ 所生成的串一定到达状态 I_2 , I_1 或 I_3 .
- (3) $Follow(S) = \{\$, (,)\}.$ 根据题三,左结合的活前缀只能是S(*S?)?,故 SLR 分析表如下所示:

	action		goto	
状态	()	\$ (<u> </u>
0	s1/r3	(r3)	r3	2
1	s1/r3	r3	<u>r3</u>	3
2	s1/r3	(r3)	acc/r3	4
3	s1/r3	s5/r3	<u>r3</u>	4
4	r1/s1/r3	/r1/r3	r1/r3	4
5	r2	r2	r2	

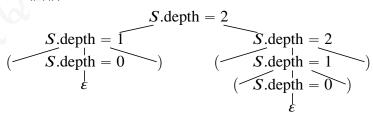
(4) 语句"()()"的分析过程如下所示:

剩余串	分析栈	分析动作
()()\$	0	shift
)()\$	0(1	reduce $S \to \varepsilon$
)()\$	0(1S3	shift
()\$	0(1S3)5	reduce $S \to (S)$
()\$	0.52	shfit
)\$	0S2(1	reduce $S \to \varepsilon$
)\$	0S2(1S3	shift
\$	0S2(1S3)5	reduce $S \to (S)$
\$	0S2S4	reduce $S \to SS$
\$	0.52	accept

五、(1)

$$rac{\dot{\mathcal{F}} \, \, \dot{\mathcal{E}} \, \, \dot{\mathcal{E}} \, \, \dot{\mathcal{M}} \, \, \dot{\mathcal{M}}}{S \to S_1 \, S_2} \quad \dot{\mathcal{S}}. \text{depth} = \max(S_1. \text{depth}, S_2. \text{depth})}{S \to (S_1)} \quad \dot{\mathcal{S}}. \text{depth} = S_1. \text{depth} + 1} \\ S \to \varepsilon \qquad \dot{\mathcal{S}}. \text{depth} = 0$$

(2) 试画出语句"()(())"的附注语法树:



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六、

```
L1: ifnot (a < b) goto L2 | ifnot (x = 2) goto L1
if (c < d) goto L3 | L0: t1 := x + 2
ifnot (e < f) goto L3 | x := t1
L2: t0 := x + 1 | goto L1
x := t0 | L3:
if (x = 0) goto L0</pre>
```

七、 调用 printf 时,实参反向逐一计值压栈. 因此首先调用 sum(),将返回值 15 先压栈,由于 sum()的副作用修改了全局变量 n为 0,再压 n进栈时,实际上是压 0进栈.