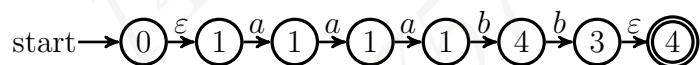


武汉大学计算机学院2013-2014学年第一学期  
2011级《编译原理》参考答案

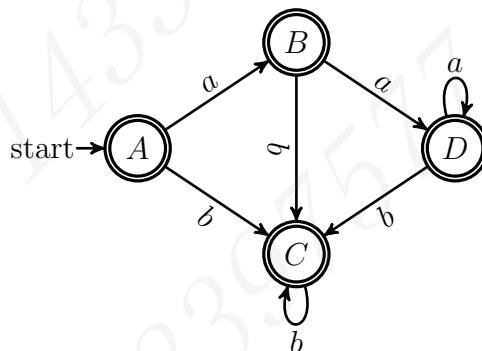
一、(1)



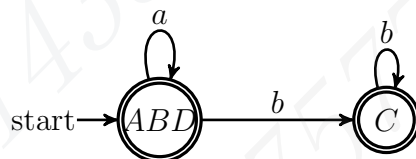
(2)

$$A = \{0, 1, 3, 4\}, B = \{1, 2, 3, 4, 5\}, C = \{3, 4\}, D = \{1, 3, 4\}.$$

状态转换图为:



(3) 最小DFA如下所示:



(4) 零个或多个连续的 $a$ 连接零个或多个连续的 $b$ .

(5)  $r = a^*b^*$ .

二、(1) 最左推导如下:

$$\begin{array}{lll} I & \xRightarrow{lm} \{L\} & \xRightarrow{lm} \{\{I\}, L\} \\ & \xRightarrow{lm} \{L, L\} & \xRightarrow{lm} \{\{n\}, L\} \\ & \xRightarrow{lm} \{I, L\} & \xRightarrow{lm} \{\{n\}, I\} \\ & \xRightarrow{lm} \{\{L\}, L\} & \xRightarrow{lm} \{\{n\}, n\} \end{array}$$

(2) 消除左递归后的文法如下:

$$\begin{array}{l} I \rightarrow \{L\} \mid n \\ L \rightarrow IL' \\ L' \rightarrow , \mid LL' \mid \varepsilon \end{array}$$

- (3)  $\text{First}(I) = \text{First}(L) = \{ \{, n \}$ ;  $\text{First}(L') = \{ , \varepsilon \}$ .  
 $\text{Follow}(I) = \{ \}, , \$$ ;  $\text{Follow}(L) = \text{Follow}(L') = \{ \}, , \}$ .
- (4) LL(1)分析表如下所示:

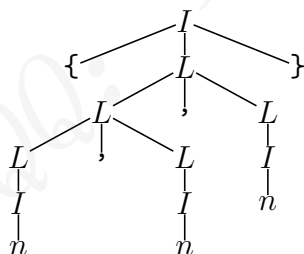
	,	$n$	{	}	\$
$I$		$I \rightarrow n$	$I \rightarrow \{L\}$		
$L$		$L \rightarrow IL'$	$L \rightarrow IL'$		
$L'$	$L' \rightarrow ,LL', L \rightarrow \varepsilon$			$L' \rightarrow \varepsilon$	

- (5) 语句“ $\{ \{n\}, n \}$ ”的分析过程如下所示:

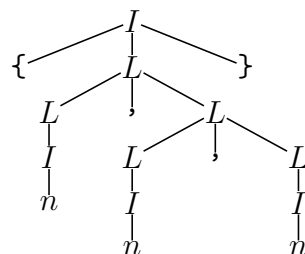
剩余串	分析栈	分析动作
$\{ \{n\}, n \} \$$	$I \$$	$I \rightarrow \{L\}$
$\{ \{n\}, n \} \$$	$\{L\} \$$	match-advance
$\{n\}, n \} \$$	$L \} \$$	$L \rightarrow IL'$
$\{n\}, n \} \$$	$IL' \} \$$	$I \rightarrow \{L\}$
$\{n\}, n \} \$$	$\{L\} L' \} \$$	match-advance
$n \}, n \} \$$	$L \} L' \} \$$	$L \rightarrow IL'$
$n \}, n \} \$$	$IL' \} L' \} \$$	$I \rightarrow n$
$n \}, n \} \$$	$nL' \} L' \} \$$	match-advance
$\}, n \} \$$	$L' \} L' \} \$$	$L' \rightarrow \varepsilon$
$\}, n \} \$$	$\} L' \} \$$	match-advance
$, n \} \$$	$L' \} \$$	$L' \rightarrow ,LL'$
$, n \} \$$	$,LL' \} \$$	match-advance
$n \} \$$	$LL' \} \$$	$L \rightarrow IL'$
$n \} \$$	$IL' L' \} \$$	$L \rightarrow IL'$
$n \} \$$	$IL' L' \} \$$	$I \rightarrow n$
$n \} \$$	$nL' L' \} \$$	match-advance
$\} \$$	$L' L' \} \$$	$L' \rightarrow \varepsilon$
$\} \$$	$L' \} \$$	$L' \rightarrow \varepsilon$
$\} \$$	$\} \$$	match-advance
$\} \$$	$\} \$$	分析成功

- 三、(1) 语句“ $\{n, n, n\}$ ”的两颗不同的语法树为:

语法树1:



语法树2:



(2) 无二义文法:

$$\begin{aligned} I &\rightarrow \{L\} \mid n \\ L &\rightarrow L, I \mid I \end{aligned}$$

四、(1) 状态 $I_3$ 的LR(0)项目集为

$$\begin{aligned} &\overline{\{I \rightarrow \{\bullet L\}\}} \\ &= \{I \rightarrow \{\bullet L\}, L \rightarrow \bullet L, L, L \rightarrow \bullet I, I \rightarrow \bullet \{L\}, I \rightarrow \bullet n\} \end{aligned}$$

(2)  $\{^*(n|\varepsilon)\}$ .

(3) SLR分析表如下所示:

状态	action					goto	
	$n$	,	{	}	\$	$I$	$L$
0	s2		s3			1	
1					acc		
2		r2		r2	r2		
3	s2		s3			4	5
4		r4		r4			
5		s6		s7			
6	s2		s3			4	8
7		r1		r1	r1		
8	<b>r3</b>			r3			

(4) 语句“ $\{n, n, n\}$ ”的分析过程如下所示:

剩余串	分析栈	分析动作
$\{n, n, n\}\$$	0	shift
$n, n, n\}\$$	0{3	shift
$, n, n\}\$$	0{3n2	reduce $I \rightarrow n$
$, n, n\}\$$	0{3I4	reduce $L \rightarrow I$
$, n, n\}\$$	0{3L5	shift
$n, n\}\$$	0{3L5, 6	shift
$, n\}\$$	0{3L5, 6n2	reduce $I \rightarrow n$
$, n\}\$$	0{3L5, 6I4	reduce $L \rightarrow L$
$, n\}\$$	0{3L5, 6L8	reduce $L \rightarrow L, L$
$, n\}\$$	0{3L5	shift
$n\}\$$	0{3L5, 6	shift
$\}\$$	0{3L5, 6n2	reduce $I \rightarrow n$
$\}\$$	0{3L5, 6I4	reduce $L \rightarrow I$
$\}\$$	0{3L5, 6L8	reduce $L \rightarrow L, L$

}\$	0{3L5	shift
\$	0{3L5}7	reduce $I \rightarrow \{L\}$
\$	0I1	分析成功

五、 (1)

产生式	语义规则
$I' \rightarrow MI$	$I.level = 0, I.limit = 1$
$M \rightarrow \varepsilon$	$count = 0$
$I \rightarrow \{LN\}$	<b>if</b> $count \% getsize(I.level) \neq 0$ <b>then</b> $error("the\ left\ brace\ is\ not\ in\ right\ pos")$ <b>if</b> $getsize(I.level) == 1$ <b>then</b> $error("the\ brace\ level\ is\ more\ than\ array\ dim")$ $L.level = I.level + 1$ $L.limit = getsize(I.level) + count$ $N.limit = getsize(I.level) + count$
$I \rightarrow n$	<b>if</b> $I.level == 0$ <b>then</b> $error("initializer\ is\ not\ in\ braces")$ <b>if</b> $count \geq I.limit$ <b>then</b> $error("the\ array\ up\ bound\ exceeds")$ $emit(getname() + "[" + 4 * count + "]=" + n.val)$ $count = count + 1$
$N \rightarrow \varepsilon$	<b>while</b> $count < N.limit$ $emit(getname() + "[" + 4 * count + "]=" + 0)$ $count = count + 1$
$L \rightarrow L_1, L_2$	$L_1.level = L.level, L_2.level = L.level$ $L_1.limit = L.limit, L_2.limit = L.limit$
$L \rightarrow I$	$I.level = L.level, I.limit = L.limit$

(2)  $a[0] = 1; a[32] = 2; [36] = 3; a[64] = 4; a[80] = 5; a[84] = 6;$

六、

```

L1:  if (a > b) goto L2      |      ifnot (m > n) goto L2
      ifnot (c > d) goto L2 |      t1 := y + 2
      ifnot (e > f) goto L0 |      y := t1
      if (i > j) goto L1     |      goto L1
L0:  t0 := x + 1             |  L2:
      x := t0

```

七、实参反向压栈，即函数main()中调用printf时，首先将最后一个实参count压栈，这时其值为0。再调用fac(5)，并将函数调用的结果120压栈。这时虽然count修改为6，但是栈中所传的实参还是0。