

# Sparse Optimization with Distributionally Robust Chance Constraint

D2 Zhicheng Zhang

@IPS, Fujisaki Lab

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# Sparse Decision-Making Under Uncertainty

## Sparse chance constraint optimization (SCCO)

Consider the sparse chance constraint optimization as

$$\begin{aligned} \min_x \quad & \|x\|_0 && (\text{sparse cost}) \\ \text{s.t.} \quad & x \in \mathcal{X} && (\text{deterministic constraint}) \\ & \mathbb{P}\{h(x, \delta) \leq 0\} \geq 1 - \epsilon && (\text{chance constraint}) \\ \Leftrightarrow \quad & \mathbb{P}\{h(x, \delta) > 0\} < \epsilon && (\text{violation}) \end{aligned}$$

- Non-convex cost function
- Uncertainty  $\delta$  is assumed to be a random variable governed by probability  $\mathbb{P}$  supported on  $\Delta \subseteq \mathbb{R}^{n_\delta}$
- Violation event does not exceed a risk level  $\epsilon \in (0, 1)$
- Multivariate integral computation and non-convex feasible set

# Sparsity

## Definition (Sparsity)

A vector  $x \in \mathbb{R}^n$  is sparse if it contains many 0's, or has  $\ell_0$  “norm”

$$\|x\|_0 = |\text{supp}(x)| \quad \Rightarrow \quad \|x\|_0 \leq s \quad (\text{s-sparse}),$$

where  $\text{supp}(x) \doteq \{I \in \{1, \dots, n\} : x_I \neq 0\}$  denotes the number of the nonzero elements in  $x$ .

## Convex Relaxation ( $\ell_1$ norm)

The  $\ell_1$  norm convex set is defined as

$$\Sigma_s := \{x \in \mathbb{R}^n : \|x\|_1 \leq s\}, \quad (\text{Lasso})$$

where the  $\ell_1$  norm constraints  $\|x\|_1 = \sum_{I=1}^n |x_I|$ .

- When the assumption  $\|x\|_\infty \leq 1$  holds, then the biconjugate function of  $\ell_0$  norm gives the result  $\|x\|_0^{**}(r) = \|r\|_1 \leq \|x\|_0$ .

# Conjugate Function

- The conjugate function of  $\|x\|_0$  is defined by

$$\|x\|_0^*(y) := \sup_x \{ \langle x, y \rangle - \|x\|_0 \} = \max \left( \sum_i |y_i|, 0 \right),$$

it is always a convex form even the  $\|x\|_0$  is non-convex.

- The biconjugate function of  $\|x\|_0$  is as follows

$$\|x\|_0^{**}(r) := \sup_y \{ \langle r, y \rangle - \|x\|_0^*(y) \} = \|r\|_1$$

- Hence,  $\ell_1$  norm is a *convex relaxation* of  $\ell_0$  norm in some sense due to the fact that  $\|x\|_0^{**}(r) \leq \|x\|_0$ .

# Scenario Approximation for Uncertainty

## Sparse Scenario Approximation (SA) [Calafiore & Campi, 05, 06]

The SA entails randomized algorithms to randomly generate  $N$  samples i.i.d. from the probability  $\mathbb{P}$ , and the uncertainty set  $\Delta$  is replaced by a finite samples  $\{\delta^{(i)}\}_{i=1}^N$ , then (SCCO) becomes sparse SA

$$J_N^* = \min_x \{ \|x\|_1 : x \in \mathcal{X}, h(x, \delta^i) \leq 0, \forall i=1, \dots, N \},$$

which may be recast as the following *epigraphic form*

$$J_N^{s*} = \min_{x,s} \left\{ s : x \in \mathcal{X}_s, \max_{i=1, \dots, N} h(x, \delta^i) \leq 0 \right\}, \text{ where } \mathcal{X}_s = \{ \|x\|_1 \leq s \cap \mathcal{X} \}$$

## Sample complexity [Campi & Garatti, 13]

Given  $\epsilon, \beta \in (0, 1)$  and  $q = \dim(\mathcal{X}_s) < n$ . The sample complexity is

$$N \geq \frac{2}{\epsilon} \left( \ln \frac{1}{\beta} + q \ln \frac{n \cdot e}{q} \right), \quad e \approx 2.718 \dots$$

then it holds that  $\mathbb{P}^N \{ V(x_N^*(s^*)) > \epsilon \} \leq \beta$ .

# Sample Average Approximation for Uncertainty

## Sample Average Approximation (SAA) [Luedke & Ahmed, 08]

The SAA is a probabilistic constraint relaxation for *out-of-sample* in SA under a significance level  $\alpha \in (0, 1)$ , that is,  $\hat{p}_N(x) \leq \alpha$ , where  $0 \leq \alpha < \epsilon$ . Therefore, the sparse SAA program is as follows

$$J_{N,\alpha}^* = \min_x \{ \|x\|_1 : x \in \mathcal{X}, \hat{p}_N(x) \leq \alpha, i = 1, \dots, N \}$$

- Chance constraint  $p(x) := \mathbb{P}(\delta : h(x, \delta) > 0) = \mathbb{E}_{\mathbb{P}}[\mathbb{I}(h(x, \delta))]$
- Estimate the “true” probability via the *discrete empirical distribution*

$$\hat{p}_N(x) = \frac{1}{N} \sum_{i=1}^N \mathbb{I}(h(x, \delta^i)), \quad i = 1, \dots, N.$$

- If  $\alpha = 0$ , then sparse SAA program reduces to sparse SA program
- Assess worst case probability of two-sided failure for given  $\hat{\epsilon}, \hat{\beta} \in (0, 1)$

$$\mathbb{P}^N \left\{ \sup_{x \in \mathcal{X}} |p(x) - \hat{p}_N(x)| > \hat{\epsilon} \right\} < \hat{\beta} \quad (\text{VC theory})$$

# Exact Sparsity via SAA

## Proposition (Exact Sparsity via SAA)

The exact  $\ell_0$  norm constraint is equivalent to a SAA of chance constraint

$$\|x\|_0 \leq s \Leftrightarrow \frac{1}{n} \sum_{l=1}^n \mathbb{I}(|x_l| \leq 0) \geq 1 - \frac{s}{n} \Leftrightarrow \frac{1}{n} \sum_{l=1}^n \mathbb{I}(|x_l| > 0) \leq \frac{s}{n}$$

- $n$  scenarios are with *equal* probability  $\frac{1}{n}$
- The  $l$ -th scenario index set is  $\mathcal{S}^l := \{x : x_l = 0, \forall l = 1, \dots, n\}$
- $\|u\|_0 \leq s$  means that at most  $s$  out of the  $n$  scenarios are violated.

## SCCO via SAA Reformulation

The problem (SCCO) can be recast as a SAA formulation in sparse cost and chance constraint, that is,

$$\min_s \left\{ s : \frac{1}{N} \sum_{i=1}^N \mathbb{I}(h(u, \delta^i)) \leq \alpha, i \in \mathcal{N}, \frac{1}{n_u} \sum_{l=1}^{n_u} \mathbb{I}(|u_l| \leq 0) \geq 1 - \frac{s}{n_u} \right\}$$

# Mixed Integer Programming

## Proposition (Exact Sparsity via MIP)

The exact  $\ell_0$  norm constraint is equivalent to a MIP formulation

$$\begin{aligned}\|x\|_0 \leq s &\Leftrightarrow |x_l| \leq M_l z_l, \sum_{l=1}^n z_l \leq s, z_l \in \{0, 1\}, l = 1, \dots, n \\ &\Leftrightarrow |x| \leq Mz, e^\top z \leq s, z \in \{0, 1\}^n, e = [1 \cdots 1]\end{aligned}$$

- Binary variables  $z_l$  is an auxiliary variable to evaluate the sparsity
- $z_l = 1$  counts the nonzero elements
- Solve big- $M$  coefficients
- Boolean convex relaxation  $z \in [0, 1]^n$  (big- $M$  free)



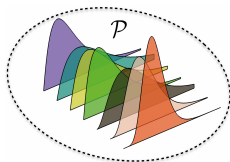
# Chance Constraint via MIP

## MIP for Chance constraint

We introduce an auxiliary binary variables  $v_i \in \{0, 1\}$  for each  $i \in \mathcal{N}$ , where  $v_i = 1$  assures that the safety or reliability event  $h(x, \delta^i) \leq 0$  holds; and otherwise  $v_i = 0$  indicates the violation event. Thus, we express a MIP form

$$\begin{aligned} h(x, \delta^i) &\leq \eta_i(1 - v_i), & (\Leftrightarrow h(x, \delta^i) + \eta_i v_i &\geq 0) \\ \sum_{i=1}^N p_i v_i &\geq 1 - \alpha, & (\Leftrightarrow \sum_{i=1}^N p_i(1 - v_i) &\leq \alpha) \\ v_i &\in \{0, 1\}, \quad \forall i = 1, \dots, N, \end{aligned}$$

where  $\eta_i \in \mathbb{R}$  and  $0 < p_i = \mathbb{P}\{\xi = \xi^i\}$  are the *non-equal* probability of possible outcomes as scenarios and satisfy  $\sum_{i=1}^N p_i = 1$ . We refer these more generic constraints as *knapsack constraints*.



## Sparse and distributionally robust optimization (SDRO)

Consider a sparse distributionally robust (chance constrained) optimization

$$\begin{aligned} \min_x \quad & \|x\|_0 \\ \text{s.t.} \quad & \inf_{\mathbb{P} \in \mathcal{P}} \mathbb{P}\{\delta \in \Delta : f(x, \delta) \leq 0\} \geq 1 - \epsilon \quad (\text{safety}) \\ & \Leftrightarrow \sup_{\mathbb{P} \in \mathcal{P}} \mathbb{P}\{\delta \in \Delta : h(x, \delta) > 0\} \leq \epsilon \end{aligned}$$

- “Ambiguity set”  $\mathcal{P}$  = a family of probability distributions.
- Moment ambiguity set  $\mathcal{P}(\mu, \Sigma)$ , and metric ambiguity set  $\mathcal{P}(\epsilon)$

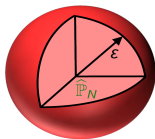
# Data-Driven Wasserstein Ambiguity Set

## Definition (Wasserstein ambiguity set)

[Gao & Kleywegt, 17; Esfahani & Kuhn, 18]

The Wasserstein ambiguity set  $\mathcal{P}^W$  can be defined as

$$\mathcal{P}^W := \mathcal{B}_\varepsilon(\hat{\mathbb{P}}_N) = \left\{ \mathbb{Q} : W_p(\hat{\mathbb{P}}_N, \mathbb{Q}) \leq \varepsilon \right\}, \quad \hat{\mathbb{P}}_N = \frac{1}{N} \sum_{i=1}^N \mathbb{I}_{\hat{\delta}(i)} \text{ (empirical)}$$



Contains every  $\mathbb{Q}$  obtainable by re-shaping  $\hat{\mathbb{P}}_N$  at a cost of at most  $\varepsilon$

- $W_p(\mathbb{Q}_1, \mathbb{Q}_2) = \inf_{\pi \in \Pi(\mathbb{Q}_1, \mathbb{Q}_2)} \left( \int_{\Delta \times \Delta} \|\delta_1 - \delta_2\|^p \pi(d\delta_1, d\delta_2) \right)^{\frac{1}{p}}$ , and  $\Pi$  is the set of couplings with marginals  $\mathbb{Q}_1$  and  $\mathbb{Q}_2$ . [Villani, 08]
- Be capable of comparing a **continuous** and a **discrete** distribution (weak convergence and convergence in the  $p$ -th moment)

# Feasibility Analysis

$$\inf_{\mathbb{P} \in \mathcal{P}^w} \mathbb{P}\{h(x, \delta) \leq 0\} \geq 1 - \epsilon \quad \Leftrightarrow \quad \sup_{\mathbb{P} \in \mathcal{B}_\epsilon(\hat{\mathbb{P}}_N)} \mathbb{P}\{h(x, \delta) > 0\} \leq \epsilon$$

## Primal Problem (Wasserstein metric)

The worst case ambiguous (violation) uncertainty quantification

$$\begin{aligned} (P) \quad J^P &= \sup_{\mathbb{P} \in \mathcal{P}^w} \mathbb{P}(h(x, \delta) > 0) = \sup_{\mathbb{P} \in \mathcal{B}_\epsilon(\hat{\mathbb{P}}_N)} \mathbb{E}_{\mathbb{P}}[\mathbb{I}(h(x, \delta))] \\ &= \begin{cases} \sup_{\Pi, \mathbb{P}} \int_{\Delta} \mathbb{I}(h(x, \delta)) \mathbb{P}(d\delta) \\ \text{s.t. } W_p(\hat{\mathbb{P}}_N, \mathbb{P}) \leq \epsilon \end{cases} \\ &= \begin{cases} \sup_{\mathbb{P}_i} \int_{\Delta} \mathbb{I}(h(x, \delta)) \mathbb{P}_i(d\delta) \\ \text{s.t. } \frac{1}{N} \sum_{i=1}^N \int_{\Delta} \|\delta - \hat{\delta}^i\|_p \mathbb{P}_i(d\delta) \leq \epsilon \end{cases} \end{aligned}$$

where the optimal value of primal problem (P) denotes  $J^P$ .

# Tractable Performance Reformulation

## Dual Problem

Using a standard Lagrangian dual variable  $\lambda \geq 0$ , the dual of primal problem (P) is as follows

$$\begin{aligned}(D) \quad J^D &= \inf_{\lambda \geq 0} \left\{ \lambda \varepsilon - \int_{\Delta} \inf_{\delta \in \Delta} \left[ \lambda \|\delta - \hat{\delta}^i\|_p - \mathbb{I}(h(x, \delta)) \right] \hat{\mathbb{P}}_N(d\hat{\delta}^i) \right\} \\ &= \inf_{\lambda \geq 0} \left\{ \lambda \varepsilon + \frac{1}{N} \sum_{i=1}^N \sup_{\delta \in \Delta} \left[ \mathbb{I}(h(x, \delta)) - \lambda \|\delta - \hat{\delta}^i\|_p \right] \right\}\end{aligned}$$

where the optimal value of dual problem (D) denotes  $J^D$ .

- Strong duality:  $J^P = J^D$
- Introduce epigraphical auxiliary variables  $\zeta_i$ ,  $\forall i \in \mathcal{N}$ , then (D) is as

$$\inf_{\lambda, \zeta_i} \left\{ \lambda \varepsilon + \frac{1}{N} \sum_{i=1}^N \zeta_i \quad \text{s.t.} \quad \sup_{\delta \in \Delta} \left[ \mathbb{I}(h(x, \delta)) - \lambda \|\delta - \hat{\delta}^i\|_p \right] \leq \zeta_i, \quad \lambda \geq 0 \right\}$$

## Reformulation

[Gao & Kleywegt, 17; Esfahani & Kuhn, 18]

For  $i \in \mathcal{N}$ , the optimal value of (P) is equal to the optimal value of (D) under the Wasserstein ambiguity set  $\mathcal{P}^W = \mathcal{B}_\epsilon(\hat{\mathbb{P}}_N)$ , that is

$$\sup_{\mathbb{P} \in \mathcal{P}^W} \mathbb{P}(h(u, \delta) > 0) = \inf_{\lambda \geq 0} \left\{ \lambda \epsilon + \frac{1}{N} \sum_{i=1}^N \sup_{\delta \in \Delta} \left[ \mathbb{I}(h(x, \delta)) - \lambda \|\delta - \hat{\delta}^i\|_p \right] \right\}$$

- Break down indicator function in infimum/supremum by discussing the condition of taking “zero” and “one”
- In general, the unsafety event can be redefined as

$$\sup_{\mathbb{P} \in \mathcal{P}^W} \mathbb{P}(x(\delta) \notin \mathcal{S}(\delta)) \leq \epsilon \quad \Rightarrow \quad x(\hat{\delta}^i) \text{ via } \hat{\mathbb{P}}_N$$

- $\mathbb{I}(h(x, \delta))$  is the violation (or unsafety) event that governed by a metric/distance between a random point and the unsafety set.

# Take Home Message

- SCCO can be approximated by scenario approximation, sample average approximation and Data-driven Wasserstein ball setting.
- SCCO can be recast as a SAA form
- SCCO can be recast as a MIP form

## Future work

- The selection of performance function  $h(x, \delta)$ .
- Wasserstein ambiguity feasible set analysis.
- Using MIP to solve sparse distributionally robust chance constrained program.