

# CSC 665: Vapnik-Chervonenkis (VC) Theory

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## 1 Infinite hypothesis classes can be PAC learnable

In the last lecture, we have seen that the size of a hypothesis class  $\mathcal{H}$  can be an important factor of sample complexity of learning from that  $\mathcal{H}$ . Specifically, if  $\mathcal{H}$  is finite, then it has a PAC sample complexity upper bound of  $O(\frac{1}{\epsilon}(\ln |\mathcal{H}| + \ln \frac{1}{\delta}))$ , and an agnostic PAC sample complexity upper bound of  $O(\frac{1}{\epsilon^2}(\ln |\mathcal{H}| + \ln \frac{1}{\delta}))$ . Does that mean that if  $\mathcal{H}$  is infinite, then  $\mathcal{H}$  is not PAC learnable?

In this section, we give a counterexample, showing that for the hypothesis class of threshold functions on the  $[0, 1]$  interval,  $\mathcal{H}$  is PAC learnable. To formalize the statement, we need some notation setup.

1. The instance domain  $\mathcal{X}$  be the  $[0, 1]$  interval,
2. The label space  $\mathcal{Y}$  be  $\{-1, +1\}$ .
3. The hypothesis class  $\mathcal{H} = \{h_t \triangleq 2\mathbf{1}(x > t) - 1 : t \in [0, 1]\}$  is the set of threshold functions over  $[0, 1]$ .  
Given classifier  $h_t$ , it will classify all examples  $x$  on the left of  $t$  as label  $-1$ , and classify all examples  $x$  on the right of  $t$  (including  $t$ ) as label  $+1$ . Note that  $\mathcal{H}$  is (uncountably) infinite.

Recall that the consistency algorithm is one that returns a classifier  $\hat{h}$  in  $\mathcal{H}$  that agrees with all training examples. We have the following theorem on the sample complexity of the consistency algorithm.

**Theorem 1.** Suppose  $D$  is a distribution over  $[0, 1]$  that is realizable with respect to  $\mathcal{H}$ . Then, for any  $\epsilon \in (0, \frac{1}{2})$ ,  $\delta \in (0, 1)$ , given  $m \geq \frac{1}{\epsilon} \ln \frac{2}{\delta}$  training examples drawn iid from  $D$ , the consistency algorithm returns a classifier  $h_{\hat{t}}$  such that with probability  $1 - \delta$ ,

$$\text{err}(h_{\hat{t}}, D) \leq \epsilon.$$

*Proof.* We will only consider the setting where  $D$  is a continuous probability distribution that has density on  $[0, 1]$ . (For a rigorous proof for general  $D$ , see Appendix A for details.)

Consider two points  $t_L$  and  $t_R$ , which are defined such that

$$\mathbb{P}(x \in [t^*, t_R]) = \epsilon,$$

$$\mathbb{P}(x \in [t_L, t^*]) = \epsilon.$$

We consider the setting where such  $t_L$  and  $t_R$  exists. See Appendix A for the proof for the general setting where  $t_L$  and  $t_R$  may not exist.

We are going to show that given a sample size  $m \geq \frac{1}{\epsilon} \ln \frac{2}{\delta}$ , there exists an event  $\bar{E}$ , such that at least one sample is in  $[t^*, t_R]$ , and at least one sample is in  $[t_L, t^*]$ . Note that if this happens, then the returned threshold  $\hat{t}$  will be inside  $[t_L, t_R]$ .

If  $\hat{t}$  is in  $[t_L, t^*]$ , then

$$\text{err}(h_{\hat{t}}, D) = \mathbb{P}(x \in [\hat{t}, t^*]) \leq \mathbb{P}(x \in [t_L, t^*]) = \epsilon.$$

Similarly, if  $\hat{t}$  is in  $[t^*, t_R]$ , then

$$\text{err}(h_{\hat{t}}, D) = \mathbb{P}(x \in [t^*, \hat{t}] \leq \mathbb{P}(x \in [t^*, t_R]) = \epsilon.$$

Now define event  $E_L$  (resp.  $E_R$ ) be such that no sample is in  $[t_L, t^*]$  (resp.  $[t^*, t_R]$ ), and define  $E = E_L \cup E_R$ . It suffices to show  $\mathbb{P}(E) \leq \delta$ . Indeed,

$$\mathbb{P}(E_L) = (1 - \mathbb{P}(x \in [t_L, t^*]))^m \leq e^{-\epsilon m} = \delta/2$$

and similarly  $\mathbb{P}(E_R) \leq \delta/2$ . This implies that  $\mathbb{P}(E) \leq \mathbb{P}(E_L) + \mathbb{P}(E_R) \leq \delta$ .  $\square$

## 2 VC dimension

We provide a more refined characterization of the complexity of a hypothesis class. Generally, if a hypothesis class is more expressive, then we may need more samples to learn from them. But how can we measure the expressiveness of a hypothesis class?

**Definition 1.** Given a hypothesis class  $\mathcal{H}$  and a set of unlabeled examples  $S = \{x_1, \dots, x_n\}$ , define the projection of  $\mathcal{H}$  to  $S$  as:

$$\Pi_{\mathcal{H}}(S) = \{(h(x_1), \dots, h(x_n)) : h \in \mathcal{H}\}.$$

Intuitively, if  $\mathcal{H}$  is more expressive, then  $|\Pi_{\mathcal{H}}(S)|$  is larger. The largest possible value of  $|\Pi_{\mathcal{H}}(S)|$  is  $2^n$ , where  $\mathcal{H}$  achieves all possible  $+1/-1$  labelings on  $S$ . In this case, we call that  $S$  is *shattered* by  $\mathcal{H}$ .

**Definition 2.** The VC dimension of  $\mathcal{H}$  (abbrev.  $\text{VC}(\mathcal{H})$ ), is the largest nonnegative integer  $d$  such that there exists  $S$  of size  $d$  that is shattered by  $\mathcal{H}$ . If no such  $d$  exists, we  $\text{VC}(\mathcal{H})$  is defined to be infinity.

We have the following more checkable definition of VC dimension:

**Lemma 1.** Suppose we are given a hypothesis class  $\mathcal{H}$  and an integer  $d$ . Then  $\text{VC}(\mathcal{H}) = d$  is equivalent to the following two holding simultaneously:

1. There exists a set of examples of size  $d$  that is shattered by  $\mathcal{H}$ .
2. Any set of examples of size  $d + 1$  are not shattered by  $\mathcal{H}$ .

Examples of VC dimension:

1. Thresholds in  $\mathbb{R}$ .  $\mathcal{H} = \{h_t(x) \triangleq \mathbf{21}(x > t) - 1 : t \in [0, 1]\}$ . It can be seen that  $\{0.5\}$  is shattered by  $\mathcal{H}$ . However, consider any set  $S = \{x_1, x_2\}$ . Suppose  $x_1 \leq x_2$ . Then it is impossible to find  $h_t$  in  $\mathcal{H}$  such that  $h_t(x_1) = +1$  and  $h_t(x_2) = -1$ . Therefore,  $\text{VC}(\mathcal{H}) = 1$ .
2. Intervals in  $\mathbb{R}$ .  $\mathcal{H} = \{h_{a,b}(x) \triangleq \mathbf{21}(a \leq x \leq b) - 1 : t \in [0, 1]\}$ . It can be seen that  $\{0.2, 0.5\}$  is shattered by  $\mathcal{H}$ . However, consider any set  $S = \{x_1, x_2, x_3\}$ . Suppose  $x_1 \leq x_2 \leq x_3$ . Then it is impossible to find  $h_{a,b}$  in  $\mathcal{H}$  such that  $h_{a,b}(x_1) = +1$ ,  $h_{a,b}(x_2) = -1$ ,  $h_{a,b}(x_3) = +1$ . The reason is that:  $h_{a,b}(x_1) = +1$  implies that  $a \leq x_1$ ;  $h_{a,b}(x_3) = +1$  implies that  $x_3 \leq b$ . However, this would imply that  $x_2 \in [a, b]$ , therefore  $h_{a,b}(x_2) = +1$ , and  $h_{a,b}(x_2) = -1$  is impossible. Therefore,  $\text{VC}(\mathcal{H}) = 2$ .
3. Homogeneous linear classifiers in  $\mathbb{R}^d$ .  $\mathcal{H} = \{h_w(x) \triangleq \mathbf{21}(w \cdot x > 0) - 1 : w \in \mathbb{R}^d\}$ . It can be seen that the canonical basis vectors  $\{e_1, \dots, e_d\}$  (or more generally, any set of linearly independent examples) is shattered by  $\mathcal{H}$ . To see this, note that given a set of linearly independent examples  $x_1, \dots, x_m$ , consider the matrix  $M \in \mathbb{R}^{m \times d}$  whose rows are the  $x_i$ 's. Note that  $M$  has rank  $m$ , therefore its columns also spans the whole  $\mathbb{R}^m$ . Hence, for any vector  $l$  in  $\mathbb{R}^m$ , there is a vector  $w$  in  $\mathbb{R}^d$ , such that

$$Mw = \begin{bmatrix} \langle w, x_1 \rangle \\ \dots \\ \langle w, x_d \rangle \end{bmatrix} = l.$$

This immediately implies that for any labeling in  $\{-1, +1\}^m$ , there is a linear classifier in  $\mathbb{R}^m$  that achieves that labeling. Hence,  $\text{VC}(\mathcal{H}) \geq d$ .

However, consider any set  $S = \{x_1, \dots, x_{d+1}\}$ . We now show that  $S$  is not shatterable.

First,  $x_1, \dots, x_{d+1}$  are  $d+1$  vectors in  $\mathbb{R}^d$ , therefore they must be linearly dependent. Thus, there exists  $\alpha_1, \dots, \alpha_{d+1}$  not all zero, such that

$$\sum_{i=1}^{d+1} \alpha_i x_i = 0. \quad (1)$$

Furthermore, there exists  $\alpha_1, \dots, \alpha_{d+1}$ , such that there exists  $i^*$  in  $\{1, \dots, d+1\}$ ,  $\alpha_{i^*} > 0$ ,

$$\sum_{i=1}^{d+1} \alpha_i x_i = 0.$$

The reason is as follows: if there already exists a positive  $\alpha_i$  in Equation 1, then we are done; otherwise, we can flip the sign of the  $\alpha_i$ 's and ensuring at least one positive  $\alpha_i$ .

Now consider the following labeling  $(l_1, \dots, l_{d+1})$ , where  $l_i = \begin{cases} +1, \alpha_i > 0 \\ -1, \alpha_i \leq 0 \end{cases}$ . Can a linear classifier achieve such labeling? Suppose there is a  $w$  that achieves so. Then, for all  $i$  in  $\{1, \dots, d+1\}$ ,

$$\begin{cases} w \cdot x_i > 0, \alpha_i > 0 \\ w \cdot x_i \leq 0, \alpha_i \leq 0 \end{cases}$$

Thus, for all  $i$ ,  $\alpha_i \langle w, x_i \rangle \geq 0$ . Specifically, for index  $i^*$ ,  $\alpha_{i^*} \langle w, x_{i^*} \rangle > 0$ . Summing over all  $i$ 's, this implies that

$$\sum_{i=1}^{d+1} \alpha_i \langle w, x_i \rangle > 0.$$

This contradicts Equation 1, which would imply that

$$\sum_{i=1}^{d+1} \alpha_i \langle w, x_i \rangle = 0.$$

4. Non-homogeneous linear classifiers in  $\mathbb{R}^d$ .  $\mathcal{H} = \{h_w(x) \triangleq 2\mathbf{1}(w \cdot x + w_0 > 0) - 1 : (w, w_0) \in \mathbb{R}^{d+1}\}$ .

Using the same reasoning as above, it can be shown that the VC dimension of  $\mathcal{H}$  is  $d+1$ . We leave the proof to you as an exercise.

Finally, we define the notion of growth function, which measures the largest possible number of for  $\mathcal{H}$  to datasets of fixed size  $n$ .

**Definition 3.** Define the growth function  $\mathcal{S}(\mathcal{H}, n)$  as the maximum number of labelings one can generate on a dataset of size  $n$ , formally,

$$\mathcal{S}(\mathcal{H}, n) = \max_{S: |S|=n} |\Pi_{\mathcal{H}}(S)|.$$

We also have the simple observation for finite classes.

**Lemma 2.** If  $\mathcal{H}$  is finite, then  $\mathcal{S}(\mathcal{H}, n) \leq |\mathcal{H}|$  and  $\text{VC}(\mathcal{H}) \leq \log_2 |\mathcal{H}|$ .

*Proof.* The first statement is trivial as  $|\Pi_{\mathcal{H}}(S)| \leq |\mathcal{H}|$ . For the second statement, suppose  $\mathcal{H}$  shatters  $S$ . Then,  $2^{|S|} \leq |\Pi_{\mathcal{H}}(S)| \leq |\mathcal{H}|$ , implying that  $|S| \leq \log_2 |\mathcal{H}|$ . Therefore,  $\text{VC}(\mathcal{H})$ , the maximum sizes of a dataset shatterable by  $\mathcal{H}$  is at most  $\log_2 |\mathcal{H}|$ .  $\square$

### 3 Sauer's Lemma: bounding the growth function

Suppose we have a hypothesis class  $\mathcal{H}$  of VC dimension  $d$ , and a set of  $m$  examples  $\{x_1, \dots, x_m\}$ . We already know that when  $m \leq d$ ,  $\mathcal{S}(\mathcal{H}, n)$  can be as large as  $2^m$ . Can we give a good characterization of  $\mathcal{S}(\mathcal{H}, n)$  when  $m > d$  (other than the trivial upper bound of  $2^m - 1$ )? We have the following important combinatorial lemma, discovered independently by several authors (including Sauer, Shelah, Perles, Vapnik and Chervonenkis) in the 70s.

**Theorem 2** (Sauer's Lemma). *Suppose  $\mathcal{H}$  is a nonempty hypothesis class, and  $S = \{x_1, \dots, x_n\}$  is a set of  $m$  unlabeled examples. Then,*

$$|\Pi_{\mathcal{H}}(S)| \leq |\{T \subseteq S : \mathcal{H} \text{ shatters } T\}|.$$

Consequently, if  $\text{VC}(\mathcal{H}) = d$ , then

$$\mathcal{S}(\mathcal{H}, m) \leq \sum_{i=0}^d \binom{m}{i}.$$

(The right hand side is often abbreviated as  $\binom{m}{\leq d}$ .) Here we use the convention that  $\mathcal{H}$  always shatters an empty set.

**Remark.** We see that the growth function, as a function of  $m$ , has the following behavior on its upper bound: when  $m \leq d$ , the upper bound grows exponentially with  $m$ ; however, when  $m > d$ , the upper bound grows as a polynomial of  $m$ , which is substantially slower. We will see in the next section why this type of growth is useful for establishing uniform convergence guarantees.

*Proof.* We will show the first claim by induction on the size of sample  $m$ .

- Base case. If  $m = 1$ , then there are two subcases to consider: if  $\mathcal{H}$  classifies  $x_1$  in both  $+1$  and  $-1$  labels, then the left hand side is 2, and the right hand side is also 2. Otherwise,  $\mathcal{H}$  classifies  $x_1$  in only one label, then both sides are equal to 1.
- Inductive case. Before proceeding, we need the following important definition. Define a modification of the original hypothesis class  $\mathcal{H}$ : for every labeling  $(l_1, \dots, l_m)$  in  $\Pi_{\mathcal{H}}(S)$ , we select one representative classifier  $h$  in  $\mathcal{H}$  that achieves the labeling; we call the collection of the classifiers selected  $\mathcal{H}_S$ . Note that  $|\mathcal{H}_S| = |\Pi_{\mathcal{H}}(S)|$ . In addition, define  $S' = \{x_1, \dots, x_{m-1}\}$ .

Now, given  $\mathcal{H}_S$ , let us decompose it to two hypothesis classes,  $\mathcal{H}_1$  and  $\mathcal{H}_2$ , in the following manner. Consider a labeling  $(l_1, \dots, l_{m-1})$  achieved by  $\mathcal{H}_2$  on examples  $S'$ .

- If both  $(l_1, \dots, l_{m-1}, +1)$  and  $(l_1, \dots, l_{m-1}, -1)$  are achievable by  $\mathcal{H}_S$ , then we allocate the pair of classifiers such that one of them goes to  $\mathcal{H}_1$ , and the other goes to  $\mathcal{H}_2$ .
- Otherwise, only one of  $(l_1, \dots, l_{m-1}, +1)$  and  $(l_1, \dots, l_{m-1}, -1)$  is achievable by  $\mathcal{H}_S$ , then we send the classifier that achieves that labeling to  $\mathcal{H}_1$ .

See Tables 1, 2 and 3 for an example.

Classifier	$x_1$	$x_2$	$x_3$	$x_4$
$h_1$	–	–	–	–
$h_2$	–	–	–	+
$h_3$	–	+	–	+
$h_4$	+	–	–	–
$h_5$	+	–	–	+

Table 1: An example with  $m = 4$  and  $|\mathcal{H}_S| = 5$ . The matrix shows  $\mathcal{H}_S$ 's labelings on  $\{x_1, x_2, x_3, x_4\}$ .

Classifier	$x_1$	$x_2$	$x_3$	$x_4$
$h_1$	−	−	−	−
$h_3$	−	+	−	+
$h_5$	+	−	−	+

Classifier	$x_1$	$x_2$	$x_3$	$x_4$
$h_2$	−	−	−	+
$h_4$	+	−	−	−

Table 2:  $\mathcal{H}_2$ 's labelings on  $\{x_1, x_2, x_3, x_4\}$ . Table 3:  $\mathcal{H}_1$ 's labelings on  $\{x_1, x_2, x_3, x_4\}$ .

By construction, we have the following three simple but important observations:

- Claim 1.** 1.  $|\mathcal{H}_1| = |\Pi_{\mathcal{H}_1}(S')|$ ,  $|\mathcal{H}_2| = |\Pi_{\mathcal{H}_2}(S')|$ .  
2. If  $T \subset S'$  and  $\mathcal{H}_1$  shatters  $T$ , then it is also achieved by  $\mathcal{H}_S$ .  
3. If  $T \subset S'$  and  $\mathcal{H}_2$  shatters  $T$ , then  $\mathcal{H}_S$  shatters  $T \cup \{x_m\}$ .

Now, let us upper bound the size of  $\mathcal{H}_S$ :

$$\begin{aligned}
|\mathcal{H}_S| &= |\mathcal{H}_1| + |\mathcal{H}_2| \\
&= |\Pi_{\mathcal{H}_1}(S')| + |\Pi_{\mathcal{H}_2}(S')| \\
&\leq |\{T \subseteq S' : T \text{ shattered by } \mathcal{H}_1\}| + |\{T \subseteq S' : T \text{ shattered by } \mathcal{H}_2\}| \\
&\leq |\{T \subseteq S' : T \text{ shattered by } \mathcal{H}\}| + |\{T \subseteq S' : T \cup x_m \text{ shattered by } \mathcal{H}\}| \\
&= |\{T \subseteq S : x_m \notin T, T \text{ shattered by } \mathcal{H}\}| + |\{T \subseteq S : x_m \in T, T \text{ shattered by } \mathcal{H}\}| \\
&= |\{T \subseteq S : T \text{ shattered by } \mathcal{H}\}|
\end{aligned}$$

For the second statement, observe that all subsets  $T$  shatterable by  $\mathcal{H}$  is of size at most  $d$ . The right hand side is exactly the number of subsets of size at most  $d$ .  $\square$

*Proof of Claim 3.* We show the three items respectively.

1. The first statement is trivial, as by construction, for every labeling  $(l_1, \dots, l_{m-1})$ , there is at most one classifier in  $\mathcal{H}_1$  (resp.  $\mathcal{H}_2$ ).
2. The second statement is also trivial, as  $\mathcal{H}_1$  is a subset of  $\mathcal{H}_S$ .
3. Suppose some classifier  $h$  in  $\mathcal{H}_2$  achieves certain labeling  $(b_1, \dots, b_{|T|})$  on  $T$ . Suppose  $h$ 's full labeling on  $S'$  is  $(l_1, \dots, l_{m-1})$  (which is consistent with  $(b_1, \dots, b_{|T|})$ ). Then by construction, both  $(l_1, \dots, l_{m-1}, +1)$  and  $(l_1, \dots, l_{m-1}, -1)$  are achieved by  $\mathcal{H}_S$ . This implies that  $\mathcal{H}_S$  achieves labelings  $(b_1, \dots, b_{|T|}, +1)$  and  $(b_1, \dots, b_{|T|}, -1)$  on  $T \cup \{x_m\}$ . Therefore, if  $\mathcal{H}_2$  achieves all  $2^{|T|}$  labelings on  $T$ , then  $\mathcal{H}_S$  achieves all  $2^{|T|+1}$  labelings on  $T \cup \{x_m\}$ .

**Example.** Consider the example in Tables 1 and 3. Observe that  $\mathcal{H}_2$  shatters  $T = \{x_1\}$  with classifiers  $h_2$  and  $h_4$ . It can be seen that  $\mathcal{H}_S$  also shatters  $T \cup \{x_4\} = \{x_1, x_4\}$  with classifiers  $h_1, h_2, h_4$  and  $h_5$ .  $\square$

**Remark.** The growth function bound  $\binom{m}{\leq d}$  can further be upper bounded by  $(m+1)^d$  or  $(\frac{em}{d})^d$ .

## A A rigorous proof of Theorem 1

*Proof.* As  $D$  is realizable wrt  $\mathcal{H}$ , there exists a classifier  $h_{t^*}$  that has zero error on  $D$ .

Let us consider two critical thresholds  $t_L$  and  $t_R$ , defined as follows:

$$t_L = \sup \{t \in [0, 1] : \mathbb{P}(t \leq x \leq t^*) \geq \epsilon\}$$

If  $\mathbb{P}(0 \leq x \leq t^*) < \epsilon$ , then  $t_L$  is defined as 0.

$$t_R = \inf \{t \in [0, 1] : \mathbb{P}(t^* < x \leq t) \geq \epsilon\}$$

If  $\mathbb{P}(t^* < x \leq 1) < \epsilon$ , then  $t_R$  is defined as 1.

Suppose for the moment that both  $t_L$  and  $t_R$  are in  $(0, 1)$ . Our plan is to show the following:

1. With probability  $1 - \delta$ , the returned threshold  $\hat{t}$  lies in  $[t_L, t_R]$ .
2. Wherever  $\hat{t}$  lies in  $[t_L, t_R]$ ,  $h_{\hat{t}}$  has error at most  $\epsilon$ .

We show the two items respectively:

1. By Lemma 3, we have that

$$\mathbb{P}(t^* < x \leq t_R) \geq \epsilon, \quad \mathbb{P}(t_L \leq x \leq t^*) \geq \epsilon.$$

Now, consider event  $E_L$  (resp.  $E_R$ ) as the one that for all  $i$ , none of  $x_i$  are in  $[t_L, t^*]$  (resp.  $(t^*, t_R]$ ). In addition, define  $E = E_L \cup E_R$ .

Observe that

$$\mathbb{P}(E_L) = \mathbb{P}(\text{for all } i, x_i \notin [t_L, t^*]) \leq (1 - \epsilon)^m \leq e^{-m\epsilon} \leq \delta/2.$$

Similarly,  $\mathbb{P}(E_R) \leq \delta/2$ . By union bound,  $\mathbb{P}(E) \leq \mathbb{P}(E_L) + \mathbb{P}(E_R) \leq \delta$ . Therefore, in the event  $\bar{E}$  (which happens with probability  $1 - \delta$ ), there is an  $x_i$  (resp.  $x_j$ ) in  $[t_L, t^*]$  (resp.  $(t^*, t_R]$ ). Note that  $x_i$  has label  $-1$  and  $x_j$  has label  $+1$ . Thus, the consistency algorithm will return a threshold  $\hat{t}$  between  $[t_L, t_R]$  (Note that  $\hat{t}$  cannot be  $t_R$ , as this would misclassify  $x_j$ ).

2. Suppose  $\bar{E}$  happens. We show that the generalization error of the returned threshold classifier  $h_{\hat{t}}$  is at most  $\epsilon$ .

- (a) Suppose  $\hat{t} < t^*$ . As argued above,  $\hat{t} \geq t_L$ . Therefore,

$$\text{err}(h_{\hat{t}}, D) = \mathbb{P}(\hat{t} < x \leq t^*) \leq \mathbb{P}(t_L < x \leq t^*) \leq \epsilon,$$

where the inequality is from item 1 of Lemma 3.

- (b) Suppose  $\hat{t} \geq t^*$ . As argued above,  $\hat{t} < t_R$ . Therefore,

$$\text{err}(h_{\hat{t}}, D) = \mathbb{P}(t^* < x \leq \hat{t}) \leq \mathbb{P}(t^* < x < t_R) \leq \epsilon,$$

where the inequality is from item 2 of Lemma 3.

Now for the general case, where  $t_L$  can be 0 or  $t_R$  can be 1. Note that both cannot happen at the same time. Suppose  $t_L = 0$ , then by the exact same reasoning, we can show that with probability  $1 - \delta$ ,  $\hat{t}$  is in  $[t^*, t_R]$  or  $[0, t^*]$ . In the former case, as have been argued before,

$$\text{err}(h_{\hat{t}}, D) \leq \mathbb{P}(t^* \leq x \leq \hat{t}) \leq \mathbb{P}(t^* \leq x < t_R) \leq \epsilon.$$

In the latter case,

$$\text{err}(h_{\hat{t}}, D) = \mathbb{P}(\hat{t} < x \leq t^*) \leq \mathbb{P}(0 \leq x \leq t^*) \leq \epsilon.$$

In summary, with probability  $1 - \delta$ ,  $\text{err}(h_{\hat{t}}, D) \leq \epsilon$ . The case of  $t_R = 1$  is symmetric and is left as exercise.  $\square$

The following lemma crucially uses the continuity property of probability measure, that is, If  $A_1 \subset \dots \subset A_n \subset \dots$ , and  $A = \bigcup_{n=1}^{\infty} A_n$  (abbrev.  $A_n \uparrow A$ ), then  $\lim_{n \rightarrow \infty} \mathbb{P}(A_n) = \mathbb{P}(A)$ .

**Lemma 3.** 1. Suppose  $\mathbb{P}(0 \leq x \leq t^*) \geq \epsilon$ . Consider

$$t_L \triangleq \sup \{t \in [0, 1] : \mathbb{P}(t \leq x \leq t^*) \geq \epsilon\}.$$

Then,

$$\mathbb{P}(t_L \leq x \leq t^*) \geq \epsilon,$$

$$\mathbb{P}(t_L < x \leq t^*) \leq \epsilon.$$

2. Suppose  $\mathbb{P}(t^* < x \leq 1) \geq \epsilon$ . Consider

$$t_R \triangleq \inf \{t \in [0, 1] : \mathbb{P}(t^* < x \leq t) \geq \epsilon\}.$$

Then,

$$\mathbb{P}(t^* < x \leq t_R) \geq \epsilon,$$

$$\mathbb{P}(t^* < x < t_R) \leq \epsilon.$$

*Proof.* We only show the first item. The second item is left as an exercise.

First, by the definition of  $t_L$ , for all  $t < t_L$ ,  $\mathbb{P}(t \leq x \leq t^*) \geq \epsilon$ . As events  $\{t_L - \frac{1}{n} \leq x \leq t^*\} \downarrow \{t_L \leq x \leq t^*\}$  as  $n \rightarrow \infty$ , this implies that

$$\mathbb{P}(t_L \leq x \leq t^*) = \lim_{n \rightarrow \infty} \mathbb{P}(t_L - \frac{1}{n} \leq x \leq t^*) \geq \epsilon.$$

Second, by the definition of  $t_L$ , for all  $t > t_L$ ,  $\mathbb{P}(t \leq x \leq t^*) < \epsilon$ . As events  $\{t_L + \frac{1}{n} \leq x \leq t^*\} \downarrow \{t_L < x \leq t^*\}$  as  $n \rightarrow \infty$ , this implies that

$$\mathbb{P}(t_L < x \leq t^*) = \lim_{n \rightarrow \infty} \mathbb{P}(t_L + \frac{1}{n} \leq x \leq t^*) \leq \epsilon. \quad \square$$