

CSC 665: Midterm

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November 10, 2019

Please complete the following set of problems. You must do the exercises completely on your own (no collaboration allowed this time). The exam is due **on Oct 24, 12:30pm, on Gradescope**. You are free to cite existing theorems from the textbooks and course notes.

Problem 1

Define $\mathcal{H} = \{\text{sign}(p(x)) : p \text{ is a polynomial of } x \text{ of degree } \leq n\}$ (where $x \in \mathbb{R}$). Here $\text{sign}(z) = 2\mathbf{1}(z > 0) - 1$. What is the VC dimension of \mathcal{H} ?

Solution

We show $\text{VC}(\mathcal{H}) = n + 1$.

1. We first show $\text{VC}(\mathcal{H}) \geq n + 1$, i.e. there exists $n + 1$ points that are shattered by \mathcal{H} . Pick $n + 1$ distinct numbers x_1, \dots, x_{n+1} in \mathbb{R} . By a standard fact in analysis, for any values y_1, \dots, y_{n+1} in \mathbb{R} , there exists a unique degree-at-most- n polynomial p that passes all points $(x_i, y_i)_{i=1}^{n+1}$ (this can be shown using Cramer's rule or Lagrange polynomials).

Specifically, for any y_1, \dots, y_{n+1} in $\{\pm 1\}^{n+1}$, there exists p of degree of at most n such that $p(x_i) = y_i$; as a consequence, $\text{sign}(p)$ in \mathcal{H} also achieves the labeling of (y_1, \dots, y_{n+1}) on (x_1, \dots, x_{n+1}) . Therefore, \mathcal{H} shatters x_1, \dots, x_{n+1} , a set of size $n + 1$.

2. We now show $\text{VC}(\mathcal{H}) \leq n + 1$. Note that an alternative way of writing \mathcal{H} is $\{\text{sign}(\langle a, \phi(x) \rangle) : a \in \mathbb{R}^{n+1}\}$, where $\phi(x) = (1, x, \dots, x^n)$. If there are $n + 2$ points x_1, \dots, x_{n+2} shattered by \mathcal{H} , this also means that $\phi(x_1), \dots, \phi(x_{n+2})$ is shattered by $\mathcal{F} = \{\text{sign}(\langle a, x \rangle) : a \in \mathbb{R}^{n+1}\}$. But in class, we have seen that \mathcal{F} has VC dimension $n + 1$, contradiction.

Problem 2

Suppose we have an algorithm \mathcal{B} that learns hypothesis class \mathcal{H} in the following sense. There exists a function $m(\epsilon)$, such that for any $\epsilon > 0$, suppose \mathcal{B} draws $m \geq m(\epsilon)$ training examples from a distribution D realizable by \mathcal{H} , then with probability $\geq \frac{1}{2}$, \mathcal{B} returns a classifier \hat{h} with error at most ϵ on D .

Now, given \mathcal{B} and the ability of drawing fresh training examples, how can you design an algorithm \mathcal{A} that (ϵ, δ) -PAC learns \mathcal{H} for any ϵ, δ ? What is its sample complexity? (You may want to run \mathcal{B} multiple times.)

Solution

Consider the following algorithm $\mathcal{A}(\epsilon, \delta)$:

1. Repeated $\mathcal{B}(\epsilon/2)$ for k times, getting classifiers $h_1 \dots, h_k$, where $k = \lceil \log_2 \frac{2}{\delta} \rceil$.
2. $S \leftarrow \text{Sample } \frac{32}{\epsilon^2} \ln \frac{4k}{\delta}$ iid examples from D .
3. Select $\hat{h} = \arg \min_{h \in \mathcal{F}} \text{err}(h, S)$, where $\mathcal{F} = \{h_1, \dots, h_k\}$.

We show that \mathcal{A} outputs a classifier with error ϵ with probability $1 - \delta$. Define event $E_1 = \{\exists i, \text{err}(h_i, D) \leq \epsilon/2\}$, and event $E_2 = \{\forall h \in \mathcal{F}, |\text{err}(h, S) - \text{err}(h, D)| \leq \epsilon/4\}$.

First,

$$\mathbb{P}(E_1) = 1 - \prod_{i=1}^k \mathbb{P}(\text{err}(h_i, D) > \epsilon/2) \geq 1 - \frac{1}{2^k} \geq 1 - \delta/2.$$

Second, by Hoeffding's inequality, given h_1, \dots, h_k , as S is a fresh set of examples independent of h_1, \dots, h_k ,

$$\mathbb{P}(E_2 | h_1, \dots, h_k) \geq 1 - \sum_{i=1}^k \mathbb{P}(|\text{err}(h_i, S) - \text{err}(h_i, D)| > \epsilon/4) = 1 - 2k \cdot e^{-2m \cdot \frac{\epsilon^2}{16}} \geq 1 - \delta/2.$$

Therefore, $\mathbb{P}(E_2) \geq 1 - \delta/2$, and consequently, by union bound, $\mathbb{P}(E_1 \cap E_2) \geq 1 - \delta$.

Now on event E_2 , we have that

$$\text{err}(\hat{h}, D) \leq \text{err}(\hat{h}, S) + \epsilon/4 = \min_{i=1}^k \text{err}(h_i, S) + \epsilon/4 \leq \min_{i=1}^k \text{err}(h_i, D) + \epsilon/2.$$

Therefore, on event $E_1 \cap E_2$,

$$\text{err}(\hat{h}, D) \leq \min_{i=1}^k \text{err}(h_i, D) + \epsilon/2 \leq \epsilon/2 + \epsilon/2 = \epsilon.$$

As can be seen from the description of \mathcal{A} , it has a sample complexity of

$$km(\epsilon/2) + \frac{32}{\epsilon^2} \ln \frac{4k}{\delta} = \lceil \log_2 \frac{2}{\delta} \rceil \cdot m(\epsilon/2) + \frac{32}{\epsilon^2} \ln \frac{4k}{\delta}.$$

Problem 3

Suppose X_1, \dots, X_n is a sequence of n iid random variables, and let $\sigma^2 = \text{var}(X_i)$ and $\mu = \mathbb{E}(X_i)$. Suppose $n = mk$ for some integer m and odd integer $k \geq 20 \ln \frac{1}{\delta}$. Denote by

$$\hat{\mu} = \text{median}(\hat{\mu}_1, \dots, \hat{\mu}_k),$$

where $\hat{\mu}_i = \frac{1}{m} \sum_{j=(i-1)m+1}^{im} X_j$.

1. Show that for every j ,

$$\mathbb{P}(|\hat{\mu}_j - \mu| \leq \frac{2\sigma}{\sqrt{m}}) \geq \frac{3}{4}.$$

2. Show that

$$\mathbb{P}(|\hat{\mu} - \mu| \leq \frac{2\sigma}{\sqrt{m}}) \geq 1 - \delta.$$

Solution

1. $\text{var}(\hat{\mu}_j) = \frac{1}{m^2} \sum_{j=(i-1)m+1}^{im} \text{var}(X_j) = \frac{1}{m^2} \cdot m \cdot \sigma^2 = \frac{\sigma^2}{m}.$

By Chebyshev's Inequality,

$$\mathbb{P}(|\hat{\mu}_j - \mu| > \frac{2\sigma}{\sqrt{m}}) \leq \frac{\delta(\hat{\mu}_j)}{(\frac{2\sigma}{\sqrt{m}})^2} \leq \frac{1}{4}.$$

The stated result follows by taking the complement of the event considered.

2. Denote by

$$Z_j = \mathbf{1}(|\hat{\mu}_j - \mu| \leq \frac{2\sigma}{\sqrt{m}}).$$

Note that Z_1, \dots, Z_k are independent, and has mean p at least $\frac{3}{4}$. By Hoeffding's inequality, when $k \geq 20 \ln \frac{1}{\delta}$,

$$\mathbb{P}(\sum_{j=1}^k Z_j \geq \frac{k}{2}) \geq 1 - \exp\left(2k(p - \frac{1}{2})^2\right) \geq 1 - \delta.$$

Now, consider the event $E = \left\{\sum_{j=1}^k Z_j \geq \frac{k}{2}\right\}$. We claim that under E , $\hat{\mu}$ would be inside interval $I = [\mu - \frac{2\sigma}{\sqrt{m}}, \mu + \frac{2\sigma}{\sqrt{m}}]$. Indeed, if $\hat{\mu}$ is outside I , then two cases would happen:

- $\hat{\mu} < \mu - \frac{2\sigma}{\sqrt{m}}$. As $\hat{\mu}$ is the median of μ_i 's, at least half of μ_i 's would also be smaller than $\mu - \frac{2\sigma}{\sqrt{m}}$, contradiction to the fact that E happens.
- $\hat{\mu} > \mu + \frac{2\sigma}{\sqrt{m}}$. Symmetrically, this would also contradict with the fact that E happens.

In summary, in event E , which happens with probability $1 - \delta$, $\hat{\mu} \in [\mu - \frac{2\sigma}{\sqrt{m}}, \mu + \frac{2\sigma}{\sqrt{m}}]$.

Remark. You may wonder: *what is the motivation of this question?* The estimator $\hat{\mu}$ is interesting in that it gives a better mean estimator than naive sample mean for heavy-tailed random variables. Sample mean can sometimes have bad concentration properties; see [1, 2] for discussions.

References

- [1] Jean-Yves Audibert, Olivier Catoni, et al. Robust linear least squares regression. *The Annals of Statistics*, 39(5):2766–2794, 2011.
- [2] Daniel Hsu and Sivan Sabato. Loss minimization and parameter estimation with heavy tails. *The Journal of Machine Learning Research*, 17(1):543–582, 2016.