CSC 665: Homework 2

Chicheng Zhang

October 11, 2019

Please complete the following set of exercises. You must write down your solutions **on your own**. If you have discussed with your classmates on any of the questions, please indicate so in your solutions. The homework is due **on Oct 10, 12:30pm, on Gradescope**. You are free to cite existing theorems from the textbook and course notes.

Problem 1

Do Exercise 2.3 in (Shalev-Shwartz and Ben-David, 2014). For item 2, you can assume that the joint distribution of (X_1, X_2) is continuous over $\mathbb{R}^2$.

Problem 2

1. Show the following inequality: for positive a, b and x, if $x > 2a\ln(2a) + 2b$, then $x > a\ln x + b$.
2. Show the following basic inequality: for n, d such that $n \geq 2$ and $n \geq d$, $\binom{n}{\leq d} \leq n^{d+1}$.
3. Consider l hypothesis classes $\mathcal{H}_1, \mathcal{H}_2, \mathcal{H}_l$, where $\mathrm{VC}(\mathcal{H}_i) = v \geq 1$. Define $\mathcal{H} \triangleq \cup_{i=1}^{l} \mathcal{H}_i$. Show that there exists some constant $c > 0$ such that
$$\mathrm{VC}(\mathcal{H}) \leq c \cdot \big(v\ln(v) + \ln(l)\big).$$
4. Let $\mathcal{H} = \big\{\mathrm{sign}(\langle w, x\rangle) : w \in \mathbb{R}^d, |\{i : w_i \neq 0\}| = k\big\}$ be the set of k-sparse homogenenous linear classifiers in $\mathbb{R}^d$, where $k \leq d$. Show that there exists some constant $c > 0$ such that
$$\mathrm{VC}(\mathcal{H}) \leq c \cdot (k\ln d).$$
5. Consider l hypothesis classes $\mathcal{H}_1, \ldots, \mathcal{H}_l$, where $\mathrm{VC}(\mathcal{H}_i) = d_i \geq 1$. Suppose f is a function from $\{\pm 1\}^l$ to $\{\pm 1\}$ (for example, the majority function $f(z_1, \ldots, z_l) = \mathrm{sign}(\sum_{i=1}^{l} z_i)$ or the parity function $f(z_1, \ldots, z_l) = \prod_{i=1}^{l} z_i$). Define $\mathcal{H} \triangleq \big\{f(h_1(x), \ldots, h_l(x)) : h_1 \in \mathcal{H}_1, \ldots, h_l \in \mathcal{H}_l\big\}$. Show that there exists some constant $c > 0$ such that
$$\mathrm{VC}(\mathcal{H}) \leq c\left(\sum_{i=1}^{l} d_i\right)\ln\left(\sum_{i=1}^{l} d_i\right).$$

Problem 3

In this exercise, we will show that, under the *realizable setting*, with hypothesis class $\mathcal{H}$ having VC dimension d, ERM (in fact, the consistency algorithm) will have a PAC sample complexity of $O\left(\frac{1}{\epsilon}(d\ln\frac{1}{\epsilon}+\ln\frac{1}{\delta})\right)$. Suppose $S=\{Z_1,\ldots,Z_m\}$ a set of m training examples drawn iid from distribution D, where each $Z_i=(X_i,Y_i)$ is a labeled example. In addition, $\mathcal{F}=\left\{\mathbf{1}(h(x)\neq y): h\in\mathcal{H}\right\}$ is the zero-one loss function class. Our proof will mostly follow the steps for showing agnostic PAC sample complexity given in the lecture.

1. **Double Sampling Trick.** Fix a training set S. Suppose $\mathbb{E}_S f(Z)=0$ and $\mathbb{E}_D f(Z)\geq\epsilon$. Show that for a fresh set of random examples S' of size m ($m\geq\frac{16}{\epsilon}$) sampled iid from D:

$$\mathbb{P}_{S'\sim D^m}\left(\mathbb{E}_{S'}f(Z)\geq\frac{\epsilon}{2}\right)\geq\frac{1}{2}.$$

2. **Conditioning.** Denote by events

$$E'\triangleq\left\{\text{there exists } f\in\mathcal{F}, \mathbb{E}_S f(Z)=0, \mathbb{E}_{S'}f(Z)\geq\frac{\epsilon}{2}\right\},$$

$$E\triangleq\left\{\text{there exists } f\in\mathcal{F}, \mathbb{E}_S f(Z)=0, \mathbb{E}_D f(Z)\geq\epsilon\right\}.$$

Show $\mathbb{P}_{S,S'\sim D^m}(E'|E)\geq\frac{1}{2}$, and conclude that $\mathbb{P}_{S\sim D^m}(E)\leq 2\mathbb{P}_{S,S'\sim D^m}(E')$.

3. **Symmetrization.** Introduce $\sigma=(\sigma_1,\ldots,\sigma_m)$ where each $\sigma_i\in\{\pm1\}$. Denote by

$$(W_i,W_i')=\begin{cases}(Z_i,Z_i') & \sigma_i=+1,\\ (Z_i',Z_i) & \sigma_i=-1.\end{cases}$$

Show that

$$\mathbb{P}_{S,S'\sim D^m}(E')=\mathbb{P}_{S,S'\sim D^m,\sigma\sim\mathrm{R}^m}\left(\text{exists } f\in\mathcal{F}, \sum_{i=1}^m f(W_i)=0, \sum_{i=1}^m f(W_i')\geq\frac{m\epsilon}{2}\right),$$

where R is the Rademacher distribution, i.e. uniform in $\{\pm1\}$.

4. **The randomness in Rademacher random variables.** Fix two size m training sets S and S'. Show that for a fixed classifier f in $\mathcal{F}$,

$$\mathbb{P}_{\sigma\sim\mathrm{R}^n}\left(\sum_{i=1}^m f(W_i)=0, \sum_{i=1}^m f(W_i')\geq\frac{m\epsilon}{2}\right)\leq\exp\left(-\frac{m\epsilon}{4}\right).$$

5. Use the above items to conclude that for $m\geq\frac{16}{\epsilon}$,

$$\mathbb{P}_{S\sim D^m}(\text{there exists } f\in\mathcal{F}, \mathbb{E}_S f(Z)=0, \mathbb{E}_D f(Z)\geq\epsilon)\leq 2\mathcal{S}(\mathcal{F},2m)\exp\left\{-\frac{m\epsilon}{4}\right\}.$$

In addition, show that ERM has a PAC sample complexity of $O\left(\frac{1}{\epsilon}(d\ln\frac{1}{\epsilon}+\ln\frac{1}{\delta})\right)$.

Problem 4

In this exercise, we develop sample complexity lower bounds for *realizable* PAC learning using Le Cam's method and Assouad's method. Suppose hypothesis class $\mathcal{H}$ has VC dimension $d \geq 2$, and it shatters examples $z_0, z_1, \ldots, z_{d-1}$. In additon, suppose $\epsilon, \delta \in (0, \frac{1}{8})$ are target error and target failure probability. A learning algorithm $\mathcal{A}$ is a mapping from training set S to $\{\pm 1\}$. In the following, you can use the elementary fact that for $x \in (0, \frac{1}{2})$, $e^{-x} \geq 1 - x \geq e^{-2x}$.

1. Consider D_{-1} and D_{+1} as follows: for every i in $\{\pm 1\}$,

$$D_i(x, y) = \begin{cases} 1 - 2\epsilon, & (x, y) = (z_0, -1), \\ 2\epsilon, & (x, y) = (z_1, i), \\ 0 & \text{otherwise.} \end{cases}$$

   Note that $\min_{h \in \mathcal{H}} \mathrm{err}(h', D_i) = 0$ for both $i \in \{\pm 1\}$. For every j in $\{\pm 1\}$, denote by $P_j((x_i, y_i)_{i=1}^m) = \prod_{i=1}^m D_j(x_i, y_i)$ the distribution over training sets (observations). Use Le Cam's method to show that for any hypothesis tester f, there exists an i in $\{\pm 1\}$, such that

$$\mathbb{P}_i(f(S) \neq i) > \frac{1}{2}(1 - 4\epsilon)^m.$$

2. Conclude that for any learning algorithm $\mathcal{A}$, if sample size $m \leq \frac{1}{4\epsilon} \ln \frac{1}{4\delta}$, then there exists an i in $\{\pm 1\}$,

$$\mathbb{P}_i(\mathrm{err}(\hat{h}, D_i) > \epsilon) > \delta.$$

3. For every $\tau \in \{\pm 1\}^{d-1}$, consider D_τ as follows:

$$D_\tau(x, y) = \begin{cases} 1 - 4\epsilon, & (x, y) = (z_0, -1), \\ \frac{4\epsilon}{d-1}, & (x, y) = (z_i, \tau_i) \text{ for some} i \in \{1, \ldots, d-1\}, \\ 0 & \text{otherwise.} \end{cases}$$

   Note that $\min_{h \in \mathcal{H}} \mathrm{err}(h', D_\tau) = 0$ for all $\tau \in \{\pm 1\}^{d-1}$. For every τ, denote by $P_\tau((x_i, y_i)_{i=1}^m) = \prod_{i=1}^m D_\tau(x_i, y_i)$ the distribution over training sets (observations).

   Use Assouad's method to show that for any hypothesis tester $f_1, \ldots, f_{d-1}$, there exists $\tau \in \{\pm 1\}^{d-1}$,

$$\mathbb{E}_\tau \left[\sum_{j=1}^{d-1} \mathbf{1}(f_j(S) \neq \tau_j) \right] > \frac{d-1}{2} \left(1 - \frac{4\epsilon}{d-1} \right)^m.$$

4. Conclude that for any learning algorithm $\mathcal{A}$, suppose that sample size $m \leq \frac{d-1}{128\epsilon}$, then there exists a $\tau \in \{\pm 1\}^d$, such that

$$\mathbb{P}_\tau(\mathrm{err}(\hat{h}, D_\tau) > \epsilon) > \frac{1}{4}.$$

Hints

2.1 Use the elementary fact that $\ln(z) \leq z - 1$ for $z = \frac{x}{2a}$.

2.2 Use the elementary fact that $\binom{n}{i} \leq n^i$.

2.3 (1) consider S of size n shattered by $\mathcal{H}$. We know that $|\Pi_{\mathcal{H}}(S)| = 2^n$. Use Sauer's Lemma to obtain an upper bound on $|\Pi_{\mathcal{H}}(S)|$ in terms of v. (2) consider using the contrapositive of item 1.

2.4 Write $\mathcal{H}$ as a union of $\binom{d}{k}$ hypothesis classes, each of which has VC dimension k, then apply item 3.)

3.1 Use Chernoff bound for Bernoulli distributions (the version with exponent $-\frac{mp\mu^2}{4}$).

3.4 Consider three cases: (1) there exists some i, $(f(Z_i), f(Z'_i)) = (1,1)$; (2) $\sum_{i=1}^m f(Z_i) + f(Z'_i) < \frac{m\epsilon}{2}$; (3) both (1) and (2) are not satisfied. Observe that in the first two cases, the probability is identically zero.

4.1 Consider observation $S = ((z_0, -1), \ldots, (z_0, -1))$. Show that $\mathbb{P}_{-1}(S) = \mathbb{P}_{+1}(S)$.

4.2 Define an appropriate hypothesis tester f that depends on $\mathcal{A}$.

4.3 Define $A_j = \left\{S = (x_i, y_i)_{i=1}^m : x_i \neq z_j \text{ for all } i\right\}$. Show that for every $\tau \overset{j}{\sim} \tau'$, $\mathbb{P}_\tau(S) = \mathbb{P}_{\tau'}(S)$ for all S in A_j. In addition, for $\sigma \in \left\{\tau, \tau'\right\}$, $\mathbb{P}_\sigma(S \in A_j) > \frac{7}{8}$. Intuitively, seeing only examples other than z_j does not help determining the optimal classifier's labeling on z_j.

4.4 First show that $\sum_{j=1}^{d-1} \mathbf{1}(f_j(S) \neq \tau_j) > \frac{d-1}{2}$ with probability $> \frac{1}{4}$. Then define an appropriate hypothesis tester $f = (f_1, \ldots, f_{d-1})$ that depends on $\mathcal{A}$.

Problem 5 (No need to submit)

In this problem, we develop an alternative proof of Sauer's Lemma: any hypotheis class $\mathcal{H}$ with VC dimension d can have at most $\binom{n}{\leq d}$ labelings on any dataset $S = \{z_1, \ldots, z_n\}$. Throughout, we will be using the notation that

$$\binom{\{1,\ldots,n\}}{d+1} \triangleq \left\{(i_1, \ldots, i_{d+1}) : 1 \leq i_1 < \ldots < i_{d+1} \leq n\right\}$$

to denote the set of $(d+1)$-tuples whose entries are distinct. Note that $\left|\binom{\{1,\ldots,n\}}{d+1}\right| = \binom{n}{d+1}$.

1. Show that for any indices $I = (i_1, \ldots, i_{d+1}) \in \binom{\{1,\ldots,n\}}{d+1}$, there exists a string $s_I \in \{\pm 1\}^{d+1}$, such that none of the labelings in
$$L_I = \left\{b \in \{\pm 1\}^n : (b_{i_1}, \ldots, b_{i_{d+1}}) = s_I\right\}$$
are achievable by classifiers in $\mathcal{H}$.

2. Show the following basic facts:
   (a) For a finite set A and a function f, denote by $f(A) = \left\{f(a) : a \in A\right\}$. Then $|f(A)| \leq |A|$, where $|B|$ denotes the cardinality of set B.
   (b) Suppose $\mathcal{I}$ is a set of indices. Given a collection of sets $\{L_I\}_{I\in\mathcal{I}}$ and a function f,
$$\left|\bigcup_{I\in\mathcal{I}} f(L_I)\right| \leq \left|\bigcup_{I\in\mathcal{I}} L_I\right|. \tag{1}$$

3. Use the above two facts to conclude that
$$\left|\bigcup_{I\in\binom{\{1,\ldots,n\}}{d+1}} L_I\right| \geq \sum_{i=d+1}^{n} \binom{n}{i}.$$
(Hint: consider functions $f_1, \ldots, f_n$, where $f_i(s_1, \ldots, s_n) = (s_1, \ldots, s_{i-1}, -1, s_{i+1}, \ldots, s_n)$ is the function that sets a length n string's i-th entry to -1. Iteratively applying Equation (1) for $f_1, \ldots, f_n$, what do you get?)

4. Use item 3 to conclude that $|\Pi_{\mathcal{H}}(S)| \leq \binom{n}{\leq d}$.