Improved algorithms for efficient active learning halfspaces with Massart and Tsybakov noise

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Abstract

We give an efficient PAC active halfspace learning algorithm that has improved noise-tolerance and label efficiency under benign noise conditions, given that the unlabeled data distribution satisfies certain structural properties [DKKTZ20]. Specifically:

- 1. Under Massart noise, it achieves optimal label complexity; such efficient and label-optimal results were previously only known when the unlabeled data distribution is uniform [YZ17].
- 2. Under two subfamilies of Tsybakov noise, it achieves improved label complexities compared to passive learning algorithms.

Problem: efficient active learning halfspaces with benign noise

• (x, \rightarrow) drawn from a distribution D

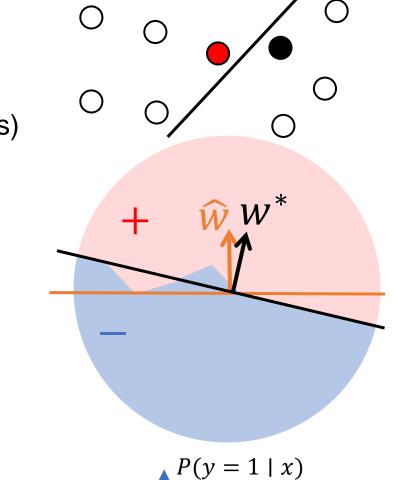
features interactive label queries

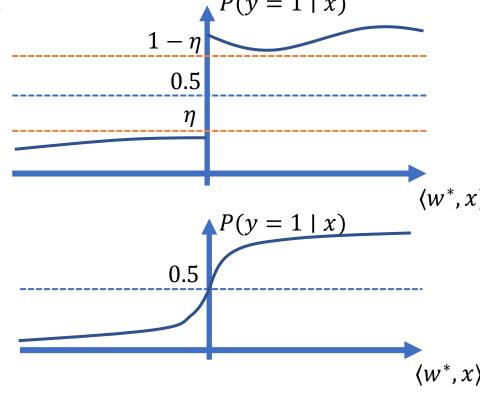


- Linear classifiers: $H = \{ sign(w \cdot x) : w \in \mathbb{R}^d \}$
- Error $err(w) = P(y \neq sign(w \cdot x))$
- Optimal linear classifier $w^* = \operatorname{argmin}_w \operatorname{err}(w)$
- Goal: computationally efficient algorithm that returns a vector \widehat{w} , such that $\operatorname{err}(\widehat{w}) \operatorname{err}(w^*) \leq \epsilon$, using a few label queries



- η -Massart [MN06]: for all x, $\eta(x) \le \eta < \frac{1}{2}$
- α -Tsybakov [T04] for $\alpha \in (0,1)$: for all t, $P_{\rm D}(1/2 \eta(x) \le t) \le O(t^{\alpha/(1-\alpha)})$
- α -Geometric Tsybakov [e.g., CN08]: for all $x, \frac{1}{2} \eta(x) \ge |w^* \cdot x|^{\frac{1-\alpha}{\alpha}}$





Main result: Massart noise

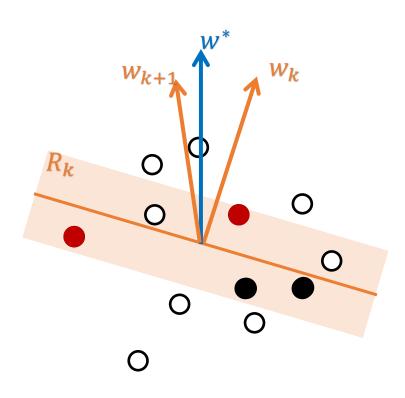
Algorithm	Efficient?	Label complexity in $\widetilde{0}$
[BL13]	No	$\frac{d}{(1-2\eta)^2}\operatorname{polylog}(1/\epsilon)$
[ZSA20]	Yes	$\frac{d}{(1-2\eta)^4} \text{ polylog}(1/\epsilon)$
This work	Yes	$\frac{d}{(1-2\eta)^2}$ polylog $(1/\epsilon)$

Main result: Tsybakov noise

Algorithm	Efficient?	Label complexity in 0
[BL13]	No	$d\left(\frac{1}{\epsilon}\right)^{2-2\alpha}$
[DKKTZ20]	Yes	$d\left(\frac{-}{\epsilon}\right)$ $\operatorname{poly}(d)\left(\frac{1}{\epsilon}\right)^{O(1/\alpha)}$
This work $(\alpha \in \left(\frac{1}{2}\right))$,1]) Yes	$d\left(\frac{1}{\epsilon}\right)_{2-2\alpha}^{\frac{2-2\alpha}{2\alpha-1}}$
This work (Geome Tsybakov)	tric Yes	$d\left(\frac{1}{\epsilon}\right)^{\frac{2-2\alpha}{\alpha}}$

Algorithm skeleton

 $w_1 \leftarrow \text{Initialize}()$. //Acute Initialization In phases $k=1,2,\ldots,k_0=\log(1/\epsilon)$: $w_{k+1} \leftarrow \text{Refine}(w_k,2^{-(k+1)})$. // Refinement Return w_{k_0+1} .



Refine: design challenges and related work

A series of prior works combine margin-based sampling with loss minimization techniques to design Refine:

- [BL13]:computationally inefficient (0-1 loss minimization)
- [ABHU15, ABHZ16]: analysis only tolerates $\eta \leq$ small constant, or requires high label complexity
- [ZSA20]: specialized to Massart noise (needs to know η)

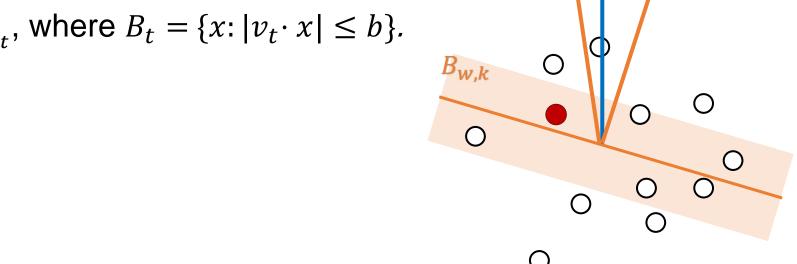
Refine: our design

For t = 1, 2, ..., T:

- 1. **Sample:** $(x_t, y_t) \leftarrow \text{example drawn from } D|_{B_t}, \text{ where } B_t = \{x: |v_t \cdot x| \leq b\}.$
- 2. **Update:** $v_{t+1} \leftarrow v_t \alpha g_t$, where $g_t = -y_t x_t$



Key difference from [ZSA20]: simpler definition of g_t leads to broader noise tolerance $^{\circ}$ Algorithmically similar to ``nonconvex optimization'' view [GCB09, DKTZ20], but analysis very different (see next)



Analysis: key ideas

Theorem: If $\theta(v_1, w^*) \leq 2\theta$, then with high probability, Refine (v_1, θ) returns a vector v with $\theta(v, w^*) \leq \theta$, if T is of order:

- $\frac{d}{(1-2\eta)^2}$, under η -Massart noise;
- $d\left(\frac{1}{\theta}\right)^{\frac{2-2\alpha}{2\alpha-1}}$, under α -Tsybakov noise with $\alpha \in \left(\frac{1}{2}, 1\right]$;
- $\psi_b(v) = \mathbb{E}[(1 2\eta(x)) | w^* \cdot x | \mid |v \cdot x| \le b]$

Key observation: Refine optimizes the following

`proximity function" in a nonstandard way:

Idea: rewriting OGD's regret guarantees over g_t 's:

$$\frac{1}{T} \sum_{t=1}^{T} \langle -w^*, g_t \rangle \le \frac{1}{T} \sum_{t=1}^{T} \langle -v_t, g_t \rangle + O\left(\frac{1}{\sqrt{T}}\right)$$

- $d\left(\frac{1}{\theta}\right)^{\frac{2-2\alpha}{\alpha}}$, under α -Geometric Tsybakov noise.
- Concentrates to $\frac{1}{T}\sum_{t=1}^{T}\psi_b(v_t)$

Can be made small by tuning b, T

The ``proximity function'' ψ_b

Lemma (simplified): For ``structured' D, under one of the three noise conditions, $\psi_b(v)$ is lower bounded by an increasing function of $\theta(v, w^*)$.

Consequently, optimizing $\psi_b(v) \Rightarrow$ optimizing $\theta(v, w^*)$

