

CSC 665: Homework 2

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Please complete the following set of exercises. You must write down your solutions **on your own**. If you have discussed with your classmates on any of the questions, please indicate so in your solutions. The homework is due **on Oct 15, 12:30pm, on Gradescope**. You are free to cite existing theorems from the textbook and course notes.

Problem 1

Do Exercise 2.3 in (Shalev-Shwartz and Ben-David, 2014). For item 2, you can assume that the joint distribution of (X_1, X_2) is continuous over \mathbb{R}^2 .

Problem 2

1. Show the following inequality: for positive a, b and x , if $x > 2a \ln(2a) + 2b$, then $x > a \ln x + b$.
2. Show the following basic inequality: for n, d such that $n \geq 2$ and $n \geq d$, $\binom{n}{\leq d} \leq n^{d+1}$.
3. Consider l hypothesis classes $\mathcal{H}_1, \mathcal{H}_2, \dots, \mathcal{H}_l$, where $\text{VC}(\mathcal{H}_i) = v \geq 1$. Define $\mathcal{H} \triangleq \cup_{i=1}^l \mathcal{H}_i$. Show that there exists some constant $c > 0$ such that

$$\text{VC}(\mathcal{H}) \leq c \cdot (v \ln(v) + \ln(l)).$$

4. Let $\mathcal{H} = \{\text{sign}(\langle w, x \rangle) : w \in \mathbb{R}^d, |\{i : w_i \neq 0\}| = k\}$ be the set of k -sparse homogenous linear classifiers in \mathbb{R}^d , where $k \leq d$. Show that there exists some constant $c > 0$ such that

$$\text{VC}(\mathcal{H}) \leq c \cdot (k \ln d).$$

5. Consider l hypothesis classes $\mathcal{H}_1, \dots, \mathcal{H}_l$, where $\text{VC}(\mathcal{H}_i) = d_i \geq 1$. Suppose f is a function from $\{\pm 1\}^l$ to $\{\pm 1\}$ (for example, the majority function $f(z_1, \dots, z_l) = \text{sign}(\sum_{i=1}^l z_i)$ or the parity function $f(z_1, \dots, z_l) = \prod_{i=1}^l z_i$). Define $\mathcal{H} \triangleq \{f(h_1(x), \dots, h_l(x)) : h_1 \in \mathcal{H}_1, \dots, h_l \in \mathcal{H}_l\}$. Show that there exists some constant $c > 0$ such that

$$\text{VC}(\mathcal{H}) \leq c \left(\sum_{i=1}^l d_i \right) \ln \left(\sum_{i=1}^l d_i \right).$$

Problem 3

In this exercise, we will show that, under the *realizable setting*, with hypothesis class \mathcal{H} having VC dimension d , ERM (in fact, the consistency algorithm) will have a PAC sample complexity of $O\left(\frac{1}{\epsilon}(d \ln \frac{1}{\epsilon} + \ln \frac{1}{\delta})\right)$. Suppose $S = \{Z_1, \dots, Z_m\}$ a set of m training examples drawn iid from distribution D , where each $Z_i = (X_i, Y_i)$ is a labeled example. In addition, $\mathcal{F} = \{\mathbf{1}(h(x) \neq y) : h \in \mathcal{H}\}$ is the zero-one loss function class. Our proof will mostly follow the steps for showing agnostic PAC sample complexity given in the lecture.

1. **Double Sampling Trick.** Fix a training set S . Suppose $\mathbb{E}_S f(Z) = 0$ and $\mathbb{E}_D f(Z) \geq \epsilon$. Show that for a fresh set of random examples S' of size m ($m \geq \frac{16}{\epsilon}$) sampled iid from D :

$$\mathbb{P}_{S' \sim D^m} \left(\mathbb{E}_{S'} f(Z) \geq \frac{\epsilon}{2} \right) \geq \frac{1}{2}.$$

2. **Conditioning.** Denote by events

$$E' \triangleq \left\{ \text{there exists } f \in \mathcal{F}, \mathbb{E}_S f(Z) = 0, \mathbb{E}_{S'} f(Z) \geq \frac{\epsilon}{2} \right\},$$

$$E \triangleq \left\{ \text{there exists } f \in \mathcal{F}, \mathbb{E}_S f(Z) = 0, \mathbb{E}_D f(Z) \geq \epsilon \right\}.$$

Show $\mathbb{P}_{S, S' \sim D^m}(E' | E) \geq \frac{1}{2}$, and conclude that $\mathbb{P}_{S \sim D^m}(E) \leq 2\mathbb{P}_{S, S' \sim D^m}(E')$.

3. **Symmetrization.** Introduce $\sigma = (\sigma_1, \dots, \sigma_m)$ where each $\sigma_i \in \{\pm 1\}$. Denote by

$$(W_i, W'_i) = \begin{cases} (Z_i, Z'_i) & \sigma_i = +1, \\ (Z'_i, Z_i) & \sigma_i = -1. \end{cases}$$

Show that

$$\mathbb{P}_{S, S' \sim D^m}(E') = \mathbb{P}_{S, S' \sim D^m, \sigma \sim R^m} \left(\text{exists } f \in \mathcal{F}, \sum_{i=1}^m f(W_i) = 0, \sum_{i=1}^m f(W'_i) \geq \frac{m\epsilon}{2} \right),$$

where R is the Rademacher distribution, i.e. uniform in $\{\pm 1\}$.

4. **The randomness in Rademacher random variables.** Fix two size m training sets S and S' . Show that for a fixed classifier f in \mathcal{F} ,

$$\mathbb{P}_{\sigma \sim R^n} \left(\sum_{i=1}^m f(W_i) = 0, \sum_{i=1}^m f(W'_i) \geq \frac{m\epsilon}{2} \right) \leq \exp\left(-\frac{m\epsilon}{4}\right).$$

5. Use the above items to conclude that for $m \geq \frac{16}{\epsilon}$,

$$\mathbb{P}_{S \sim D^m}(\text{there exists } f \in \mathcal{F}, \mathbb{E}_S f(Z) = 0, \mathbb{E}_D f(Z) \geq \epsilon) \leq 2\mathcal{S}(\mathcal{F}, 2m) \exp\left\{-\frac{m\epsilon}{4}\right\}.$$

In addition, show that ERM has a PAC sample complexity of $O\left(\frac{1}{\epsilon}(d \ln \frac{1}{\epsilon} + \ln \frac{1}{\delta})\right)$.

Problem 4

In this exercise, we develop sample complexity lower bounds for *realizable* PAC learning using Le Cam's method and Assouad's method. Suppose hypothesis class \mathcal{H} has VC dimension $d \geq 2$, and it shatters examples z_0, z_1, \dots, z_{d-1} . In addition, suppose $\epsilon, \delta \in (0, \frac{1}{8})$ are target error and target failure probability. A learning algorithm \mathcal{A} is a mapping from training set S to $\{\pm 1\}$. In the following, you can use the elementary fact that for $x \in (0, \frac{1}{2})$, $e^{-x} \geq 1 - x \geq e^{-2x}$.

1. Consider D_{-1} and D_{+1} as follows: for every i in $\{\pm 1\}$,

$$D_i(x, y) = \begin{cases} 1 - 2\epsilon, & (x, y) = (z_0, -1), \\ 2\epsilon, & (x, y) = (z_1, i), \\ 0 & \text{otherwise.} \end{cases}$$

Note that $\min_{h \in \mathcal{H}} \text{err}(h', D_i) = 0$ for both $i \in \{\pm 1\}$. For every j in $\{\pm 1\}$, denote by $P_j((x_i, y_i)_{i=1}^m) = \prod_{i=1}^m D_j(x_i, y_i)$ the distribution over training sets (observations). Use Le Cam's method to show that for any hypothesis tester f , there exists an i in $\{\pm 1\}$, such that

$$\mathbb{P}_i(f(S) \neq i) > \frac{1}{2}(1 - 4\epsilon)^m.$$

2. Conclude that for any learning algorithm \mathcal{A} , if sample size $m \leq \frac{1}{4\epsilon} \ln \frac{1}{4\delta}$, then there exists an i in $\{\pm 1\}$,

$$\mathbb{P}_i(\text{err}(\hat{h}, D_i) > \epsilon) > \delta.$$

3. For every $\tau \in \{\pm 1\}^{d-1}$, consider D_τ as follows:

$$D_\tau(x, y) = \begin{cases} 1 - 4\epsilon, & (x, y) = (z_0, -1), \\ \frac{4\epsilon}{d-1}, & (x, y) = (z_i, \tau_i) \text{ for some } i \in \{1, \dots, d-1\}, \\ 0 & \text{otherwise.} \end{cases}$$

Note that $\min_{h \in \mathcal{H}} \text{err}(h', D_\tau) = 0$ for all $\tau \in \{\pm 1\}^{d-1}$. For every τ , denote by $P_\tau((x_i, y_i)_{i=1}^m) = \prod_{i=1}^m D_\tau(x_i, y_i)$ the distribution over training sets (observations).

Use Assouad's method to show that for any hypothesis tester f_1, \dots, f_{d-1} , there exists $\tau \in \{\pm 1\}^{d-1}$,

$$\mathbb{E}_\tau \left[\sum_{j=1}^{d-1} \mathbf{1}(f_j(S) \neq \tau_j) \right] > \frac{d-1}{2} \left(1 - \frac{4\epsilon}{d-1} \right)^m.$$

4. Conclude that for any learning algorithm \mathcal{A} , suppose that sample size $m \leq \frac{d-1}{128\epsilon}$, then there exists a $\tau \in \{\pm 1\}^d$, such that

$$\mathbb{P}_\tau(\text{err}(\hat{h}, D_\tau) > \epsilon) > \frac{1}{4}.$$

Hints

2.1 Use the elementary fact that $\ln(z) \leq z - 1$ for $z = \frac{x}{2a}$.

2.2 Use the elementary fact that $\binom{n}{i} \leq n^i$.

- 2.3 (1) consider S of size n shattered by \mathcal{H} . We know that $|\Pi_{\mathcal{H}}(S)| = 2^n$. Use Sauer's Lemma to obtain an upper bound on $|\Pi_{\mathcal{H}}(S)|$ in terms of v . (2) consider using the contrapositive of item 1.
- 2.4 Write \mathcal{H} as a union of $\binom{d}{k}$ hypothesis classes, each of which has VC dimension k , then apply item 3.)
- 3.1 Use Chernoff bound for Bernoulli distributions (the version with exponent $-\frac{m p \mu^2}{4}$).
- 3.4 Consider three cases: (1) there exists some i , $(f(Z_i), f(Z'_i)) = (1, 1)$; (2) $\sum_{i=1}^m f(Z_i) + f(Z'_i) < \frac{m\epsilon}{2}$; (3) both (1) and (2) are not satisfied. Observe that in the first two cases, the probability is identically zero.
- 4.1 Consider observation $S = ((z_0, -1), \dots, (z_0, -1))$. Show that $\mathbb{P}_{-1}(S) = \mathbb{P}_{+1}(S)$.
- 4.2 Define an appropriate hypothesis tester f that depends on \mathcal{A} .
- 4.3 Define $A_j = \{S = (x_i, y_i)_{i=1}^m : x_i \neq z_j \text{ for all } i\}$. Show that for every $\tau \stackrel{j}{\sim} \tau'$, $\mathbb{P}_{\tau}(S) = \mathbb{P}_{\tau'}(S)$ for all S in A_j . In addition, for $\sigma \in \{\tau, \tau'\}$, $\mathbb{P}_{\sigma}(S \in A_j) > \frac{7}{8}$. Intuitively, seeing only examples other than z_j does not help determining the optimal classifier's labeling on z_j .
- 4.4 First show that $\sum_{j=1}^{d-1} \mathbf{1}(f_j(S) \neq \tau_j) > \frac{d-1}{2}$ with probability $> \frac{1}{4}$. Then define an appropriate hypothesis tester $f = (f_1, \dots, f_{d-1})$ that depends on \mathcal{A} .

Problem 5 (No need to submit)

In this problem, we develop an alternative proof of Sauer's Lemma: any hypothesis class \mathcal{H} with VC dimension d can have at most $\binom{n}{\leq d}$ labelings on any dataset $S = \{z_1, \dots, z_n\}$. Throughout, we will be using the notation that

$$\binom{\{1, \dots, n\}}{d+1} \triangleq \{(i_1, \dots, i_{d+1}) : 1 \leq i_1 < \dots < i_{d+1} \leq n\}$$

to denote the set of $(d+1)$ -tuples whose entries are distinct. Note that $\left| \binom{\{1, \dots, n\}}{d+1} \right| = \binom{n}{d+1}$.

1. Show that for any indices $I = (i_1, \dots, i_{d+1}) \in \binom{\{1, \dots, n\}}{d+1}$, there exists a string $s_I \in \{\pm 1\}^{d+1}$, such that none of the labelings in

$$L_I = \{b \in \{\pm 1\}^n : (b_{i_1}, \dots, b_{i_{d+1}}) = s_I\}$$

are achievable by classifiers in \mathcal{H} .

2. Show the following basic facts:

- (a) For a finite set A and a function f , denote by $f(A) = \{f(a) : a \in A\}$. Then $|f(A)| \leq |A|$, where $|B|$ denotes the cardinality of set B .
- (b) Suppose \mathcal{I} is a set of indices. Given a collection of sets $\{L_I\}_{I \in \mathcal{I}}$ and a function f ,

$$\left| \bigcup_{I \in \mathcal{I}} f(L_I) \right| \leq \left| \bigcup_{I \in \mathcal{I}} L_I \right|. \quad (1)$$

3. Use the above two facts to conclude that

$$\left| \bigcup_{I \in \binom{\{1, \dots, n\}}{d+1}} L_I \right| \geq \sum_{i=d+1}^n \binom{n}{i}.$$

(Hint: consider functions f_1, \dots, f_n , where $f_i(s_1, \dots, s_n) = (s_1, \dots, s_{i-1}, -1, s_{i+1}, \dots, s_n)$ is the function that sets a length n string's i -th entry to -1 . Iteratively applying Equation (1) for f_1, \dots, f_n , what do you get?)

4. Use item 3 to conclude that $|\Pi_{\mathcal{H}}(S)| \leq \binom{n}{\leq d}$.