

### University College London

DEPARTMENT OF ELECTRONIC AND ELECTRICAL ENGINEERING

# Performance characterisation of 8-bit RISC and OISC architectures

Author:
Mindaugas

JARMOLOVIČIUS
zceemja@ucl.ac.uk

Supervisor:
Prof. Robert
KILLEY
r.killey@ucl.ac.uk

Second Assessor:
Dr. Ed
ROMANS
e.romans@ucl.ac.uk

A BEng Project Final Report

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#### 1 Abstract

- 2 Introduction
- 3 Goals and Objective
- Theory and Analytical Full list of RISC instructions are listed in 4 Bases

### Decided design criteria:

- Minimal instruction size
- Minimalistic design

#### 5 Technical Method

This section describes methods and design choices used to construct two processors.

#### Machine Code 5.1

#### 5.1.1 RISC

As the aim of instruction size to be as minimal as possible, RISC instruction decided to be 8bits with optional additional immediate value from 1 to 3 bytes. Immediate values are explained in section 5.4.

Decision was made to have instruction compose of operation code two operands - source/destination and source, which is similar to x86 architecture rather than MIPS. Three possible combinations of register address sizes are possible in such case from one to three bits. Two was selected as it allow having four general purpose registers which is sufficient for most applications, and allow four bits for operation code - allowing up to 16 instructions.

Due to small amount of available operation codes and not all instructions requiring two operands (for example JUMP instruction may not need any operands or could use one operand to have address offset), other two type instructions are added to the design - with one and zero operands. See figure 5.1.1. This enabled processor to have 45 different instructions while maintaining minimal instruction size. Final design has:

- 2-operand instructions
- 1-operand instructions
- 0-operand instructions

table ?? in Appendix section.

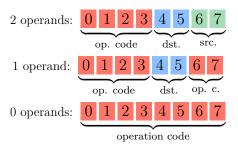


Figure 5.1.1: RISC instructions composition. Number inside box represents bit index. Destination (dst.) bits represents of source and destination register address.

#### 5.1.2OISC

As OISC requires only a single instruction, composition of instruction mainly requires two parts - source and destination. To allow higher instruction flexibility a immediate bit has been added to replace source address by immediate value. Composition of finalised machine code is shown in figure 5.1.2.

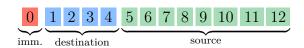


Figure 5.1.2: OISC instruction composition. Number inside box represents bit index.

Decision was made to have source address to be eight bits to allow it be replaced with immediate value. Destination address was chosen to be as minimal as possible, leaving only four bits or 16 possible destinations. Final design has 15 destination and 41 source addresses. This is not the most space efficient design as 41 source addresses would require only six bits for address, wasting two bits every time non-immediate source is used.

Full list of OISC sources and destinations are listed in table 2 in Appendix section.

## 5.2 Arithmetic Logic Unit

# 5.3 Memory

This section describes how instruction memory (ROM) is implemented for both processors.

### 5.4 Instruction decoding

This section describes RISC and OISC differences between instruction decoding and immediate value handling.

### 5.4.1 RISC

Already described in previous section 5.3, instruction from memory comes as 4 bytes. Least significant byte is sent to control block, other three bytes are sent to immediate override block (IMO) shown in figure 5.4.1. These three bytes are labelled as **immr**.

IMO block is a solution to change immediate value which enabled dynamically calculated memory pointers, branches dependant on register value or any other function that needs instruction immediate value been replaced by calculated register value. IMO is controlled by control block and cdi.imoctl signal, which is changed by CIO, CI1 and CI2 instructions. When signal is 0h, this block is transparent connecting immr directly to imm. When any of CI in-

structions executed, one of IMO register is overridden by reg1 value from register file. In order to override two or three bytes of immediate, CI instructions need to be executed in order. Only for one next instruction after last CI will have immediate bytes changed depending on what are values in IMO registers.

This circuit has two disadvantages:

- 1. Overriding immediate bytes takes one or more clock cycles,
- 2. At override, **immr** bytes are ignored therefore they are wasting instruction memory space.

Point 2 can be resolved by designing a circuit that would subtract the amount of overridden IMO bytes from pc\_off signal (program counter offset that is dependant on i-size value) at the program counter, thus effectively saving instruction memory space. This solution however would introduce a complication with the assembler as additional checks would need to be done during compiling to check if IMO instruction are used.

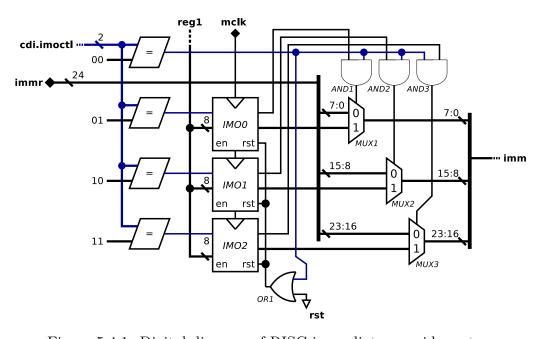


Figure 5.4.1: Digital diagram of RISC immediate override system

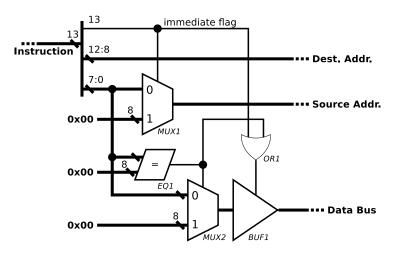


Figure 5.4.2: Digital diagram of OISC instruction decoder

### 5.4.2 OISC

OISC immediate value is set in instruction decoder shown in figure 5.4.2. Decoder operation is simple - instruction machine code is split into three parts as described in 5.1.2. If instruction source address is 00h, connect data bus with constant 0 via MUX2. If immediate bit is 1, set source address to 00h (to make sure no other buffer source connects to data bus), and connect instruction source address (immediate value) to databus via MUX2 and BUF1.

### 6 Results and Analysis

### 6.1 Benchmark Programs

### 6.1.1 Number of instructions

### 6.1.2 Instruction composition

Function composition was executed with following code:

Listing 1: RISC assembly frame for executring tests

Listing 2: OISC assembly frame for executring tests

## 6.2 Maximum clock frequency

6.3

## 7 Conclusion

## 8 References

# 9 Appendix

## 9.1 Processor instruction set tables

Table 1: Instruction set for RISC processor. \* Required immediate size in bytes

Instr.	Description	I-size *		
	2 register instructions			
MOVE	Copy value from one register to other	0		
ADD	Arithmetical addition	0		
SUB	Arithmetical subtraction	0		
AND	Logical AND	0		
OR	Logical OR	0		
XOR	Logical XOR	0		
MUL	Arithmetical multiplication	0		
DIV	Arithmetical division (inc. modulus)	0		
	1 register instructions			
COPY0	Copy intimidate to a register 0	1		
COPY1	Copy intimidate to a register 1	1		
COPY2	Copy intimidate to a register 2	1		
COPY3	Copy intimidate to a register 3	1		
ADDC	Arithmetical addition with carry bit	0		
ADDI	Arithmetical addition with immediate	1		
SUBC	Arithmetical subtraction with carry bit	0		
SUBI	Arithmetical subtraction with immediate	1		
ANDI	Logical AND with immediate	1		
ORI	Logical OR with immediate	1		
XORI	Logical XOR with immediate	1		
CI0	Replace intimidate value byte 0 for next instruction	1		
CI1	Replace intimidate value byte 1 for next instruction	1		
CI2	Replace intimidate value byte 2 for next instruction	1		
SLL	Shift left logical	1		
SRL	Shift right logical	1		
SRA	Shift right arithmetical	1		
LWHI	Load word (high byte)	3		
SWHI	Store word (high byte, reg. only)	0		
LWLO	Load word (low byte)	3		
SWLO	Store word (low byte, stores high byte reg.)	3		
INC	Increase by 1	0		
DEC	Decrease by 1	0		
GETAH	Get ALU high byte reg. (only for MUL & DIV & ROL &	0		
	ROR)			
GETIF	Get interrupt flags	0		
PUSH	Push to stack	0		
POP	Pop from stack	0		
COM	Send/Receive to/from com. block	1		
BEQ	Branch on equal	3		
BGT	Branch on greater than	3		

Table 1: Instruction set for RISC processor. \* Required immediate size in bytes

Instr.	Description	I-size *
BGE	Branch on greater equal than	3
BZ	Branch on zero	2
0 register instructions		
CALL	Call function, put return to stack	2
RET	Return from function	0
JUMP	Jump to address	2
RETI	Return from interrupt	0
INTRE	Set interrupt entry pointer	2

Table 2: Instructions for OISC processor.

Name	Description		
Destination Addresses			
ACC0	Set ALU source A accumulator		
ACC1	Set ALU source B accumulator		
BR0	Set Branch pointer register (low byte)		
BR1	Set Branch pointer register (high byte)		
BRZ	If source value is 0, set program counter to branch pointer		
STACK	Push value to stack		
MEM0	Set Memory pointer register (low byte)		
MEM1	Set Memory pointer register (middle byte)		
MEM2	Set Memory pointer register (high byte)		
MEMHI	Save high byte to memory at memory pointer		
MEMLO	Save low byte to memory at memory pointer		
COMA	Set communication block address register		
COMD	Send value to communication block		
REG0	Set general purpose register 0		
REG1	set general purpose register 1		
	Source Addresses		
NULL	Get constant 0		
ALU0	Get value at ALU source A accumulator		
ALU1	Get value at ALU source B accumulator		
ADD	Get Arithmetical addition of ALU sources		
ADDC	Get Arithmetical addition carry		
ADC	Get Arithmetical addition of ALU sources and carry		
SUB	Get Arithmetical subtraction of ALU sources		
SUBC	Get Arithmetical subtraction carry		
SBC	Get Arithmetical subtraction of ALU sources and carry		
AND	Get Logical AND of ALU sources		
OR	Get Logical OR of ALU sources		
XOR	Get Logical XOR of ALU sources		
SLL	Get ALU source A shifted left by source B		
SRL	Get ALU source A shifted right by source B		
ROL	Get rolled off value from previous SLL instance		
ROR	Get rolled off value from previous SRL instance		

Table 2: Instructions for OISC processor.

Name	Description
MULLO	Get Arithmetical multiplication of ALU sources (low byte)
MULHI	Get Arithmetical multiplication of ALU sources (high byte)
DIV	Get Arithmetical division of ALU sources
MOD	Get Arithmetical modulus of ALU sources
EQ	Check if ALU source A is equal to source B
GT	Check if ALU source A is greater than source B
GE	Check if ALU source A is greater or equal to source B
NE	Check if ALU source A is not equal to source B
LT	Check if ALU source A is less than source B
LE	Check if ALU source A is less or equal to to source B
BR0	Get Branch pointer register value (low byte)
BR1	Get Branch pointer register value (high byte)
PC0	Get Program counter value (low byte)
PC1	Get Program counter value (high byte)
MEM0	Get Memory pointer register value (low byte)
MEM1	Get Memory pointer register value (middle byte)
MEM2	Get Memory pointer register value (high byte)
MEMHI	Load high byte from memory at memory pointer
MEMLO	Load low byte from memory at memory pointer
STACK	Pop value from stack
ST0	Get stack address value (low byte)
ST1	Get stack address value (high byte)
COMA	Get communication block address register value
COMD	Read value from communication block
REG0	Get value from general purpose register 0
REG1	Get value from general purpose register 1