LAGRANGIAN RELAXATION WITH GAMS

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Abstract. This document describes an implementation of Lagrangian Relaxation using GAMS.

1. Introduction

Lagrangian Relaxation techniques [2, 3] form an important and popular tool in discrete optimization. We will show how Lagrangian Relaxation with subgradient optimization can be implemented in a GAMS environment.

2. Lagrangian Relaxation

We consider the Mixed Integer Programming model:

MIP minimize
$$z = c^T x$$

$$Ax \ge b$$

$$Bx \ge d$$

$$x \ge 0$$

$$x_j \in \{0, 1, \dots, n\} \text{ for } j \in J$$

There are two sets of linear constraints. We assume the set $Ax \geq b$ are the complicating constraints: if we relax the problem by removing these constraints, the remaining problem

is relatively easy to solve.

We can form the Lagrangian Dual:

(2)
$$L(u) = \min c^{T} x + u^{T} (b - Ax)$$
$$Bx \ge d$$
$$x \ge 0$$
$$x_{j} \in \{0, 1, \dots, n\} \text{ for } j \in J$$

For any $u \ge 0$, L(u) forms a lower bound on problem MIP, as $u^T(b-Ax) \le 0$. I.e. we have $L(u) \le z$.

The task is to find

$$\max_{u \ge 0} L(u)$$

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which can provide a better bound than a linear programming relaxation. I.e.

$$(4) z_{LP} \le \max_{u \ge 0} L(u) \le z$$

where z_{LP} is the optimal objective of the linear programming relaxation.

3. Subgradient Optimization

It can be shown that L(u) is a piecewise linear function. Solving (3) is therefore a nondifferentiable optimization problem. A successful technique for this problem is *Subgradient Optimization*.

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Following the notation in [5], the subgradient algorithm can be summarized as:
{Input}
An upper bound L^*
An initial value u^0 \ge 0
{Initialization}
\theta_0 := 2
{Subgradient iterations}
for j := 0, 1, ... do
   \gamma^j := g(x^j) {gradient of L(u^j)}
  \begin{aligned} &f_j := \theta_j(L^* - L(u^j))/||\gamma^j||^2 \text{ {step size}} \\ &u^{j+1} := \max\{0, u^j + t_j \gamma^j\} \\ &\text{ if } ||u^{j+1} - u^j|| < \varepsilon \text{ then} \end{aligned}
       Stop
   if no progress in more than K iterations then
       \theta_{j+1} := \theta_j/2
       \theta_{j+1} := \theta_j
   end if
   j := j + 1
end for
```

There are many variants possible with respect to the calculation of the step size and the updating of the parameter θ [1, 4].

4. Example

In the GAMS model below we illustrate the technique described above using a generalized assignment problem:

(5)
$$\min \sum_{i} \sum_{j} c_{i,j} x_{i,j}$$
$$\sum_{j} x_{i,j} = 1 \ \forall i$$
$$\sum_{i} a_{i,j} x_{i,j} \leq b_{j} \ \forall j$$

The assignment constraint $\sum_{j} x_{i,j} = 1$ will be dualized. The resulting obtained bound of 12.5 is tighter than the LP relaxation bound of 6.4. To be complete, the optimal MIP solution is 18.

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$ontext
  Lagrangian Relaxation
  using a Generalized Assignment Problem
  LP Relaxation : 6.4286
  Lagrangian Relaxation: 12.5
  Optimal Integer Solution: 18
  Erwin Kalvelagen, Amsterdam Optimization
  Reference:
    Richard Kipp Martin, Large Scale Linear and Integer Optimization,
Kluwer, 1999
$offtext
set i 'tasks' /i1*i3/;
set j 'servers' /j1*j2/;
parameter b(j) 'available resources' /
   j1 13
   j2 11
table c(i,j) 'cost coefficients'
      j1
           j2
i1
               2
        9
i2
i3
        3
               8
table a(i,j) 'resource usage'
       j1
             j2
i1
        6
               8
i2
               5
i3
               6
* standard MIP problem formulation
* solve as RMIP to get initial values for the duals
variables
  cost 'objective variable'
x(i,j) 'assignments'
binary variable x;
equations
  obj
                'objective'
   assign(i) 'assignment constraint'
   resource(j) 'resource limitation constraint'
                   cost =e= sum((i,j), c(i,j)*x(i,j));
obj..
assign(i)..
                   sum(j, x(i,j)) =e= 1;
sum(i, a(i,j)*x(i,j)) =l= b(j);
resource(j)..
option optcr=0;
model genassign /obj,assign,resource/;
solve genassign minimizing cost using rmip;
```

 $^{^{1} \}verb|http://www.amsterdamoptimization.com/models/lagrel.gms|$

```
* Let assign be the complicating constraint
parameter u(i);
variable bound;
equation LR 'lagrangian relaxation';
LR.. bound === sum((i,j), c(i,j)*x(i,j))
+ sum(i, u(i)*[1-sum(j,x(i,j))]);
model ldual /LR,resource/;
* subgradient iterations
set iter /iter1*iter50/;
scalar continue /1/;
parameter stepsize;
scalar theta /2/;
scalar noimprovement /0/;
scalar bestbound /-INF/;
parameter gamma(i);
scalar norm;
scalar upperbound;
parameter uprevious(i);
scalar deltau;
parameter results(iter,*);
* initialize u with relaxed duals
u(i) = assign.m(i);
display u;
* an upperbound on L
parameter initx(i,j) / i1.j1 1, i2.j2 1, i3.j2 1 /;
upperbound = sum[(i,j), c(i,j)*initx(i,j)];
display upperbound;
loop(iter$continue,
* solve the lagrangian dual problem
    option optcr=0;
    option limrow = 0;
option limcol = 0;
    ldual.solprint = 0;
    solve Idual minimizing bound using mip;
results(iter,'dual obj') = bound.1;
if (bound.1 > bestbound,
        bestbound = bound.1;
        display bestbound;
       noimprovement = 0;
       noimprovement = noimprovement + 1;
        if (noimprovement > 1,
           theta = theta/2;
           noimprovement = 0;
       );
    results(iter, 'noimprov') = noimprovement;
```

```
results(iter,'theta') = theta;
* calculate step size
    gamma(i) = 1-sum(j,x.l(i,j));
    norm = sum(i,sqr(gamma(i)));
    stepsize = theta*(upperbound-bound.1)/norm;
    results(iter,'norm') = norm;
    results(iter,'step') = stepsize;
* update duals u
    uprevious(i) = u(i);
    u(i) = max(0, u(i)+stepsize*gamma(i));
    display u;
* converged ?
   deltau = smax(i,abs(uprevious(i)-u(i)));
  results(iter,'deltau') = deltau;

if( deltau < 0.01,

    display "Converged";
      continue = 0;
);
display results;
```

References

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