# Primal and Dual Decomposition: Theory and Distributed Algorithms

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### Outline

- Distributed Algorithms
- Primal Decomposition
- Lagrange Dual Decomposition

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### Distributed Algorithms

In previous lectures, we have studied optimization algorithms like

- Descent algorithm, e.g., Gradient method, Newton method
- Interior point algorithm
- Cutting plane algorithm
- Fixed-point iteration algorithm, e.g., power method

Distributed algorithms are preferred because:

- It's scalable
- It's robust
- Centralized command is **not** feasible or is too costly

#### Key issues:

- Local computation vs. global communication
- Scope, scale, and physical meaning of communication overhead
- Theoretical issues: Convergence? Optimality? Speed?
- Practical issues: Robustness? Synchronization? Complexity? Stability?
- Problem separability structure for decomposition: vertical and horizontal

### Decomposition: LP Example

LP with variables u, v:

maximize 
$$c_1^T u + c_2^T v$$
  
subject to  $A_1 u \leq b_1$   
 $A_2 v \leq b_2$   
 $F_1 u + F_2 v \leq h$ 

Coupling constraint:  $F_1u + F_2v \leq h$ . Otherwise, separable into two LP

## Primal Decomposition

Introduce variable z and rewrite coupling constraint as

$$F_1u \leq z, \quad F_2v \leq h-z$$

LP decomposed into a master problem and two subproblems:

$$minimize_z\phi_1(z) + \phi_2(z)$$

where

$$\phi_1(z) = \inf_{u} \{ c_1^T u | A_1 u \leq b_1, Fu \leq z \}$$

$$\phi_2(z) = \inf_{v} \{ c_2^T v | A_2 v \le b_2, F_2 v \le h - z \}$$

Subgradient of function  $f: \mathbf{R}^n \to \mathbf{R}$  at x is a vector g such that

$$f(y) \ge f(x) + g^T(y - x), \ \forall y$$

## Primal Decomposition

#### For each iteration *t*:

- 1. Solve two separate LPs to obtain optimal u(t), v(t) and associated dual variables  $\lambda_1(t), \lambda_2(t)$ 
  - 2. Subgradient update:  $g(t) = -\lambda_1(t) + \lambda_2(t)$
  - 3. Mater algorithm update:  $z(t+1) = z(t) \alpha(t)g(t)$  where  $\alpha(t) \geq 0$ ,  $\lim_{t \to \infty} \alpha_t = 0$  and  $\sum_{t=1}^{\infty} \alpha(t) = \infty$

#### Interpretation:

- ullet z fixes allocation of resources between two subproblems and master problem iteratively finds best allocation of resources
  - More of each resource is allocated to the subproblem with larger Lagrange multiplier at each step

## **Dual Decomposition**

#### Form partial Lagrangian:

$$L(u, v, \lambda) = c_1^T u + x_2^T v + \lambda^T (F_1 u + F_2 v - h)$$
  
=  $(F_1^T \lambda + c_1)^T u + (F_2^T + c_2)^T v - \lambda^T h$ 

#### **Dual function:**

$$q(\lambda) = \inf_{u,v} \{ L(u,v,\lambda) | A_1 u \leq b1, A_2 v \leq b_2 \}$$
  
=  $-\lambda^T h + \inf_{u:A_1 u \leq b_1} (F_1^T \lambda + c_1)^T u + \inf_{v:A_2 u \leq b_2} (F_2^T \lambda + c_2)^T v$ 

### Dual problem:

maximize 
$$q(\lambda)$$
 subject to  $\lambda \succeq 0$ 

### **Dual Decomposition**

Solve the following LP in u, with minimizer  $u^*(\lambda(t))$ 

minimize 
$$(F_1^T \lambda(t) + c_1)^T u$$
  
subject to  $A_1 u \leq b_1$ 

Solve the following LP in v, with minimizer  $v^*(\lambda(t))$ 

minimize 
$$(F_2^T \lambda(t) + c_2)^T v$$
  
subject to  $A_2 v \leq b_2$ 

Use the following subgradient (to -q) to update  $\lambda$ :

$$g(t) = -F_1 u^*(\lambda(t)) - F_2 v^*(\lambda(t)) + h, \quad \lambda(t+1) = [\lambda(t) - \alpha(t)g(t)]^+$$

#### Interpretation:

Master algorithm adjusts prices  $\lambda$ , which regulates the separate solutions of two subproblems

## Summary

- Decouple a coupling constraint: primal or dual decomposition
- Dual decomposition algorithm: a large-scale engineering application in the next lecture

### Reading assignment:

• Sections 3.4, 7.5 in D. P. Bertsekas and J. N. Tsitsiklis, <u>Parallel and Distributed</u> Computation: Numerical Methods, Athena Scientific 1999.