

## Relaxed plan

Set of actions to reach goal state

from current state in relaxed problem

**Definition (Relaxed Plan Heuristic).** A heuristic function is called a *relaxed plan heuristic*, denoted  $h^{FF}$ , if, given a state  $s$ , it returns  $\infty$  if no relaxed plan exists, and otherwise returns  $\sum_{a \in RPlan} c(a)$  where  $RPlan$  is the action set returned by relaxed plan extraction on a closed well-founded best-supporter function for  $s$ .

→ Recall: If a relaxed plan exists, then there also exists a closed well-founded best-supporter function, see previous slide.

## Best Supporter Functions

notation:  $bs(p)$  = best supporter for fact  $p$

Action that makes fact  $p$  true which has "lowest accumulated cost" in approximated relaxation problem.

i.e. action with lowest cost + lowest

cost of achieving prerequisites

can use  $h^{add}$  and  $h^{max}$  to find best supporters  
 $bs^{add}$   $bs^{max}$

$bs(p)$  is only defined for actions that are not true in the initial state and not defined for facts that are unreachable from current state in relaxed problem

**Definition (Best-Supporters from  $h^{max}$  and  $h^{add}$ ).** Let  $\Pi = (F, A, c, I, G)$  be a STRIPS planning task, and let  $s$  be a state.

The  $h^{max}$  supporter function  $bs_s^{max} : \{p \in F \mid 0 < h^{max}(s, \{p\}) < \infty\} \mapsto A$  is defined by  $bs_s^{max}(p) := \arg \min_{a \in A, p \in add_a} c(a) + h^{max}(s, pre_a)$ .

The  $h^{add}$  supporter function  $bs_s^{add} : \{p \in F \mid 0 < h^{add}(s, \{p\}) < \infty\} \mapsto A$  is defined by  $bs_s^{add}(p) := \arg \min_{a \in A, p \in add_a} c(a) + h^{add}(s, pre_a)$ .

Finding Best supporters using  $h^{add}/h^{max}$

While performing Bellman Ford, maintain array/dictionary

for storing BS for each action. Initialise each BS to NONE

- In each iteration, apply every possible action given preconditions. If the action updates the value of a fact  $p$  (makes true or reduces value), set  $BS(p) = \text{action}$

## Ideal Properties of $bs(p)$

**Closed:**  $bs$  defined for every fact  $p$

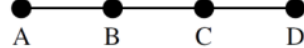
which may contribute to reaching goal  
(except those already true in initial state)

**Well-Sounded:** Paths using best supporters don't  
run into cycles.

**Support Graph:** Directed graph showing connections  
between facts and best supporter actions

Well founded  $\iff$  Support graph is acyclic

### Support Graphs and Prerequisite (C) in "Logistics"



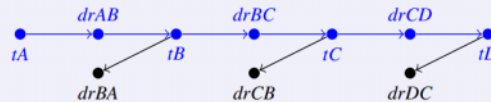
- Initial state:  $tA$ .
- Goal:  $tD$ .
- Actions:  $drXY$ .

#### How to do it (well-founded)

Best-supporter function:

$p$	$bs(p)$
$t(B)$	$dr(A, B)$
$t(C)$	$dr(B, C)$
$t(D)$	$dr(C, D)$

Yields support graph backchaining:



#### How to NOT do it (not well-founded)

Best-supporter function:

$p$	$bs(p)$
$t(B)$	$dr(C, B)$
$t(C)$	$dr(B, C)$
$t(D)$	$dr(C, D)$

Yields support graph backchaining:



Using  $h^{add}$  and  $h^{max}$  to derive best supporters guarantees  
closure and wellfoundedness

**Calculating  $h^{ss}$**

#### Relaxed Plan Extraction for state $s$ and best-supporter function $bs$

```

Open := G \ s; Closed := ∅; RPlan := ∅
while Open ≠ ∅ do:
  select g ∈ Open
  Open := Open \ {g}; Closed := Closed ∪ {g};
  RPlan := RPlan ∪ {bs(g)}; Open := Open ∪ (prebs(g) \ (s ∪ Closed))
endwhile
return RPlan

```

→ Starting with the top-level goals, iteratively close open singleton sub-goals by selecting the best supporter.

**This is fast!** Number of iterations bounded by  $|P|$ , each near-constant time.

- Similar to BFS / DFS
- Initialise open set with all goal facts not true

in initial state, closed set with all initial facts,  
and empty action set

• While open set is not empty:

pop any element from open set. Add to closed set

Add best supporter to action set. Add all preconditions

for action not in closed set to open set

$h^{ss} = \text{number of actions in action set (unique actions)}$

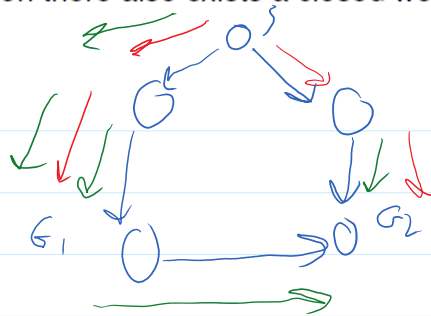
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$$h^{ss} \geq h^*$$

• May be inadmissible

$$h^{ss}(s) = \infty \iff h^*(s) = \infty$$



**Definition (Helpful Actions)** Let  $h^{FF}$  be a relaxed plan heuristic, let  $s$  be a state, and let *RPlan* be the action set returned by relaxed plan extraction on the closed well-founded best-supporter function for  $s$  which underlies  $h^{FF}$ . Then an action  $a$  *applicable* to  $s$  is called *helpful* if it is *contained in RPlan*.

#### Remarks

- Initially introduced in FF [Hoffmann and Nebel (2011)], restricting Enforced Hill-Climbing to use **only** the helpful actions
- Expanding only helpful actions does not guarantee completeness.
- Other planners use helpful actions as preferred operators, expanding first nodes resulting from helpful actions.

## WIDTH

Sort of measure of Problem Difficulty

• low width = easy

Single fact in goal set

or if multiple goals, easy to achieve one at a time

• High width = hard

Many many goal conditions

or goals are hard to achieve one at a time  
(solving one makes others harder)

## Novelty

**Key definition:** the **novelty**  $w(s)$  of a **state**  $s$  is the size of the smallest subset of atoms in  $s$  that is true for the first time in the search.

- e.g.  $w(s) = 1$  if there is **one** atom  $p \in s$  such that  $s$  is the first state that makes  $p$  true.
- Otherwise,  $w(s) = 2$  if there are **two** different atoms  $p, q \in s$  such that  $s$  is the first state that makes  $p \wedge q$  true.
- Otherwise,  $w(s) = 3$  if there are **three** different atoms...

Pseudocode:

- Initialise empty set to store seen facts (seen set)
- While performing search algorithm:
  - 1) Add PowerSet of facts to seen fact set  
i.e. if state  $= \{A, B, C\}$  then add:  
 $\{A, B, C, (A, B), (A, C), (B, C), (A, B, C)\}$
  - 2) if new single fact added, then novelty = 1
  - 3) else if new pair of facts added, then novelty = 2
  - 4) else if new triple, novelty = 3 etc
  - 5) If nothing new added, novelty = 1 + num facts in state

## Iterated Width (IW)

### Algorithm

- $IW(k)$  = breadth-first search that **prunes** newly generated states whose  $novelty(s) > k$ .
- $IW$  is a **sequence of calls**  $IW(k)$  for  $i = 0, 1, 2, \dots$  over problem  $P$  until problem solved or  $i$  exceeds number of variables in problem

### Properties

$IW(k)$  expands at most  $O(n^k)$  states, where  $n$  is the number of atoms.

- For  $IW(k)$ , don't need to calculate novelty if it is guaranteed to be bigger than  $k$ , can just prune

Dealing with multiple goal facts?

Use Serialised iterative width

- IW from initial state to first goal
- Then IW from first goal to next goal  
etc

