Lecture 13 Bin Packing

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December 7, 2022

Outline

- The First Fit Algorithm
- Karmarkar and Karp's Algorithm

The First Fit Algorithm

Problem Statement

1-d bin packing

- Input: n items $a_1, a_2, \cdots, a_n$, each wit size $\in (0, 1]$, and an infinite number of unit-size bins
- Output: A feasible packing with the least number of bins used
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- feasible packing: all items are packed and no bins accept items summing over 1

NP-complete to decide if two bins suffice to accommodate given items.

Asymptotic Analysis

- An algorithm A has asymptotic approximation ratio at most α if

$$A(I) \leq \alpha OPT(I) + k$$

for any instance I and some "constant" k.
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- No algorithm with absolute approximation ratio $< \frac{3}{2}$ unless P=NP
- Conventional packing rules: Next-Fit, First-Fit, Best-Fit, Harmonic Fit,...

Simple Packing Rules

Next Fit (NF)

- Pack items one by one and keep at a time one bin open.
- Pack the current item into the opened bin if it has enough space; otherwise close the bin and open a new one for the item.
- $NL(I) \leq 2OPT(I) - 1$.

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Any Fit

- Pack items one by one and never close a bin before the game is over.
- Never open a new bin as long as there exists an opened bin in which the current item fits.

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Specific Fit Rules

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Performance

- WF can work as badly as NF.
- The others have an asymptotic ratio of 1.7.

FF

More precisely, we can show that

$$\forall I,\ FF(I) \leq 1.7OPT(I) + \frac{4}{5}.$$

Analysis of the FF algorithm

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- $FF(I) \le w(I) + 4/5 \quad \Longleftarrow \quad \forall B_j, w(B_j) \ge 1$ in an "average" way

Weighting Function

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We call the items, respectively, tiny, small, medium and big based on the above classification.

Weights of Optimal Bins

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- If B_j^* does not contain any big item, $w(B_j^*) \leq 1.2 + 0.3 < 1.7$.
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- If B_j contains two medium items, $w(B_j) \geq \frac{6}{5} \cdot \frac{2}{3} + 0.2 = 1$.

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Sort the FF bins in the order that they appear. Remove all those bins having weights at least one. Now consider the left bins (keeping the order). Note that each of such bins

- does not contain a large item;
- contains at most one medium item;
- has a content less than $\frac{5}{6}$.

Further Observations

Let k be the number of bins left. Without loss of generality, assume $k \geq 2$.

♣ Only one item in a bin

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- All but at most one have a content larger than $2/3$;
- If there exists a bin whose content is less than $2/3$, it must be either the last (if B_p does not exist) or the second last bin B_q (if B_p exists).

Analysis

Put B_p aside. Consider any two adjacent bins B and C:

- Recall that $2/3 \leq c(B) < 5/6$. Suppose $c(B) = 5/6 - z, z \in (0, 1/6]$.

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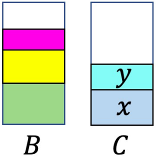

$$x > 1 - (\tfrac{5}{6} - z) = \tfrac{1}{6} + z \in (\tfrac{1}{6}, \tfrac{1}{3}], \;\; y > \tfrac{1}{6} + z \in (\tfrac{1}{6}, \tfrac{1}{3}]$$

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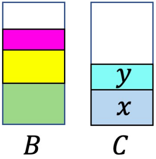

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$$v(x) > \tfrac{3}{5}(\tfrac{1}{6} + z - \tfrac{1}{6}) = \tfrac{3}{5}z, \quad v(y) > \tfrac{3}{5}z$$

- $\frac{6}{5}c(B) + v(C) \ge \frac{6}{5}(\frac{5}{6} - z) + v(x) + v(y)$

$$\ge \tfrac{6}{5}(\tfrac{5}{6} - z) + \tfrac{3}{5}z + \tfrac{3}{5}z = 1$$

Analysis

No matter if B_p and B_q exist, we apply the above analysis from the first bin to the second last bin.

- For $i = 1, 2, \ldots, k-2$, $\frac{6}{5}c(B_i) + v(B_{i+1}) \geq 1$
- $c(B_{k-1}) + c(B_k) > 1$
- $$\begin{aligned} \sum_{i=1}^{k} w(B_i) &\geq \sum_{i=1}^{k-2}(\frac{6}{5}c(B_i) + v(B_{i+1})) + \frac{6}{5}(c(B_{k-1}) + c(B_k)) \\ &\geq (k-2) + \frac{6}{5} = k - \frac{4}{5} \end{aligned}$$

Analysis

At this moment, we have shown

Summary

- Each optimum bin has a weight at most 1.7.
- With a total loss of 4/5, the average weight of FF bins is at least one.

which implies that

- $W(I) \leq 1.7OPT(I)$
- $W(I) \geq FF(I) - 4/5$

So, $FF(I) \leq 1.7OPT(I) + \frac{4}{5}$.

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So, $FF(I) \leq 1.7OPT(I) + \frac{4}{5}$.

A tight analysis shows $FF(I) \leq 1.7OPT(I)$ (Dosa and Sgall 2013).

Better Results

- First Fit Decreasing (FFD): apply First Fit after sorting the items in non-increasing order of their sizes

  .
- $FFD(I) \leq \frac{3}{2}OPT(I)$ (best possible assuming that $P \neq NP$).
- $FFD(I) \leq \frac{11}{9}OPT(I) + \frac{6}{9}$ (Dosa 2007).

Even Better Results

Bin packing can be approximated to

- $OPT + O(\log^2 OPT)$ (Karmarkar and Karp 1982)
- $OPT + O(\log OPT * \log\log OPT)$ (Rothvoss 2013)
- $OPT + O(\log OPT)$ (Rothvoss and Hoberg 2017)

Karmarkar and Karp's Algorithm

Packing Configurations

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- m : the number of different item sizes
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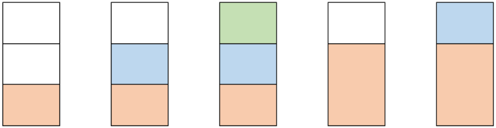
Configuration: a feasible packing into a bin



Configuration LP

N : the number of different configurations
t_{ij} : the number of items of size s_i in configuration j

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\begin{aligned}
\min \quad & \sum_{j=1}^{N} x_j & & \text{(Configuration LP)} \\
\text{s.t.} \quad & \sum_{j=1}^{N} t_{ij} x_j \geq b_i & i = 1, \ldots, m & \\
& x_j \geq 0 & j = 1, \ldots, N &
\end{aligned}
$$

Configuration LP

Theorem

The configuration LP can be solved within an additive error of at most 1 in time polynomial in m and $\log(n/s_m)$, where s_m is the size of smallest item in the instance.

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The configuration LP can be solved within an additive error of at most 1 in time polynomial in m and $\log(n/s_m)$, where s_m is the size of smallest item in the instance.

- Ensure that $s_m \geq 1/SIZE(I)$, where $SIZE(I) = \sum s_i$.
- Those items smaller than $1/SIZE(I)$ can be packed by FF or NF without increasing the additive error much.
- Therefore, the configuration LP can be solved in polynomial time.

Rounding the LP Solution

Rounding Scheme

There are at most m non-zero variables in an extreme point. Directly rounding up can only guarantee a feasible packing with $OPT(I) + m$ bins. To do better we round down the optimal LP solution x^*:

- Pack $\lfloor x_j^* \rfloor$ bins according to the configuration, for $j = 1, 2, \ldots, m$.
- Denote the set of items already packed as I_z.
- $SIZE(I - I_z) \leq m$.
- Recurse on the remaining items $I - I_z$.

Claim

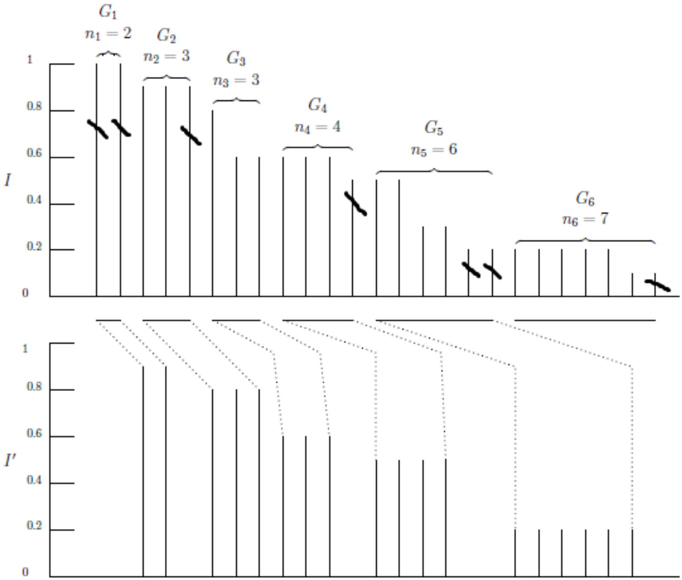
$LP(I - I_z) + LP(I_z) \leq LP(I)$.

Proof. $LP(I - I_z) + LP(I_z) \leq \sum_j (x_j^* - \lfloor x_j^* \rfloor) + \sum_j \lfloor x_j^* \rfloor = LP(I)$.

Harmonic Grouping

- Sort the items in non-increasing order.
- Following the order we form a group whenever the total size is at least 2, and start a new group with the next item.
- Let r be the number of groups, where G_i is the i-th group with n_i items, for $i = 1, 2, \ldots, r$.
- $n_i \geq n_{i-1}$, $i = 2, \ldots, r-1$
- Discard group G_1 and G_r.
- For $i = 2, \ldots, r-1$, discard $n_i - n_{i-1}$ smallest items in G_i, and round the remaining n_{i-1} items to the size of the largest item in G_i.
- The remaining items form a new instance I'.
- $LP(I) \geq LP(I')$.

Illustration of the Grouping



Key Issues

- $I \to I'$ by harmonic grouping;
- The number m of distinct item sizes in I' is at most $SIZE(I)/2$;
- The total size of all discarded items is $O(\log SIZE(I))$.

Algorithm

BIN PACK(I)

$k = 1$

if $SIZE(I) < 10$ **then**

  Pack remaining items using First Fit

else

  Apply **harmonic grouping scheme** to create instance I'; pack discarded items in $O(\log SIZE(I))$ bins using First Fit

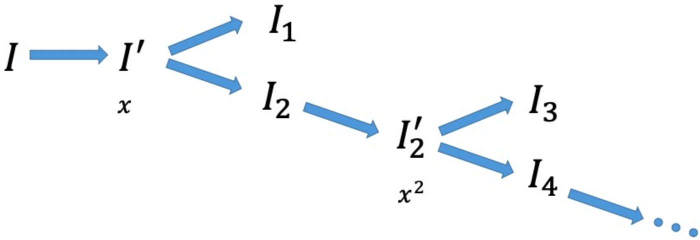
  Let x be optimal solution to configuration LP for instance I'

  Pack $\lfloor x_j \rfloor$ bins in configuration T_j for $j = 1, \cdots, N$; call the packed items instance I_{2k-1}

  Let I_{2k} be remaining items from I'

  Pack I_{2k} via BIN PACK(I_{2k}); $k = k + 1$

Algorithm



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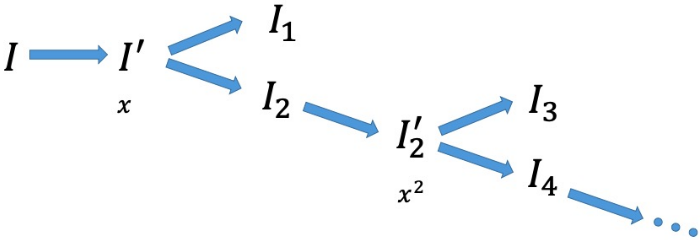

The harmonic grouping scheme can be used to design an approximation algorithm that always finds a packing with $OPT_{LP}(I) + O(\log^2(SIZE(I)))$ bins.

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- $OPT_{LP}(I_1) + OPT_{LP}(I_2) \leq OPT_{LP}(I') \leq OPT_{LP}(I)$;

Bounding the Discarded Items

- $I \to I'$ by harmonic grouping;
- The number of distinct item sizes in I' is at most $SIZE(I)/2$;
  (Each of these groups has size at least 2)
- The total size of all discarded items is $O(\log SIZE(I))$;
  (Each of these groups has size at most 3)
  $n_1 \leq n_2 \leq \cdots n_{r-1}$
  $r \leq \lceil \frac{SIZE(I)}{2} \rceil$
  $SIZE(G_1) + SIZE(G_r) \leq 6$

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  $n_1 \leq n_2 \leq \cdots n_{r-1}$
  $r \leq \lceil \frac{SIZE(I)}{2} \rceil$
  $SIZE(G_1) + SIZE(G_r) \leq 6$
  Average size of the discarded items in group G_i is at most $3/n_i$

Bounding the Discarded Items

- $I \rightarrow I'$ by harmonic grouping;
- The number of distinct item sizes in I' is at most $SIZE(I)/2$;
(Each of these groups has size at least 2)
- The total size of all discarded items is $O(\log SIZE(I))$;
(Each of these groups has size at most 3)
$n_1 \leq n_2 \leq \cdots n_{r-1}$
$r \leq \lceil \frac{SIZE(I)}{2} \rceil$
$SIZE(G_1) + SIZE(G_r) \leq 6$
Average size of the discarded items in group G_i is at most $3/n_i$

$SIZE(I_d) = SIZE(G_1) + SIZE(G_r) + \sum_{i=2}^{r-1}(n_i - n_{i-1})\frac{3}{n_i} \leq$
$6 + \sum_{i=1}^{r-1}\frac{3}{i} = O(\log SIZE(I))$