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IMPLEMENTATION OF COPY ON WRITE ON FORK IN XV6

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ABSTRACT

This project explores the integration of Copy-On-Write (COW) functionality into the fork in XV6, aiming to optimize management and memory resource The implementation enhances performance scalability, system empowering users to tailor XV6 to their unique requirements. Through meticulous exploration of kernel configuration options and system call modifications, the project understanding of elevates memory management techniques in operating system design.

INTRODUCTION

In operating systems, efficient memory management is paramount for optimizing system performance and resource utilization. However, traditional methods of memory allocation and process creation pose challenges in terms of overhead and redundancy.

During fork, a common issue is the need to duplicate the parent's memory space for the child, consuming both time and resources. Hence we use Copy-On-Write (COW), a technique devised to alleviate the inefficiencies of memory duplication during process creation.

XV6 RISC-V OPERATING SYSTEM

- Xv6 is a Unix-based operating system that provides students with a structured yet accessible environment to delve into the inner workings of operating systems.
- It offers an invaluable platform for comprehending critical OS components like process management, memory allocation, and file systems.
- The RISC-V iteration of xv6 signifies a notable adaptation, aligning this educational OS with the RISC-V instruction set architecture (ISA).
- XV6, regardless of its traditional or RISC-V variant, deliberately maintains a minimalist approach, eschewing many features found in commercial operating systems.

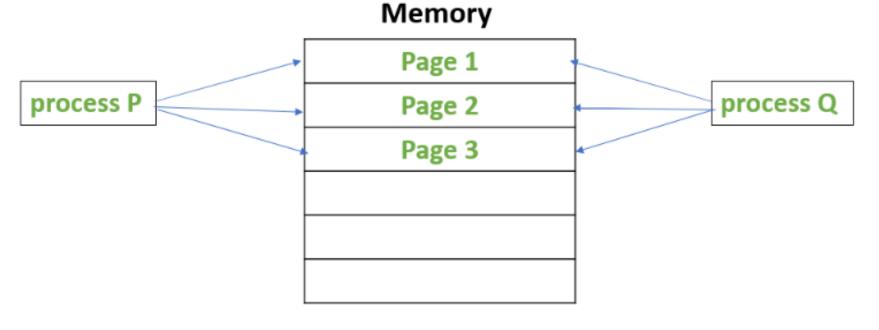
TOOLS/ API USED

- xv6 Operating System
- System Calls (fork(), sbrk(),brk())
- Memory Management (Page Allocation, Deallocation and Page Tables)
- Concurrency Control (Locks, Semaphores)
- Testing Framework (Tools and Scripts)
- QEMU

COPY ON WRITE

- 1. Copy-On-Write (COW), a technique devised to alleviate the inefficiencies of memory duplication during process creation.
- 2. Operates on the principle of deferred copying, postponing the actual duplication of memory pages until a write operation is attempted.
- 3. Rather than immediately duplicating memory pages upon fork, COW allows both processes to share the same memory pages initially.
- 4. When either process attempts to modify a ensuring data integrity while minimising unnecessary copying and resource

shared page, only then is a copy created, consumption



Before process P modifies Page 3

APPLICATIONS

Database cloning:

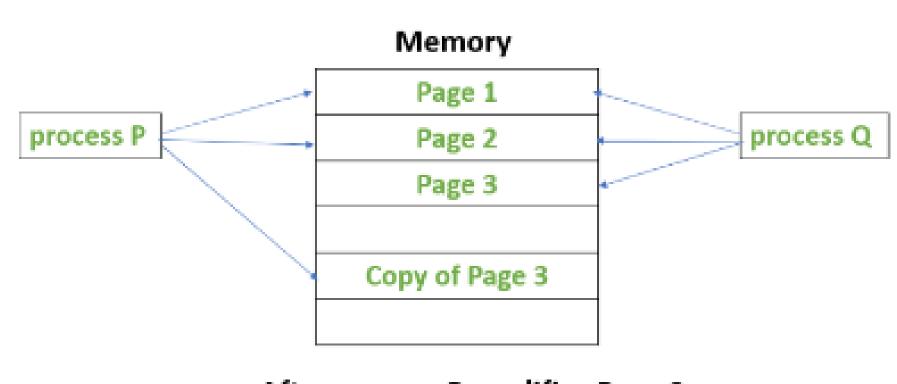
COW allows for creating database clones efficiently. Only data differences are copied, reducing cloning time and resource usage

Process Forking:

When a process creates a child process (e.g., using fork in Unix-like systems), both initially share the same memory pages using COW. This saves memory and reduces process creation time, as copying occurs only when necessary (on write).

Snapshots and Versioning:

File systems can leverage COW to create snapshots of directories or files. This allows for reverting to previous versions, recovering from accidental changes, or managing file versions efficiently.



After process P modifies Page 3

METHODOLOGY



MAP PARENT'S PAGES

Update uvmcopy() to map the parent's physical pages into the child instead of allocating new ones. Clear the writable bit (PTE_W) in the PTEs for both child and parent on shared pages.

*pte &= ~PTE W;

*pte |= PTE_COW;

pa = PTE2PA(*pte);

acquire(&reflock);

release(&reflock);

flags = PTE_FLAGS(*pte);

reference counter[pa/PGSIZE]+=1;

HANDLE COW PAGE FAULTS

Recognize page faults in usertrap(). On a write page fault for a COW page, allocate a new page, copy the content, and update the PTE with PTE_W set. Maintain read-only status for read only pages

IMPLEMENT REFERENCE COUNTING

Set reference count to one when allocating with kalloc(). Increment for shared pages during fork(), and decrement during page table drops. Modify kfree() to free only when the reference count is zero.

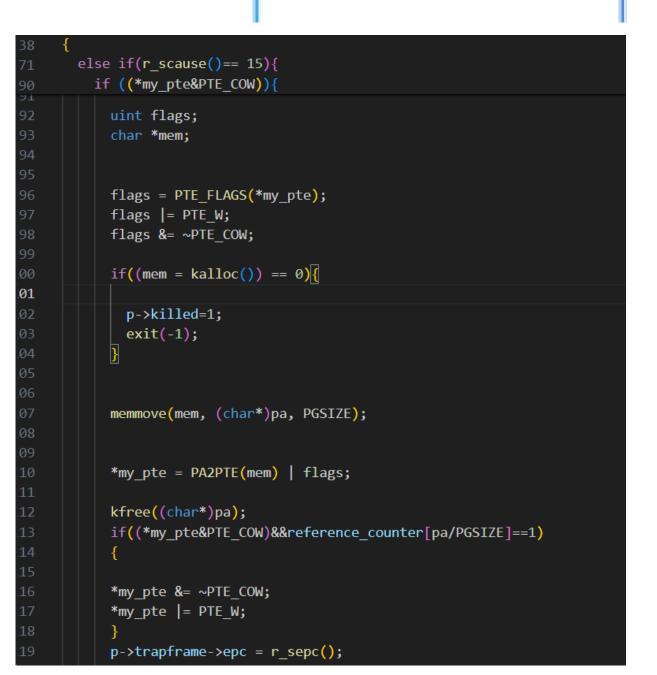
UPDATE COPYOUT()

Modify copyout() to

handle COW pages similarly to page faults. Allocate a new page, copy the content, and update the PTE with PTE_W if the original page was writable.

RECORD COW MAPPING IN PTE

Use RISC-V PTE's reserved bits to record COW mappings in each PTE. Manage COW mappings during page faults and copying out data.



OUTPUT

The successful execution of the cowtest confirms the correct implementation of Copy-On-Write (CoW) memory management in the xv6 operating system.

> xv6 kernel is booting hart 2 starting hart 1 starting init: starting sh \$ cowtest ok COW TEST PASSED

CONCLUSION

In conclusion, the introduction of Copy-on-Write (CoW) presents a demonstrably effective technique for optimising memory management.

This mechanism hinges on the principle of delaying memory duplication until a write attempt is detected. This approach not only minimises memory consumption during the initial fork event, but also extends these benefits throughout the processes' lifecycles.

The implementation of CoW in xv6 serves as a compelling illustration of the benefits leveraging architectural derived from as RISC-V's memory management capabilities, efficient process creation.

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