# Data Mining and Analysis

Graph Mining: 1

CSCE 676 :: Fall 2019

Texas A&M University

Department of Computer Science & Engineering

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#### Resources

Networks, Crowds, and Markets. Chapter 2 and Chapter 3

MMDS Chapter 10.1, 10.2, 10.4 (ignore 10.4.4)

Louvain method (Wikipedia)

DMTT Chapter 17.4

#### Agenda

Today

Basics, Frequent itemsets —> graph mining, Finding Important Nodes

Wednesday

Social Networks, Community Detection

Friday

**Community Detection** 

Later in the semester —> graph/node embeddings!

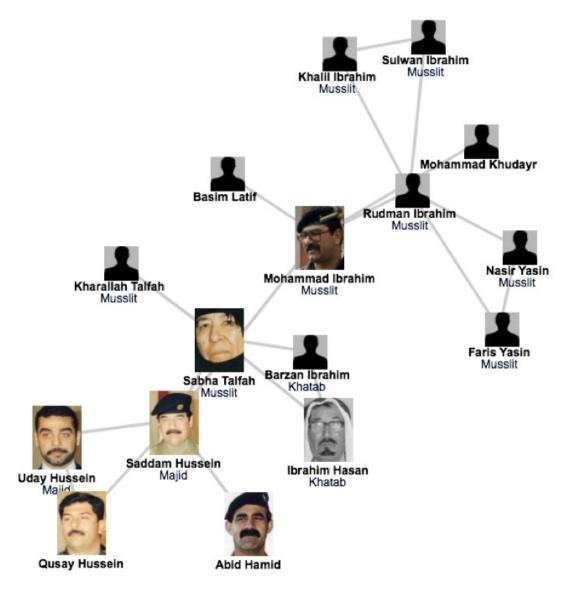
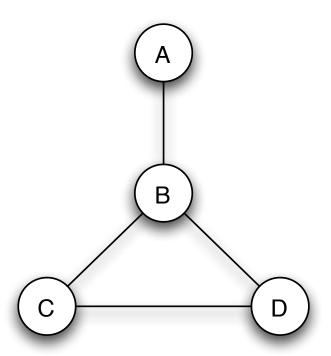


Image 1.2b
The network of Saddam Hussein.

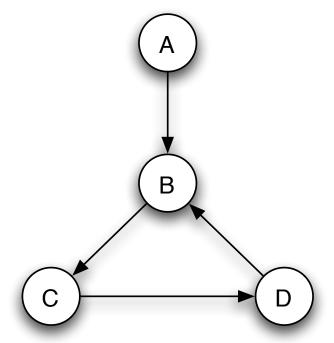
The Social Network. A small region of the social network reconstructed by the US forces in the process of searching for Saddam Hussein. The map represents the relationship between individuals in Saddam's inner circle.

#### Basic concepts

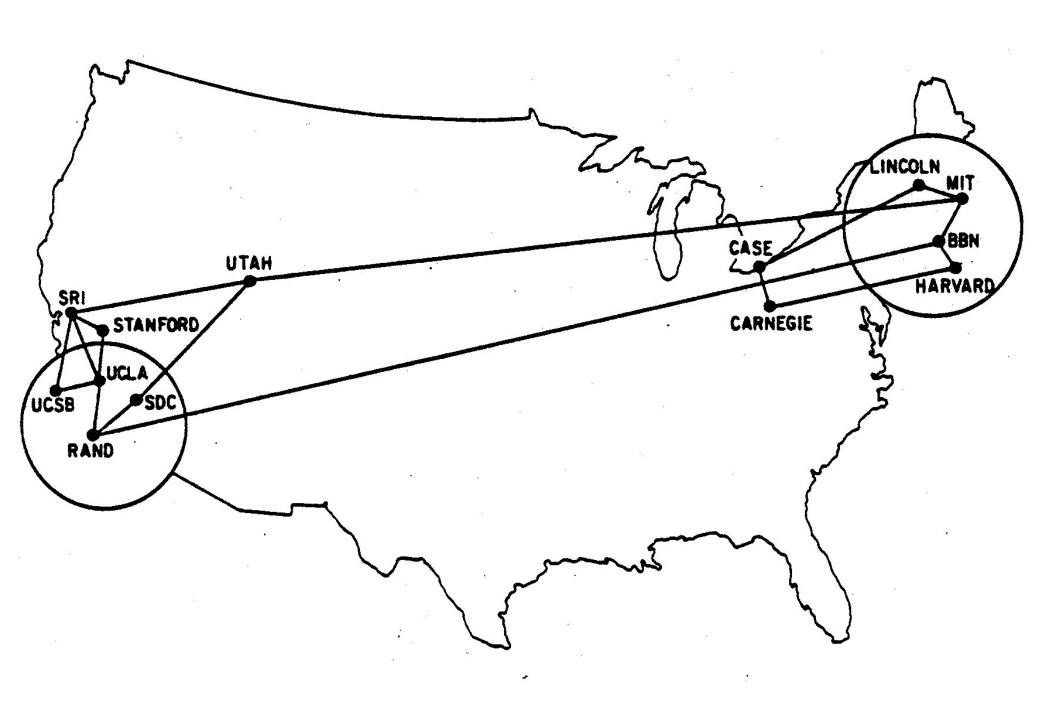
```
nodes
edges (directed or undirected)
paths
cycles
components (and giant
components)
```

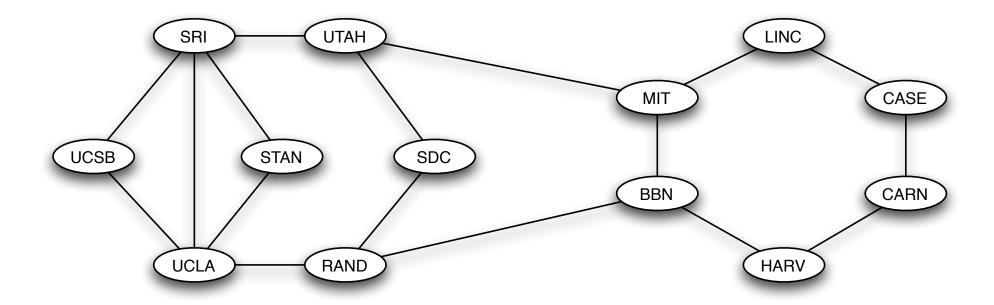


(a) A graph on 4 nodes.



(b) A directed graph on 4 nodes.





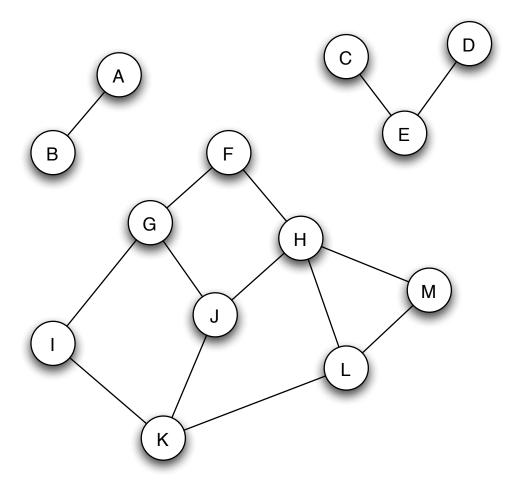


Figure 2.5: A graph with three connected components.

In practice, most graphs have a single giant component. But consider the Western hemisphere vs Europe (at the age of exploration) ... human diseases evolved independently, technology

#### **Network Datasets**

Collaboration graphs
Who-talks-to-whom graphs
Information linkage graphs
Technological networks
Networks in the natural world

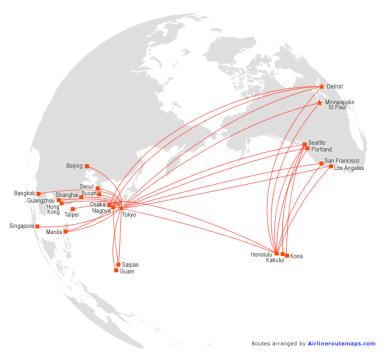
#### Some Network Resources

https://snap.stanford.edu/data/index.html

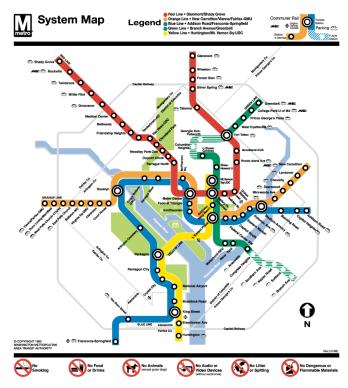
https://aws.amazon.com/datasets/ marvel-universe-social-graph/

http://www-personal.umich.edu/~mejn/netdata/

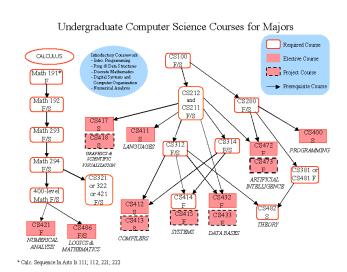
https://networkdata.ics.uci.edu/ index.html



(a) Airline routes



(b) Subway map



(c) Flowchart of college courses



(d) Tank Street Bridge in Brisbane

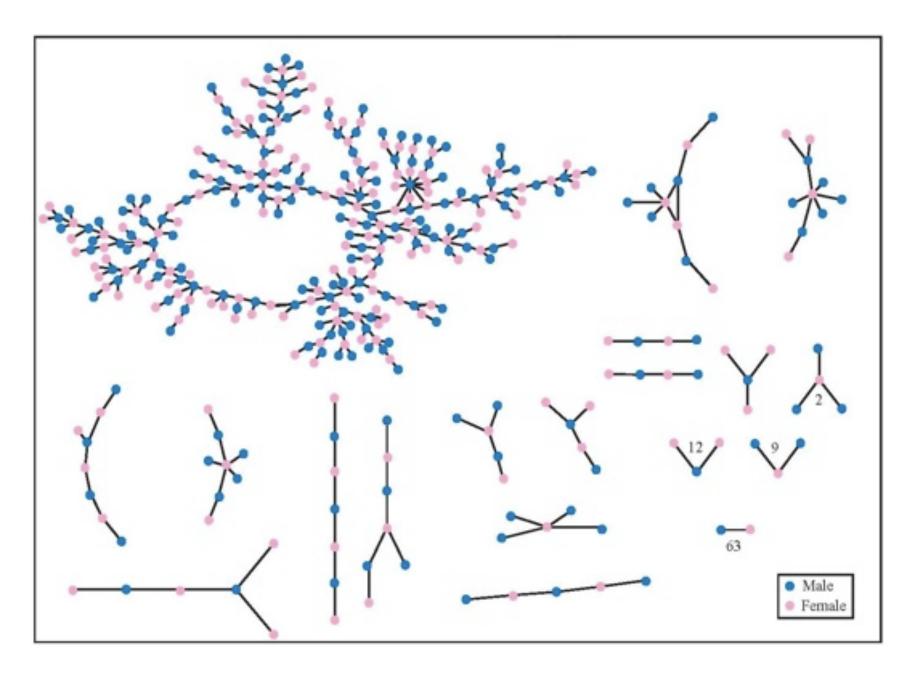
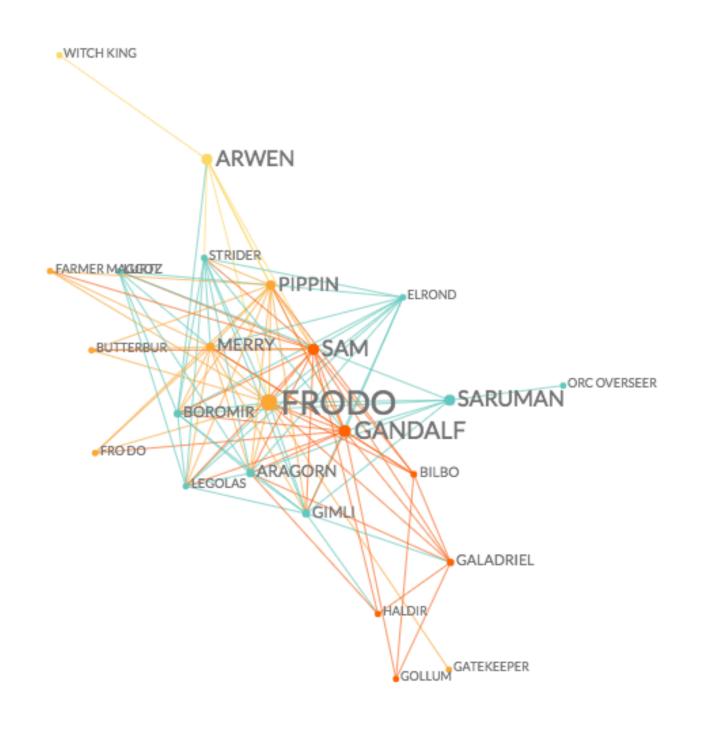
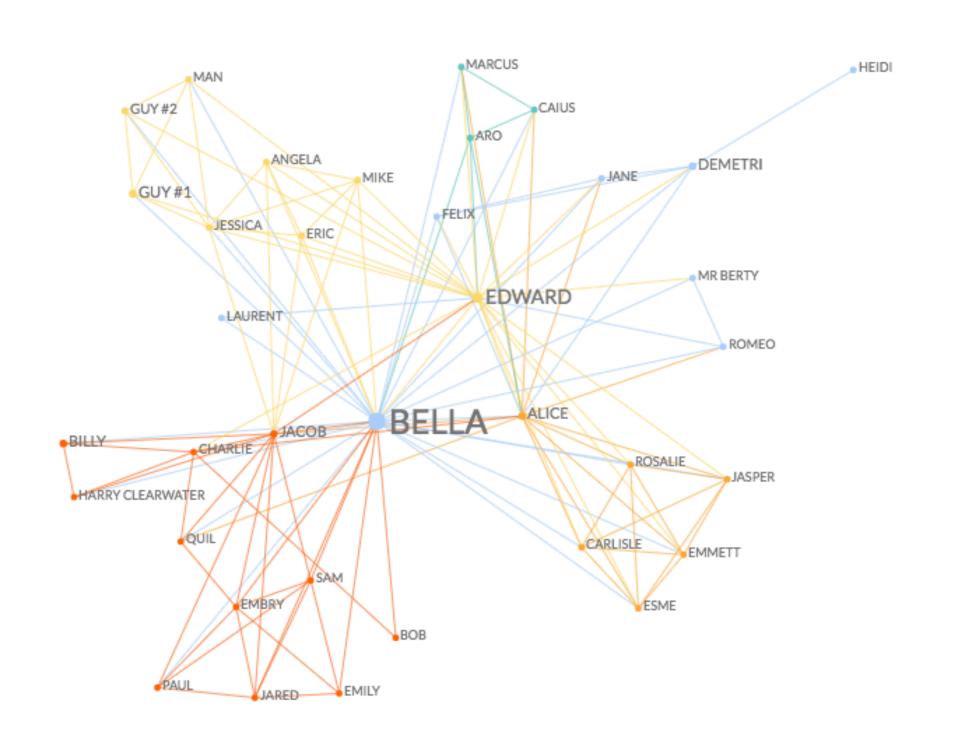
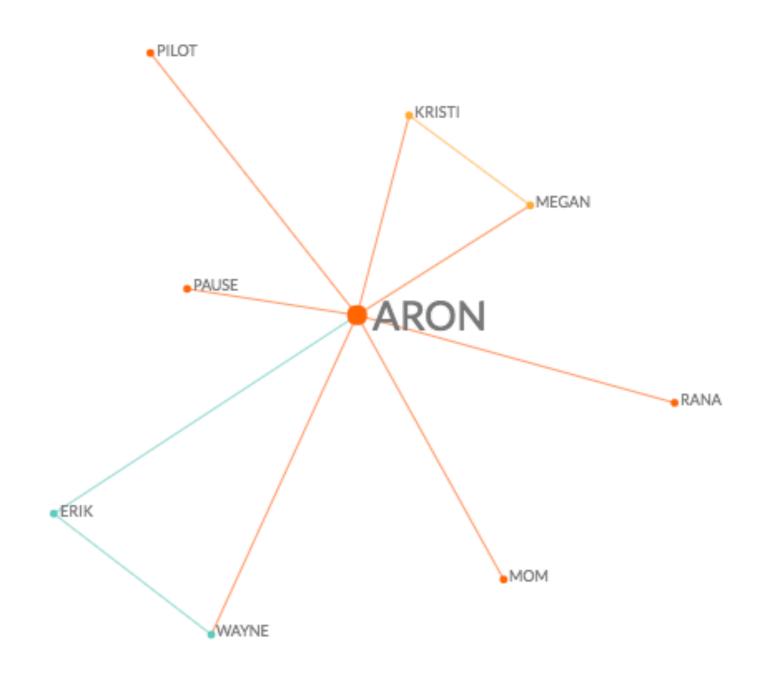
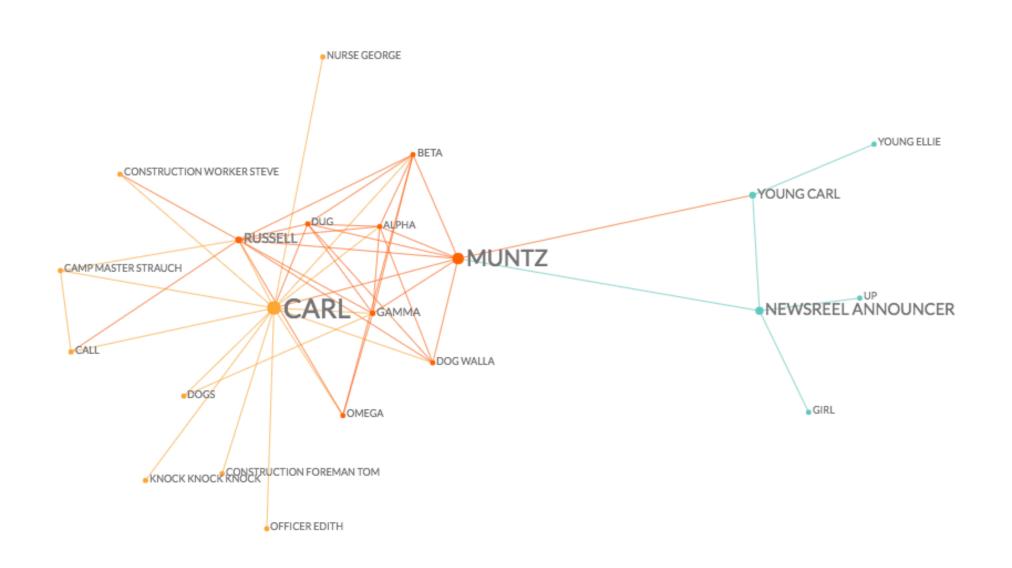


Figure 2.7: A network in which the nodes are students in a large American high school, and an edge joins two who had a romantic relationship at some point during the 18-month period in which the study was conducted [49].







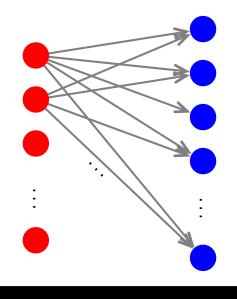


# Connecting Graph Mining to Frequent Item Sets

#### Idea 1: Trawling [Kumar '99]

Searching for small communities in the Web graph

What is the signature of a community / discussion in a Web graph?



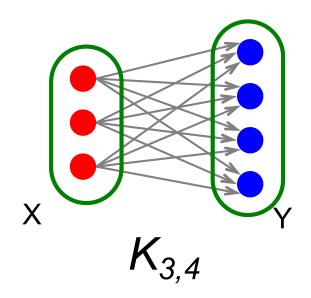
Intuition: Many people all talking about the same things

Dense 2-layer graph

### Searching for Small Communities

A more well-defined problem: Enumerate complete bipartite subgraphs K<sub>s,t</sub>

Where  $K_{s,t}$ : s nodes on the "left" where each links to the same t other nodes on the "right"



$$|X| = s = 3$$
  
 $|Y| = t = 4$ 

Fully connected

#### Frequent Itemset Enumeration

Market Basket Analysis!

Universe U of n items

Baskets: m subsets of U: S<sub>1</sub>, S<sub>2</sub>, ..., S<sub>m</sub> ⊆ U

(S<sub>i</sub> is a set of items one person bought)

Support: Frequency threshold f

Goal:

Find all subsets T s.t.  $T \subseteq S_i$  of at least f sets  $S_i$  (items in T were bought together at least f times)

What's the connection between the itemsets and complete bipartite graphs?

#### From Itemsets to Bipartite Ks,t

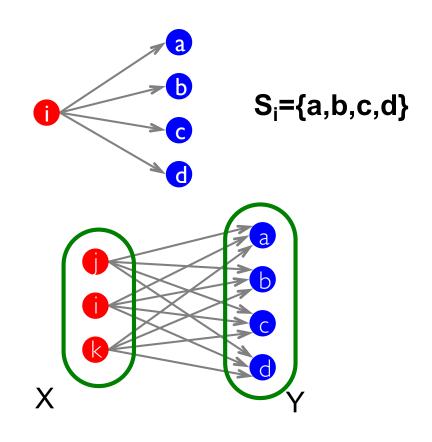
Frequent itemsets = complete bipartite graphs!

How?

View each node i as a set Si of nodes i points to

 $K_{s,t} = a \text{ set } Y \text{ of size } t \text{ that } occurs in s sets <math>S_i$ 

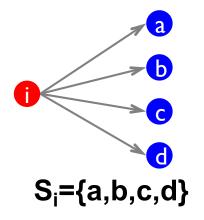
Looking for K<sub>s,t</sub> —> set of frequency threshold to s and look at layer t — all frequent sets of size t



**s** ... minimum support (|X|=s) **t** ... itemset size (|Y|=t)

#### From Itemsets to Bipartite Ks,t

View each node i as a set S<sub>i</sub> of nodes i points to



Find frequent itemsets:

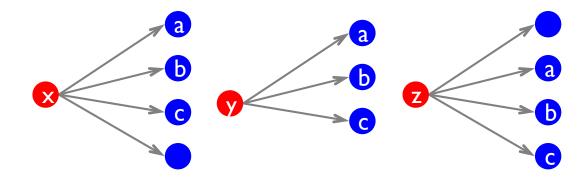
s ... minimum suppor

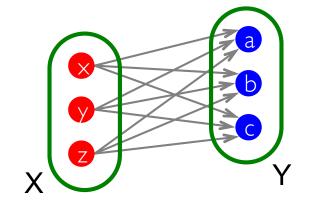
t ... itemset size

We found Ks,t!

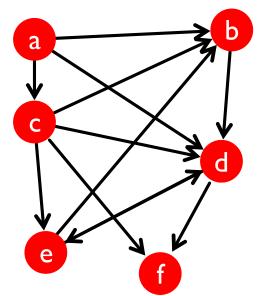
 $K_{s,t}$  = a set Y of size t that occurs in s sets  $S_i$ 

Say we find a frequent itemset Y={a,b,c} of supp s So, there are s nodes that link to all of {a,b,c}:





### Example



#### Itemsets:

$$a = \{b,c,d\}$$

$$b = \{d\}$$

$$c = \{b,d,e,f\}$$

$$d = \{e,f\}$$

$$e = \{b,d\}$$

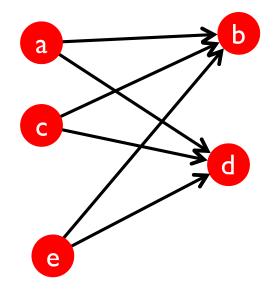
$$f = \{\}$$

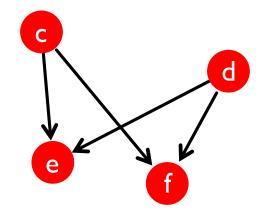
Support threshold s=2

{b,d}: support 3

{e,f}: support 2

And we just found 2 bipartite subgraphs:





#### Example

#### A community of Australian fire brigades

Authorities	Hubs
NSW Rural Fire Service Internet Site	New South Wales Firial Australian Links
NSW Fire Brigades	Feuerwehrlinks Australien
Sutherland Rural Fire Service	FireNet Information Network
CFA: County Fire Authority	The Cherrybrook Rurre Brigade Home Page
"The National Centeted Children's Ho	New South Wales Firial Australian Links
CRAFTI Internet Connexions-INFO	Fire Departments, F Information Network
Welcome to Blackwoo Fire Safety Serv	The Australian Firefighter Page
The World Famous Guestbook Server	Kristiansand brannvdens brannvesener
Wilberforce County Fire Brigade	Australian Fire Services Links
NEW SOUTH WALES FIRES 377 STATION	The 911 F,P,M., Firmp; Canada A Section
Woronora Bushfire Brigade	Feuerwehrlinks Australien
Mongarlowe Bush Fire – Home Page	Sanctuary Point Rural Fire Brigade
Golden Square Fire Brigade	Fire Trails "1ghters around the
FIREBREAK Home Page	FireSafe - Fire and Safety Directory
Guises Creek Voluntfficial Home Page	Kristiansand Firededepartments of th

## Idea 2: Frequent Subgraph Mining

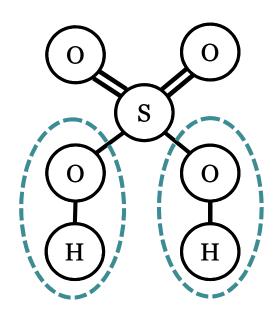
Instead of finding frequent itemsets, lets look for frequent subgraphs

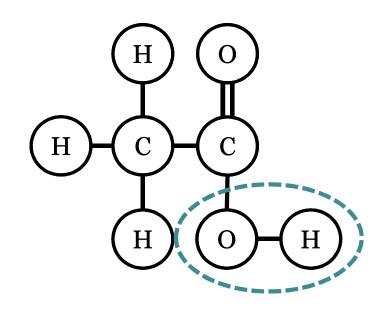
Frequent subgraph mining:

Discovery of graph structures that occur a significant number of times across a set of graphs

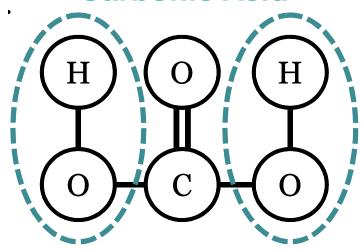
#### **Acetic Acid**

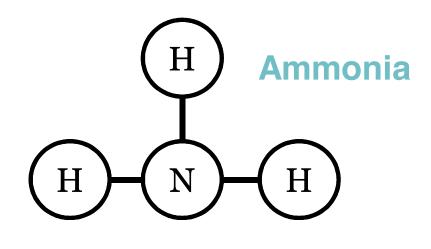
#### **Sulfuric Acid**

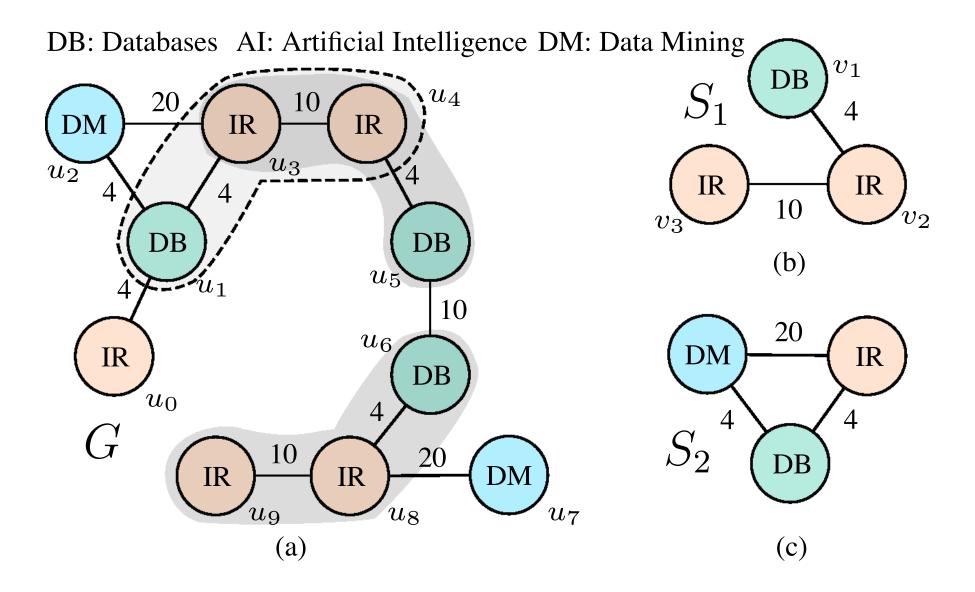




#### **Carbonic Acid**



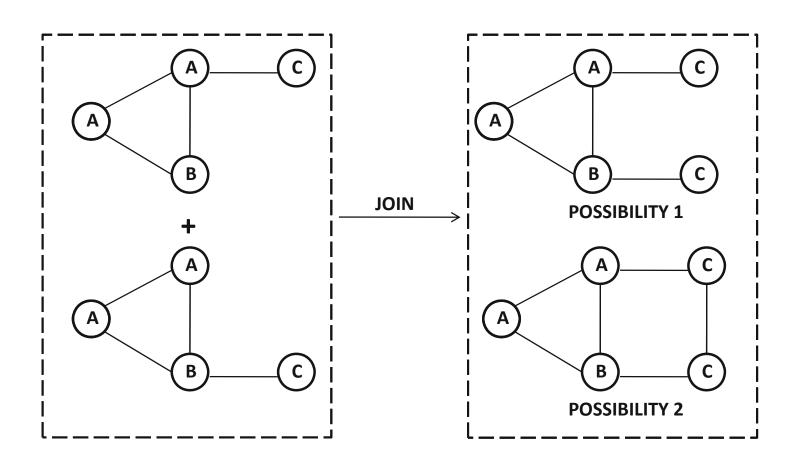




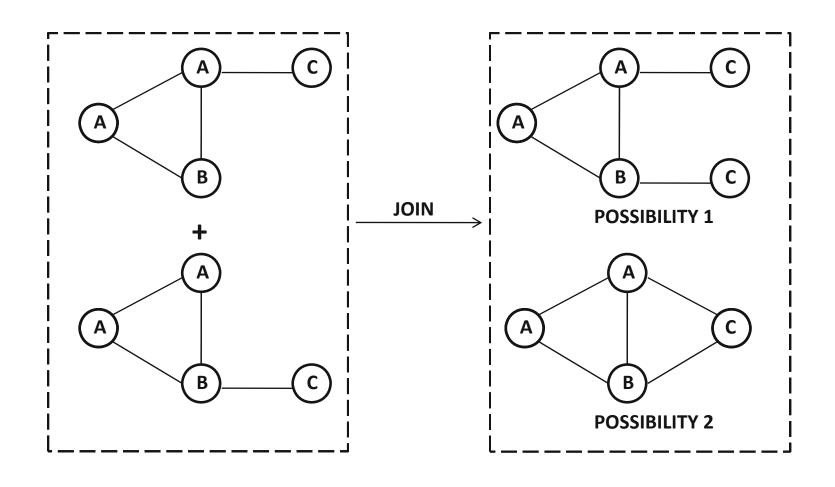
GRAMI: Frequent Subgraph and Pattern Mining in a Single Large Graph, VLDB 2014

```
Algorithm GraphApriori(Graph Database: \mathcal{G},
           Minimum Support: minsup);
begin
 \mathcal{F}_1 = \{ \text{ All Frequent singleton graphs } \};
 k = 1;
 while \mathcal{F}_k is not empty do begin
  Generate \mathcal{C}_{k+1} by joining pairs of graphs in \mathcal{F}_k that
       share a subgraph of size (k-1) in common;
 Prune subgraphs from C_{k+1} that violate downward closure;
 Determine \mathcal{F}_{k+1} by support counting on (\mathcal{C}_{k+1},\mathcal{G}) and retaining
           subgraphs from C_{k+1} with support at least minsup;
   k = k + 1;
 end;
 \mathbf{return}(\cup_{i=1}^k \mathcal{F}_i);
end
```

## Node-based join



## Edge-based Join



# Finding Important Nodes

### Which nodes are important?

Why would we care?

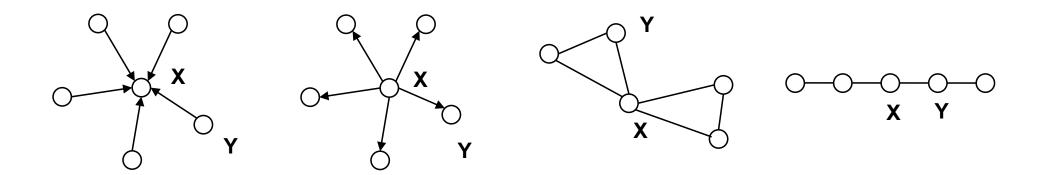
. . .

How would we find?

. . .

### Q: How to Measure Centrality?

In each of the following networks, X has higher centrality than Y according to a particular measure



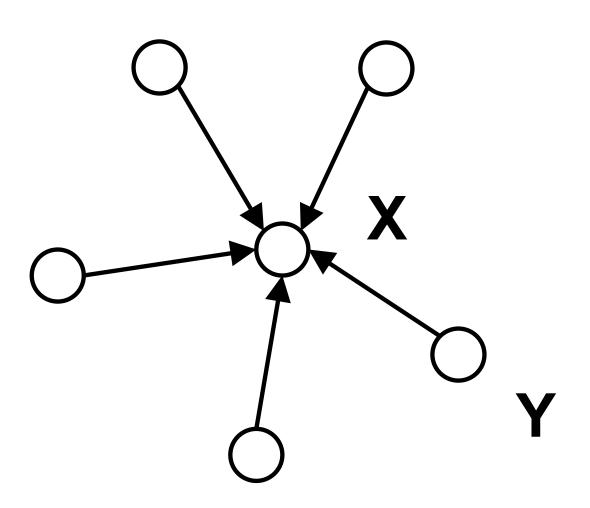
indegree

outdegree

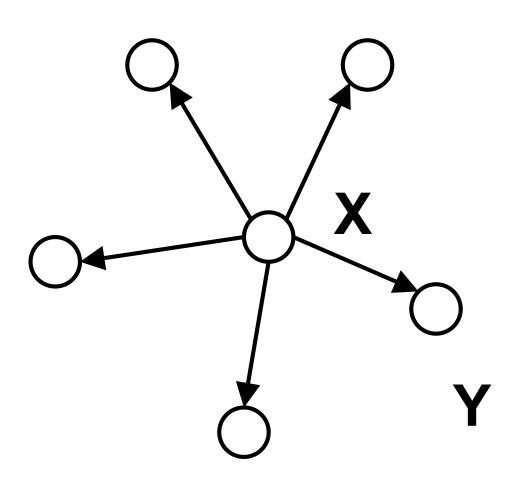
betweenness

closeness

## In-degree centrality

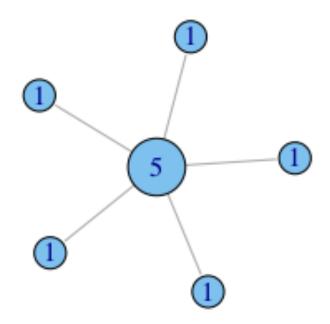


## Out-degree centrality



### Undirected degree centrality

Nodes with more friends are more central

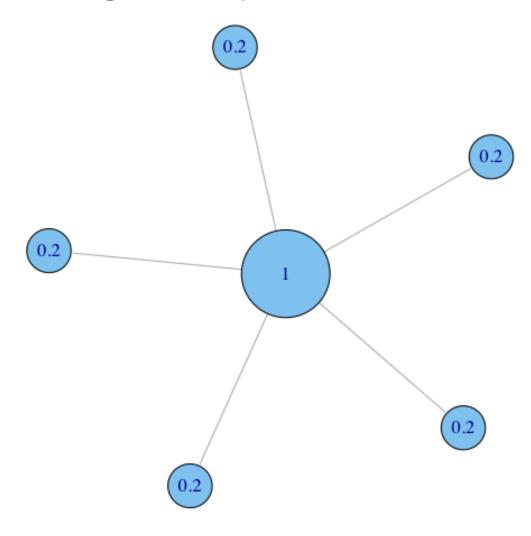


Key assumption: the connections that your friends have are unimportant; all that matters is what your friends can do directly for you

Examples?

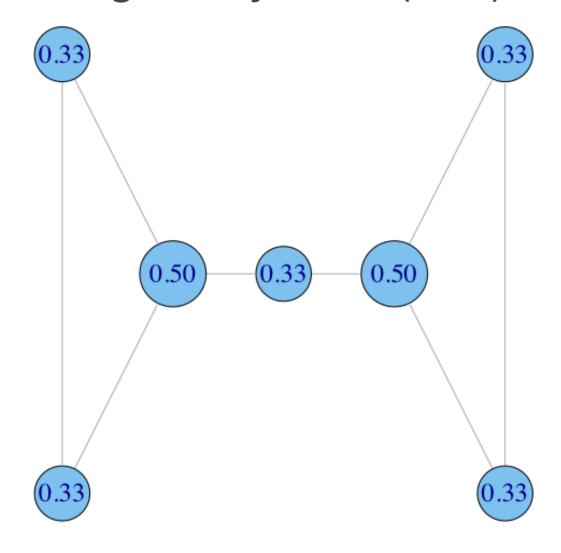
#### Normalization

divide degree by max (N-1)

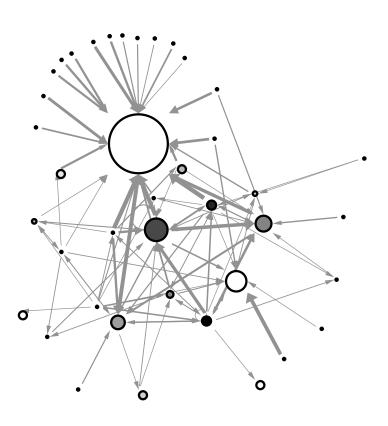


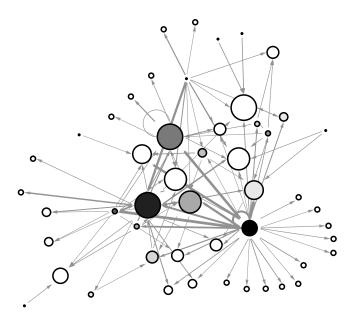
#### Normalization

divide degree by max (N-1)



### Example: financial trading





# Q: What are these degree-based centrality measures missing?

Brokerage!

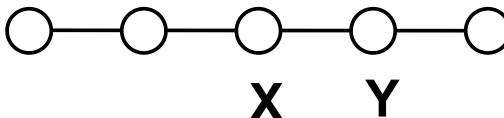
Connecting me to others

friends-of-friends, ... Y

#### Betweenness centrality

Betweenness: Capturing brokerage in a centrality measure

Intuition: how many pairs of individuals would have to go through you in order to reach one another in the minimum number of hops?

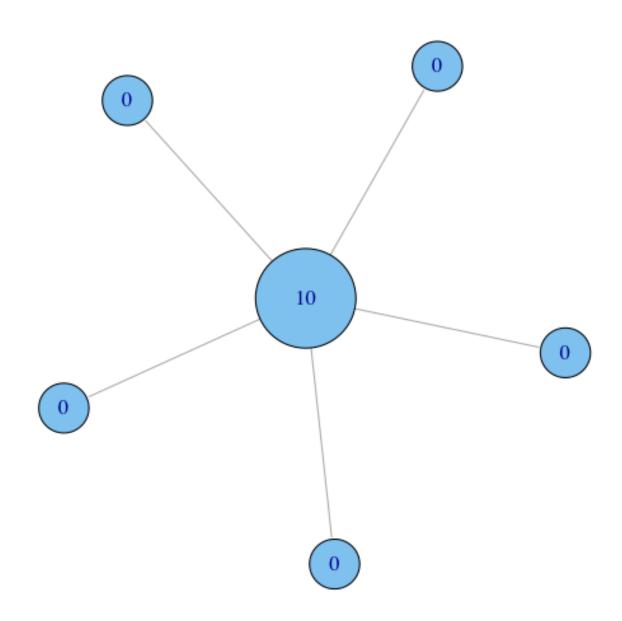


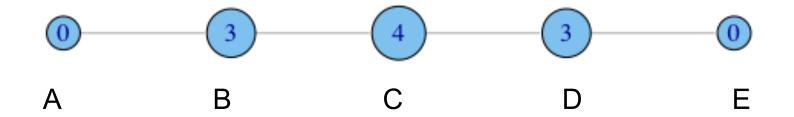
#### Betweenness centrality

$$g(v) = \sum_{s 
eq v 
eq t} rac{\sigma_{st}(v)}{\sigma_{st}}$$

where  $\sigma_{st}$  is the total number of shortest paths from node s to node t and  $\sigma_{st}(v)$  is the number of those paths that pass through v.

Usually normalized by total number of possible vertex pairs (excluding itself)





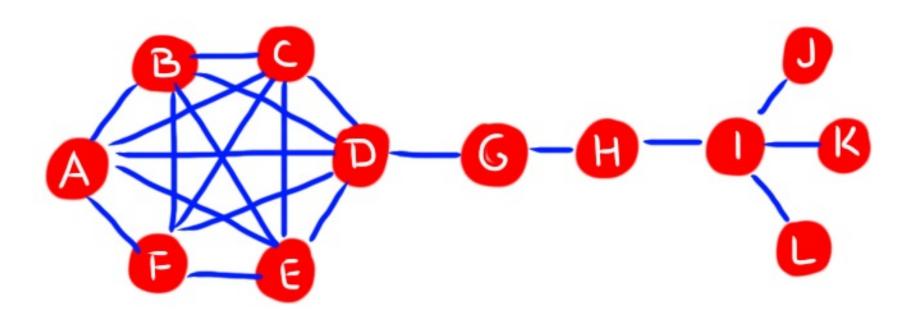
A lies between no two other vertices

B lies between A and 3 other vertices: C, D, and E

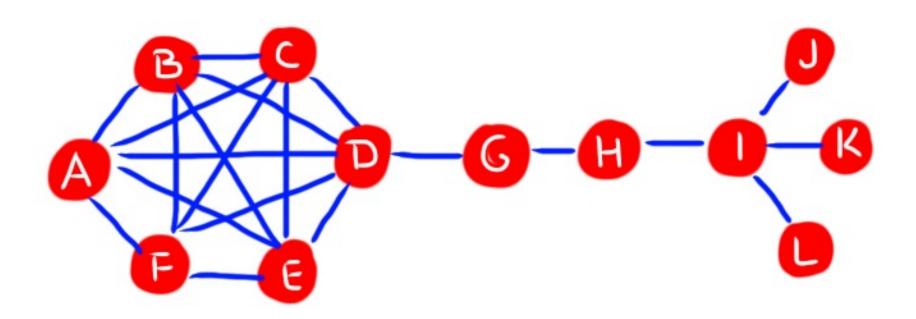
C lies between 4 pairs of vertices (A,D),(A,E), (B,D),(B,E)

note that there are no alternate paths for these pairs to take, so C gets full credit

## Q: Find a node with high betweenness but low degree



# Q: Find a node with low betweenness but high degree

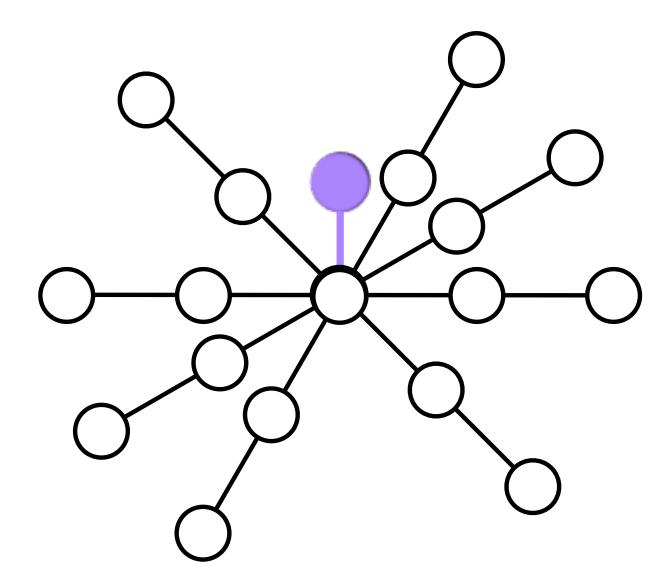


#### Closeness centrality

What if it's not so important to have many direct friends?

Or be "between" others

But one still wants to be in the "middle" of things, not too far from the center



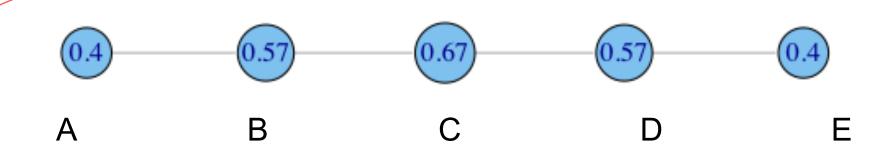
#### Closeness centrality

Closeness is based on the length of the average shortest path between a node and all other nodes in the network

$$C_c(i) = \left[\sum_{j=1}^N d(i,j)\right]^{-1}$$

Normalized:

$$C_C'(i) = (C_C(i))/(N-1)$$



$$C'_{c}(A) = \left[\frac{\sum_{j=1}^{N} d(A,j)}{N-1}\right]^{-1} = \left[\frac{1+2+3+4}{4}\right]^{-1} = \left[\frac{10}{4}\right]^{-1} = 0.4$$

### Q: node with high degree but low closeness?

