### Stack ISA - Example of assembly program: A = B \* (C+D)

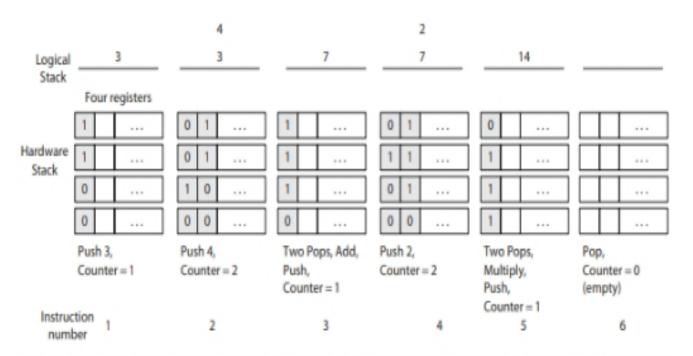


FIGURE 8.4 An illustration of stack content when computing the reverse polish notation CD + B \* = A; it is assumed that (B) = 2, (C) = 3, and (D) = 4.

```
Instruction
number
         PUSH (C) //stack + M[C]
  1:
         PUSH (D) //stack ← M[D]
  2:
                  //stack ← (C) + (D), values popped, added,
         ADD
  3:
                  //result pushed
         PUSH (B) //stack ← M[B]
  4:
         MUL //stack ← ((C) + (D)) * (B), values popped, added,
  5:
             //result pushed
         POP (A) //M[A] ← (((C) + (D)) * (B)), value is popped
  6:
                  //and stored in memory
```

Acc ISA - Example of assembly program: A = B \* (C+D)

```
1. LD (C) // ACC \leftarrow M[C]

2. ADD (D) // ACC \leftarrow ACC + M[D]

3. MUL (B) // ACC \leftarrow ACC * M[B]

4. ST(A) // M[A] \leftarrow ACC
```

CISC-ISA: Example of assembly program: A = B \* (C+D)B = 4; C = 5; D = 10

The value in R1 after execution of instruction No. 1 is 5.

The value of R1 after execution of instruction No. 3 is 60

```
1. LD R1, (C) // R1 \leftarrow M[C]
2. ADD R1, (D) // R1 \leftarrow R1 + M[D]
3. MUL R1, (B) // R1 \leftarrow R1 * M[B]
4. ST (A), R1 // M[A] \leftarrow R1
```

### RISC-ISA: Example of assembly program: A = B \* (C + D)

```
1. LD R1, (C)  // R1 \leftarrow M[C]

2. LD R2, (D)  // R2 \leftarrow M[D]

3. Add R3, R1, R2  // R3 \leftarrow R1 + R2

4. LD R4, (B)  // R4 \leftarrow M[B]

5. MUL R5, R3, R4  // R5 \leftarrow R3 * R4

6. ST (A), R5  // M[A] \leftarrow R5
```

Computation is performed by a RISC ISA. A = B \* (C + D). What is the value in R5 after the execution of code line # 6: (B = 5; C = 10; D = 15) ie: Code line # 6 has been completed. (20 pts)

R5?

# Addressing Modes and Syntax Examples

- o Immediate
  - E.g., Add R1, 9;
- o Direct
  - E.g., ADD R1, (M[9]);
- o Register
  - E.g., ADD R1, R2;
- o Register direct
  - E.g., ADD R1, (R2);
- o Register indexed
  - E.g., ADD R1, R2, (M[9]);

Operand Notation	Addressing Mode
V	I, immediate: V is an immediate input operand, a 2's complement number.
(V)	D, direct: V is a memory address and (V) indicates the content of memory address V (i.e., M[V]).
R	R, register: Indicates an input data register source or a destination register or both
R, (V)	X, indexed: $V$ is a memory address and $R+V$ is the address of the next data item in memory (i.e., $M[R+V]$ ).

TABLE 8.1 Examples of Addressing Modes

```
Immediate E.g., Add R1, 9;
//
Direct E.g., ADD R1, (M[9]);
Register E.g., ADD R1, R2;
Register direct E.g., ADD R1, (R2);
Register indexed E.g., ADD R1, R2, (M[9]);
```

## oRegister indexed E.g., ADD R1, R2, (M[9]);

• What is the value in R1?

 $RTN\#1: R1 \leftarrow R1 + M[R2 + 9]$ 

- Example 8.2 Assembly code listing of an Acc-ISA assembly language program for the high level (c) program in Example code 1

#### **Example Instruction format for Acc-ISA** .code ST (sum) ST (j) Example code #1 L1: CMP JGT 7 L2 int array[8]; MVX int i, sum; X(array) LD sum = 0;ADD (sum) for (i = 0; i < 8; i++)(sum) LD sum = sum + array[i]; ADD L2: .data RB 16 RB 2 RB 2 array: sum: Example 8.2. The listing of an Acc-ISA assembly language program for the program in Example code 1.

.code L1: The listing of an Acc-ISA assembly language program for the program in

```
LD
            0
      ST
      ST
L1:
      CMP 7
      JGT
             L2
      MVX
      LD
            X(array)
      ADD
            (sum)
      ST
            (sum) )
      LD
            (i)
```

.code //start program code

```
ADD 1
ST (i)
JMP L1 //loop back (End of for loop)
```

# Program level Translation:

Now, here is an example of a real C If-Then-Else:

```
if(x == 10)
{
    x = 0;
}
else
{
    x++;
}
```

Which gets translated into the following assembly/machine code:

```
X = 5; Ox 05 = 5 in decimal;
X = 0xA which is equal decimal 10.
Mov eax, $x
Cmp eax, 0x0A; 0x0A = 10
```

Jne else

Mov eax, 0

Jmp end

Else:

Inc eax

End;

Mov \$x, eax