## **Network Security**

AA 2020/2021 Security Protocols

## Examples

- IPSec
- WLAN Security
- DNS



## Typical Attacks to IPv4

- Lack of confidentiality (stealing credentials)
- Lack of source authentication (spoofing, DOS)
- Source routing (spoofing and redirection)



## IP Security Objectives

- Application level:
  - Transparent to applications and users (below transport layer)
- Host Level
  - Provide security for individual hosts
- Router Level
  - router or neighbor advertisements come from authorized routers
  - redirect message come from routers to which the initial packet was sent
  - A routing update is not forged



#### **IPSec**

- A set of security protocols
- A general framework that allows a pair of communicating entities (IP addresses!) to choose the appropriate crypto for the communication.
- IPSec service
  - Connectionless integrity
  - Data origin authentication
  - Rejection of replayed packets
  - Confidentiality (encryption)
  - Limited traffic flow confidentiallity

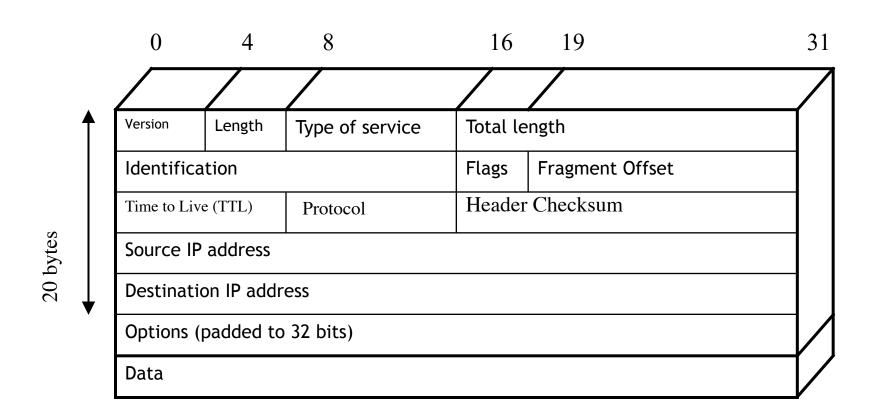


#### **IPsec Basic Features**

- Two basic modes of use:
  - "Transport" mode: for IPsec-aware hosts as endpoints.
  - "Tunnel" mode: for IPsec-unaware hosts, established by intermediate gateways or host OS.
- Provides authentication and/or confidentiality services for data.
  - AH and ESP protocols.
- Provides flexible set of key establishment methods:
  - IKE, IKEv2.

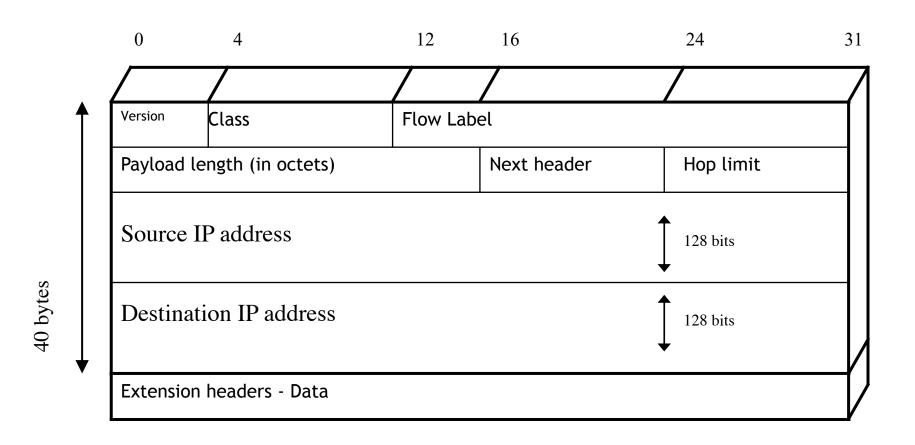


## IPv4 Header



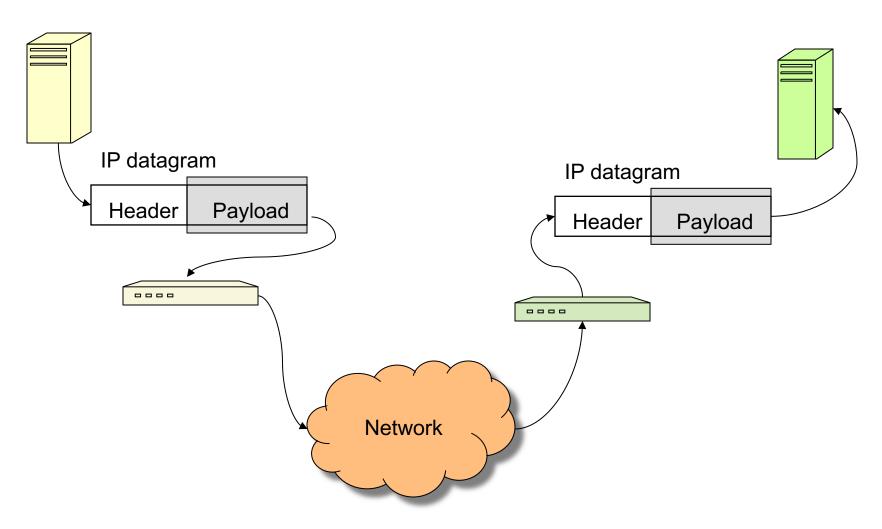


## IPv6 Header





# IPSec Transport Mode



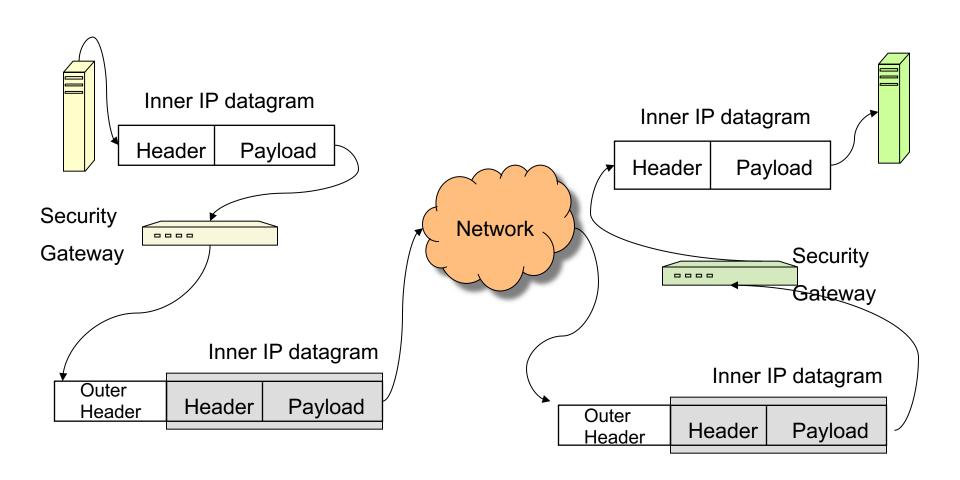


## IPsec Transport Mode

- Protection for upper-layer protocols.
- Protection covers IP datagram payload (and selected header fields).
  - Could be TCP packet, UDP, ICMP message,....
- Host-to-host (end-to-end) security:
  - IPsec processing performed at endpoints of secure channel.
  - Endpoint hosts must be IPsec-aware.



## **IPsec Tunnel Mode**





#### **IPsec Tunnel Mode**

- Protection for entire IP datagram.
- Entire datagram plus security fields treated as new payload of 'outer' IP datagram.
- Original 'inner' IP datagram encapsulated within 'outer' IP datagram.
- IPsec processing performed at security gateways on behalf of endpoint hosts.
  - Gateway could be perimeter firewall or router.
  - Gateway-to-gateway rather than end-to-end security.
  - Hosts need not be IPsec-aware.
- Inner IP datagram not visible to intermediate routers:
  - Even original source and destination addresses encapsulated and so 'hidden'.



#### Protocols

- AH: Authentication Header for authentication and integrity
- ESP: Encapsulating Security Payload for confidentiality and authentication

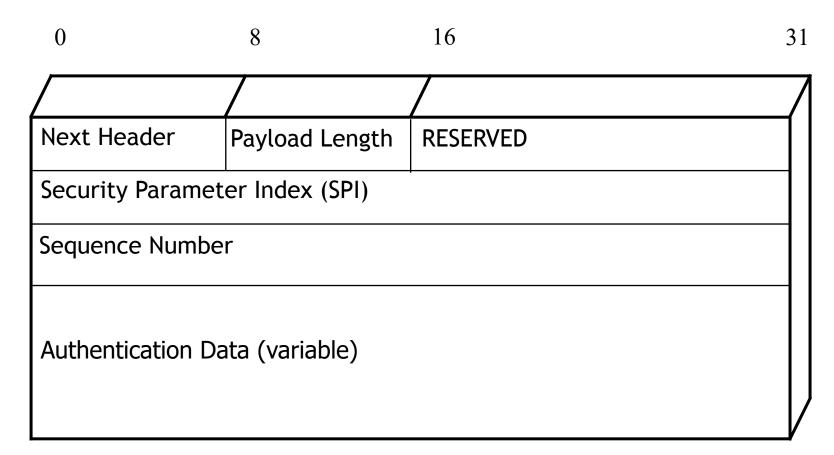


#### **AH Protocol**

- AH = Authentication Header (RFC 2402).
- Provides data origin authentication and data integrity.
- AH authenticates whole payload and most of header.
- Prevents IP address spoofing.
  - Source IP address is authenticated.
- Creates stateful channel.
  - Use of sequence numbers.
- Prevents replay of old datagrams.
  - AH sequence number is authenticated.
- Uses MAC and secret key shared between endpoints.



## Authentication Header (RFC 2402)



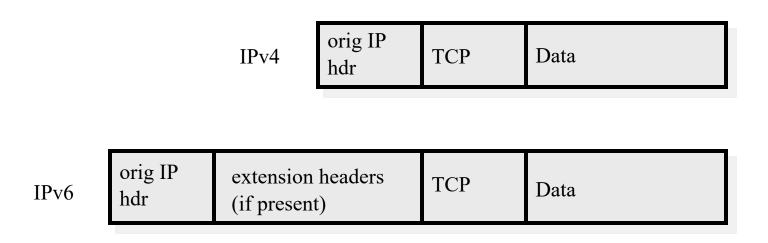


#### AH Protocol

- AH specifies a header added to IP datagrams.
- Fields in header include:
  - Payload length
  - SPI = Security Parameters Index
    - Identifies which algorithms and keys are to be used for IPSec processing (more later).
  - Sequence number
  - Authentication data (the MAC value)
    - Calculate over immutable IP header fields (so omit TTL) and payload or inner IP datagram.

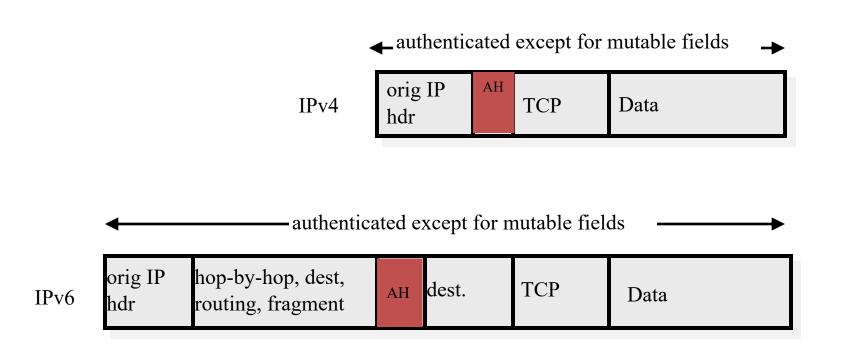


# Before applying AH



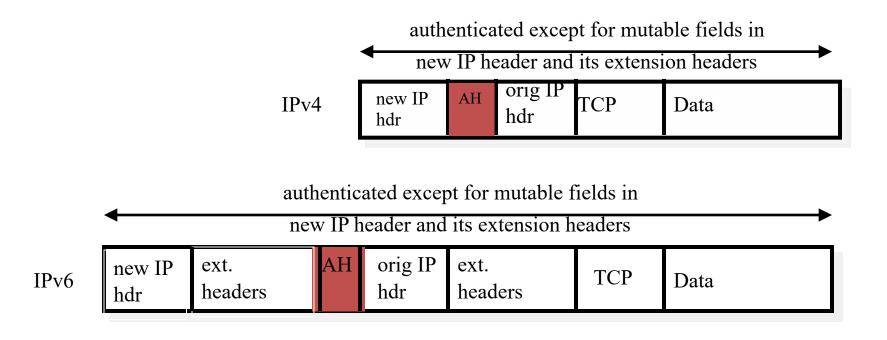


## Transport Mode (AH Authentication)





# Tunnel Mode (AH Authentication)





### **ESP Protocol**

- ESP = Encapsulating Security Payload (RFC 2406).
- Provides one or both:
  - Confidentiality for payload/inner datagram; sequence number not protected by encryption.
  - Authentication of payload/inner datagram; but <u>not</u> of any header fields (original header or outer header).
- Traffic-flow confidentiality in tunnel mode.
- Uses symmetric encryption and MACs based on secret keys shared between endpoints.

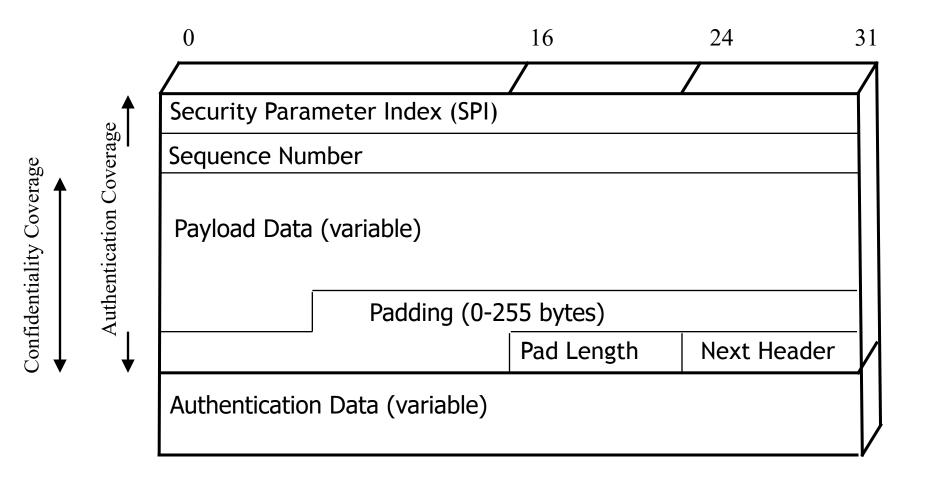


### **ESP Protocol**

- ESP specifies a header and trailing fields to be added to IP datagrams.
- Header fields include:
  - SPI (Security Parameters Index): identifies which algorithms and keys are to be used for IPsec processing (more later).
  - Sequence number.
- Trailer fields include:
  - Any padding needed for encryption algorithm (may also help disguise payload length).
  - Padding length.
  - Authentication data (if any) the MAC value.

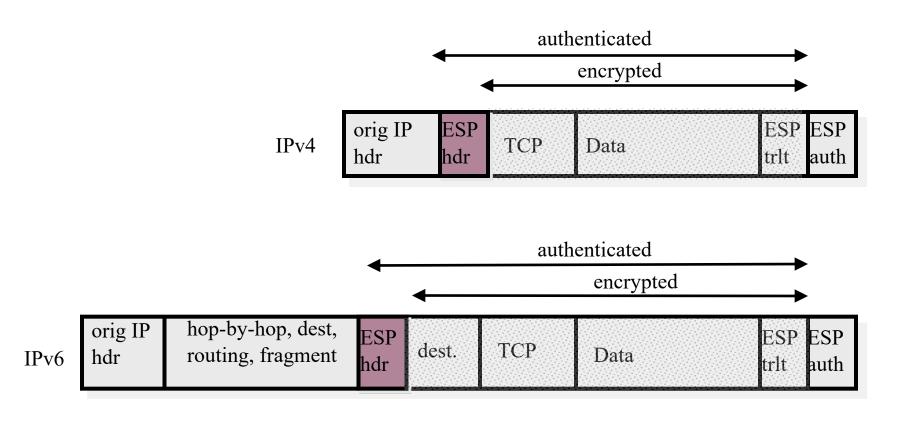


## **Encapsulating Security Payload**



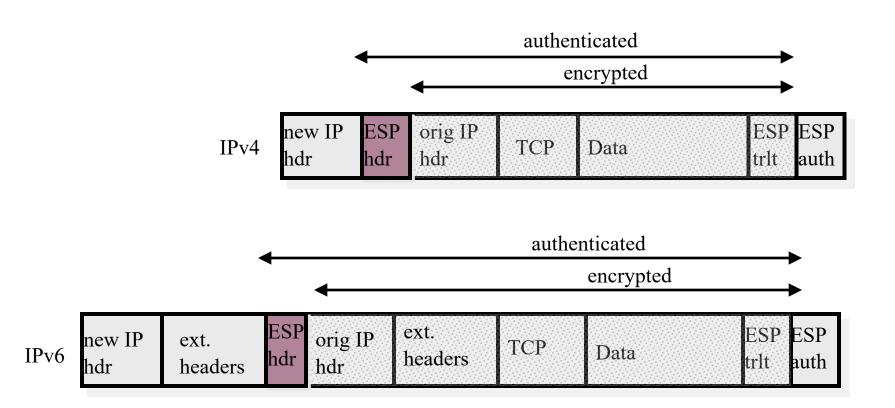
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# ESP Encryption and Authentication (Transport)



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# ESP Encryption and Authentication (Tunnel)





## Combining MAC and ENC

Encryption key  $K_F$  MAC key =  $K_T$ 

Option 1: MAC-then-Encrypt (SSL)

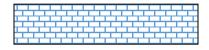
 $MAC(M,K_T)$ 

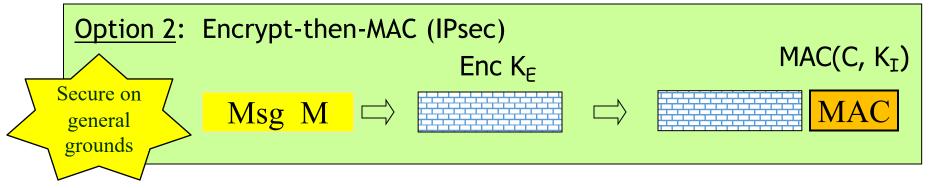
Enc K<sub>F</sub>











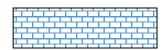
Option 3: Encrypt-and-MAC (SSH)

Enc K<sub>F</sub>

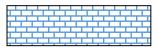
 $MAC(M, K_T)$ 

Msg M











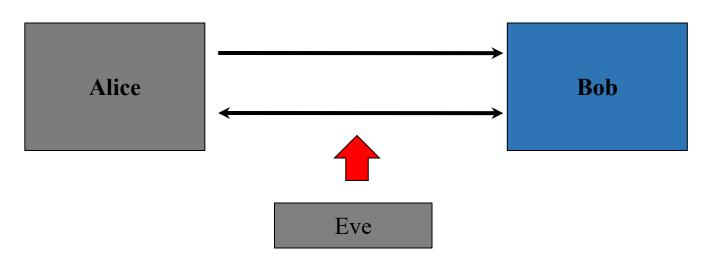
## IPSec Key Management

- IPSec is a heavy consumer of symmetric keys:
  - One key for each SA.
  - Potentially, different SAs for every combination from:
     {ESP,AH} x {tunnel,transport} x {sender, receiver} x {protocol} x {port}.
- Where do these SAs and keys come from?
- Two sources:
  - Manual keying.
    - Fine for small number of nodes and testing purposes.
    - Hopeless for reasonably sized networks of IPSec-aware hosts.
  - IKE: Internet Key Exchange, RFC 2409.
    - RFC documentation hard to follow.
    - Algorithms and parameters negotiation
    - Protocols have many options and parameters.
  - IKEv2
    - Addresses problems and complexities of IKE (i.e. DoS).



## Diffie-Hellman Protocol

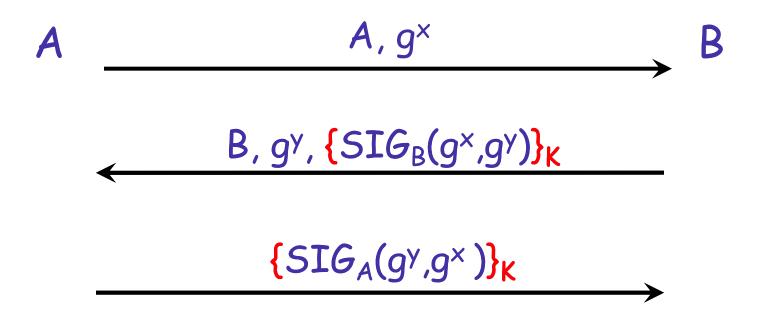
- Simple public-key algorithm for key exchange
- Based on Discrete Logarithm Problem
- Secure against eavesdropping only





## Authenticated DH: STS

Use signature and proof of knowledge



Note: power modulo p

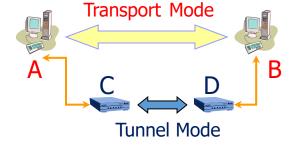
p large prime and q is primitive root module p

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# SPD and SADB Example

SADB: Security Associations DB

SPD: Security Policies DB



A's SADB

#### A's SPD

From	То	Protocol	Port	Policy
Α	В	Any	Any	AH[HMAC-MD5]
From	То	Protocol	SPI	SA Record
Δ	R	ΔЦ	12	HMAC-MD5 key

From	То	Protocol	Port	Policy	Tunnel Dest
		Any	Any	ESP[3DES]	D

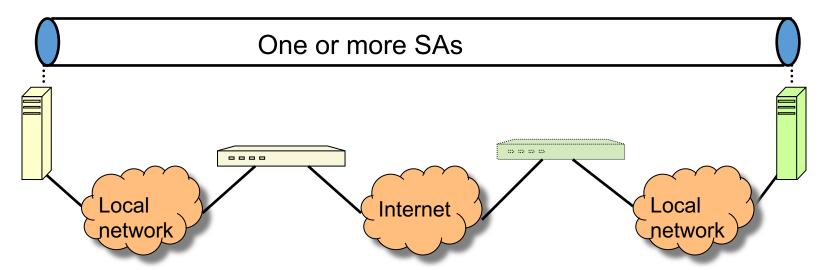
C's SPD
---------

From	То	Protocol	SPI	SA Record
		ESP	14	3DES key

C's SADB



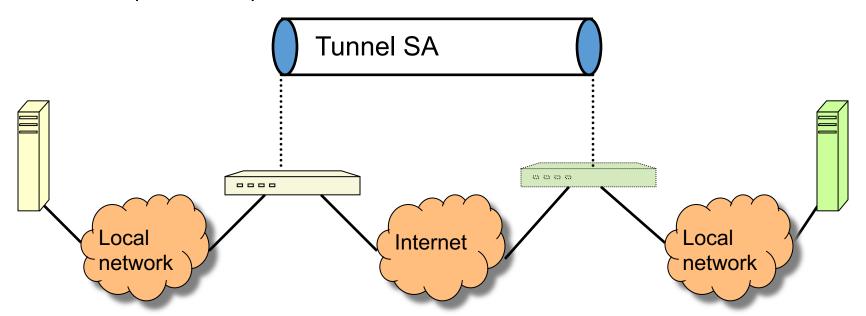
- 1. End-to-end application of IPsec between IPsec-aware hosts; one or more SAs, one of the following combinations:
  - AH in transport
  - ESP in transport
  - AH followed by ESP, both transport
  - Any of the above, tunnelled inside AH or ESP.





#### 2. Gateway-to-gateway only:

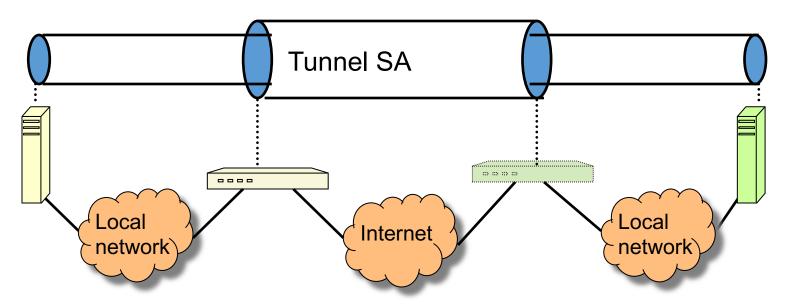
- No IPsec at hosts.
- Simple Virtual Private Network (VPN).
- Single tunnel SA supporting any of AH, ESP (conf only) or ESP (conf+auth).





#### 3. A combination of 1 and 2 above:

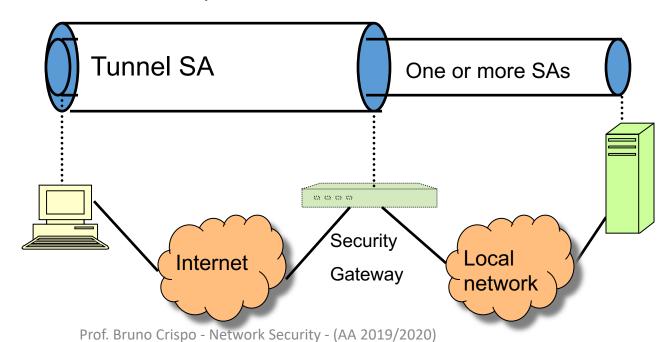
- Gateway-to-gateway tunnel as in 2 carrying host-to-host traffic as in 1.
- Gives additional, flexible security on local networks (between gateways and hosts)
- E.g., ESP in tunnel mode carrying AH in transport mode.





#### 4. Remote host support:

- Single gateway (typically firewall).
- Remote host uses Internet to reach firewall, then gain access to server behind firewall.
- Traffic protected in inner tunnel to server as in case 1 above.
- Outer tunnel protects inner traffic over Internet.





#### Final Notes on IPSec

- IPSec and firewalls have problems working together.
  - Authentication of source IP addresses in AH is the issue.
  - Some firewalls change these addresses on out-bound datagrams (NAT).
- IPSec support for ICMP is somewhat complicated.
- Managing IPSec policy and deployments is tricky.
  - Getting it wrong can mean losing connectivity, e.g. by making exchanges of routing updates unreadable.
  - Getting it wrong can mean loss of security.
  - Many, many IPSec options, rather poor documentation.



## IPSec documents:

- RFC 2401: An overview of security architecture
- RFC 2402: Description of a packet authentication extension to IPv4 and IPv6
- RFC 2406: Description of a packet encryption extension to IPv4 and IPv6
- RFC 2408: Specification of key managament capabilities
- and many more...



#### HTTP

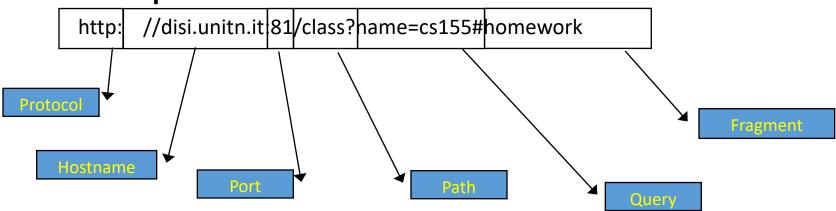
- Main protocol on which the www works
- Based on the notion that client can either request or submit data to a server
- Two methods
  - GET → Requests data from a specified resource
    - GET /test/demo\_form.asp?name1=value1&name2=value2 HTTP/1.1
  - POST → Submits data to be processed to a specified resource
    - POST /test/demo\_form.asp HTTP/1.1 Host: w3schools.com
       name1=value1&name2=value2
- HTTP is stateless
  - HTTP cookies enable statefulness



#### **URLs**

Global identifiers of network-retrievable documents

#### Example:



- Special characters are encoded as hex:
  - %0A = newline
  - %20 or + = space, %2B = + (special exception)



## HTTP GET Request

```
Method File Parameters HTTP version Headers

GET /index.php&user=luca&password=1234 HTTP/1.1

Accept: image/gif, image/x-bitmap, image/jpeg, */*

Accept-Language: en

Connection: Keep-Alive
User-Agent: Mozilla/1.22 (compatible; MSIE 2.0; Windows 95)

Host: www.example.com

Referer: http://www.google.com?q=example
```

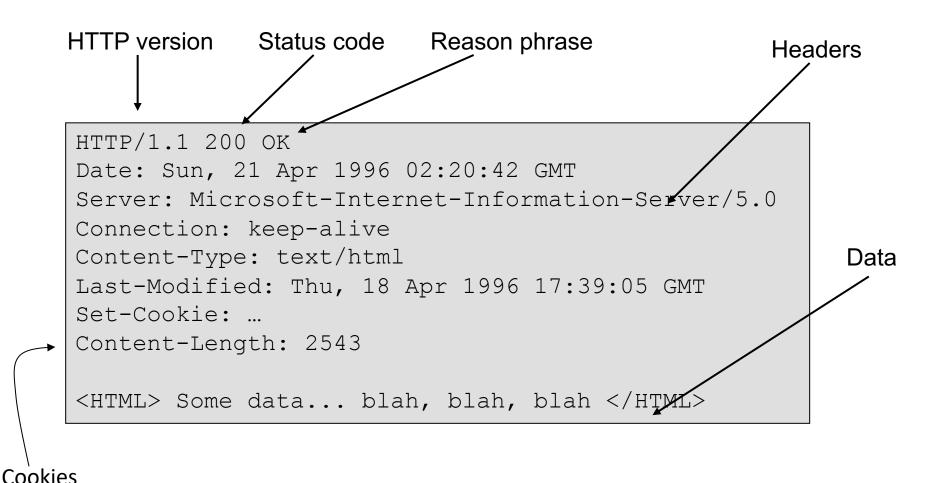


## HTTP POST Request





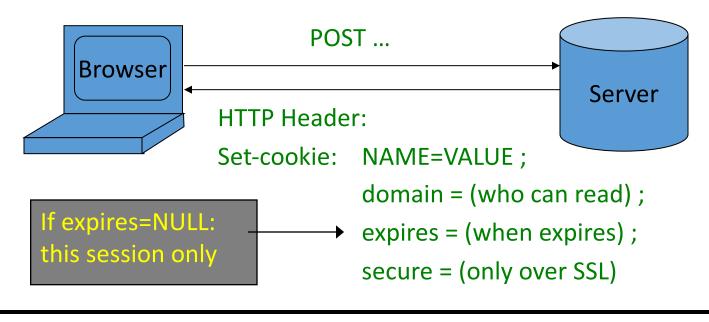
## HTTP Response

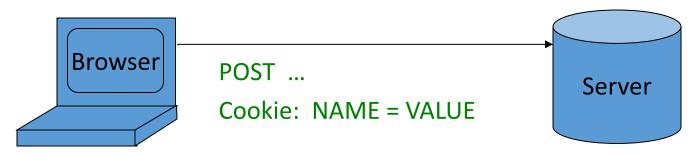




## Cookies

Used to store state on user's machine

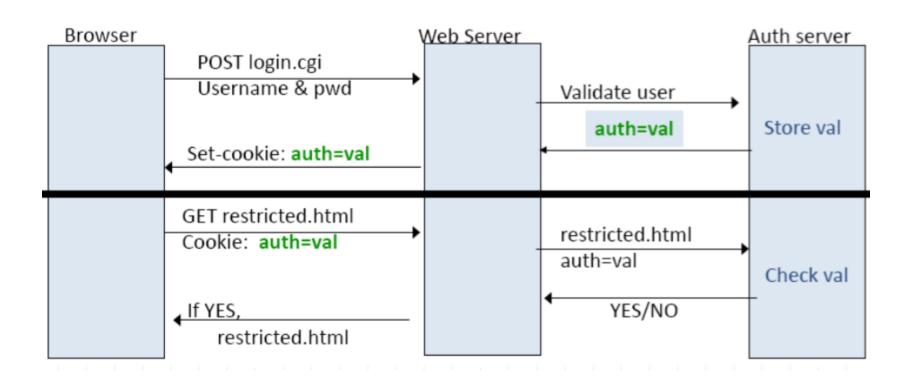




HTTP is stateless protocol; cookies add state

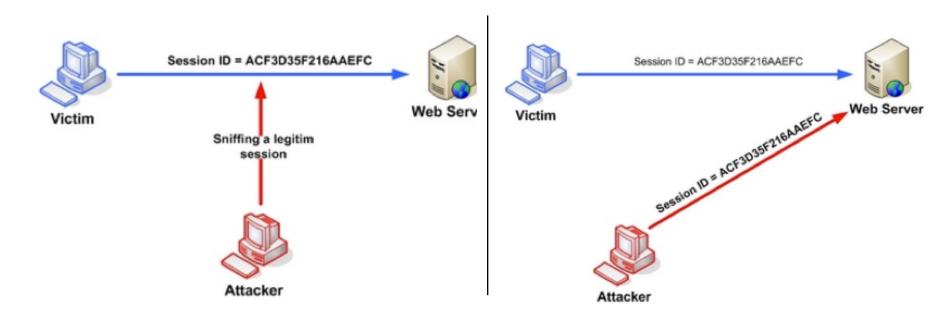


# Cookie example: authentication





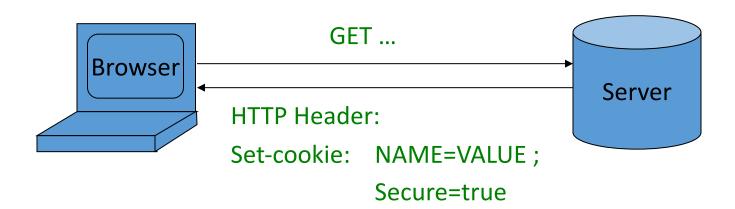
## Attack example: HTTP session hijacking



- Session ID used by webserver to authenticate client "victim"
  - Send over cookie in-the-clear
- Attacker can read the session ID cookie and spoof the victim's identity
  - e.g. access to personal webpages/accounts (e.g. Facebook until 2011)
- https://www.owasp.org/index.php/Session\_hijacking\_attack



### Secure Cookies



- Provides confidentiality against network attacker
  - Browser will only send cookie back over encrypted channels
- ... but no integrity
  - Can rewrite secure cookies over HTTP
     ⇒⇒ network attacker can rewrite secure cookies



## Suggested reading

- Bykova, Marina, and Shawn Ostermann. "Statistical analysis of malformed packets and their origins in the modern Internet." Proceedings of the 2nd ACM SIGCOMM Workshop on Internet measurment. ACM, 2002.
- Hao Yang; Osterweil, E.; Massey, D.; Songwu Lu; Lixia Zhang. Deploying Cryptography in Internet-Scale Systems: A Case Study on DNSSEC. IEEE Transactions on Dependable and Secure Computing. Vol 8, Issue 5.
- Internet Census 2012. Port scanning /0 using insecure embedded devices.
  - http://internetcensus2012.bitbucket.org/paper.html
- Blackert, W. J., et al. "Analyzing interaction between distributed denial of service attacks and mitigation technologies." DARPA information survivability conference and exposition, 2003. Proceedings. Vol. 1. IEEE, 2003.
- S. M. Bellovin. 1989. Security problems in the TCP/IP protocol suite. SIGCOMM Comput. Commun. Rev. 19, 2 (April 1989), 32-48. DOI=http://dx.doi.org/10.1145/378444.378449