Fiborracci Series

0, 1, 1, 2, 3, 5, 8, 13 ---- n^{th} fib term

Base case — 0th term — 1st term

$n = 6$

$t_n = t_{n-1} + t_{n-2}$

?

$n = 3$

2^0 — 1 — f(3)

2^1 — 2 — f(2) f(1)

2^2 — 4 — f(1) f(0) | Base case

Base case

```
int fib(int n){
if(n==0||n==1)
    return n
return fib(n-1) + fib(n-2);
}
```

$Sum = \frac{a(r^n - 1)}{(r-1)}$

T.C = total calls + w.p in each call

$[2^0 + 2^1 + 2^2 + 2^3 + \cdots 2^{n-1}]$ — G.P

$\frac{2^0 * (2^n - 1)}{2-1} = \frac{1 * (2^n - 1)}{1} = [(2^n - 1)]$

$(2^n - 1) + c$

$[T.C = O(2^n)]$

S.C = Depth of R.T * Memory

$= n * 1$

$[S.C = O(n)]$

→ Check if Array is Sorted

0	1	2	3	4
1	2	3	4	5

Ascending order

we will check for idx i

arr[i] >= arr[i-1]

↓ comparison base

arr, n=5

Compare arr[n-1] >= a[n-2]

we will go from the end

1	2	3	4	5
			n-2	n-1

solved sorted

will go on until base case hit

arr, n=4

arr, n=3

arr, n=2

arr, n=1

arr, n=0

Base case → return True, already sorted

```
IsSorted(arr, n){
    if(n==0 || n==1) return True
    return arr[n-1] >= arr[n-2]
        && isSorted(arr, n-1)
}
```

f(4)
↓
f(3)
↓
f(2)
↓
f(1)
↓
f(0)

T.C = n * O(1)

T.C = O(n)

S.C = O(n)

Binary Search

$$mid = st + \frac{(e-s)}{2}$$

-1	0	3	5	9	12
st = 0		mid			end = 5

```
int bs(arr, tar, start, end){
 if(st <= end){
   mid = (st + e-s)/2
 ① if(a[mid] == tar){
       return mid
   }
 ② if(a[mid] <= tar){
       2nd Half
   return bs(arr, tar, mid+1, end)
 ③ else 1st Half
   return bs(arr, tar, start, mid-1)
   }
   return -1
   }
```

$T.C = \log_2 * O(1)$

$T.C = \log_2$

$\boxed{T.C = O(\log n)}$

$\boxed{S.C = O(\log n)}$

$\frac{n}{2^{1c}} = 1$

$n = 2^{1c}$

$\log_2 n = 1c$

$f(6) \rightarrow f(3) \rightarrow f(1)$

$n \rightarrow n/2^1 \rightarrow n/2^2 \rightarrow n/2^3 \rightarrow \frac{1}{2^k}$