

Fibonacci Series

0, 1, 1, 2, 3, 5, 8, 13

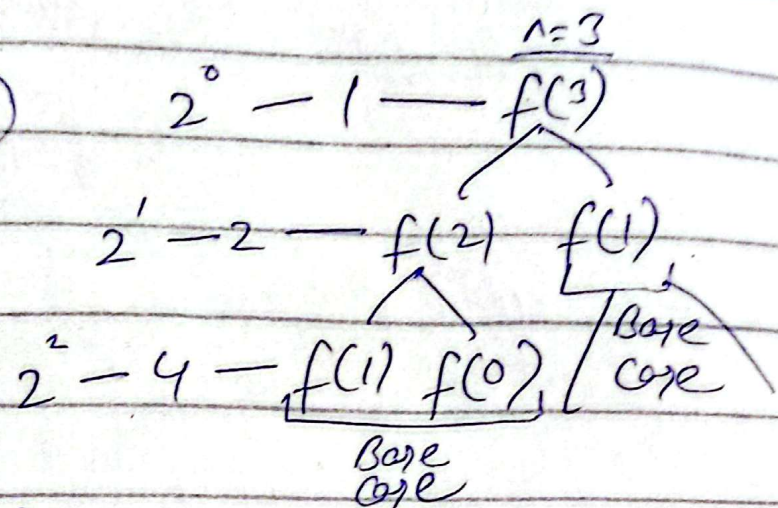
Base case: 0th term, 1st term

n^{th} fib term

$n=6$

$$t_n = (t_{n-1}) + (t_{n-2})$$

?



```
int fib(int n) {
    if (n == 0 || n == 1)
        return n;
```

```
    return fib(n-1) + fib(n-2);
}
```

$$Sum = \frac{a(r^n - 1)}{(r - 1)}$$

T.C = total calls + w.d in each call

$$[2^0 + 2^1 + 2^2 + 2^3 + \dots + 2^{n-1}] - G.P$$

$$\frac{2^0 (2^n - 1)}{2 - 1} = \frac{1 (2^n - 1)}{2 - 1} = (2^n - 1)$$

$$T.C = O(2^n)$$

S.C = Depth of R.T + Memory

= $n + 1$

$$S.C = O(n)$$

→ Check if Array is Sorted

0 1 2 3 4
[1|2|3|4|5]

Ascending order

we will check -
for idx i
 $arr[i] \geq arr[i-1]$
↓
comparison one

arr, n=5
Compare $arr[n-1] \geq arr[n-2]$
we will go from the end

[1|2|3|4|5]
 n-2 n-1 } solved
 sorted

will go on until base case hit

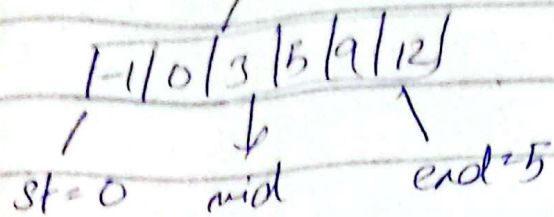
arr, n=4
arr, n=3
arr, n=2
arr, n=1 } Base Case → return True
arr, n=0 } already sorted

isSorted(arr, n) {
 if (n==0 || n==1) return True
 return $arr[n-1] \geq arr[n-2]$
 && isSorted(arr, n-1)
}

f(4)
↓
f(3)
↓
f(2)
↓
f(1)
↓
f(0)

T.C = n * O(1)
T.C = O(n)
S.C = O(n)

→ Binary Search



$$mid = st + \frac{(e-s)}{2}$$

```
int bs(arr, tar, start, end) {
```

```
    if (st == end) {
```

$$mid = (st + e - s) / 2$$

```
    ① if (arr[mid] == tar) {
```

```
        return mid
```

```
    ② if (arr[mid] < tar) {
```

2nd Half

```
        return bs(arr, tar, mid + 1, end)
```

```
    ③ else 1st Half
```

```
        return bs(arr, tar, start, mid - 1)
```

```
    }
```

```
    return -1
```

```
}
```

$$T.C = \log_2 * O(1)$$

$$T.C = \log_2$$

$$T.C = O(\log n)$$

$$S.C = O(\log n)$$

$$\frac{n}{2^k} = 1$$

$$n = 2^k$$

$$\log_2 n = k$$

