Parallel Fringe Search

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ETH Zürich

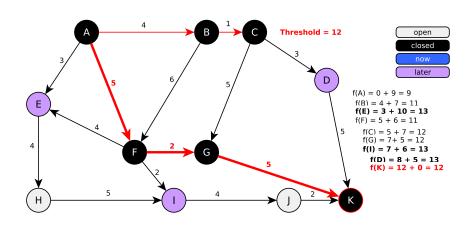
Design of Parallel and High-Performance Computing

December 15, 2013

Overview

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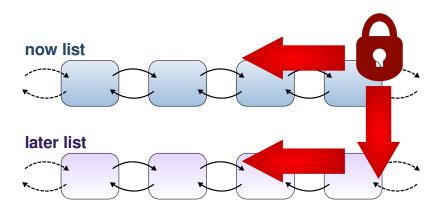
Fringe Search

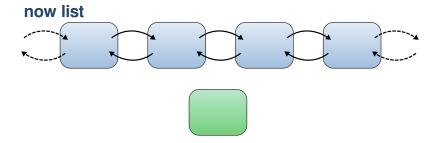


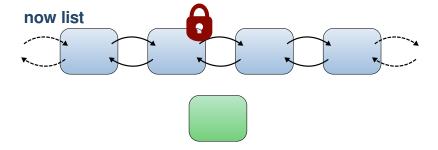
What we have done

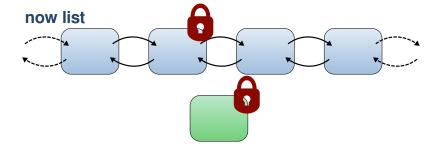
- Serial implementation of fringe search (much faster than Boost A*)
- Parallel implementation with Open MP
 - 2 different locking concepts
 - Locks implemented using inline assembly (faster than Open MP locks)
- Benchmarking
 - Strong scaling
 - Weak scaling
 - Path quality

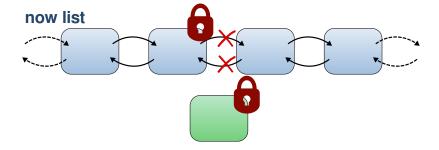
Locking concept: Deadlock prevention

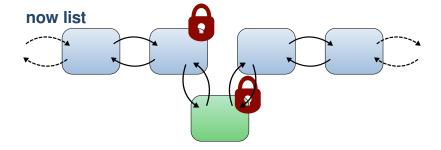


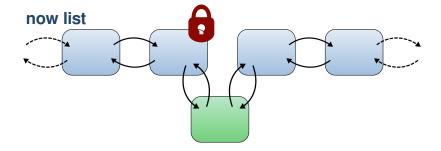


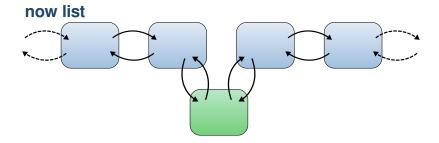


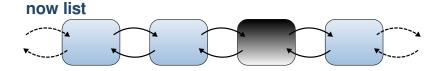


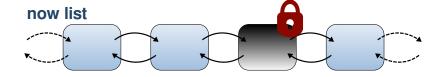


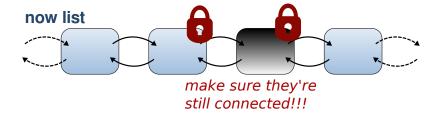


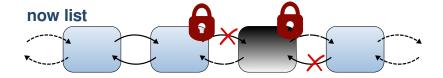


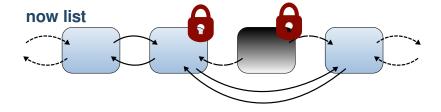


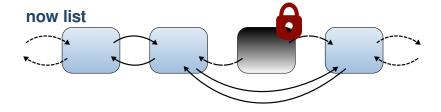


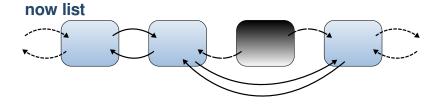












Locking concept: 2 concepts for removing nodes

Normal:

Lock node and predecessor as shown before and remove it right away

Lazy locking:

- Don't lock anything and just mark the node as removed
- Other threads will clean up and remove it later

Benchmarking

Setup:

Each of the following boxplots was generated from data from **50 runs** on **1 node** of kanifushi.inf.ethz.ch.

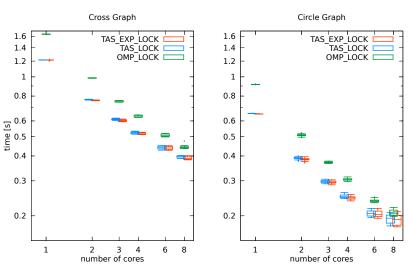
Specifications of kanifushi.inf.ethz.ch:

- NUMA model with 32 CPUs on 4 nodes
- 8 CPUs per node
- Intel(R) Xeon(R) CPU E7- 4830 @ 2.13GHz
- per CPU: 32KB L1 cache, 256KB L2 cache
- per node: 24MB L3 cache, 16GB memory

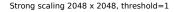
The code is written in C++ / Open MP and it has been compiled with g++ v. 4.6.1 using O1 optimization.

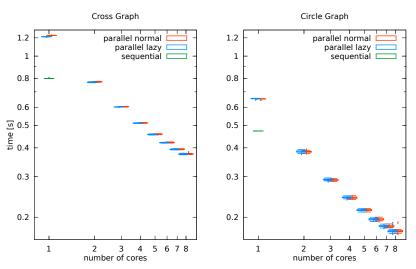
Benchmarking: Locks

Strong scaling 2048 x 2048, threshold=1

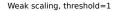


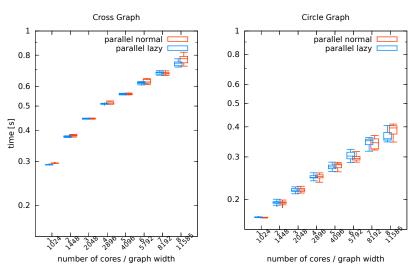
Benchmarking: Strong scaling





Benchmarking: Weak scaling

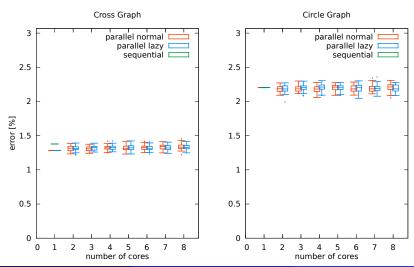




Benchmarking: Path length $\leftrightarrow \#$ cores

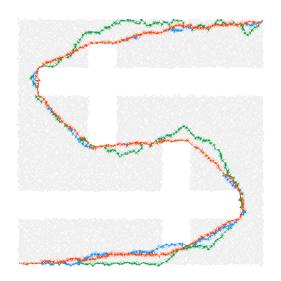
compared to A* from Boost Graph Library

relative Error, 2048 x 2048, threshold=1



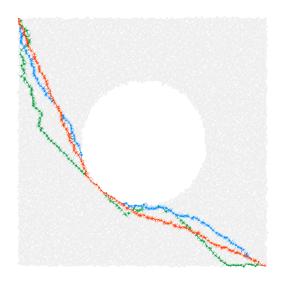
Threshold update

- threshold += 0.1
- threshold +=1
- threshold += 10



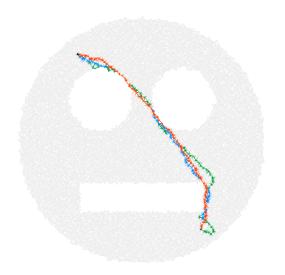
Threshold update

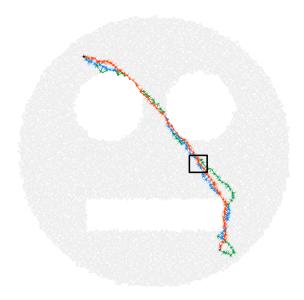
- threshold += 0.1
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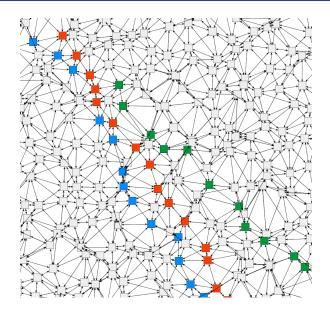


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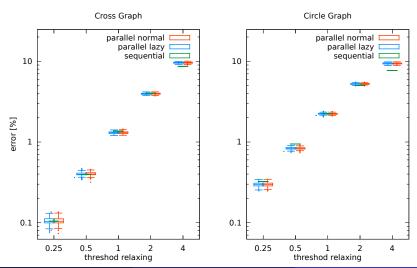






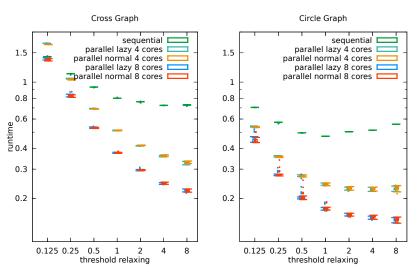
compared to A* from Boost Graph Library

relative Error, 2048 x 2048



Benchmarking: Run time \leftrightarrow Threshold update

runtime with different thresholds, 2048 x 2048



Conclusions

In general Fringe Search is a good single source shortest path algorithm, that can be very well implemented in parallel.

- Path quality not dependent of # cores
- Good strong scaling
- Weak scaling is not perfect
- ullet quality \leftrightarrow runtime trade-off can be tuned for desired result

The End