

Solve minimum window substring problem

In this post we are gonna discuss how to solve minimum window substring problem using Sliding window principle





Desriable Window- Has all the characters from t.

Problem Statement:

Given a string S and a string T, find the minimum window in S which will contain all the characters in T in complexity O(n).

Example:

Input: S = "ADOBECODEBANC", T = "ABC"

Output: "BANC"

 $Note: If there is no such window in S \ that \ covers \ all \ characters \ in \ T, \ return \ the \ empty \ string \ ``.'$

If there is such window, you are guaranteed that there will always be only one unique minimum window in S.

Step by step solution to the problem:

As in this problem we have to *find the minimum window*, first paradigm which comes to our mind to solve this is Sliding window.

Base Case:

if either s or t is null then we can't find any common window between them.

```
if(s==null \mid \mid t==null) return "";
```

Main Algorithm:

 $first \ lets \ keep \ all \ the \ character \ count \ of \ t \ in \ a \ map \ like \ below \ code \ so \ that \ we \ can \ check \ a \ character's \ availability \ from \ this \ map \ easily \ in \ O(1).$

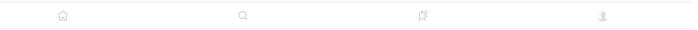
```
int[] map=new int[128];
for(Character c:t.toCharArray())map[c]++;
```

Let's take two pointers start and end to iterate through s and initialize them with 0 along with some other variables.

```
int start = 0, end = 0, minStart = 0, minLen = Integer.MAX VALUE, counter = t.length();
```

Now we iterate over s using end pointer and validate if current character is present in t or not by checking map[arr[end]] > 0. If it's present decrement the counter to signify that we got one character match and (count-1) character to be matched. Also we decrement map[arr[end]] to signify that we have traversed current character using end pointer.

 $In this process if we get {\bf count} = {\bf =0}, that means we got an desired window, so replace this window as result window if this is the minimum window. \\$



```
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             while(end<arr.length){
                  if (map[arr[end]] > 0) counter--;
                  map[arr[end]]--;
                  while(counter==0){
                       if((end-start)<minLen){</pre>
                          minLen=end-start
                          minStart=start;
10
11
                       map[arr[start]]++;
12
                      if(map[arr[start]]>0)counter++;
13
                      start++;
14
15
MinimumWindowSubstring.java hosted with 💝 by GitHub
                                                                                             view raw
```

Two important things to note here is :

```
1. if (map[arr[end]] > 0) counter - ;
  end++;
```

this signifies we got match of char in s with char in t, so we have counter-1 char to match in the current window.

```
2. if(map[arr[start]]>0)counter++;
    start++;
```

this signifies that after restoring the actual char count in map by traversing through start pointer, if count>0 then increment counter and start as we are removing this char from current window, and moving the window forward.

 $The complete code for this problem can be found in \underline{https://github.com/GyanTech877/algorithms/blob/master/slidingwindow/SlidingWindowMaximum.java$

Happy Learning 😇

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