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## **Set Matrix Zeros**

Set Matrix Zeros -Interview Problem



Difficulty: Medium

Asked in: Amazon, Google

## **Understanding the Problem**

**Problem Description:** Given a matrix, **A** of size **M x N** of **0's and 1's**. If an element is **0**, set its entire row and column to **0**.

#### For Example:

#### Input:

[[1, 0, 1],

[1, 1, 1],

[1, 1, 1]]

#### Output:

[[0, 0, 0],

[1, 0, 1],

[1, 0, 1]]

#### Input:

[[1, 0, 1],

[1, 1, 1],

[1, 0, 1]]

#### Output:

[[0, 0, 0],

[1, 0, 1], [0, 0, 0]]

#### Possible follow-up questions to ask the interviewer:-

- 1. What are the constraints on the size of the matrix? (**Ans:**  $1 \le N$ ,  $M \le 1000$ )
- 2. Do we need to change the rows and columns of elements which we set to zero? (**Ans:** No, only the zeroes present in the original matrix)

#### **Solutions**

We are discussing three ways to solve this problem:

- 1. Brute Force Approach: Using nested loops and extra space
- 2. **Using Hash Table:** Storing the status of rows and columns in the Hash Table.
- 3. **Using In-Place Hashing:** Storing hash in the matrix's first row and column.

### 1. Brute Force Approach

#### Common mistake

A simple solution that comes to mind for this problem is that we can traverse the matrix and for every '0' we encounter, we update all values in its row and column to 0. This brute force solution seems quite easy but may actually lead to a wrong answer. (How? Think!)

#### Correct Solution idea

One way would be to create an auxiliary matrix of the same size filled with 1's. We could make all changes required in this auxiliary matrix and then copy it to the original matrix in the end. We search for zeroes in the original matrix and update rows and columns in the auxiliary matrix.

#### **Solution Steps**

- 1. First, create a temporary matrix of the same size M x N and initialize its all places with 1.
- 2. Now scan the original matrix and if A[i][j] == 0 then set all the positions of row i and col j with 0 in new matrix.
- 3. Copy all elements from the temporary matrix to the original matrix.

#### Pseudo-Code

```
void setMatrixZeros(int A[][], int M, int N)
   int temp[M][N]
   for (i = 0 \text{ to } M-1)
          for (j = 0 \text{ to } N-1)
               temp[i][j] = 1
   for(i = 0 \text{ to } M - 1)
     {
          for (j = 0 \text{ to } N - 1)
               if (A[i][j] == 0)
               {
                    for (k = 0 \text{ to } M - 1)
                         temp[k][j] = 0
                    for (k = 0 \text{ to } N - 1)
                         temp[i][k] = 0
               }
          }
```

```
}
for (i = 0 to M-1)
    for (j = 0 to N-1)
        arr[i][j] = temp[i][j]
}
```

#### **Complexity Analysis**

For every cell, we are traversing its row and column and there are a total of N\*M cells.

Time Complexity: O(N\*M\*(N+M)) (**Think!**)

Space Complexity: O(N\*M), for storing the auxiliary matrix.

#### Critical ideas to think!

- How could we implement it in place in order to optimize space complexity?
- How can we reduce the time complexity of the solution?

### 2. Using Hash Table

Instead of using an auxiliary matrix, we could just store the rows and columns which need to be updated (Think!).

We could use a Hash Table to do so. We shall create a hash table for all rows and columns and set them to false. If the value of any row or column in hash is later updated to true, it's a signal for us to update that particular row or column's value to 0.

#### **Solution Steps**

- 1. Create 2 Hash Tables row and col respectively of size M and N.
- 2. Now initialize all the values of row[] and col[] to false.
- 3. Now iterate over the matrix and for every A[i][j] == 0, set row[i] = true and col[j] = true.
- 4. After completion of step 3, again iterate through matrix A and for any A[i][j], if row[i] or col[j] is true then update A[i][j] to 0.

#### Pseudo Code

```
void setMatrixZeros(int A[][], int M, int N)
     bool row[M], col[N]
    for(i = 0 \text{ to } M-1)
          row[i] = false
     for(j = 0 \text{ to } N-1)
          col[j] = false
    for(i = 0 \text{ to } M-1)
     {
          for(j = 0 \text{ to } N-1)
          {
               if(A[i][j] == 0)
                    row[i] = true
                    col[j] = true
               }
          }
     }
     for(i = 0 \text{ to } M-1)
          for(j = 0 \text{ to } N-1)
```

```
{
    if(row[i] == true or col[j] == true)
        A[i][j] = 0
}
}
```

#### **Complexity Analysis**

Time Complexity: Creating and filling values in 2 Hash Tables + Traversing the matrix + Updating the matrix

```
= O(N+M) + O(N*M) + O(N*M)
= O(N*M)
```

Space Complexity: O(M + N), for storing hash tables.

#### Critical ideas to think!

- How can we optimize the space complexity?
- What other matrix problems could be solved using a similar type of hashing approach?



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## 3. Using In-Place Hashing

We are using extra space to store Hash Tables of rows and columns. We can optimize this space by trying to store hash for row and column in the matrix itself. (How?)

We can use the first row and first column of the matrix to store the status of the rows and columns respectively. The problem will occur for the status of first row and column which can be handled using two variables which will store the status of the first row and column **(Think!)** 

#### **Solution Steps**

- 1. Initialize **firstRow** and **firstCol** to false. These two variables will store the status of the first row and first column.
- 2. Now use the first row and first column as your hash which stores the status of that row and column.
- 3. Now iterate over the matrix and for every A[i][j] == 0, set A[i][0] = 0 and col[0][j] = 0.
- 4. Now update the values of the matrix except first row and first column to 0 if A[i][0] = true or A[0][i] = true for A[i][j].
- 5. Now update the values of the first row and first column.

```
void setMatrixZeros(int A[][], int M, int N)
     for(j = 0 \text{ to } N-1)
          if(A[0][j] == 0)
               firstRow = true
     for(i = 0 \text{ to } M-1)
          if(A[i][0] == 0)
               firstCol = true
     for(i = 0 \text{ to } M-1)
     {
          for(j = 0 \text{ to } N-1)
          {
               if(A[i][j] == 0)
                    A[i][0] = 0
                    A[0][j] = 0
               }
          }
     }
   for (i = 1 \text{ to } M-1)
     {
          for (j = 1 \text{ to } N-1)
```

```
{
             if(A[0][j] == 0 || A[i][0] == 0)
                 A[i][j] = 0
         }
    }
   if(firstRow == true)
    {
         for( i = 0 to N-1 )
             A[0][i] = 0
    }
   if(firstCol == true)
    {
         for(i = 0 \text{ to } M-1)
             A[i][0] = 0
    }
}
```

#### **Complexity Analysis**

Time Complexity: Checking the first row and first column + Traversing and updating the matrix + Updating first row and first column.

```
= O(M) + O(N) + O(M*N) + O(M) + O(N)
= O(M*N)
Space Complexity: O(1)
```

#### Critical ideas to think!

- Why are the first row and first column updated separately?
- Can you use other data structures to efficiently solve this problem?

## **Comparison of different solutions**

| Approach                        | Time Complexity | Space Complexity |
|---------------------------------|-----------------|------------------|
| Brute Force: Using nested loops | O(M*N*(M+N))    | O(M*N)           |
| Using Hash Table                | O(M*N)          | O(M+N)           |
| Using In-Place Hashing          | O(M*N)          | O(1)             |

## Suggested problems to solve

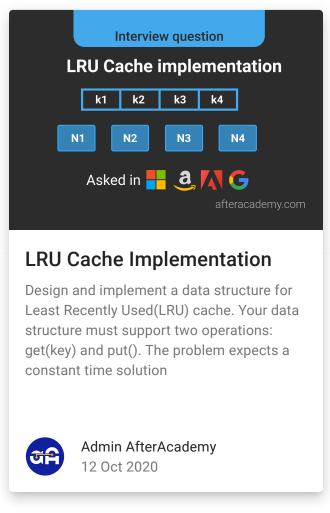
- Given a boolean matrix, find k such that all elements in the kth row are 0 and all elements in the kth column are 1
- Print unique rows in a given Boolean matrix
- Check if matrix can be converted to another matrix by transposing square sub-matrices
- Print a given matrix in spiral form
- Rotate a matrix by 90 degree

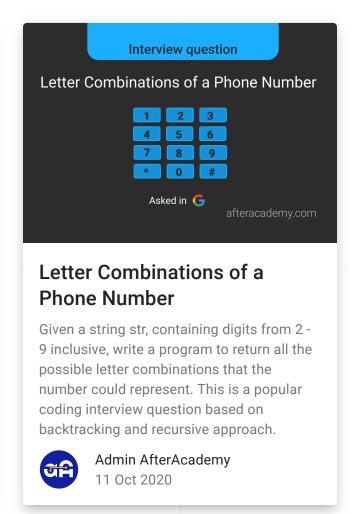
If you have any more approaches or you find an error/bug in the above solutions, please comment down below.

Happy Coding! Enjoy Algorithms!

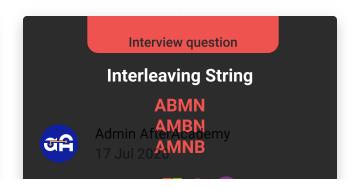
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Reverse a Stack Using

## Reverse a Stack Using Recursion

Given a stack of integers st, write a program to reverse the stack using recursion. This problem will clear the concept of recursion.



#### **Interleaving Strings**

Given three strings S1, S2 and S3, write a program which checks whether S3 is an interleaving of S1 and S2. The problem is a typical dynamic programming problem.



#### **Surrounded regions**

Given a 2-D matrix board where every element is either 'O' or 'X', write a program to replace 'O' with 'X' if surrounded by 'X'. It is a famous problem based on the concept of flood fill algorithms



Admin AfterAcademy 22 Sep 2020



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