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7.4.3. Loops in Assembly

Like if statements, loops in assembly are also implemented using jump instructions. However, loops enable instructions to be *revisited* based on the result of an evaluated condition.

The sumUp function shown in the following example sums up all the positive integers from 1 to a user-defined integer. This code is intentionally written suboptimally to illustrate a while loop in C.

```
int sumUp(int n) {
    //initialize total and i
    int total = 0;
    int i = 1;

while (i <= n) { //while i is less than or equal to n
        total += i; //add i to total
        i++; //increment i by 1
    }
    return total;
}</pre>
```

Compiling this code and disassembling it using GDB yields the following assembly code:

```
Dump of assembler code for function sumUp:
0x400526 <+0>:
                 push
                        %rbp
0x400527 <+1>:
                 mov
                        %rsp,%rbp
0x40052a <+4>:
                        %edi,-0x14(%rbp)
                 mov
0x40052d <+7>:
                 mov
                        $0x0, -0x8(%rbp)
0x400534 <+14>:
                         $0x1, -0x4(%rbp)
                 mov
0x40053b <+21>:
                 jmp
                        0x400547 < sumUp + 33 >
0x40053d <+23>:
                 mov
                        -0x4(%rbp),%eax
0x400540 <+26>:
                 add
                        %eax, -0x8(%rbp)
0x400543 <+29>:
                         $0x1,-0x4(%rbp)
                 add
0x400547 <+33>:
                        -0x4(%rbp),%eax
                 mov
0x40054a <+36>:
                         -0x14(%rbp), %eax
                 cmp
0x40054d <+39>:
                         0x40053d < sumUp + 23 >
                 jle
0x40054f <+41>:
                         -0x8(%rbp),%eax
                 mov
0x400552 <+44>:
                        %rbp
                 pop
0x400553 <+45>: retq
```

Again, we will not draw out the stack explicitly in this example. However, we encourage readers to draw the stack out themselves.

The First Five Instructions

The first five instructions of this function set the stack up for function execution and set up temporary values for function execution:

Recall that stack locations store *temporary variables* in a function. For simplicity we will refer to the location marked by $\protect\mbox{"rbp-0x8}$ as total and $\protect\mbox{"rbp}$ - $\protect\mbox{0x4}$ as i. The input parameter to sumUp (n) is moved to stack location $\protect\mbox{"rbp-0x14}$. Despite the placement of temporary variables on the stack, keep in mind that the stack pointer has not changed after the execution of the first instruction (i.e., push $\protect\mbox{"rbp}$).

The Heart of the Loop

The next seven instructions in the sumUp function represent the heart of the loop:

```
0x40053b <+21>:
                        0x400547 < sumUp+33> # qoto < sumUp+33>
                 jmp
0x40053d <+23>:
                        -0x4(%rbp),%eax
                                             # copy i to %eax
                mov
0x400540 <+26>:
                        %eax,-0x8(%rbp)
                                             # add i to total (total += i)
                add
0x400543 <+29>:
                        $0x1,-0x4(%rbp)
                                             # add 1 to i (i += 1)
                add
0x400547 <+33>:
                mov
                        -0x4(%rbp),%eax
                                             # copy i to %eax
0x40054a <+36>: cmp
                        -0x14(%rbp),%eax
                                            # compare i to n
0x40054d <+39>:
                jle
                        0x40053d < sumUp+23 > # if (i <= n) goto < sumUp+23 >
```

- The first instruction is a direct jump to <sumUp+33>, which sets the instruction pointer (%rip) to address 0x400547.
- The next instruction that executes is mov 0x4(%rbp), %eax, which places the value of i in register %eax. Register %rip is updated to 0x40054a.
- The cmp instruction at <sumUp+36> compares i to n and sets the appropriate condition code registers. Register %rip is set to 0x40054d.

The jle instruction then executes. The instructions that execute next depend on whether or not the branch is taken.

Suppose that the branch is taken (i.e., $i \le n$ is true). Then the instruction pointer is set to 0x40053d and program execution jumps to sumUp+23 > 0. The following instructions then execute in sequence:

- The mov instruction at <sumUp+23> copies i to register %eax.
- The add %eax, -0x8(%rbp) adds i to total (i.e., total += i).
- The add instruction at <sumUp+29> then adds 1 to i (i.e., i += 1).
- The mov instruction at <sumUp+33> copies the updated value of i to register %eax.
- The cmp instruction then compares i to n and sets the appropriate condition code registers.
- Next, jle executes. If i is less than or equal to n, program execution once again jumps to <sumUp+23> and the loop (defined between <sumUp+23> and <sumUp+39>) repeats.

If the branch is *not* taken (i.e., i is *not* less than or equal to n), the following instructions execute:

These instructions copy total to register %eax, restore %rbp to its original value, and exit the function. Thus, the function returns total upon exit.

Table 1 shows the assembly and C goto forms of the sumUp function:

Table 1. Translating sumUp into goto C form.

Assembly

Translated goto Form

```
<sumUp>:
  <+0>:
           push
                  %rbp
   <+1>:
                  %rsp,%rbp
           mov
   <+4>: mov
                  %edi, -0x14(%rbp)
   <+7>:
                  $0x0, -0x8(%rbp)
          mov
   <+14>:
                  $0x1, -0x4(%rbp)
          mov
                  0x400547
   <+21>:
           jmp
<sumUp+33>
   <+23>: mov
                  -0x4(%rbp),%eax
   <+26>: add
                 %eax, -0x8(%rbp)
   <+29>: add
                  $0x1, -0x4(%rbp)
   <+33>: mov
                 -0x4(%rbp),%eax
   <+36>: cmp
                  -0x14(%rbp), %eax
   <+39>: jle
                 0x40053d
<sumUp+23>
   <+41>: mov
                  -0x8(%rbp),%eax
   <+44>: pop
                  %rbp
   <+45>: retq
```

```
int sumUp(int n) {
    int total = 0;
    int i = 1;
    goto start;

body:
    total += i;
    i += 1;

start:
    if (i <= n) {
        goto body;
    }
    return total;
}</pre>
```

The preceding code is also equivalent to the following C code without goto statements:

```
int sumUp(int n) {
   int total = 0;
   int i = 1;
   while (i <= n) {
      total += i;
      i += 1;
   }
   return total;
}</pre>
```

for Loops in Assembly

The primary loop in the sumUp function can also be written as a for loop:

```
int sumUp2(int n) {
  int total = 0;  //initialize total to 0
```

This version yields assembly code identical to our while loop example. We repeat the assembly code below and annotate each line with its English translation:

```
Dump of assembler code for function sumUp2:
                        %rbp
0x400554 <+0>:
                                               #save %rbp
                 push
0x400555 <+1>:
                 mov
                        %rsp,%rbp
                                               #update %rpb (new stack
frame)
0x400558 <+4>:
                 mov
                        %edi,-0x14(%rbp)
                                               #copy %edi to %rbp-0x14 (n)
                        $0x0,-0x8(%rbp)
0x40055b <+7>:
                                               #copy 0 to %rbp-0x8 (total)
                 movl
0x400562 <+14>:
                 movl
                        9x1, -0x4(%rbp)
                                               #copy 1 to %rbp-0x4 (i)
0x400569 <+21>:
                 jmp
                        0x400575 < sumUp2+33>
                                               #goto <sumUp2+33>
0x40056b <+23>:
                 mov
                       -0x4(%rbp),%eax
                                               #copy i to %eax [loop]
0x40056e <+26>:
                        %eax,-0x8(%rbp)
                                               #add i to total (total+=i)
                 add
0x400571 <+29>:
                                               #add 1 to i (i++)
                 addl
                        0x1, -0x4(%rbp)
0x400575 <+33>:
                        -0x4(%rbp),%eax
                 mov
                                               #copy i to %eax [start]
0x400578 <+36>:
                cmp
                        -0x14(%rbp), %eax
                                               #compare i with n
0x40057b <+39>:
                        0x40056b < sumUp2+23>
                                               #if (i <= n) goto loop
                jle
0x40057d <+41>:
                        -0x8(%rbp),%eax
                                               #copy total to %eax
                 mov
0x400580 <+44>:
                                               #prepare to leave the
                pop
                        %rbp
function
0x400581 < +45>: retq
                                               #return total
```

To understand why the for loop version of this code results in identical assembly to the while loop version of the code, recall that the for loop has the following representation:

```
for ( <initialization>; <boolean expression>; <step> ){
      <body>
}
```

and is equivalent to the following while loop representation:

Since <u>every for loop can be represented by a while loop</u>, the following two C programs are equivalent representations for the previous assembly:

Table 2. Equivalent ways to write the sumUp function.

For loop While loop

```
С
int sumUp2(int n) {
                                        int sumUp(int n){
    int total = 0;
                                            int total = 0;
    int i = 1;
                                            int i = 1;
    for (i; i <= n; i++) {
                                            while (i \le n) {
        total += i;
                                                total += i;
                                                i += 1;
    return total;
                                            return total;
}
                                        }
```

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7.4.3. Loops in Assembly

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