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rzaharia Updated all readme files to contain links to the next step

2 years ago



284 lines (231 loc) · 8.66 KB

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# Part 50: Mopping Up, part 1

We have definitely reached the "mopping up" phase, as in this part of our compiler writing journey I don't introduce any major feature. Instead, I fix a couple of problems and add a couple of minor functions.

## Consecutive Cases

At present, the compiler can't parse

```
switch(x) {  
  case 1:  
  case 2: printf("Hello\n");  
}
```



because the parser expects a compound statement after the ':' token. In

switch\_statement() in stmt.c :

```
// Scan the ':' and get the compound expression  
match(T_COLON, ":");  
left= compound_statement(1); casecount++;  
...  
// Build a sub-tree with the compound statement as the left child  
casetail->right= mkastunary(ASTop, 0, left, NULL, casevalue);
```



What we want is to allow an empty compound statement, so that any case with a missing compound statement falls down into the next existing compound statement.

The change in `switch_statement()` is:

```
// Scan the ':' and increment the casecount
match(T_COLON, ":");
casecount++;

// If the next token is a T_CASE, the existing case will fall
// into the next case. Otherwise, parse the case body.
if (Token.token == T_CASE)
    body= NULL;
else
    body= compound_statement(1);
```



This is, however, only half the story. Now in the code generation section, we have to catch the NULL compound statement and do something about it. In `genSWITCH()` in `gen.c`:

```
// Walk the right-child linked list to
// generate the code for each case
for (i = 0, c = n->right; c != NULL; i++, c = c->right) {
    ...
    // Generate the case code. Pass in the end label for the breaks.
    // If case has no body, we will fall into the following body.
    if (c->left) genAST(c->left, NOLABEL, NOLABEL, Lend, 0);
    genfreeregs(NOREG);
}
```



So, this was a nice and simple fix. `tests/input123.c` is the test program to confirm this change works.

## Dumping the Symbol Table

While I was trying to work out why the global `Text` variable wasn't visible to the compiler, I added code in `sym.c` to dump the symbol table at the end of every source code file. There is an `-M` command line argument to enable the functionality. I won't go through the code, but here is an example of its output:

```
Symbols for misc.c
Global
-----
void exit(): global, 1 params
```



```

    int status: param, size 4
void _Exit(): global, 1 params
    int status: param, size 4
void *malloc(): global, 1 params
    int size: param, size 4
...
int Line: extern, size 4
int Putback: extern, size 4
struct symtable *Functionid: extern, size 8
char **Infile: extern, size 8
char **Outfile: extern, size 8
char *Text[]: extern, 513 elems, size 513
struct symtable *Globhead: extern, size 8
struct symtable *Globtail: extern, size 8
...
struct mkastleaf *mkastleaf(): global, 4 params
    int op: param, size 4
    int type: param, size 4
    struct symtable *sym: param, size 8
    int intvalue: param, size 4
...
Enums
-----
int (null): enumtype, size 0
int TEXTLEN: enumval, value 512
int (null): enumtype, size 0
int T_EOF: enumval, value 0
int T_ASSIGN: enumval, value 1
int T_ASPLUS: enumval, value 2
int T_ASMINUS: enumval, value 3
int T_ASSTAR: enumval, value 4
int T_ASSSLASH: enumval, value 5
...
Typedefs
-----
long size_t: typedef, size 0
char *FILE: typedef, size 0

```

## Passing Arrays as Arguments

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I made the following change, but in hindsight I realise that I probably need to rethink how I deal with arrays completely. Anyway ... when I compile `decl.c` with the compiler, I get the error:

Unknown variable:Text on line 87 of decl.c



which prompted me to write the symbol dumping code. `Text` is in the global symbol table, so why is the parser complaining that it's missing?

The answer is that `postfix()` in `expr.c`, after finding an identifier, consults the following token. If it is a `'['`, then the identifier must be an array. If there is no `'['`, then the identifier must be a variable:

```
// A variable. Check that the variable exists.  
if ((varptr = findsymbol(Text)) == NULL || varptr->stype != S_VARIABLE)  
    fatals("Unknown variable", Text);
```



This is preventing the passing of an array reference as an argument to a function. The "offending" line that prompts the error message is in `decl.c`:

```
type = type_of_typedef(Text, ctype);
```



We are passing the address of the base of `Text` as an argument. But with no following `'['`, our compiler thinks that it's a scalar variable, and complains that there is no scalar variable `Text`.

I made the change to allow `S_ARRAY` as well as `S_VARIABLE` here, but this is just the tip of a bigger problem: arrays and pointers in our compiler are not as interchangeable as they should be. I'll tackle this in the next part.

## Missing Operators

In our compiler, we've had these tokens and AST operators since part 21 of the journey:

- `||`, `T_LOGOR`, `A_LOGOR`
- `&&`, `T_LOGAND`, `A_LOGAND`

Somehow, I'd never implemented them! So, it's time to do them.

For `A_LOGAND`, we have two expressions. If both evaluate to true, we need to set a register to the rvalue of 1, otherwise 0. For `A_LOGOR`, if either evaluate to true, we need to set a register to the rvalue of 1, otherwise 0.

The `binexpr()` code in `expr.c` already parses the tokens and builds the `A_LOGOR` and `A_LOGAND` AST nodes. So we need to fix up the code generator.

In `genAST()` in `gen.c`, we now have:

```

case A_LOGOR:
    return (cglogor(leftreg, rightreg));
case A_LOGAND:
    return (cglogand(leftreg, rightreg));

```



with two corresponding functions in `cg.c`. Before we look at the `cg.c` functions, let's just see an example C expression and the assembly code that will be produced.

```

int x, y, z;
...
z= x || y;

```



when compiled, results in:

```

    movslq  x(%rip), %r10      # Load x's rvalue
    movslq  y(%rip), %r11      # Load y's rvalue
    test    %r10, %r10         # Test x's boolean value
    jne     L13                # True, jump to L13
    test    %r11, %r11         # Test y's boolean value
    jne     L13                # True, jump to L13
    movq    $0, %r10           # Neither true, set %r10 to false
    jmp     L14                # and jump to L14
L13:
    movq    $1, %r10           # Set %r10 to true
L14:
    movl    %r10d, z(%rip)     # Save boolean result to z

```



We test each expression, jump based on the boolean result and either store 0 or 1 into our output register. The assembly for `A_LOGAND` is similar, except that the conditional jumps are `je` (jump if equal to zero) and the `movq $0` and `movq $1` are swapped around.

So, without further comment, are the new `cg.c` functions:

```

// Logically OR two registers and return a
// register with the result, 1 or 0
int cglogor(int r1, int r2) {
    // Generate two labels
    int Ltrue = genlabel();
    int Lend = genlabel();

    // Test r1 and jump to true label if true
    fprintf(Outfile, "\ttest\t%s, %s\n", reglist[r1], reglist[r1]);
    fprintf(Outfile, "\tjne\tL%d\n", Ltrue);
}

```



```

// Test r2 and jump to true label if true
fprintf(Outfile, "\ttest\t%s, %s\n", reglist[r2], reglist[r2]);
fprintf(Outfile, "\tjne\tL%d\n", Ltrue);

// Didn't jump, so result is false
fprintf(Outfile, "\tmovq\t$0, %s\n", reglist[r1]);
fprintf(Outfile, "\tjmp\tL%d\n", Lend);

// Someone jumped to the true label, so result is true
cglabel(Ltrue);
fprintf(Outfile, "\tmovq\t$1, %s\n", reglist[r1]);
cglabel(Lend);
free_register(r2);
return(r1);
}

```

```

// Logically AND two registers and return a
// register with the result, 1 or 0
int cglogand(int r1, int r2) {
    // Generate two labels
    int Lfalse = genlabel();
    int Lend = genlabel();

    // Test r1 and jump to false label if not true
    fprintf(Outfile, "\ttest\t%s, %s\n", reglist[r1], reglist[r1]);
    fprintf(Outfile, "\tjne\tL%d\n", Lfalse);

    // Test r2 and jump to false label if not true
    fprintf(Outfile, "\ttest\t%s, %s\n", reglist[r2], reglist[r2]);
    fprintf(Outfile, "\tjne\tL%d\n", Lfalse);

    // Didn't jump, so result is true
    fprintf(Outfile, "\tmovq\t$1, %s\n", reglist[r1]);
    fprintf(Outfile, "\tjmp\tL%d\n", Lend);

    // Someone jumped to the false label, so result is false
    cglabel(Lfalse);
    fprintf(Outfile, "\tmovq\t$0, %s\n", reglist[r1]);
    cglabel(Lend);
    free_register(r2);
    return(r1);
}

```



The program `tests/input122.c` is the test to confirm that this new functionality works.

## Conclusion and What's Next

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So that's a few small things fixed up in this part of our journey. What I will do now is step back, rethink the array/pointer design and try to fix this up in the next part of our compiler writing journey. [Next step](#)