# Algorithms: CSE 202 — Homework 3

## Problem 1: Graph cohesiveness (KT 7.46)

In sociology, one often studies a graph G in which nodes represent people and edges represent those who are friends with each other. Let's assume for purposes of this question that friendship is symmetric, so we can consider an undirected graph.

Now suppose we want to study this graph G, looking for a "close-knit" group of people. One way to formalize this notion would be as follows. For a non-empty subset S of nodes, let e(S) denote the number of edges in S-that is, the number of edges that have both ends in S. We define the *cohesiveness* of S as e(S)/|S|. A natural thing to search for would be a set S of people achieving the maximum cohesiveness.

- 1. Give a polynomial-time algorithm that takes a rational number  $\alpha$  and determines whether there exists a set S with cohesiveness greater than  $\alpha$ .
- 2. Give a polynomial-time algorithm to find a set S of nodes with maximum cohesiveness.

### Problem 2: Remote Sensors

Devise as efficient as possible algorithm for the following problem. You have n remote sensors  $s_i$  and m < n base stations  $B_j$ . For  $1 \le j \le m$ , base station  $B_j$  is located at  $(x_j, y_j)$  in the two-dimensional plane. You are given that no two base-stations are less than 1 km apart (in standard Euclidean distance,  $\sqrt{((x_j - x_k)^2 + (y_j - y_k)^2)}$ ). All base stations have the same integer bandwidth capacity C.

For  $1 \le i \le n$ , sensor  $s_i$  is located at  $(x_i, y_i)$  in the two-dimensional plane and has an integer bandwidth requirement of  $r_i$ , which can be met by assigning bandwidth on multiple base stations. Let  $b_{i,j}$  be the amount of bandwidth assigned to sensor  $s_i$  on base station  $B_j$ . The assignment must meet the following constraints:

- No sensor may be assigned any bandwidth on a base station more than 2 km distance from it, i.e., if the distance from  $s_i$  to  $B_j$  is greater than 2,  $b_{i,j} = 0$ .
- The sum of all the bandwidth assigned to any remote sensor  $s_i$  must be at least  $r_i$ : for each  $1 \le i \le n$ ,  $\sum_i b_{i,j} \ge r_i$ .
- The sum of all bandwidth assigned on base station  $B_j$  must be at most C: for each  $1 \leq j \leq m$ ,  $\sum_{i} b_{i,j} < C$ .

Your algorithm should find a solution meeting the above constraints if possible, and otherwise output a message saying "No solution exists". Prove the correctness of your algorithm and discuss its time compplexity.

#### Problem 3: Scheduling in a medical consulting firm (KT 7.19)

You've periodically helped the medical consulting firm Doctors Without Weekends on various hospital scheduling issues, and they've just come to you with a new problem. For each of the next n days, the hospital has determined the number of doctors they want on hand; thus, on day i, they have a requirement that exactly  $p_i$  doctors be present.

There are k doctors, and each is asked to provide a list of days on which he or she is willing to work. Thus doctor j provides a set  $L_j$  of days on which he or she is willing to work. The system produced by the consulting firm should take these lists and try to return to each doctor j a list  $L'_i$  with the following properties.

- (A)  $L'_j$  is a subset of  $L_j$ , so that doctor j only works on days he or she finds acceptable.
- (B) If we consider the whole set of lists  $L'_1, \ldots, L'_k$ , it causes exactly  $p_i$  doctors to be present on day i, for  $i = 1, 2, \ldots, n$ .
  - 1. Describe a polynomial-time algorithm that implements this system. Specifically, give a polynomial-time algorithm that takes the numbers  $p_1, p_2, \ldots, p_n$ , and the lists  $L_1, \ldots, L_k$ , and does one of the following two things.
    - Return lists  $L'_1, L'_2, \dots, L'_k$  satisfying properties (A) and (B); or
    - Report (correctly) that there is no set of lists  $L'_1, L'_2, \ldots, L'_k$  that satisfies both properties (A) and (B).
  - 2. The hospital finds that the doctors tend to submit lists that are much too restrictive, and so it often happens that the system reports (correctly, but unfortunately) that no acceptable set of lists  $L'_1, L'_2, \ldots, L'_k$  exists.

Thus the hospital relaxes the requirements as follows. They add a new parameter c > 0, and the system now should try to return to each doctor j a list  $L'_{i}$  with the following properties.

- $(A^*)$   $L'_i$  contains at most c days that do not appear on the list  $L_i$ .
- (B) (Same as before) If we consider the whole set of lists  $L'_1, \ldots, L'_k$ , it causes exactly  $p_i$  doctors to be present on day i, for  $i = 1, 2, \ldots, n$ .

Describe a polynomial-time algorithm that implements this revised system. It should take the numbers  $p_1, p_2, \ldots, p_n$ , the lists  $L_1, \ldots, L_k$ , and the parameter c > 0, and do one of the following two things.

- Return lists  $L'_1, L'_2, \dots, L'_k$  satisfying properties  $(A^*)$  and (B); or
- Report (correctly) that there is no set of lists  $L'_1, L'_2, \ldots, L'_k$  that satisfies both properties  $(A^*)$  and (B).

## Problem 4: Cellular network

Consider the problem of selecting nodes for a cellular network. Any number of nodes can be chosen from a finite set of potential locations. We know the cost  $c_i \geq 0$  of establishing site i. If sites i and j are selected as nodes, then we derive the benefit  $b_{ij}$ , which is the revenue generated by the traffic between the two nodes. Both the benefits and costs are non-negative integers. Find an efficient algorithm to determine the subset of sites as the nodes for the cellular network such that the sum of the benefits provided by the edges between the selected nodes less the selected node costs is as large as possible.

Design an efficient polynomial-time algorithm.

Provide a high-level description of your algorithm, prove its correctness, and analyze its time complexity.