

Writing a competitive BZip2 encoder in Ada, from scratch, in a few days

<https://alire.ada.dev/crates/zipada>

Web links

Zip-Ada

Zip-Ada free, open-source, full Ada data compression library

Alire Crate: <https://alire.ada.dev/crates/zipada>



Home page: <https://unzip-ada.sourceforge.io/>

Sources, site #1: <https://github.com/zertovitch/zip-ada>

Sources, site #2: <https://sourceforge.net/projects/unzip-ada/>

Motivations for a BZip2 encoder:

- fun / challenge / warm-up for Advent of Code 2024
- weather
- someone did a documentation in 2016 (20 years after the software!)
- fill a gap in the Zip-Ada compatibility grid:

		Zip-Ada	
Format	Format #	Compress	Decompress
Store	0	v.22	v.1
Shrink	1	v.22	v.1
Reduce 1 .. 4	2 .. 5	v.29	v.1
Implode	6	never	v.1
Deflate	8	v.50 (v.40-49: limited)	v.1
Enhanced Deflate	9	never	v.30
BZip2	12	v.60	v.36
LZMA	14	v.51	v.47
PPMd	98		
Zstandard	93		

Expectations (low):

- BZip2 compresses few kinds of files better than, for instance, LZMA
- BZip2 compression scheme is mostly “mechanical”: on most steps, there is only one single possible encoding
- BZip2 is a weakened version of BZip1 (old patent issues).

Results: two very good surprises!

BZip2 is very simple.

1. Input: a “large” block of data (≤ 900 KB)
2. The entire block is processed “**off-line**”
 - Run Length Encoding (2x)
 - Burrows-Wheeler Transform (**b**lock-sorting)
 - Move To Front
 - Entropy coding (Huffman)
3. Output of the compressed block.

BZip2 is very simple.

→ simple source code as well (separate steps):

```

procedure Encode_Block (dyn_block_capacity : Natural_32) is
    ...
begin
    -- Data acquisition and transformation (no output):
    RLE_1;
    BWT;
    MTF_and_RLE_2;
    Entropy_Calculations;

    -- Now we output the block's compressed data:
    Put_Block_Header;
    Put_Block_Trees_Descriptors;
    Entropy_Output;
end Encode_Block;

```

Run Length Encoding #1

a	→	a	1	→	1
aa	→	aa	2	→	2
aaa	→	aaa	3	→	3
aaaa	→	aaaa[0]	4	→	5
aaaaa	→	aaaa[1]	5	→	5
aaaaaa	→	aaaa[2]	6	→	5
...			...	→	5
			259	→	5

Burrows-Wheeler Transform

Mary had a little lamb, its fleece was white as snow
 ary had a little lamb, its fleece was white as snowM
 ry had a little lamb, its fleece was white as snowMa
 y had a little lamb, its fleece was white as snowMar
 had a little lamb, its fleece was white as snowMary
 had a little lamb, its fleece was white as snowMary
 ad a little lamb, its fleece was white as snowMary h
 d a little lamb, its fleece was white as snowMary ha
 a little lamb, its fleece was white as snowMary had

...

a little lamb, its fleece was white as snowMary had
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 fleece was white as snowMary had a little lamb, it
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 was white as snowMary had a little lamb, its fleec

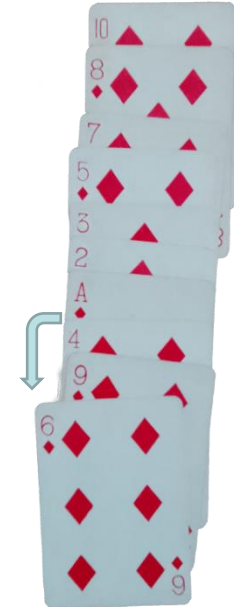
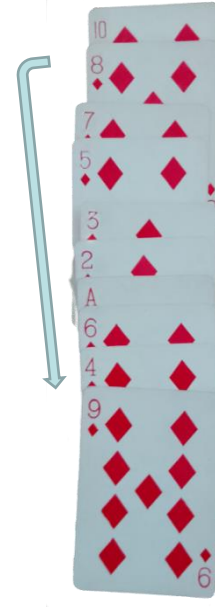
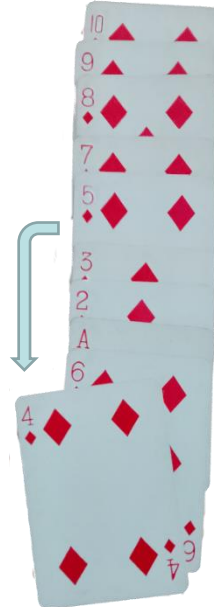
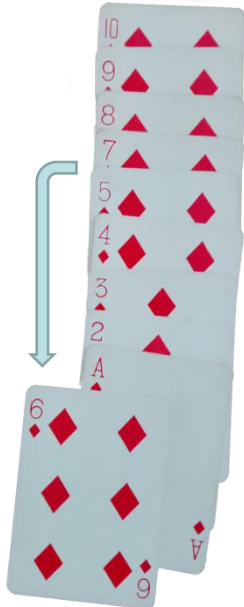
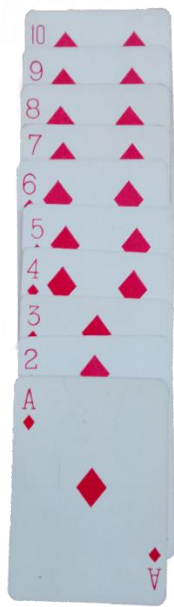
...

Sorting

Reversible!

[illegible]

Move To Front



What we have to send →

Card: “6”

Card: “4”

Card: “9”

Card: “6”

What we actually send →

Index: 6

Index: 5

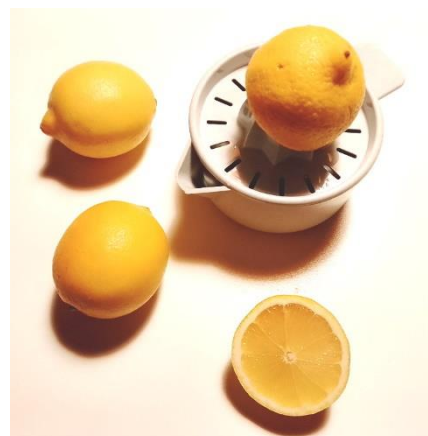
Index: 9

Index: 3

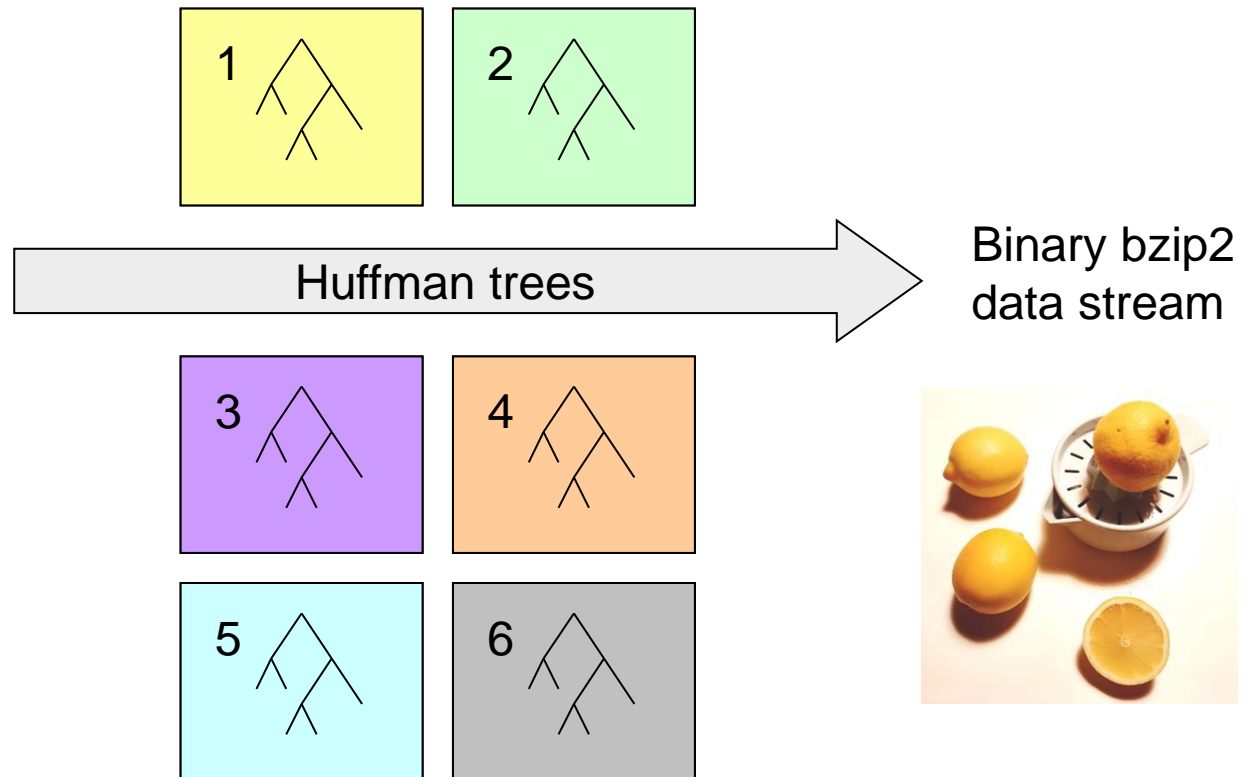
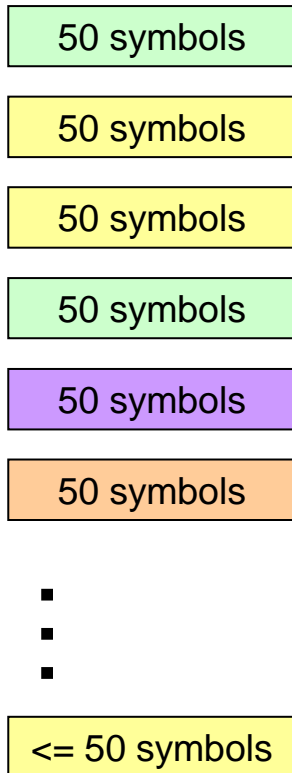
Final step: entropy coding with Huffman trees_u

Not mechanical. You have up to 6 trees, *freely* defined, that can be *freely* chosen for each string of 50 symbols (the output of Move To Front)

→ Room for **optimization!**

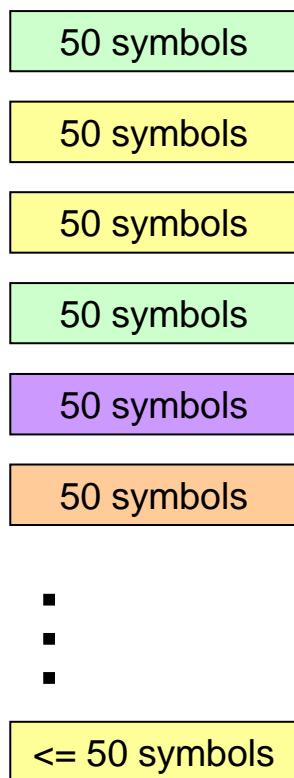


Entropy coding choices



Each symbol is of an alphabet of max 258 elements

Choice of Huffman trees



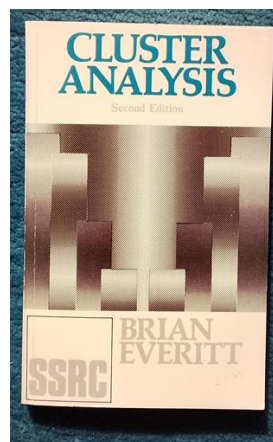
How to allocate the strings of data to the Huffman trees?

For shortest encoding, depends on the Huffman trees.

How to define convenient Huffman trees?

Depends on the data: the strings.

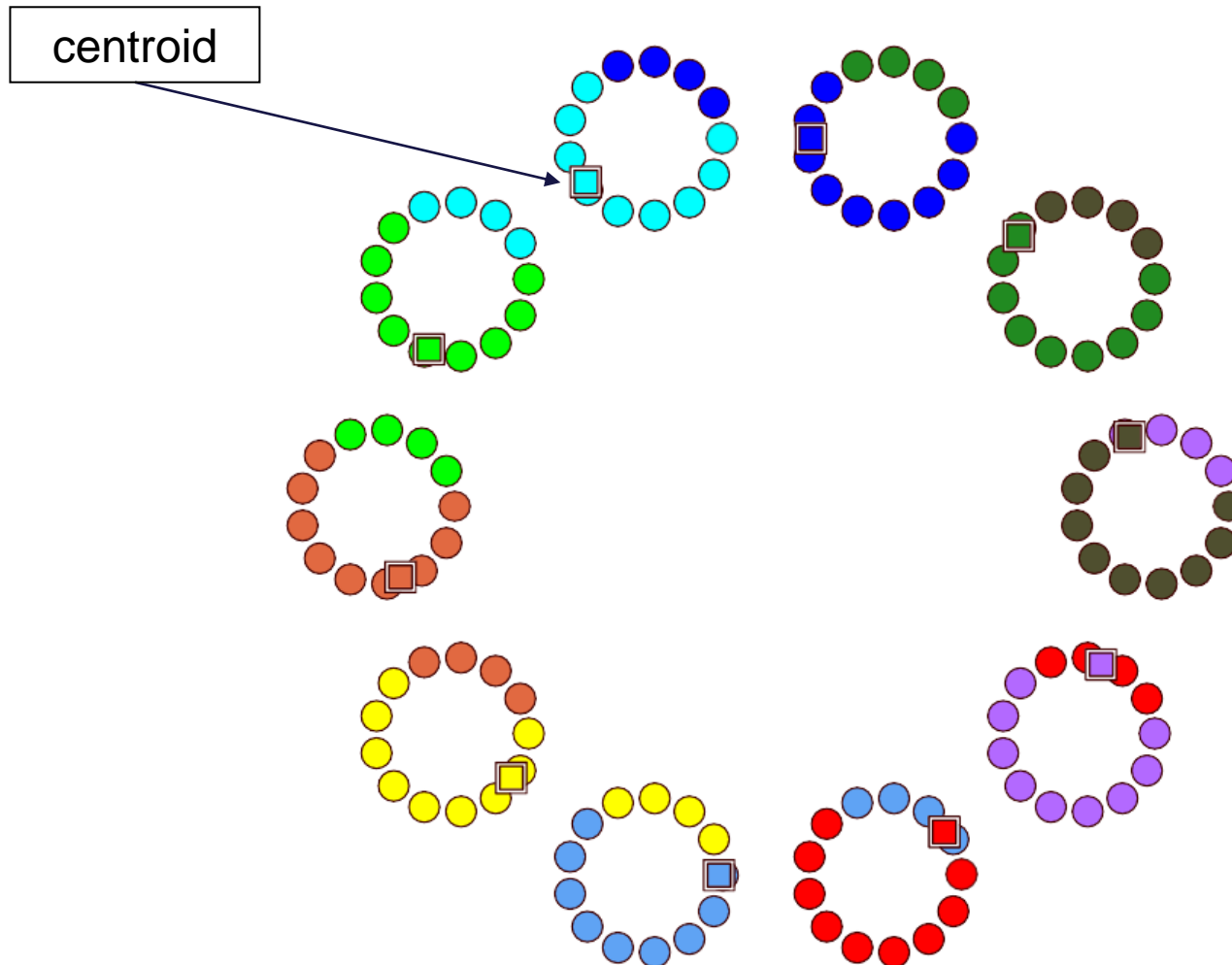
Actually, it's a



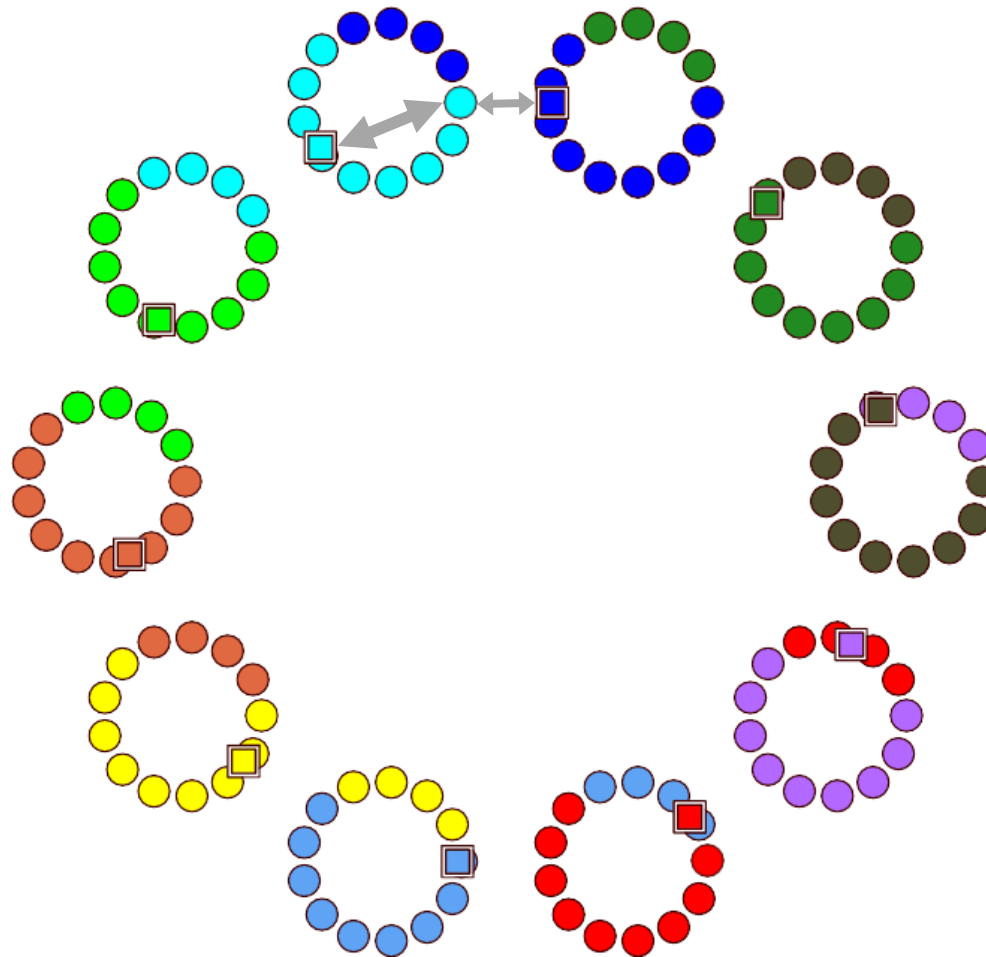
problem!

Partitioning Clustering: initial clustering and n iterations of optimization

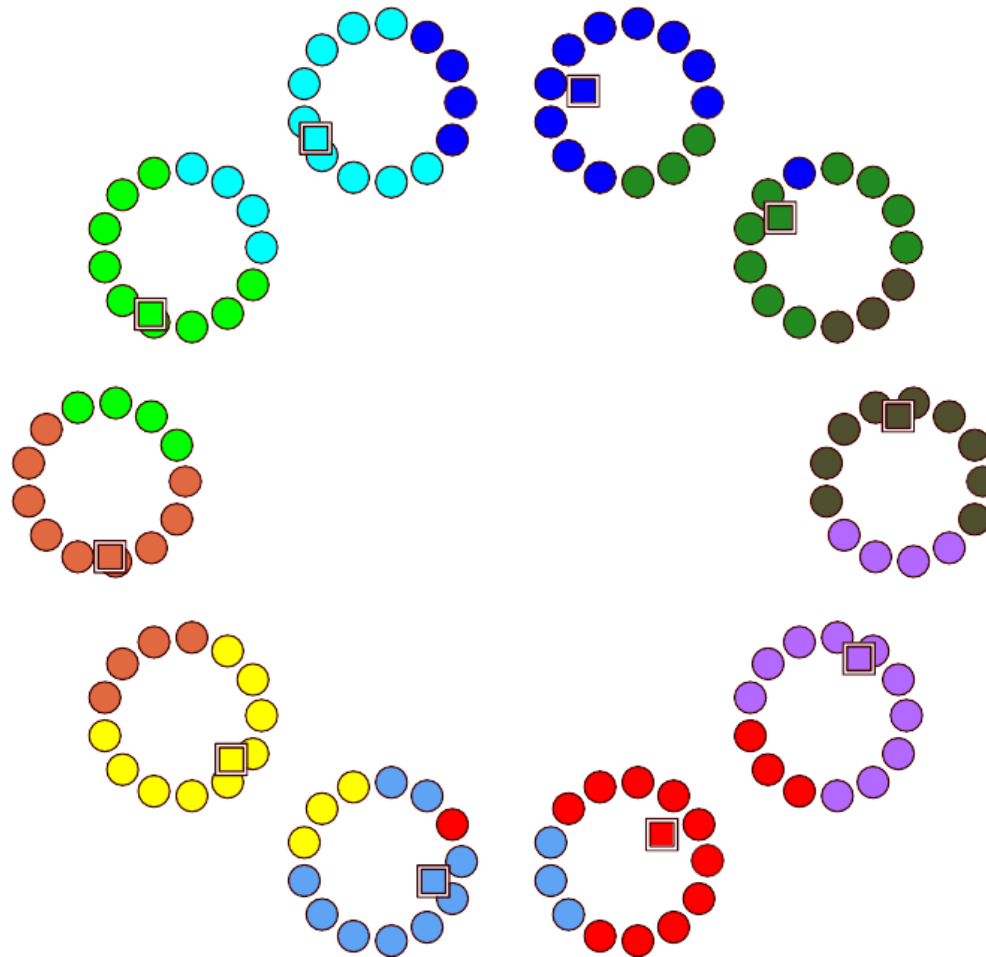
Clustering: 2D example. Initial (“imperfect”) clustering.



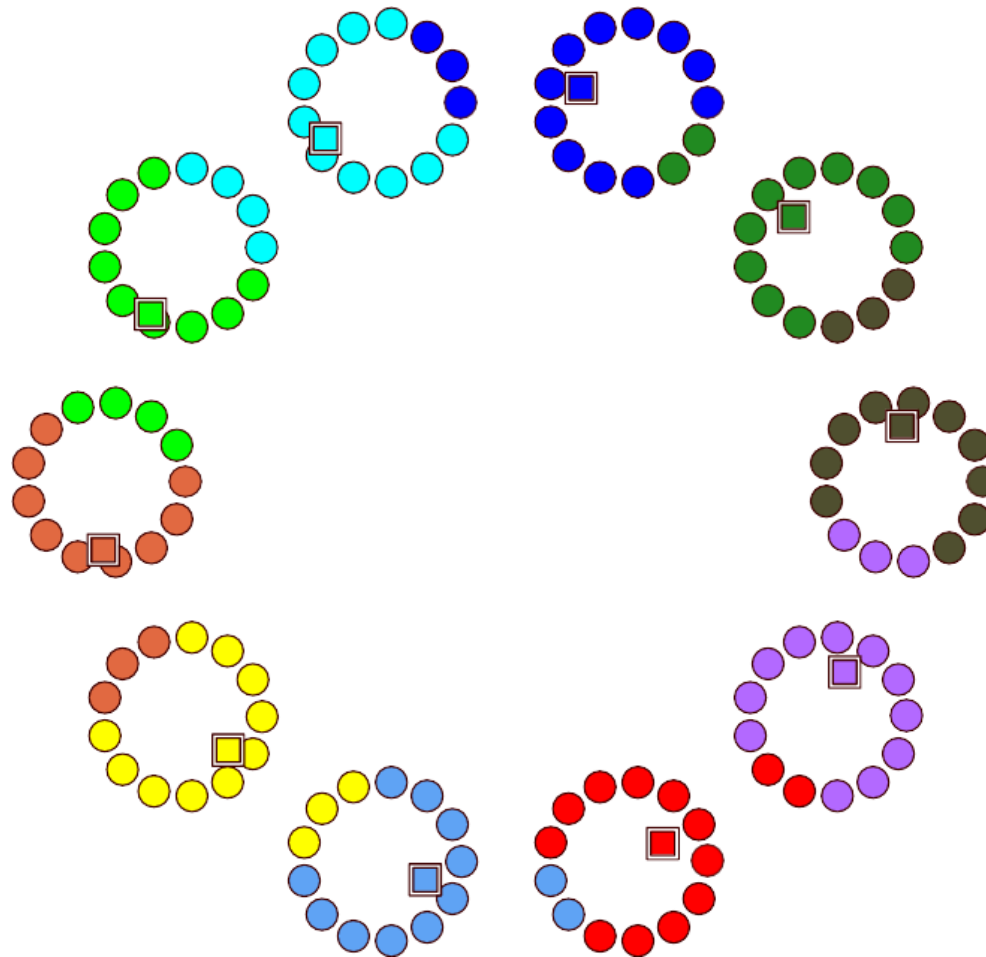
Clustering: 2D example. First reallocation round.



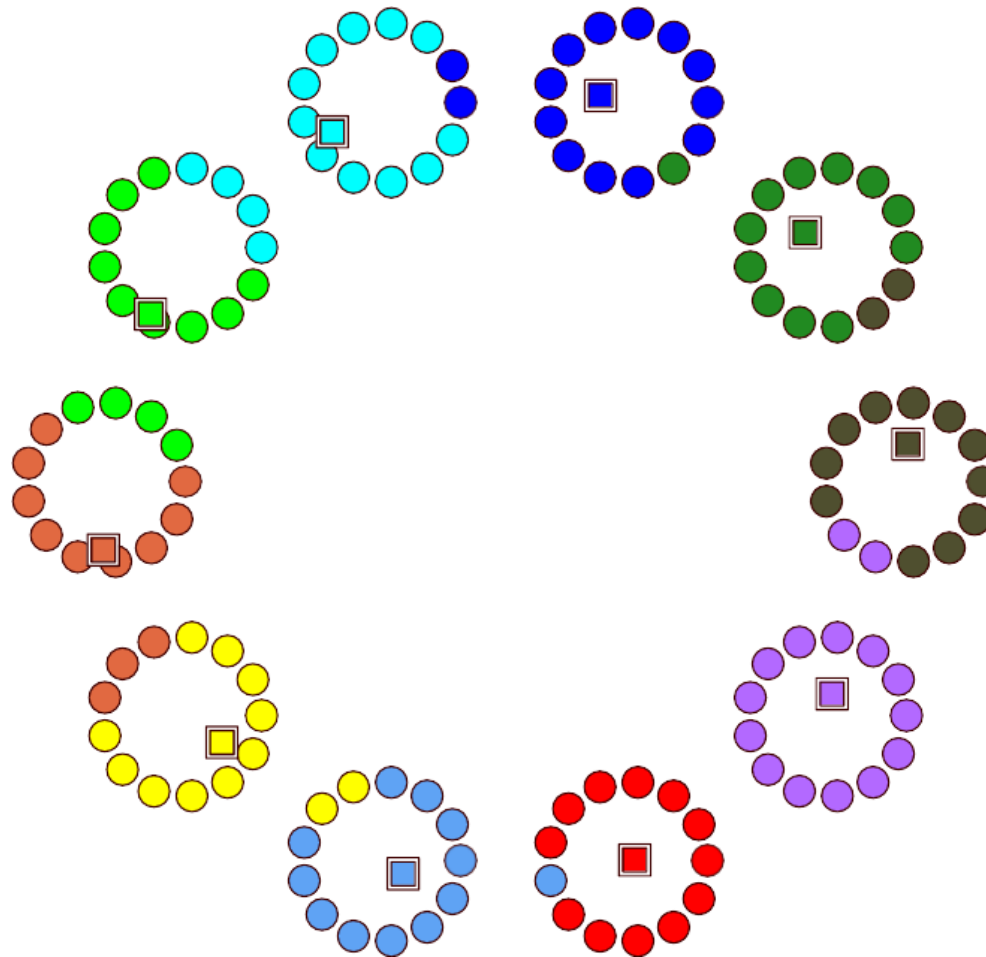
Clustering: 2D example. k-means, iteration 1



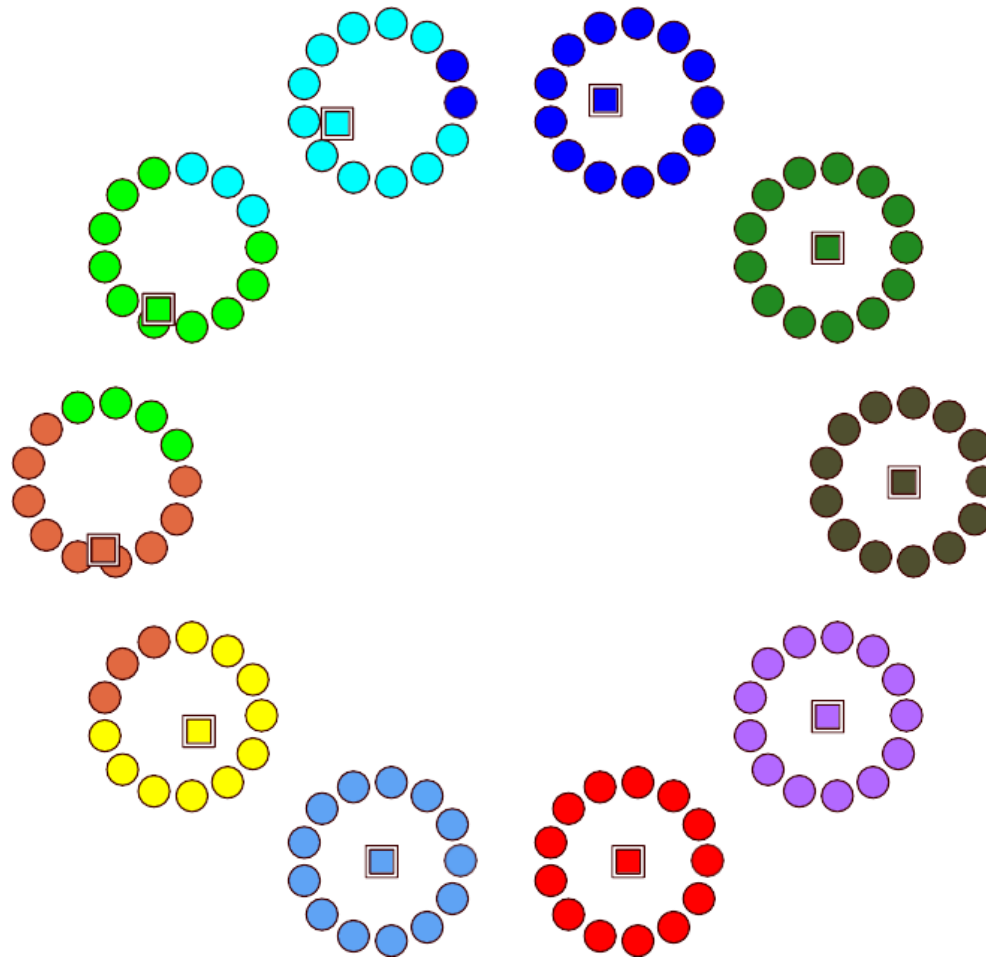
Clustering: 2D example. k-means, iteration 2



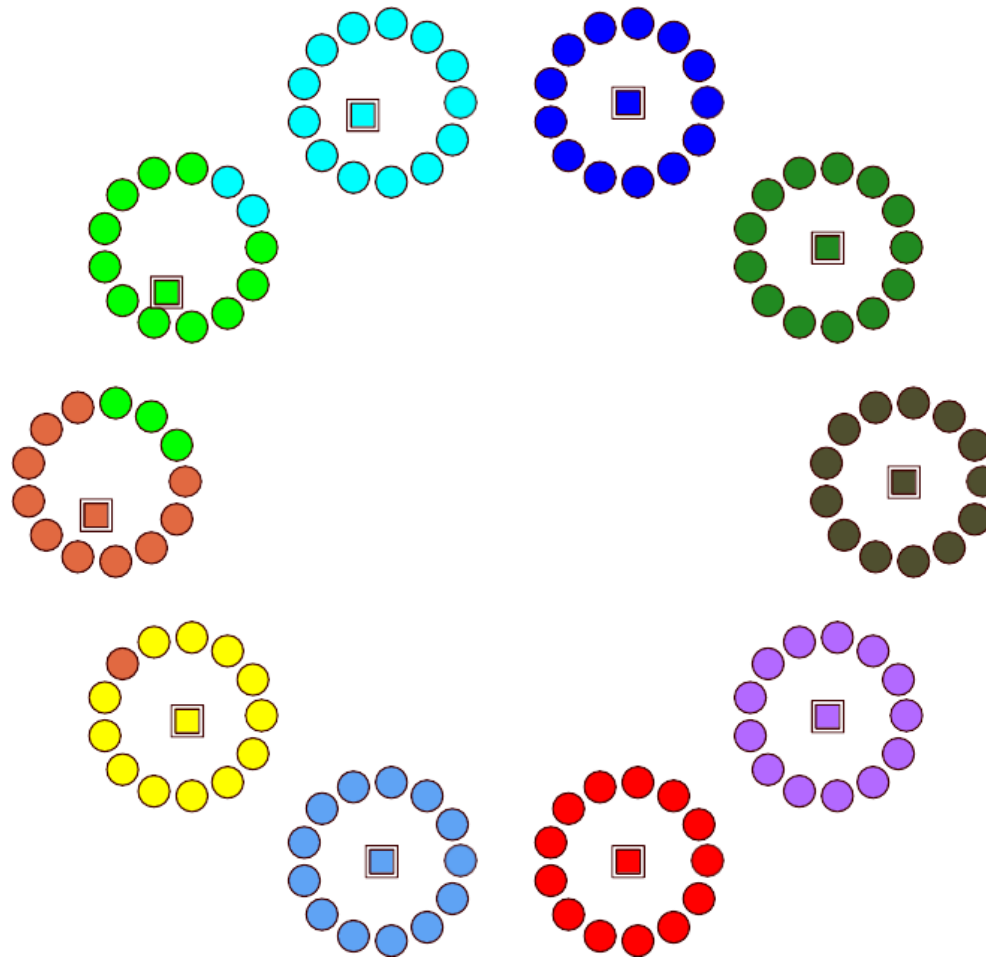
Clustering: 2D example. k-means, iteration 3



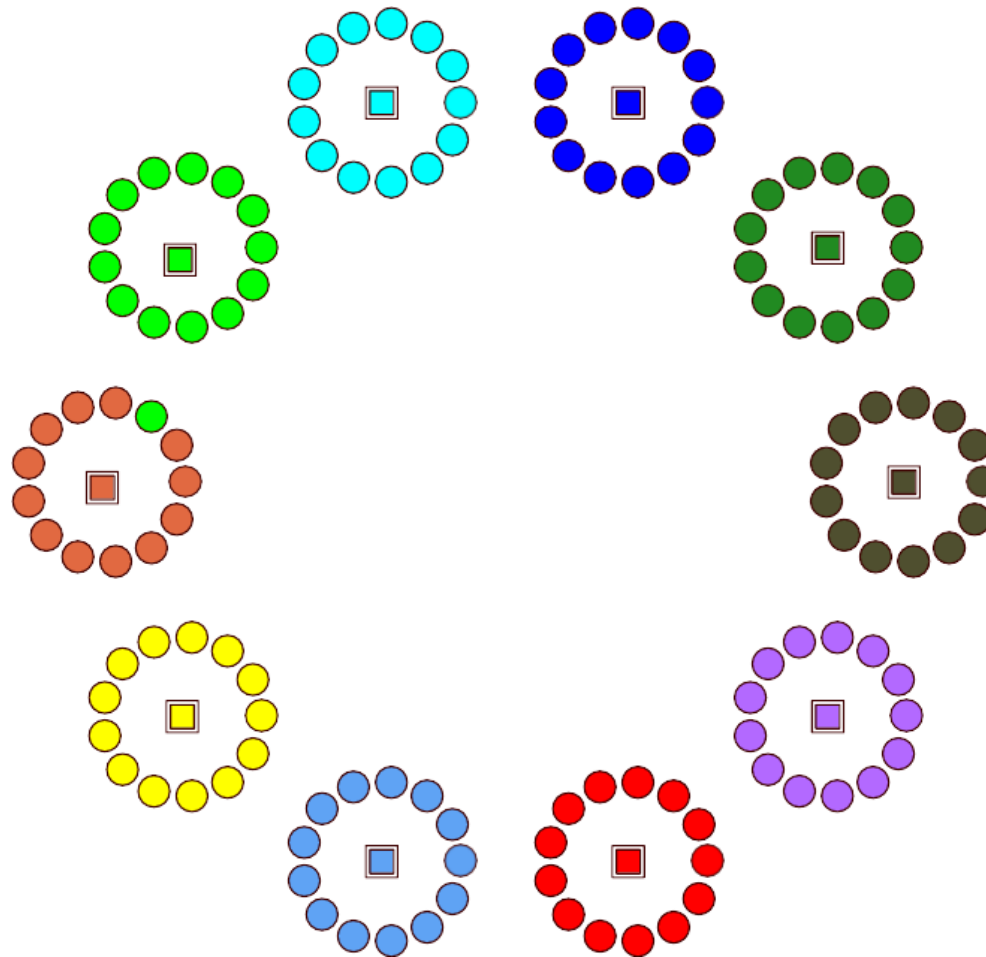
Clustering: 2D example. k-means, iteration 4



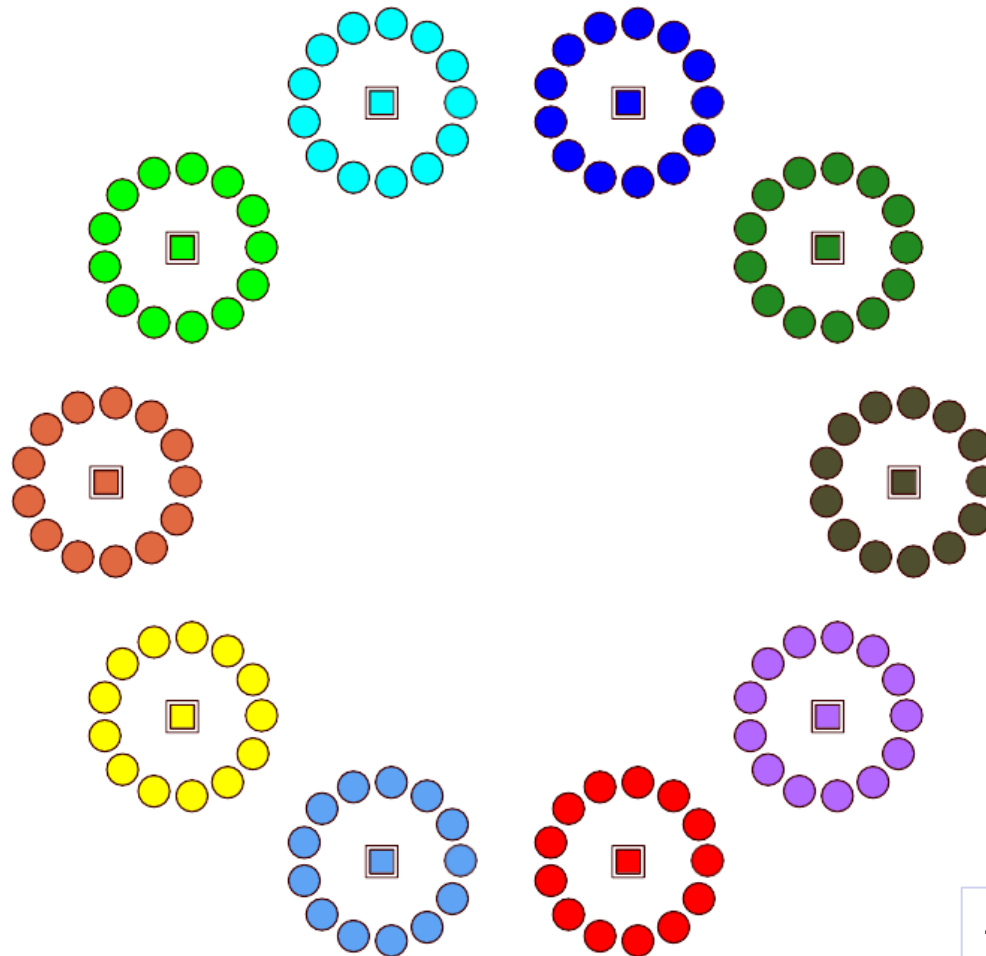
Clustering: 2D example. k-means, iteration 5



Clustering: 2D example. k-means, iteration 6

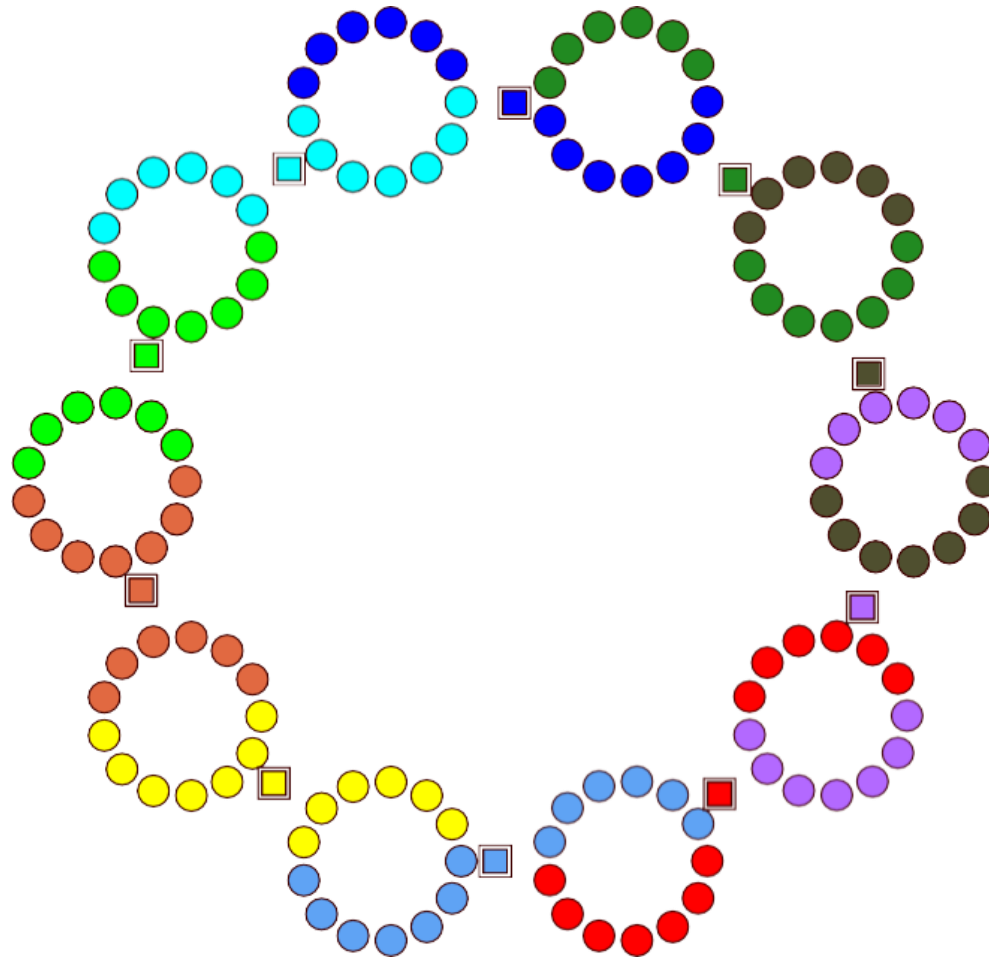


Clustering: 2D example. k-means, iteration 7 (final)

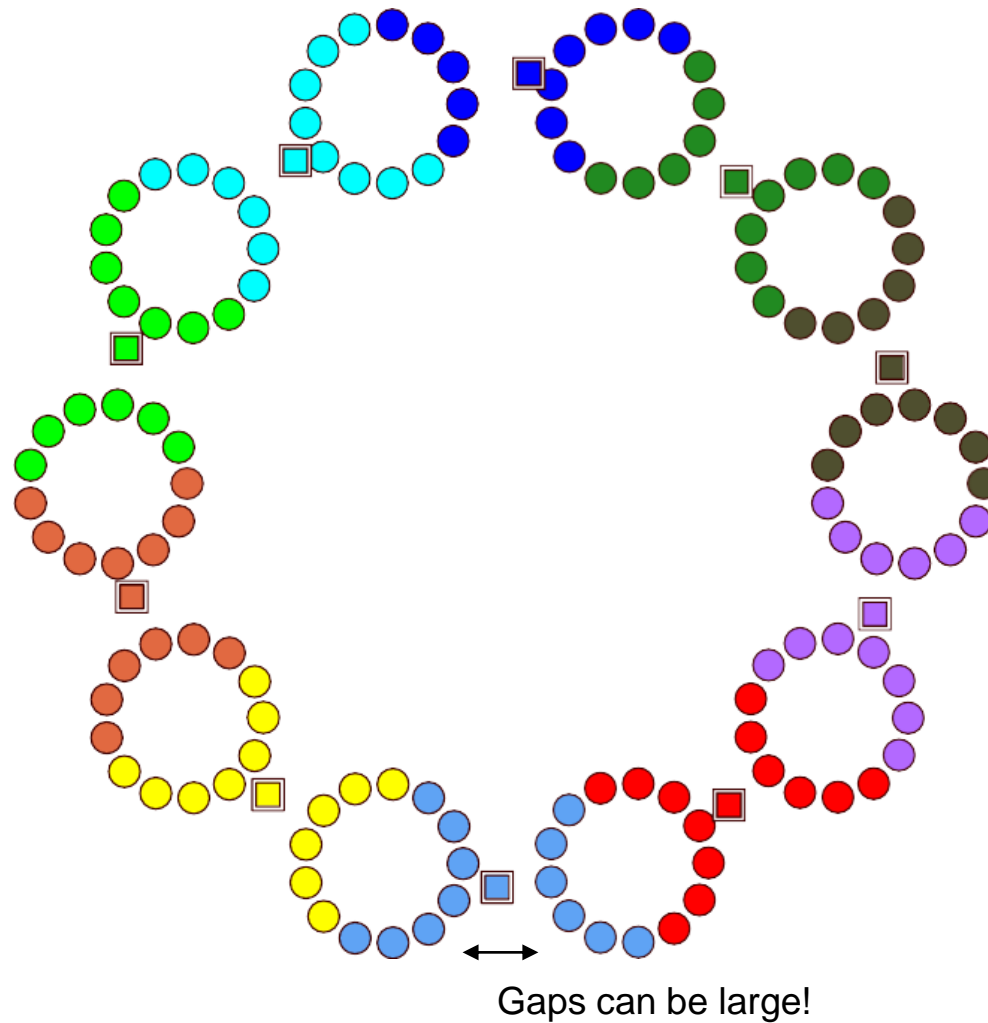


...but this example is
a “nice weather” case

The choice of the initial clustering is crucial.
If we start with:



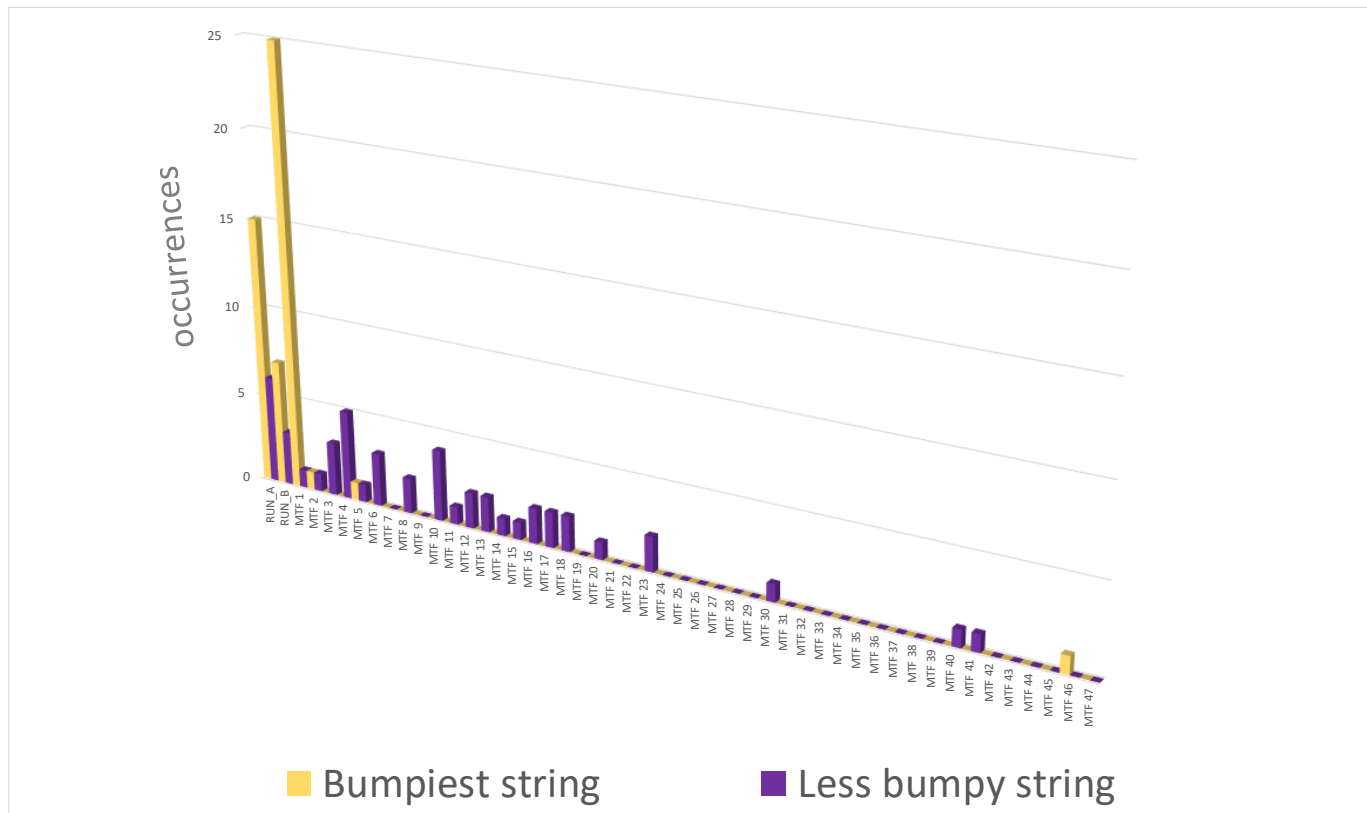
... we are stuck after two steps with this:



→ **Initial clustering** is key.

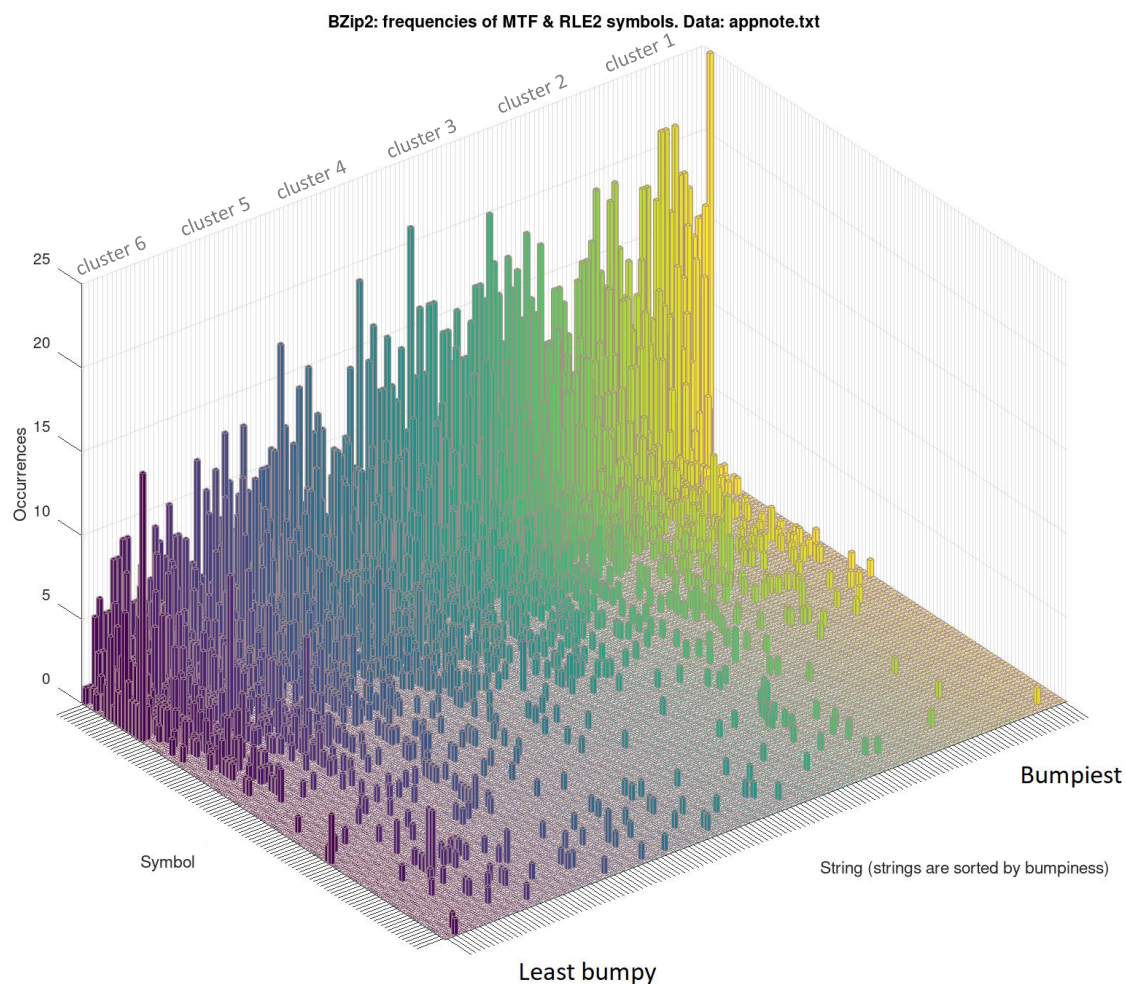
Especially, for BZip2, we
have more than 2
dimensions: up to 258 ...

Initial clustering for BZip2: sorting MTF & RLE2 symbols frequency histograms by “bumpiness”

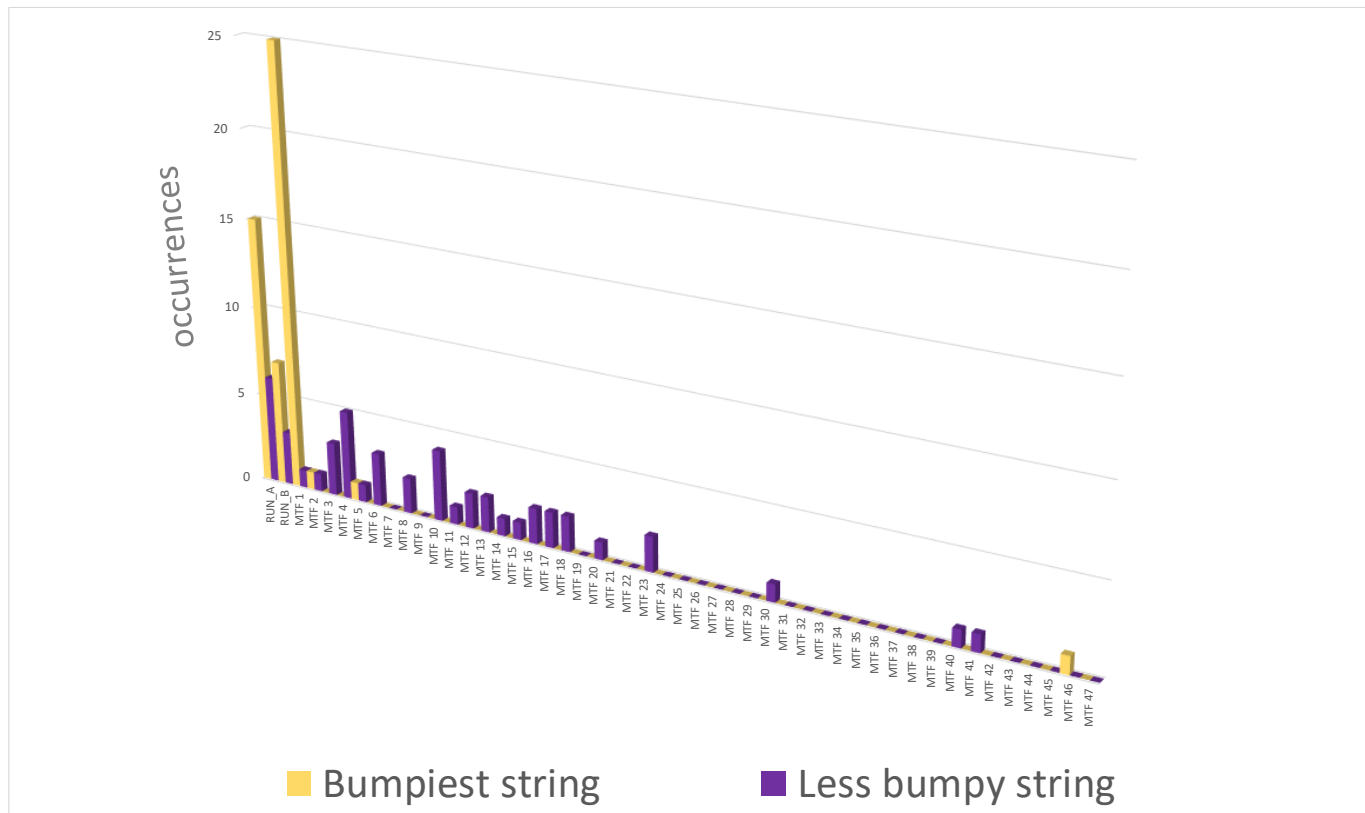


Strings with various “bumpiness” (bumpier: more **redundant MTF** data; input data: **appnote.txt**)

Initial clustering for BZip2: sorting MTF & RLE2 symbols frequency histograms by “bumpiness”



Initial clustering for BZip2: “bumpiness” function



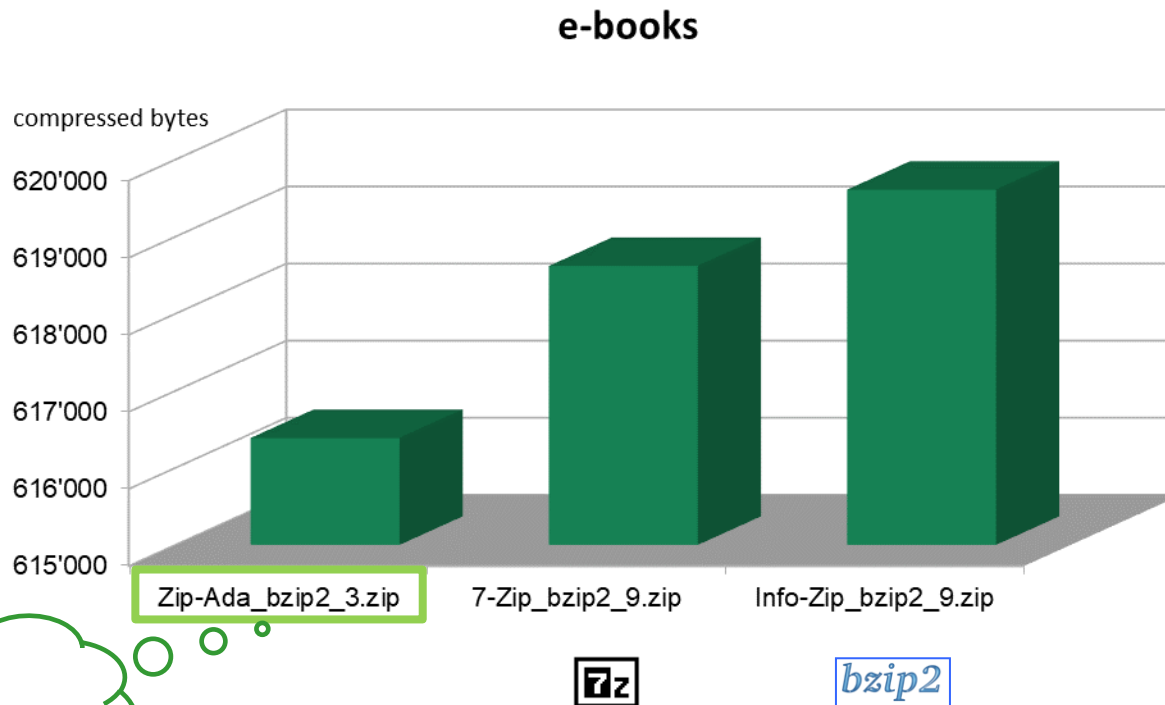
Quiz: how did I define the measure of bumpiness?

Hints:

- all strings (except at most one) have exactly 50 symbols.
- high histogram bars tend to be on the left side because of the Move-to-Front process

Results – first surprise

Zip archive, BZip2 only:



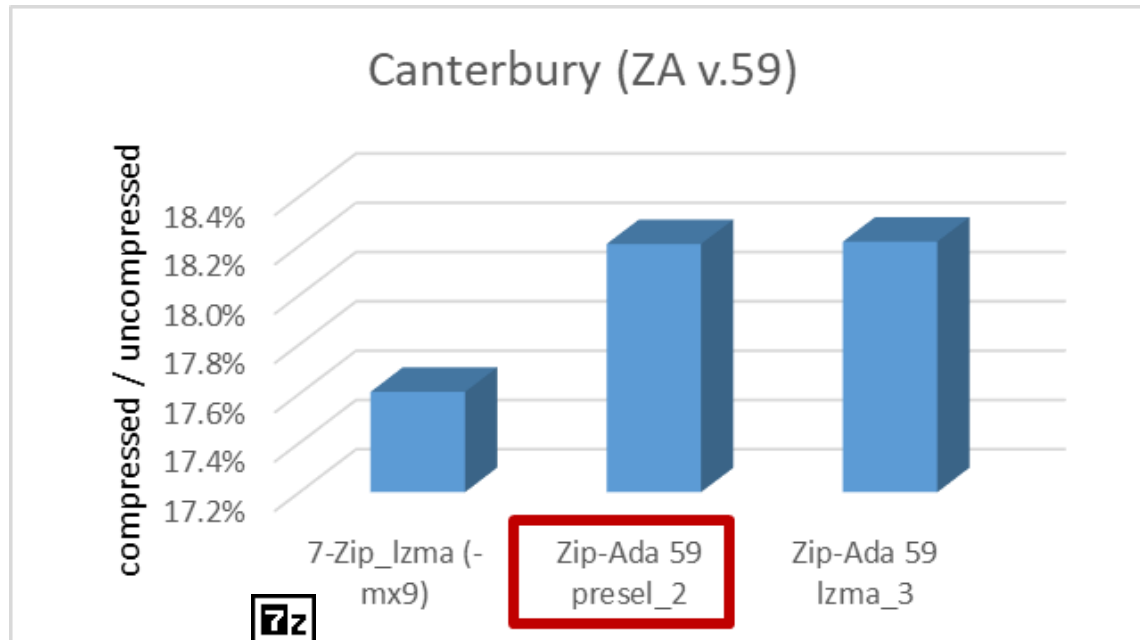
Zip-Ada
wins bigly !!!

NB: BZip2 is very good with (at least) human-written **texts** and **source code**.

Results – second surprise

Zip archive, multi-format (for Zip-Ada, Preselection_2)

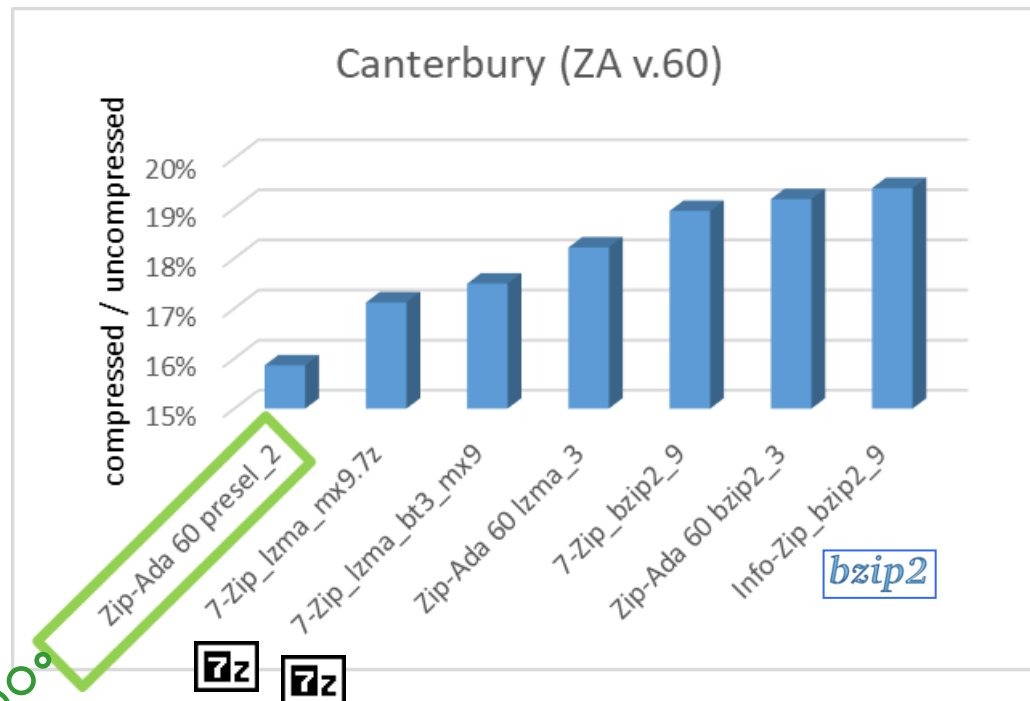
Before:



Results – second surprise

Zip archive, multi-format (for Zip-Ada, Preselection_2), or .7z archive.

After:



Zip-Ada wins
here too !!!

Benefits of Ada

Top Ada goodies:

- **Modularity**
- **Very precise typing**
- **Pointer-less programming.** No worries about **accidental deallocation**, **null pointers**, **dangling pointers**, **ownership** (the caller stays always the owner), **lifetimes**, etc.
In the BZip2 encoder, heap allocation is only used for five large arrays; pointers are used only for allocation & deallocation.

Benefits of Ada – continued

Data compression is very difficult to debug, sometimes impossible.

Ada does its best to help you doing things right the first time and avoiding stupid bugs (you only keep the clever bugs 😊). Specifically, it helps you avoiding **omissions**, **duplicates**, **confusions** between types, name **collisions**.

→ Use Ada for data compression (and the rest as well)!

Indirect benefit : you can focus on the algorithms!

Benefits of Ada – continued

Here, some **ranges** picked up from the code (bzip2-encoding.adb):

```
subtype Bit_Pos_Type is Natural range 0 .. 7;
type Buffer is array (Natural_32 range <>) of Byte;
subtype Offset_Range is Integer_32 range 0 .. block_size - 1;
subtype Max_Alphabet is Integer range 0 .. max_alphabet_size - 1;
type MTF_Array is array (Positive_32 range <>) of Max_Alphabet;

type Entropy_Coder_Range is range 1 .. max_entropy_coders;
descr : array (Entropy_Coder_Range) of
    Huffman.Encoding.Descriptor (Max_Alphabet);
entropy_coder_count : Entropy_Coder_Range
    range 2 .. Entropy_Coder_Range'Last;

subtype Alphabet_in_Use is Integer range 0 .. last_symbol_in_use;
type Huffman_Length_Array is array (Alphabet_in_Use) of Natural;
type Count_Array is array (Alphabet_in_Use) of Natural_32;

subtype Selector_Range is Positive_32 range 1 .. selector_count;
type Cluster_Attribution is array (Positive range <>) of Entropy_Coder_Range;
type Value_Array is array (Positive range <>) of Natural;

in_use_16 : array (Byte range 0 .. 15) of Boolean := (others => False);
```

Data
dependent!

