## Discrete Structures :: CS 207 :: Autumn 2019

## Practice Problem Set 13 Released: Sunday November 3, 2019

- 1. Show that the following subsets of  $\mathbb{R}$  are countably infinite:
  - (a)  $\{2^q : q \in \mathbb{Q}^+\}$
  - (b)  $\{x \in \mathbb{R} : x > 0 \land x^2 \in \mathbb{Q}\}$
- 2. Show that the following sets have the same cardinality as  $\mathbb{R}$ .
  - (a)  $(0,\infty)$
  - (b)  $[0,1] \cup \mathbb{Z}$
  - (c)  $\mathbb{R} \times \mathbb{R}$  (Hint: Flesh out the proof from the lecture.)
- 3. Suppose for each  $i \in \mathbb{N}$ ,  $S_i$  is a countable set. Using the result that  $\mathbb{N} \times \mathbb{N}$  is countable, prove that the following sets are countable (but not necessarily infinite). Be sure to handle any special cases (e.g., some  $S_i = \emptyset$ ). You may use the fact that S is countable iff  $|S| \leq |\mathbb{N}|$ .
  - (a)  $\bigcup_{i \in \mathbb{N}} S_i$
  - (b) For any  $k \in \mathbb{N}$ ,  $S_0 \times \cdots \times S_k$
- 4. Let A, B and C be sets such that |A| < |B| and  $B \subseteq C$ . Prove that |A| < |C|.
- 5.  $\mathbb{Z}[X]$  denotes the set of all polynomials in a variable X, with integer coefficients. Show that  $\mathbb{Z}[X]$  is countable.

Hint: Prove first that for every integer  $n \ge 1$  the set  $P_n$  of all polynomials of degree  $\le n$  with integer coefficients is countable. Then use the fact that the union of countably many countable sets is countable.

- 6. Consider a machine which has n bits of memory (initialized to all 0s). Its state is determined by the contents of its memory. This machine takes inputs from the set  $\Sigma = \{1, \ldots, n\}$ : On input i, it toggles its i<sup>th</sup> bit. (Toggling a bit changes its value from 0 to 1 or from 1 to 0.)
  - (a) Write down the transition function for this machine for n = 2, as table, with columns "current state," "input" and "next state." (Note that each state is labeled with 2 bits.)
  - (b) What graph does the state-diagram for this machine (for a general value of n) resemble? You can treat a pair of directed edges between two states, but pointing in opposite directions, as a single undirected edge. Justify your answer.

Hint: For this part, it may be helpful to draw a diagram for the case n=2 from the previous part.

- 7. Give deterministic finite-state acceptors for the following languages over the alphabet  $\{0,1\}$ . In each case, write down its transition function (in the form of a table) and also draw the state diagram.
  - (a)  $\{x \mid x \text{ is non-empty and begins and ends with the same symbol}\}$
  - (b)  $\{x \mid x \text{ does not contain the substring } 11\}$
  - (c)  $\{x \mid x \text{ number of zeroes is even and number of ones is odd}\}$
- 8. Give non-deterministic finite-state acceptors for the following languages over the alphabet {0,1}. In each case, write down its transition function (in the form of a table) and also draw the state diagram.
  - (a)  $\{x \mid x \text{ ends in } 00 \text{ or } 010 \}$
  - (b)  $\{x \mid x \text{ contains the substring } 11\}$
- 9. It was mentioned in class that the language consisting of *all* palindromes does not have a finite-state acceptor. In this problem we consider palindromes of a fixed length.
  - (a) Describe, as explicitly as you can, a deterministic finite-state acceptor for the language consisting of all palindromes of length exactly 100.

Hint: The first 50 input bits need to be memorized. Arrange the states in a binary tree for this.

<sup>&</sup>lt;sup>1</sup>In class, we defined a set S as countable if S is finite or  $|S| = |\mathbb{N}|$ . By the "axiom of countable choice" this is equivalent to  $|S| \leq |\mathbb{N}|$ .

- (b) Using the pigeonhole principle, argue that any such acceptor should have at least  $2^{50}$  states. Can you tighten this argument to show that your construction from above has the minimum possible number of states?
  - Hint: Consider the set of prefixes of strings in the language. Can two distinct prefixes take the machine to the same state?
- 10. Show that each of the following decision problems is in **NP**, by describing a "certificate" (for "yes" instances) that can be verified in time that is polynomial in the length of the instance.
  - (a) Is x the binary representation of a composite number?
  - (b) Is x the binary representation of an even number?
  - (c) Are the pair of graphs  $(G_1, G_2)$  (represented using their adjacency matrices) isomorphic to each other?
  - (d) Does the graph G have a Hamiltonian cycle?
  - (e) Does a polynomial p(X) with integer coefficients have an integer root? Hint: You may use the fact that all the real roots of a polynomial  $c_dX^d + \ldots + c_1X + c_0$  lie in the range [-s, s], where  $s = \max\{1, |c_0| + \ldots + |c_n|\}$ . Why do you need to use this fact?

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