

A “crash course” in nonsmooth convex optimization

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1 Introduction

These notes are supplementary material to a “crash course” I am teaching in May of 2023. The topic is *proximity operators and nonsmooth convex optimization*. These notes are not meant to be used as a standalone resource. Please cite peer-reviewed material. As a general reference text, I suggest *Convex Analysis and Monotone Operator Theory*, 2nd ed., by Bauschke and Combettes, published by Springer. Virtually all of the results in these notes also apply to real Hilbert spaces; for proofs in full-generality, read the book. If unspecified, \mathcal{H} is a real finite-dimensional vector space (e.g., \mathbb{R}^n is fine).

1.1 Optimization terminology and the extended real line

Notation 1.1 We will work with the **extended real line**, i.e., $[-\infty, +\infty] := \mathbb{R} \cup \{-\infty, +\infty\}$. Algebra on this field follows most “natural” rules one could expect (e.g., for $x \in \mathbb{R}$, $x + \infty = \infty$). However, the following quantities are **undefined**:

- Any subtraction of infinities: “ $+\infty - (+\infty)$ ”
- Zero times infinity: “ $0 \cdot (\pm\infty)$ ”
- Any quotient of infinities: “ $\pm\infty / \pm\infty$, $\pm\infty / \mp\infty$, ...”

As a result, if we work with extended-real-valued functions, we must be sure to avoid anything which is undefined (e.g., the objective function $f(x) + g(x)$ could be undefined if there exists z such that $g(z) = -\infty$ and $f(z) = \infty$.)

*Please report typos/errors found in these notes. Homework solutions should be handed in to my office ZIB 3107.

Definition 1.2 Given a real vector space \mathcal{H} , a function $f: \mathcal{H} \rightarrow [-\infty, +\infty]$, and a set $C \subset \mathcal{H}$, consider the following optimization problem.

$$\underset{x \in C}{\text{minimize}} \quad f(x) \quad (1)$$

We call f the **objective function**. We call C a **constraint**. For any $x \in C$, we say x is **feasible**. Otherwise, for $x \in \mathbb{R}^n \setminus C$, x is infeasible. If a point $x^* \in C$ satisfies

$$(\forall x \in C) \quad f(x^*) \leq f(x), \quad (2)$$

we call x^* a **solution** to the optimization problem (1).

For this class, we consider minimization; to maximize f , just use the objective function $-f$.

Definition 1.3 For $I \subset [-\infty, +\infty]$, $a \in [-\infty, +\infty]$ is a **lower bound (upper bound)** if, for every $\xi \in I$, $a \leq \xi$ ($a \geq \xi$). The **greatest lower bound**, or **infimum**, of the set I is denoted $\inf I$. Analogously, the **least upper bound**, or **supremum**, of the set I is denoted $\sup I$. In general, $\inf I, \sup I \in [-\infty, +\infty]$. If, additionally, $\inf I \in I$ ($\sup I \in I$), we call it the **minimum (maximum)**, and denote it $\min I$ ($\max I$). In these cases, we say the infimum (supremum) is *attained*.

A few things to mention:

- (i) For $I \neq \emptyset$, we have $\inf I \leq \sup I$. For the empty set, $\inf \emptyset = +\infty$ and $\sup \emptyset = -\infty$.
- (ii) While the \inf and \sup are always defined, \max and \min may not exist (e.g., consider $I = (0, 1)$ has $\inf I = 0$ and $\sup I = 1$. However, since $0, 1 \notin I$, neither $\max I$ nor $\min I$ exist.)
- (iii) Let $f: \mathbb{R}^n \rightarrow [-\infty, +\infty]$. We adopt the notation that $\inf_{x \in C} f(x) = \inf\{f(x) \mid x \in C\}$.
- (iv) It is common in optimization literature to abuse notation, and use

$$\min_{x \in C} f(x) \quad (3)$$

to describe the optimization problem (1). Technically, $\min_{x \in C} f(x)$ is not an optimization problem – it is the optimal value of the objective function at a solution, which may or may not exist.¹

Definition 1.4 Let $f: \mathcal{H} \rightarrow [-\infty, +\infty]$. We will use the following terms.

- (i) The **domain** of f is

$$\text{dom } f = \{x \in \mathcal{H} \mid f(x) < +\infty\} \quad (4)$$

¹The *Weierstraß Theorem*, loosely stated, guarantees that a solution to (1) exists if C is compact and f is lower-semicontinuous. For unconstrained functions, analytic notions of “coercivity” and “recession cones” can also yield existence results; however, they are not included in this class.

(ii) The **epigraph** of f is

$$\text{epi } f = \{(x, \xi) \in \mathcal{H} \times \mathbb{R} \mid f(x) \leq \xi\} \quad (5)$$

(iii) The function f is **proper** if $\text{dom } f \neq \emptyset$ and it never outputs the value $-\infty$ (i.e., $-\infty \notin f(\mathcal{H})$).

(iv) The function f is **lower semicontinuous** (sometimes abbreviated “lsc”) at $x \in \mathcal{H}$ if, for every sequence $(x_n)_{n \in \mathbb{N}}$ satisfying $x_n \rightarrow x$, we have $f(x) \leq \liminf f(x_n)$

For this class, we will predominantly consider proper and lsc functions. A few things to note about the lsc assumption: (1) every continuous function is lsc, and (2) lower semicontinuity basically allows for a jump-discontinuity to occur at $x \in \mathcal{H}$, but requires that f takes the lowest possible limiting value at x (cf. the figures drawn in class, or [here](#)).

1.2 Inner product and norms

Definition 1.5 Let \mathcal{H} be a real finite-dimensional vector space. A **scalar product** (sometimes called **inner product**) is a function $\langle \cdot \mid \cdot \rangle : \mathcal{H} \times \mathcal{H} \rightarrow \mathbb{R}$ which satisfies the following properties.

- (i) $(\forall x \in \mathcal{H} \setminus \{0\}) \quad \langle x \mid x \rangle > 0$
- (ii) $(\forall x, y \in \mathcal{H}) \quad \langle x \mid y \rangle = \langle y \mid x \rangle$
- (iii) $(\forall x, y, z \in \mathcal{H})(\forall \alpha \in \mathbb{R}) \quad \langle \alpha x + y \mid z \rangle = \alpha \langle x \mid z \rangle + \langle y \mid z \rangle$

Exercise 1.6 Let $0 \in \mathcal{H}$ be the zero element of \mathcal{H} . Show that, for every $x \in \mathcal{H}$, $\langle 0 \mid x \rangle = 0$.

Exercise 1.7 Consider $\mathcal{H} = \mathbb{R}^n$. For two vectors $x, y \in \mathbb{R}^n$, the *dot product* is given by $\langle x \mid y \rangle = x^\top y$. Show that the dot product on \mathbb{R}^n is a scalar product.

Exercise 1.8 Consider the vector space of matrices $\mathbb{R}^{n \times n}$. For two matrices $A = (a_{i,j})_{1 \leq i,j \leq n}$ and $B = (b_{i,j})_{1 \leq i,j \leq n}$, the *Frobenius inner product* is given by

$$\langle A \mid B \rangle = \sum_{i=1}^n \sum_{j=1}^n a_{i,j} b_{i,j} \quad (6)$$

Show (6) is an inner product.

Proposition 1.9 (Cauchy-Schwarz) For every $x, y \in \mathcal{H}$,

$$\langle x \mid y \rangle^2 \leq \langle x \mid x \rangle \langle y \mid y \rangle. \quad (7)$$

Proof. If $y = 0$, (7) holds. Now suppose that $y \neq 0$. By Definition 1.5, $\langle y | y \rangle > 0$. Set $\alpha = \langle x | y \rangle / \langle y | y \rangle$. First, we find

$$0 \leq \langle x - \alpha y | x - \alpha y \rangle \quad (8)$$

$$= \langle x | x \rangle - 2\alpha \langle x | y \rangle + \alpha^2 \langle y | y \rangle \quad (9)$$

$$= \langle x | x \rangle - 2\alpha \langle x | y \rangle + \alpha \langle x | y \rangle \quad (10)$$

$$= \langle x | x \rangle - \alpha \langle x | y \rangle. \quad (11)$$

Rearranging the inequality, we find that

$$\frac{\langle x | y \rangle^2}{\langle y | y \rangle} = \alpha \langle x | y \rangle \leq \langle x | x \rangle \quad (12)$$

$$\Leftrightarrow \langle x | y \rangle^2 \leq \langle y | y \rangle \langle x | x \rangle. \quad (13)$$

□

Definition 1.10 Let \mathcal{H} be a real finite-dimensional vector space. A function $\|\cdot\|: \mathcal{H} \rightarrow \mathbb{R}$ is a **norm** if the following hold.

- (i) $(\forall x \in \mathcal{H}) \quad \|x\| = 0 \Rightarrow x = 0$
- (ii) $(\forall x, y \in \mathcal{H}) \quad \|x + y\| \leq \|x\| + \|y\|$
- (iii) $(\forall x \in \mathcal{H})(\forall \alpha \in \mathbb{R}) \quad \|\alpha x\| = |\alpha| \|x\|$

A norm is a way to measure magnitude of vectors, or the distance from one vector to another $\|x - y\|$.

Exercise 1.11 Let \mathcal{H} be a real finite-dimensional vector space, and let $\langle \cdot | \cdot \rangle$ be a scalar product on \mathcal{H} . Show that the norm defined by

$$\|\cdot\|: \mathcal{H} \rightarrow \mathbb{R}: x \mapsto \sqrt{\langle x | x \rangle} \quad (14)$$

satisfies the properties in Definition 1.10.

The **Euclidean norm** on \mathbb{R}^n , given by $(\xi_1, \dots, \xi_n) \mapsto \sqrt{\xi_1^2 + \dots + \xi_n^2}$, arises from the dot product. Exercise 1.11 yields the following formulation of the Cauchy-Schwarz inequality

$$(\forall x, y \in \mathcal{H}) \quad \langle x | y \rangle \leq \|x\| \|y\|. \quad (\text{C-S})$$

Exercise 1.12 Let $(x_1, x_2, x_3) \in \mathbb{R}^3$. Show that

$$2x_1 - x_2^4 + 6x_3 \leq 4\sqrt{x_1^2 + x_2^8 + 9x_3^2}. \quad (15)$$

Can the coefficient 4 in (15) be reduced?

The following theorem is referenced a few times in the notes, so I will provide its statement here. Regretfully, this class does not have enough time to detail the topics of compact/closed/lsc. The following theorem is often used as a tool to ensure that a solution to an optimization problem exists.

Theorem 1.13 (Weierstraß) *Let $f: \mathcal{H} \rightarrow [-\infty, +\infty]$ be lower semicontinuous and let C be a compact subset of \mathcal{H} . Suppose that $C \cap \text{dom } f \neq \emptyset$. Then f achieves its infimum over C .*

2 Convexity

Definition 2.1 A set $C \subset \mathcal{H}$ is **convex** if, for every $x, y \in C$

$$(\forall \alpha \in]0, 1[) \quad \alpha x + (1 - \alpha)y \in C. \quad (16)$$

A function f is **convex** if $\text{epi } f$ is convex.

Proposition 2.2 $f: \mathcal{H} \rightarrow [-\infty, +\infty]$ is convex if and only if

$$(\forall x, y \in \text{dom } f) \quad (\forall \alpha \in]0, 1[) \quad f(\alpha x + (1 - \alpha)y) \leq \alpha f(x) + (1 - \alpha)f(y). \quad (17)$$

Proof. First, we note that if f is identically $+\infty$, then $\text{dom } f = \emptyset$ if and only if $\text{epi } f = \emptyset$, so (17) is vacuously true. Now assume that $\text{dom } f \neq \emptyset$. Let (x, ξ) and (y, η) be in $\text{epi } f$ and let $\alpha \in]0, 1[$.
(\Rightarrow) Assume that $\text{epi } f$ is convex. Then

$$\alpha(x, \xi) + (1 - \alpha)(y, \eta) = (\alpha x + (1 - \alpha)y, \alpha \xi + (1 - \alpha)\eta) \in \text{epi } f. \quad (18)$$

Therefore, $f(\alpha x + (1 - \alpha)y) \leq \alpha \xi + (1 - \alpha)\eta$. Taking the limit as $\xi \searrow f(x)$ and $\eta \searrow f(y)$ yields (17).
(\Leftarrow) Assume that (17) holds. By definition, $f(x) \leq \xi$ and $f(y) \leq \eta$. So, by (17),

$$f(\alpha x + (1 - \alpha)y) \leq \alpha f(x) + (1 - \alpha)f(y) \quad (19)$$

$$\leq \alpha \xi + (1 - \alpha)\eta. \quad (20)$$

Therefore, $(\alpha x + (1 - \alpha)y, \alpha \xi + (1 - \alpha)\eta) \in \text{epi } f$ which completes the proof. \square

Definition 2.3 Let $\rho > 0$ and let $x \in \mathcal{H}$. A **closed ball** of radius ρ is $B(x; \rho) = \{z \in \mathcal{H} \mid \|x - z\| \leq \rho\}$.

Definition 2.4 Let $f: \mathcal{H} \rightarrow [-\infty, +\infty]$ and let $x \in \mathcal{H}$. x is a **local minimizer** of f if there exists $\rho > 0$ such that

$$(\forall z \in \mathcal{H} \cap B(x; \rho)) \quad f(x) \leq f(z). \quad (21)$$

x is a **global minimizer** of f if

$$(\forall z \in \mathcal{H}) \quad f(x) \leq f(z). \quad (22)$$

Fact 2.5 Let f be a convex and proper function. Then every local minimizer is a global minimizer.

Proof. This is left as an exercise (easier to prove after we learn about convex subdifferentials). \square

Definition 2.6 Let $C \subset \mathcal{H}$ be nonempty.

(i) The **indicator function** of C is

$$\iota_C: \mathcal{H} \rightarrow [-\infty, +\infty] : x \mapsto \begin{cases} 0 & \text{if } x \in C \\ +\infty & \text{if } x \notin C. \end{cases} \quad (23)$$

(ii) Suppose that C is also closed. A **projection** of $x \in \mathcal{H}$ onto C is a solution to the minimization problem

$$\underset{z \in C}{\text{minimize}} \quad \|x - z\|. \quad (24)$$

A solution to (24) is a “closest” point to x which resides in C .

Fact 2.7 Let $C \subset \mathcal{H}$ and let $x \in \mathcal{H}$.

(i) Without loss of generality, constrained optimization can be rephrased as unconstrained optimization via changing the objective function:

$$\inf_{x \in C} f(x) = \inf_{x \in \mathcal{H}} f(x) + \iota_C(x). \quad (25)$$

The objective function $f + \iota_C$ on the righthand side, although a bit fancier, allows us to rephrase the constraint on the lefthand side.

(ii) C is convex if and only if its indicator function ι_C is convex.

(iii) C is closed if and only if its indicator function ι_C is lsc.

(iv) Suppose that C is closed. Then a solution to (24) exists.

(v) Suppose that C is convex. If a solution to (24) exists, it is guaranteed to be unique.

The proofs of (ii) and (iii) follow from the fact that $\text{epi } C = C \times [0, +\infty[$. Loosely speaking, the proof of (iv) follows from the Weierstraß theorem (compactness is achieved by intersecting C with $\{y \in \mathcal{H} \mid \|x - y\| \leq \eta\}$ for $\eta > 0$) and (v) follows from the fact that the norm is *strictly convex* – (a notion we have not yet defined, but the interested student could research).

Definition 2.8 Let $C \subset \mathcal{H}$ be nonempty, closed, and convex. In view of Fact 2.7(iv)–(v), for every $x \in \mathcal{H}$ there is a unique point, $\text{Proj}_C(x) \in \mathcal{H}$, which solves (24). This implicitly defines the **projection operator** of C .

$$\text{Proj}_C: \mathcal{H} \rightarrow \mathcal{H}: x \mapsto \text{Proj}_C(x) \quad (\text{solution to (24)}) \quad (26)$$

Note: if $x \in C$, then $\text{Proj}_C x = x$.

For all of the algorithms in this course, we will focus on functions from the following class

$$\Gamma_0(\mathcal{H}) = \{f: \mathcal{H} \rightarrow]-\infty, +\infty] \mid f \text{ is proper, lower semicontinuous, and convex}\}. \quad (27)$$

The following functions live in $\Gamma_0(\mathcal{H})$:

- (i) Exponentials: e^x
- (ii) Log-barriers $f(x) = \begin{cases} -\ln(x) & \text{if } x > 0 \\ +\infty & \text{otherwise.} \end{cases}$
- (iii) Any norm: $\|\cdot\|$ (e.g., $\|\cdot\|_1$ which promotes sparsity, $\|\cdot\|_{\text{nuclear}}$ which promotes low-rank)
- (iv) Hinge-Loss, ReLU, KL-Divergence, ...
- (v) Given a collection of functions $(f_i)_{i \in I}$ in $\Gamma_0(\mathcal{H})$, we can remain in $\Gamma_0(\mathcal{H})$ via the following operations.
 - (a) $\max\{f_1, \dots, f_m\}$
 - (b) Positive linear combinations: $\lambda_1 f_1 + \dots + \lambda_m f_m$, where $\{\lambda_i\}_{i=1}^m$ are positive.
 - (c) Let \mathcal{H}_1 and \mathcal{H}_2 be two finite-dimensional real vector spaces. Let $b \in \mathcal{H}_2$ and let $A: \mathcal{H}_1 \rightarrow \mathcal{H}_2$ be a linear operator (e.g., a matrix from \mathbb{R}^n to \mathbb{R}^m). If $f_1 \in \Gamma_0(\mathcal{H}_2)$, then $g(x) = f_1(Ax + b) \in \Gamma_0(\mathcal{H}_1)$.

Exercise 2.9 The **Minkowski sum** of two subsets A, B of \mathcal{H} is given by

$$A + B = \{a + b \mid a \in A \text{ and } b \in B\}. \quad (28)$$

Assume that A and B are convex. Prove that $A + B$ is convex.

Exercise 2.10 Show that the norm $\|\cdot\|$ is convex using Definition 1.10.

3 What is Differentiability?

There are a lot of ML engineers who brush off the mathematical details of what it means for a function to be differentiable. Algorithmic differentiation (sometimes misleadingly-called “automatic” differentiation) is only guaranteed to work when certain theoretical conditions about the *existence* of a gradient hold. This part of the class is dedicated to explaining that differentiability is not a freebie.

To start our discussion on differentiability, we will begin with a few preliminaries from analysis.

Definition 3.1 Let $A: \mathcal{H}_1 \rightarrow \mathcal{H}_2$. Then A is **linear** if, for every $\alpha \in \mathbb{R}$ and every $x, y \in \mathcal{H}_1$,

$$A(\lambda x) = \lambda A(x) \quad \text{and} \quad A(x + y) = A(x) + A(y). \quad (29)$$

Theorem 3.2 (Riesz-Fréchet representation) Let $A: \mathcal{H} \rightarrow \mathbb{R}$ be linear. Then there exists a unique vector $u \in \mathcal{H}$ such that, for every $x \in \mathcal{H}$, $A(x) = \langle u | x \rangle$.

Although at first-glance it looks unrelated, Theorem 3.2 is a central notion for defining the gradient. A necessary (albeit insufficient) condition for the existence of a gradient is the existence of a directional derivative, defined below.

Definition 3.3 Let $f: \mathcal{H} \rightarrow]-\infty, +\infty]$ be proper. The **directional derivative** of f at $x \in \text{dom } f$ in the direction $y \in \mathcal{H}$ is

$$f'(x; y) = \lim_{\alpha \searrow 0} \frac{f(x + \alpha y) - f(x)}{\alpha}. \quad (30)$$

From Definition 3.3, we point out a few things.

- (i) The limit in (30) might not exist.
- (ii) If f is convex, then $f'(x; y) \in [-\infty, +\infty]$.
- (iii) Even if a directional derivative exists, it might not exist in \mathbb{R} (since it could be $+\infty$ or $-\infty$).

Definition 3.4 Let $x \in \text{dom } f$. If $f'(x; \cdot)$ is linear, we say f is **differentiable at x** . In this case, the unique vector provided by Theorem 3.2 is called the **gradient** of f at x and denoted $\nabla f(x)$.

$$f'(x; \cdot) = \lim_{\alpha \searrow 0} \frac{f(x + \alpha \cdot) - f(x)}{\alpha} = \langle \nabla f(x) | \cdot \rangle \quad (31)$$

If f is differentiable at every $x \in \text{dom } f$, we say that f is **differentiable**.

Exercise 3.5 Verify that $\nabla(\frac{1}{2}\|\cdot\|^2)(x) = x$.

All of the properties we know and love about differentiability (chain rule, product rule, etc.) have to be proven. Here is an example below.

Proposition 3.6 Let $A: \mathcal{H}_1 \rightarrow \mathcal{H}_2$ be a linear operator (with adjoint denoted A^*), let $b \in \mathcal{H}_2$, and let $f: \mathcal{H} \rightarrow \mathbb{R}$ be proper and differentiable. Set $g = f(Ax + b)$. Then g is differentiable and

$$\nabla g = A^*(\nabla f(A \cdot + b)). \quad (32)$$

Proof. Since $\text{dom } f = \mathcal{H}_2$, $\text{dom } g \neq \emptyset$ so we let $x \in \text{dom } g$. By definition,

$$g'(x; y) = \lim_{\alpha \searrow 0} \frac{g(x + \alpha y) - g(x)}{\alpha} \quad (33)$$

$$= \lim_{\alpha \searrow 0} \frac{f(A(x + \alpha y) + b) - f(Ax + b)}{\alpha} \quad (34)$$

$$= \lim_{\alpha \searrow 0} \frac{f(Ax + b + \alpha Ay) - f(Ax + b)}{\alpha} \quad (35)$$

$$= f'(Ax + b; Ay). \quad (36)$$

So the directional derivative of g exists. Now, since f is differentiable,

$$g'(x; y) = f'(Ax + b; Ay) = \langle \nabla f(Ax + b) \mid Ay \rangle = \langle A^*(\nabla f(Ax + b)) \mid y \rangle. \quad (37)$$

Hence the directional derivative of g is linear and g is differentiable. The specific form of the gradient is constructed in (37) \square

Algorithmic differentiation tools use results like Proposition 3.6 to approximate a gradient of a function by reading its machine code. However, these subroutines do not check the theoretical conditions required for their theorems (e.g., *f must be differentiable*) – this must be done (and is oftentimes unjustly ignored) by the user.

Definition 3.7 Let f be proper and differentiable. f is **smooth** (“ L -smooth”) if there exists $L > 0$ such that

$$(\forall x, y \in \mathcal{H}) \quad \|\nabla f(x) - \nabla f(y)\| \leq L\|x - y\|. \quad (38)$$

Exercise 3.8 Construct a function which is differentiable and nonsmooth.

Proposition 3.9 Let $f: \mathcal{H} \rightarrow]-\infty, +\infty]$ be proper and convex. Then,

$$(\forall x \in \text{dom } f)(\forall y \in \mathcal{H}) \quad f'(x; y - x) + f(x) \leq f(y). \quad (39)$$

Proof. By Proposition 2.2, for every $\alpha \in]0, 1[$,

$$f(x + \alpha(y - x)) - f(x) = f((1 - \alpha)x + \alpha y) - f(x) \quad (40)$$

$$\leq (1 - \alpha)f(x) + \alpha f(y) - f(x) \quad (41)$$

$$= \alpha(f(y) - f(x)). \quad (42)$$

Therefore,

$$\frac{f(x + \alpha(y - x)) - f(x)}{\alpha} \leq f(y) - f(x). \quad (43)$$

Taking the limit as $\alpha \searrow 0$ implies $f'(x; y) \leq f(y) - f(x)$, which in turn yields (39). \square

Corollary 3.10 Let $f: \mathcal{H} \rightarrow]-\infty, +\infty]$ be proper and convex. If f is differentiable at an interior point x of its domain², then

$$(\forall y \in \mathcal{H}) \quad \langle y - x \mid \nabla f(x) \rangle + f(x) \leq f(y). \quad (44)$$

When the lefthand side of (44) is viewed as a function of y , we see it is the first-order Taylor series approximation of f . Therefore, it follows from (39) that a convex differentiable function always remains above its first-order Taylor approximation! This is the motivating idea in defining a (convex) subgradient.

Definition 3.11 Let $f: \mathcal{H} \rightarrow]-\infty, +\infty]$. A vector g is a **subgradient** of f at $x \in \mathcal{H}$ if

$$(\forall y \in \mathcal{H}) \quad \langle y - x \mid g \rangle + f(x) \leq f(y). \quad (45)$$

The **subdifferential** of f at x is the set of all subgradients, denoted $\partial f(x)$.

This leads to the following fundamental theorem for optimization.

Theorem 3.12 (Fermat's Rule) Let $f: \mathcal{H} \rightarrow]-\infty, +\infty]$ be proper. Then x is a minimizer of f if and only if $0 \in \partial f(x)$.

Proof. By definition,

$$0 \in \partial f(x) \Leftrightarrow (\forall y \in \mathcal{H}) \quad \langle 0 \mid y - x \rangle + f(x) \leq f(y) \quad (46)$$

$$\Leftrightarrow (\forall y \in \mathcal{H}) \quad f(x) \leq f(y). \quad (47)$$

□

Unlike differentiable functions, there are technical conditions we must check in order to get the “standard” rules one would hope for.

Theorem 3.13 (Sum rule) Let $f, g \in \Gamma_0(\mathcal{H})$ and suppose that one of the following holds:

- (i) The interior of $\text{dom } g$ intersects with $\text{dom } f$
- (ii) $\text{dom } g = \mathcal{H}$
- (iii) The relative interiors of $\text{dom } f$ and $\text{dom } g$ intersect.

Then $\partial(f + g) = \partial f + \partial g$.

Note: If f is convex and differentiable at $x \in \mathcal{H}$, then $\partial f(x) = \{\nabla f(x)\}$.

²the interior of the domain