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A-05

COSC 211

que 3

(a) hex encoding :- 0xad100000

Explanation :-

sw \$s0 0(\$t0)
 \Rightarrow sw \$16 0(\$8) offset
 \downarrow \downarrow
 rt rs

I-Format

6 bits	5 bits	5 bits	16 bits
OP	rs	rt	immediate/offset

101011	01000	10000	0000000000000000
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1010 1101 0001 0000 0000 0000 0000 0000

0xad100000

hence, proved.

(2)

(b) hex is encoded as $0 \times 1212 0007$
~~hex~~ Binary representation:

00010010 0001 0010 0000 0000 0000 0111

16 bits.

hex:

000100	rs	rt	offset
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000100	10000	10010	000000000000000111
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The offset is actually determined by counting how far is the label from the relative positions of instructions. In our example, exit is 7 instructions far from hex. Therefore our offset is 7.

Binary representation: 000900000000000111

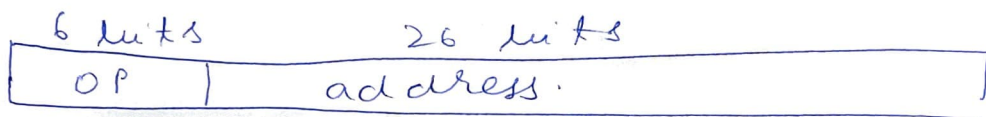
hex representation: 0×0007 .

How we use: The offset is sign extended, multiplied by 4, and added to value of PC when condition is met, PC is updated with $PC + 4 + \text{offset} * 4$ to point to label.

In our example, offset is 7. (3)
 Therefore $PC = PC + 4 + 7 * 4$.

(C) j encoded as 0×08100007 :

J-Format.



In the address we have ~~the~~ absolute address of the label divided by 4.

In our question we have: j loop

address of loop in 26 bits: ~~xxxxxxxxxxxxxxxxxxxxxxxxxxxxxxxx~~

00001000 0000 0000 0000 0001 1100

when divided by 4:

00 0001 0000 0000 0000 0000 0111



Binary :-

0000 1000 0001 0000 0000 0000 0000 0111

Hexadecimal :- 0x08100007

Hence, proved.